

# Probleme de conectivitate în grafuri neorientate



**Muchii critice**

# Muchii critice

- ▶  $G$  – graf neorientat
- ▶  $e \in E(G)$  **critică (punte, muchie de articulație)** = prin eliminarea ei crește numărul de componente conexe ale grafului

$$\text{nr componente } (G - e) > \text{nr. componente } (G)$$

- ▶ Un graf conex fără punți se numește **2-muchie conex**.

# Muchii critice

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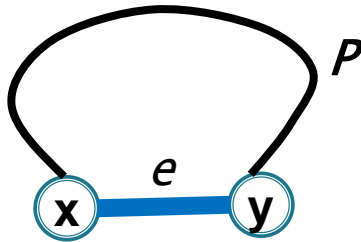
Demonstrație

# Muchii critice

- ▶ O muchie este critică  $\Leftrightarrow$  nu este conținută într-un ciclu

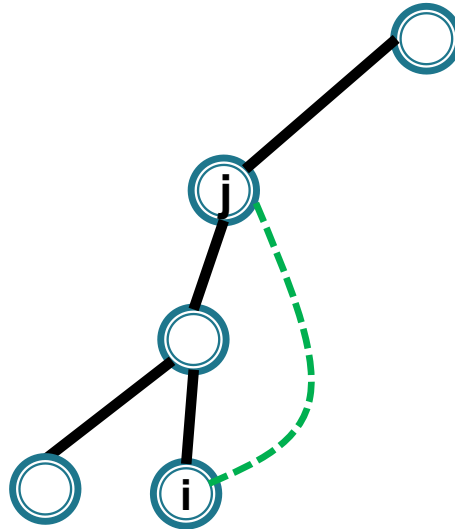
Demonstrație

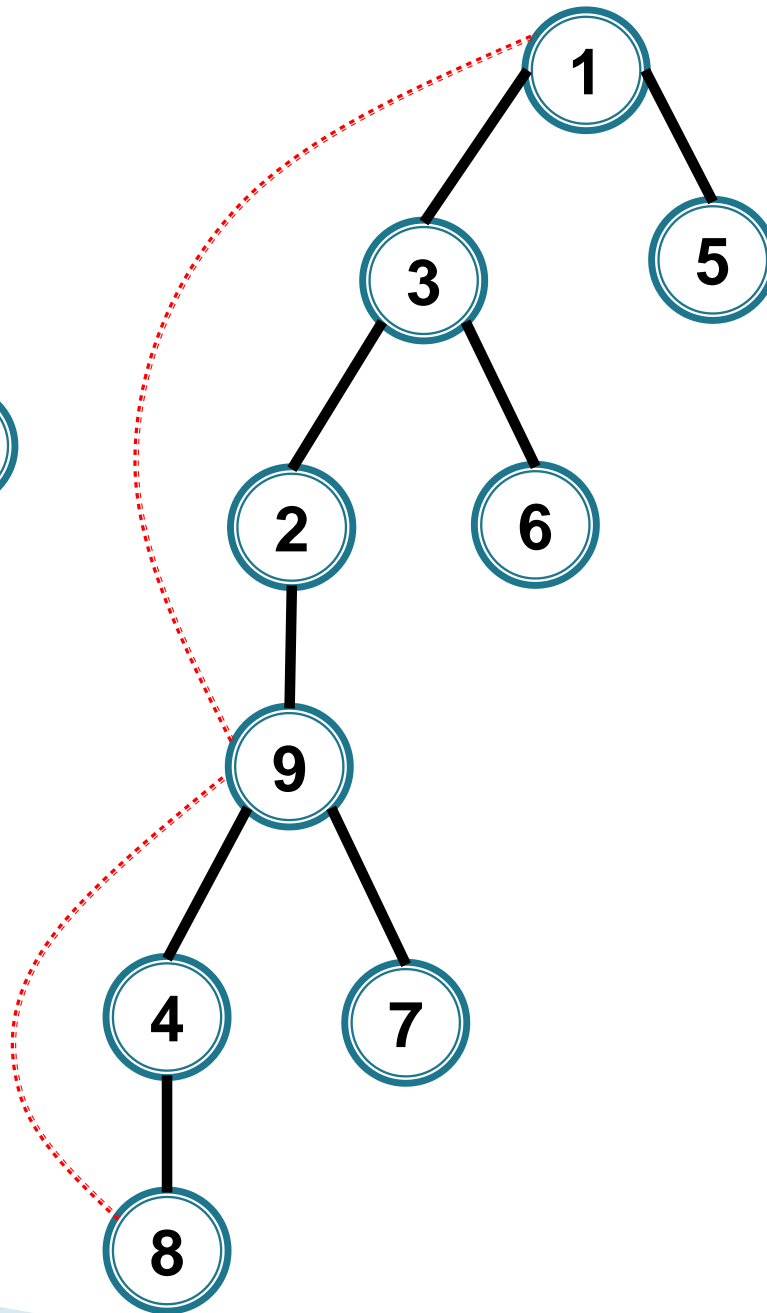
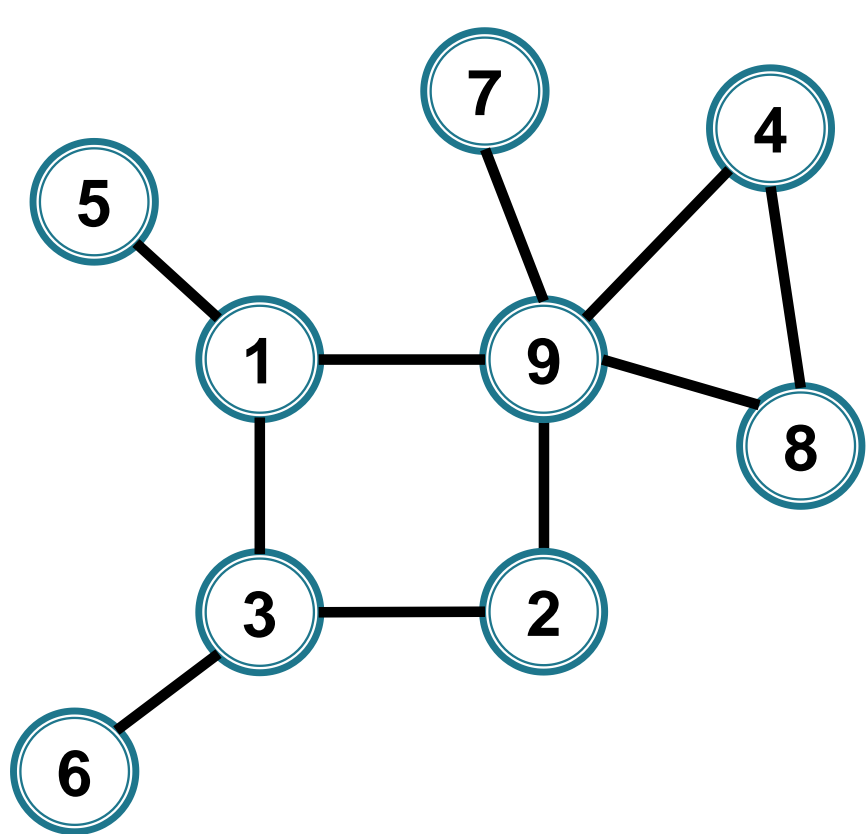
O muchie nu este critică  $\Leftrightarrow$  este conținută într-un ciclu



# Muchii critice

- ▶ Găsirea unui ciclu – parcurgere DF
  - **muchii de avansare** – ale arborelui DF (memorat cu vector tata), prin care se descoperă vârfuri noi
  - **muchii de întoarcere** – închid ciclu, nu pot fi critice

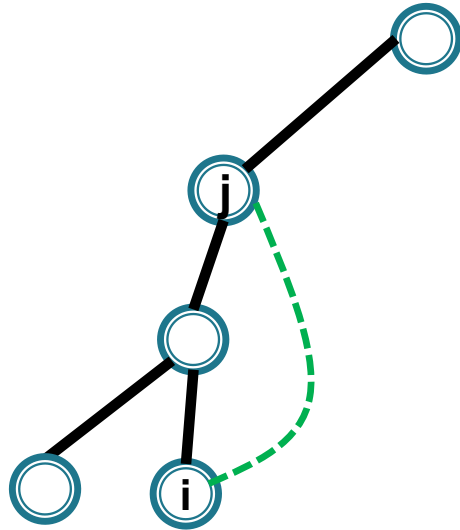






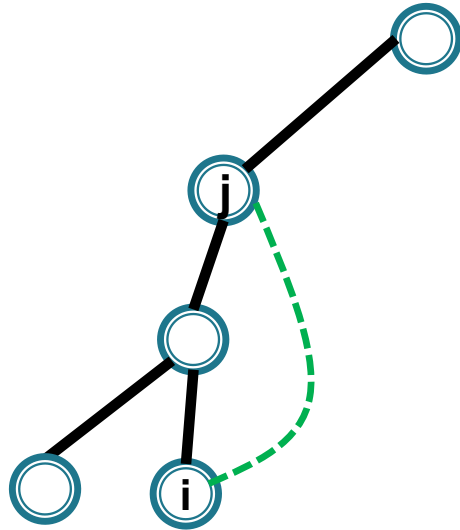
# Muchii critice

- ▶ Găsirea unui ciclu – parcurgere DF –  $O(n)$ 
  - muchii de avansare – ale arborelui DF (memorat cu vector tata), prin care se descoperă vârfuri noi
  - muchii de întoarcere – închid ciclu



# Muchii critice

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  - muchii de avansare – ale arborelui DF (memorat cu vector tata), prin care se descoperă vârfuri noi
  - muchii de întoarcere – închid ciclu, nu pot fi critice



Doar muchiile de avansare pot  
fi critice

# Muchii critice



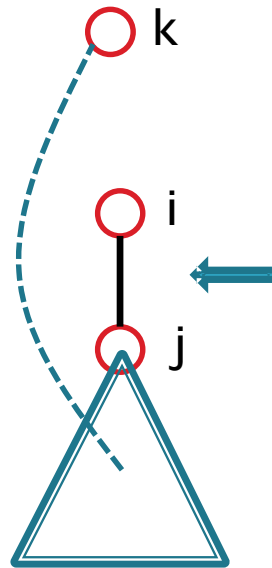
Cum testăm dacă o muchie de avansare  $(i,j)$  este critică?

# Muchii critice

Cum testăm dacă o muchie de avansare  $(i,j)$  este critică?



- nu este conținută într-un ciclu închis de o muchie de întoarcere



# Muchii critice

O muchie de avansare  $(i,j)$  este critică

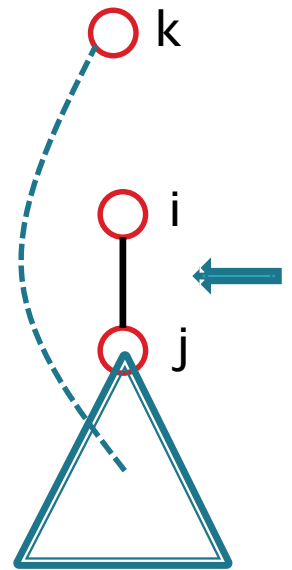
$\Leftrightarrow$

nu este conținută într-un ciclu închis de o muchie de întoarcere

$\Leftrightarrow$

nu există nicio muchie de întoarcere cu

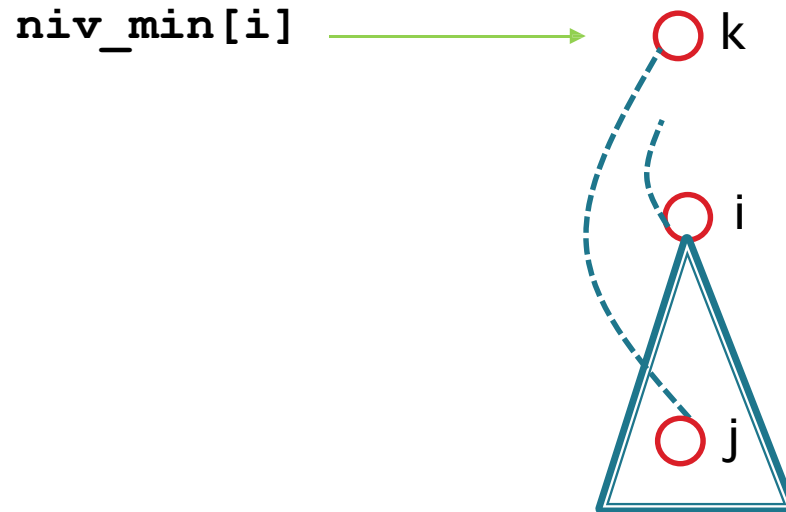
- o extremitate **în  $j$  sau într-un descendent al lui  $j$**  și
- cealaltă extremitate **în  $i$  sau într-un ascendent al lui  $i$**  (într-un vârf de pe un nivel mai mic sau egal cu nivelul lui  $i$ )



# Muchii critice

Memorăm pentru fiecare vârf  $i$ :

$\text{niv\_min}[i]$  = *intuitiv: cât de sus putem ajunge din  $i$  mergând în sensul parcurgeii DF*

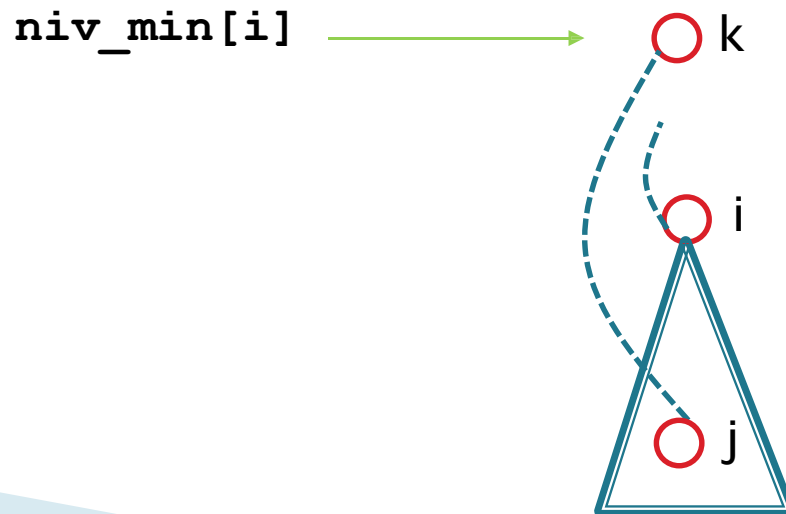


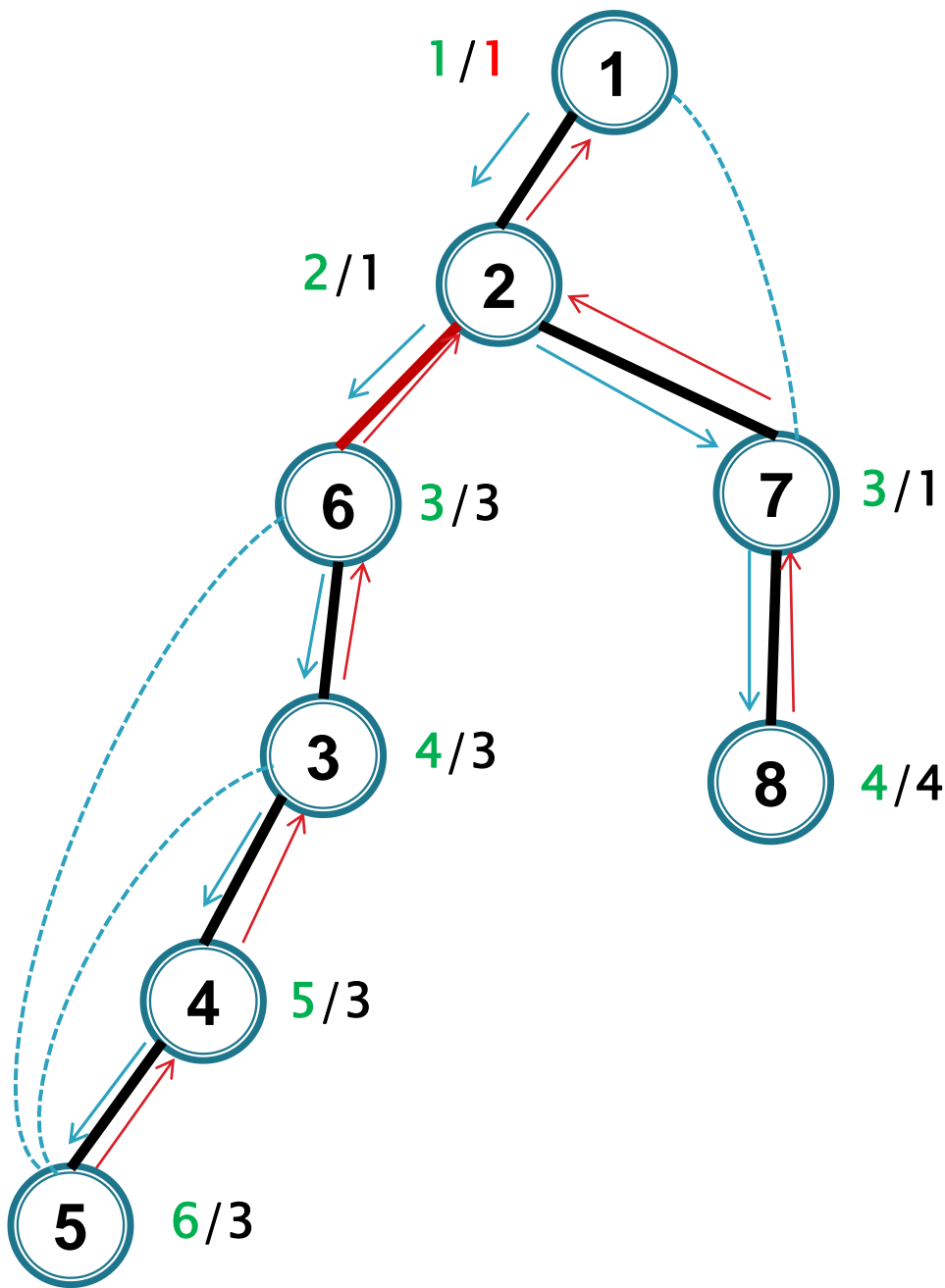
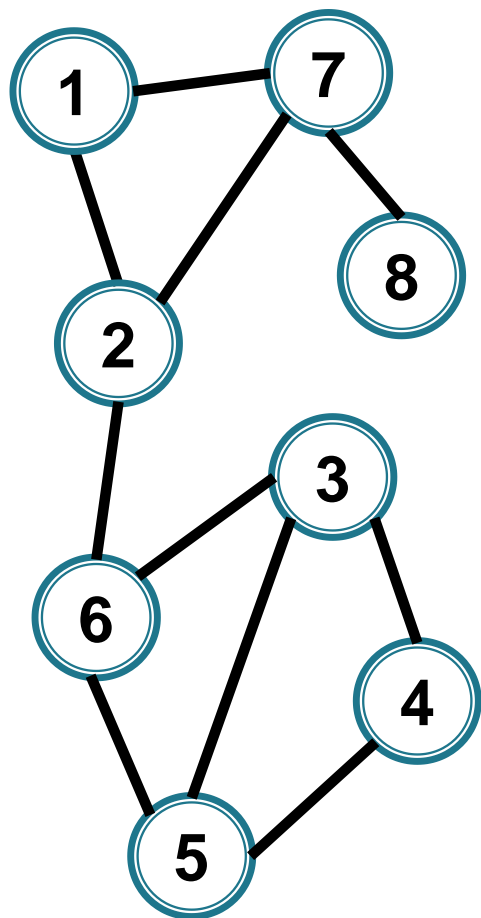
# Muchii critice

Memorăm pentru fiecare vârf  $i$ :

$niv\_min[i]$  = nivelul minim al unui vârf care este extremitate a unei muchii de întoarcere din  $i$  sau dintr-un descendent al lui  $i$

= nivelul minim la care se închide un ciclu elementar care conține vârful  $i$  (printr-o muchie de întoarcere)

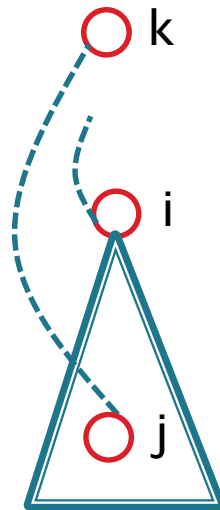






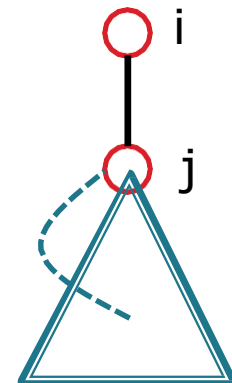
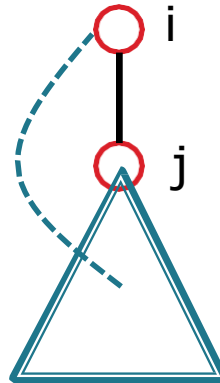
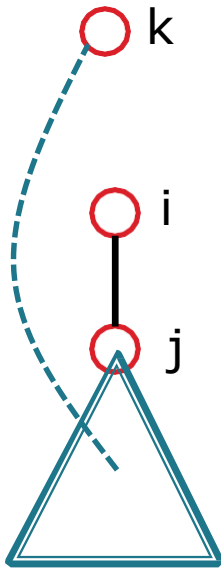
# Muchii critice

- ▶  $\text{nivel}[i] = \text{nivelul lui } i \text{ în arborele DF}$
- ▶  $\text{niv\_min}[i] = \min \{ \text{nivel}[i], A, B \}$ 
  - $A = \min \{ \text{nivel}[k] \mid i \text{ } k \text{ muchie de întoarcere} \}$
  - $B = \min \{ \text{nivel}[k] \mid j \text{ descendent al lui } i, j \text{ } k \text{ muchie de întoarcere} \}$



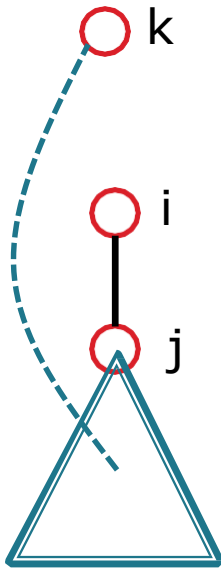
# Muchii critice

O muchie de avansare  $ij$  este critică  $\Leftrightarrow$



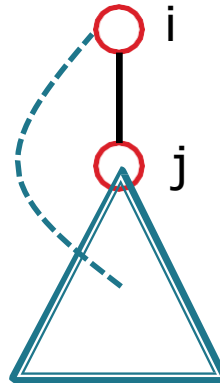
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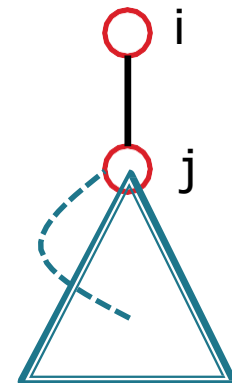
**NU este critică**

$\text{niv\_min}[j] < \text{nivel}[i]$



**NU este critică**

$\text{niv\_min}[j] = \text{nivel}[i]$

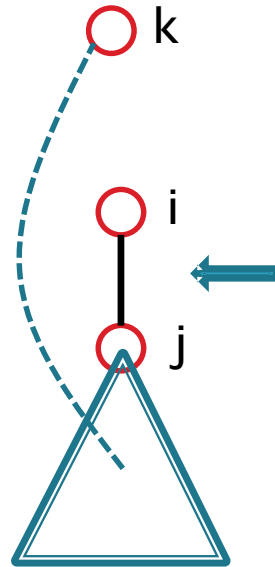


**ESTE critică**

$\text{niv\_min}[j] > \text{nivel}[i]$

# Muchii critice

O muchie de avansare  $ij$  este critică  $\Leftrightarrow \text{niv\_min}[j] > \text{nivel}[i]$



# Muchii critice



Cum calculăm eficient `niv_min[i]` ?

$$\text{niv\_min}[i] = \min \{ \text{nivel}[i], A, B \}$$
$$A = \min \{ \text{nivel}[k] \mid i \text{ k muchie de întoarcere} \}$$
$$B = \min \{ \text{nivel}[k] \mid j \text{ descendent al lui } i, \\ j \text{ k muchie de întoarcere} \}$$

# Muchii critice

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**B se poate calcula recursiv**

# Muchii critice

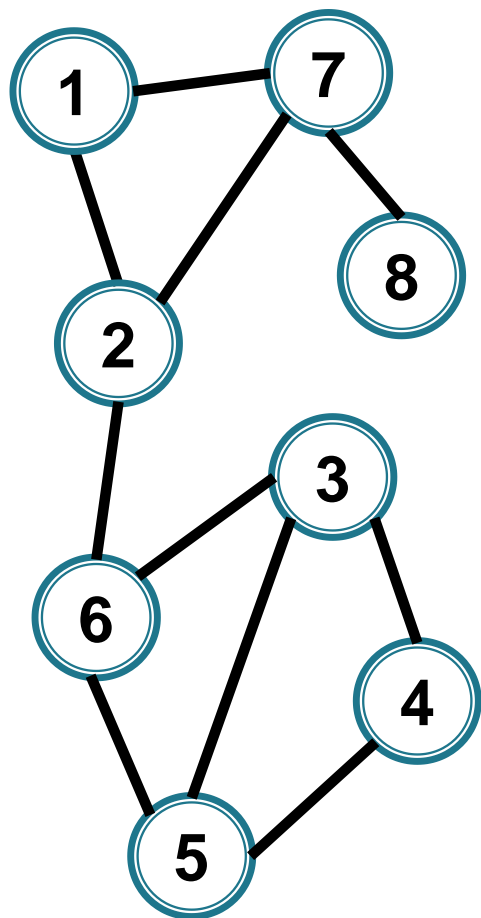
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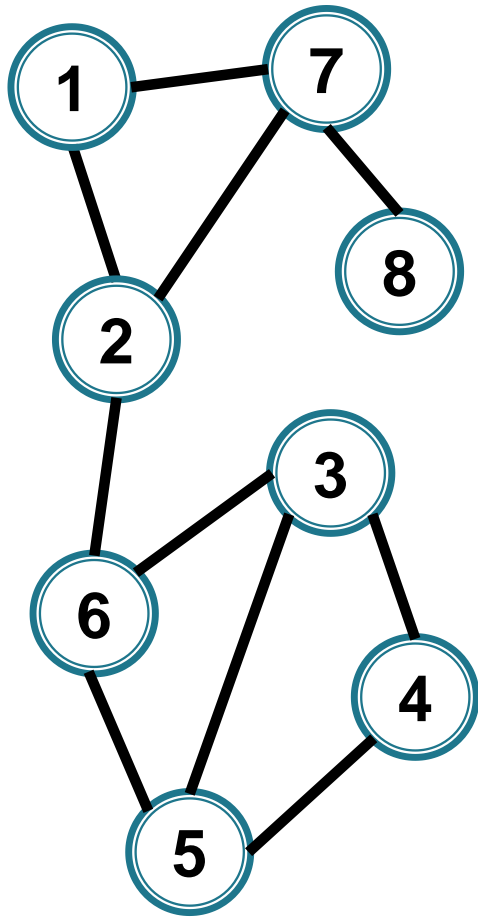
$$B = \min \{ niv\_min[j] \mid j \text{ fiu al lui } i \}$$



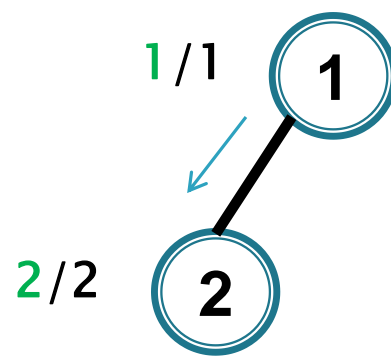
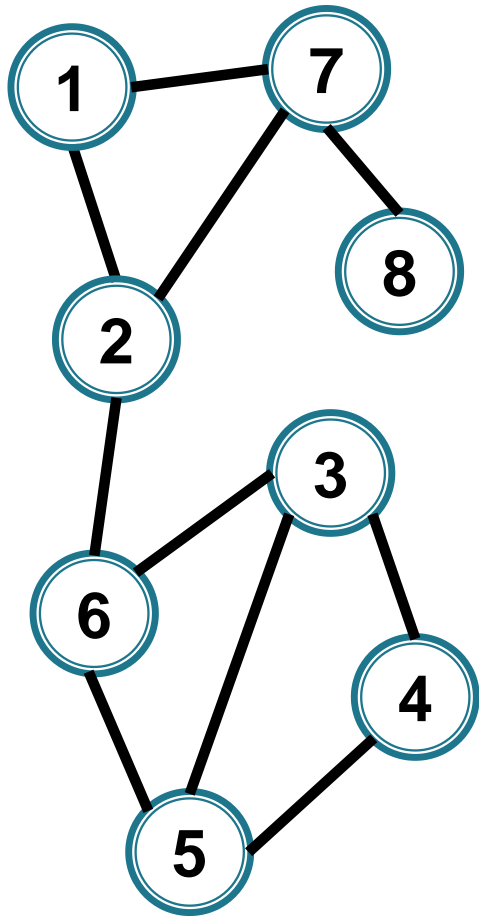


nivel/niv\_min

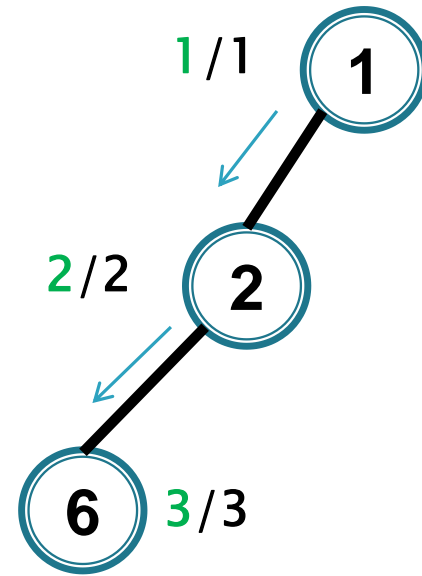
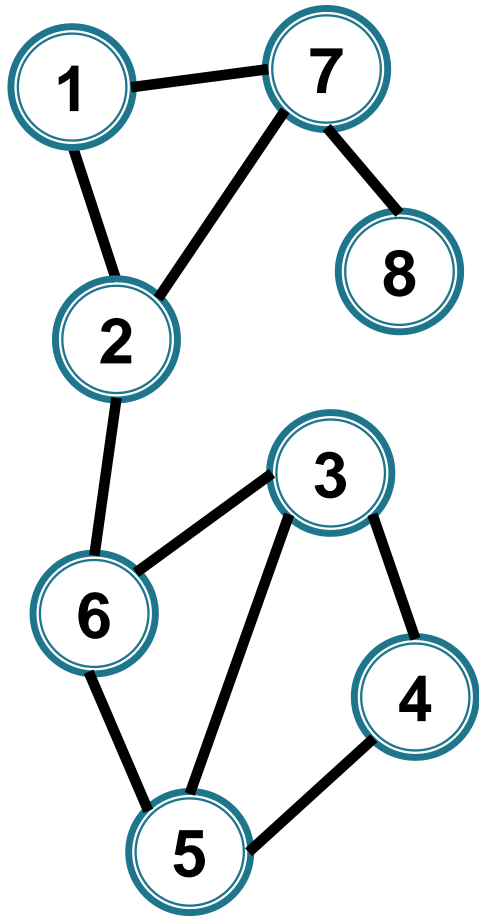
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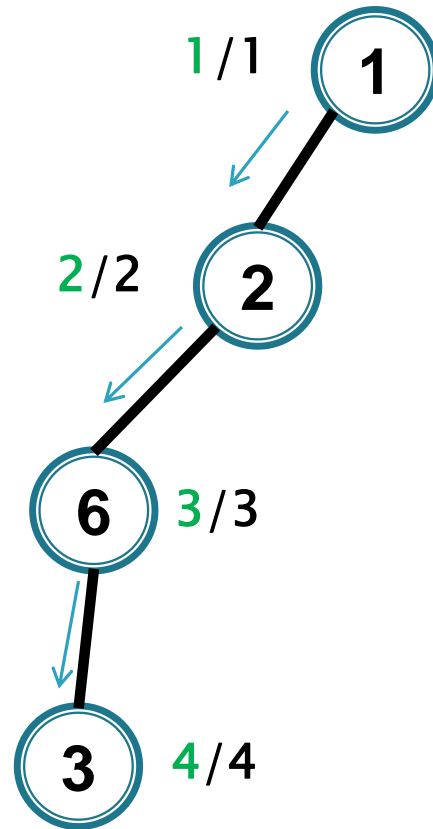
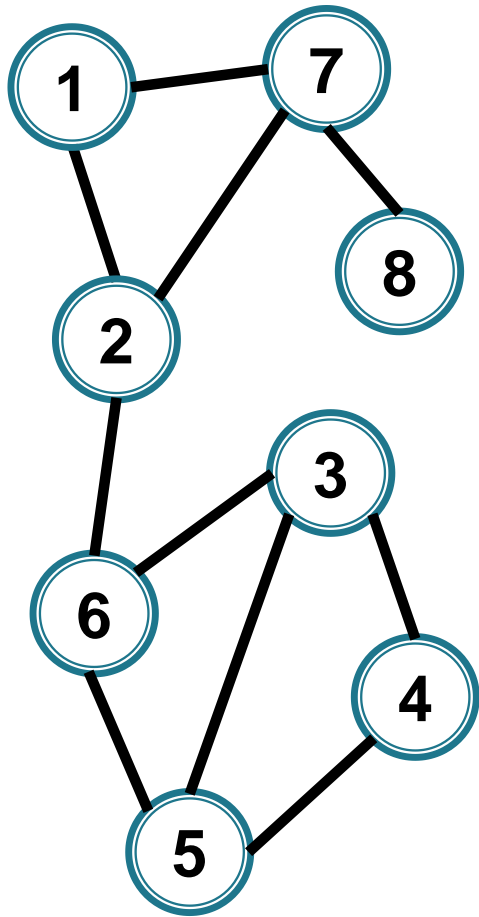
nivel/niv\_min



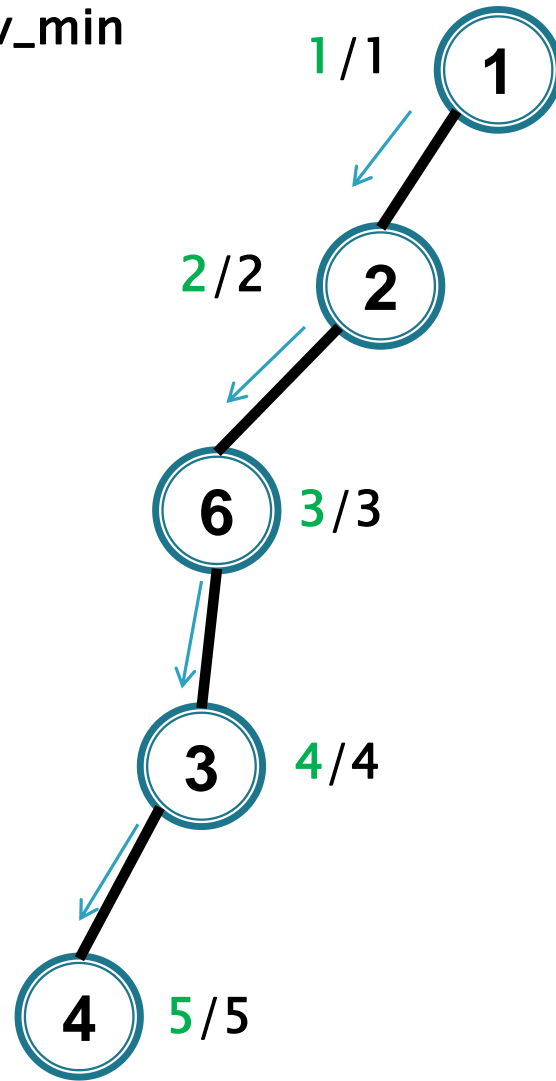
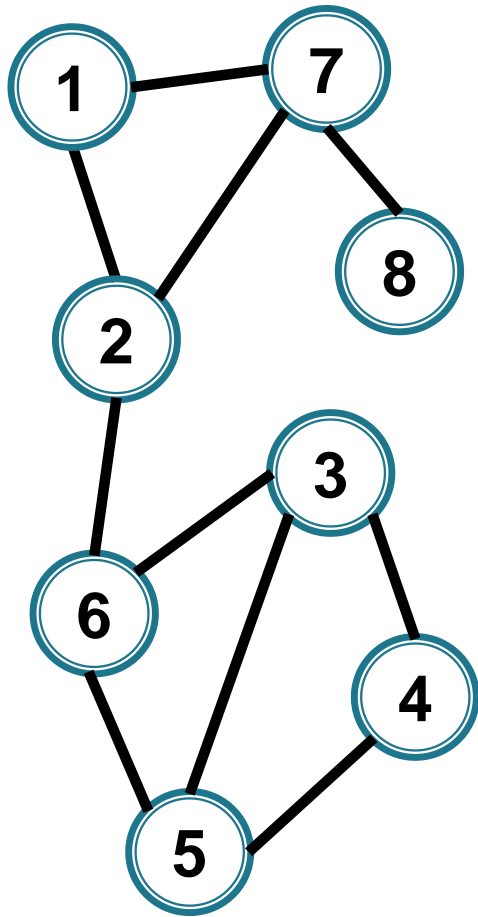
nivel/niv\_min

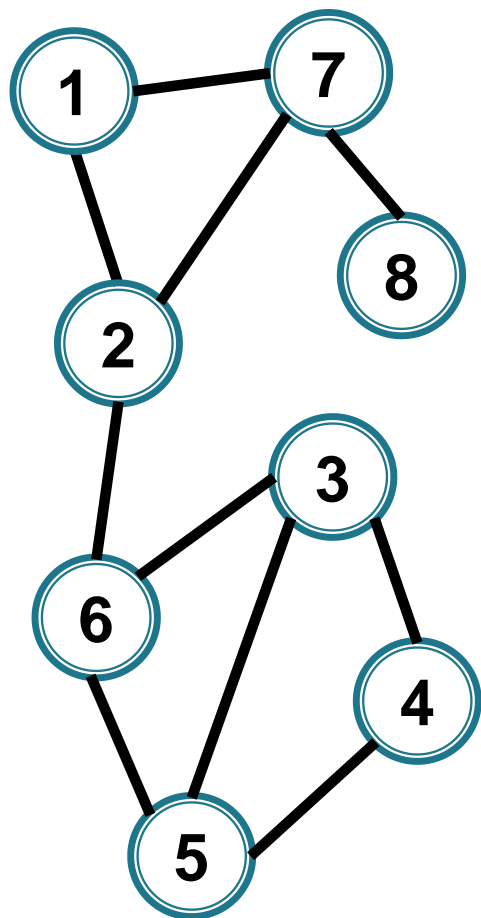


nivel/niv\_min

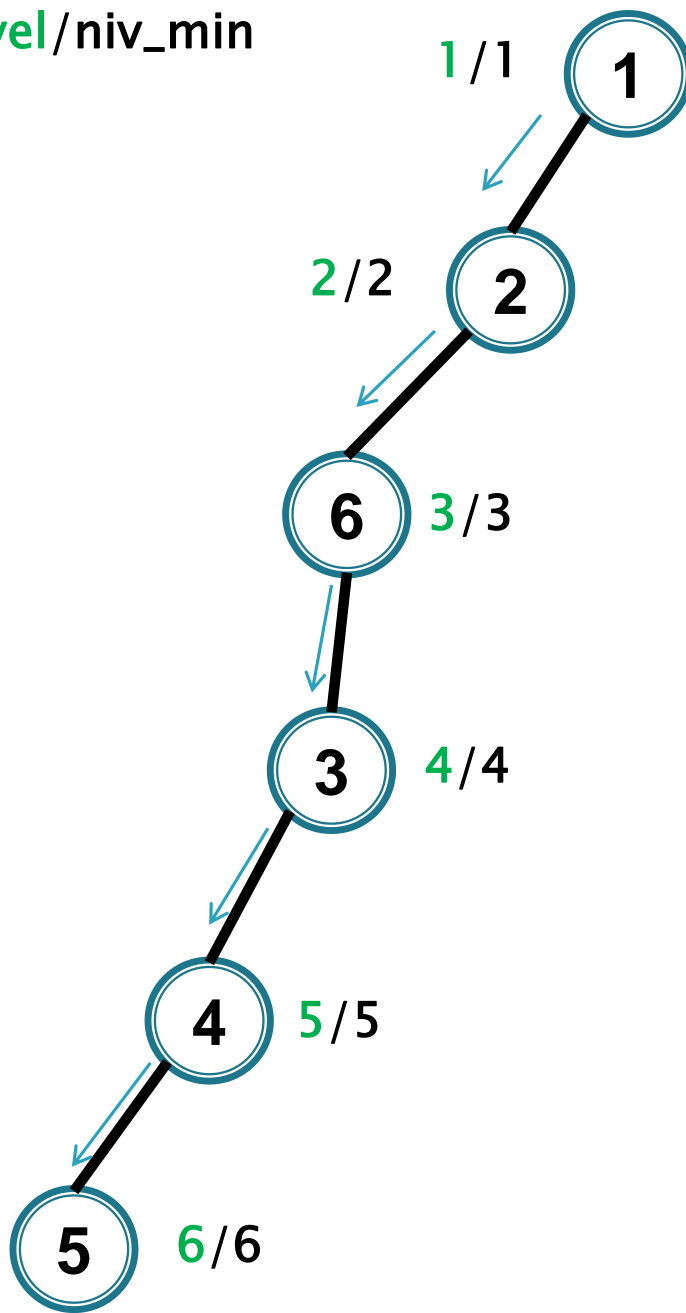


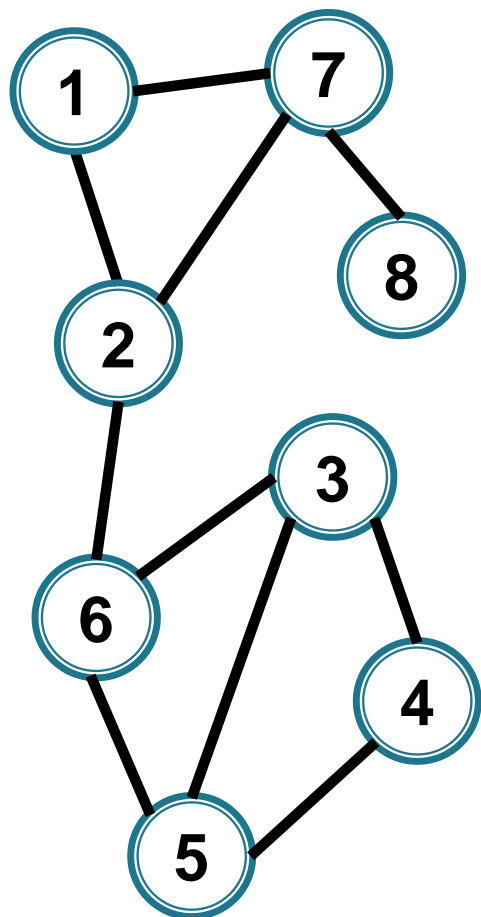
nivel/niv\_min



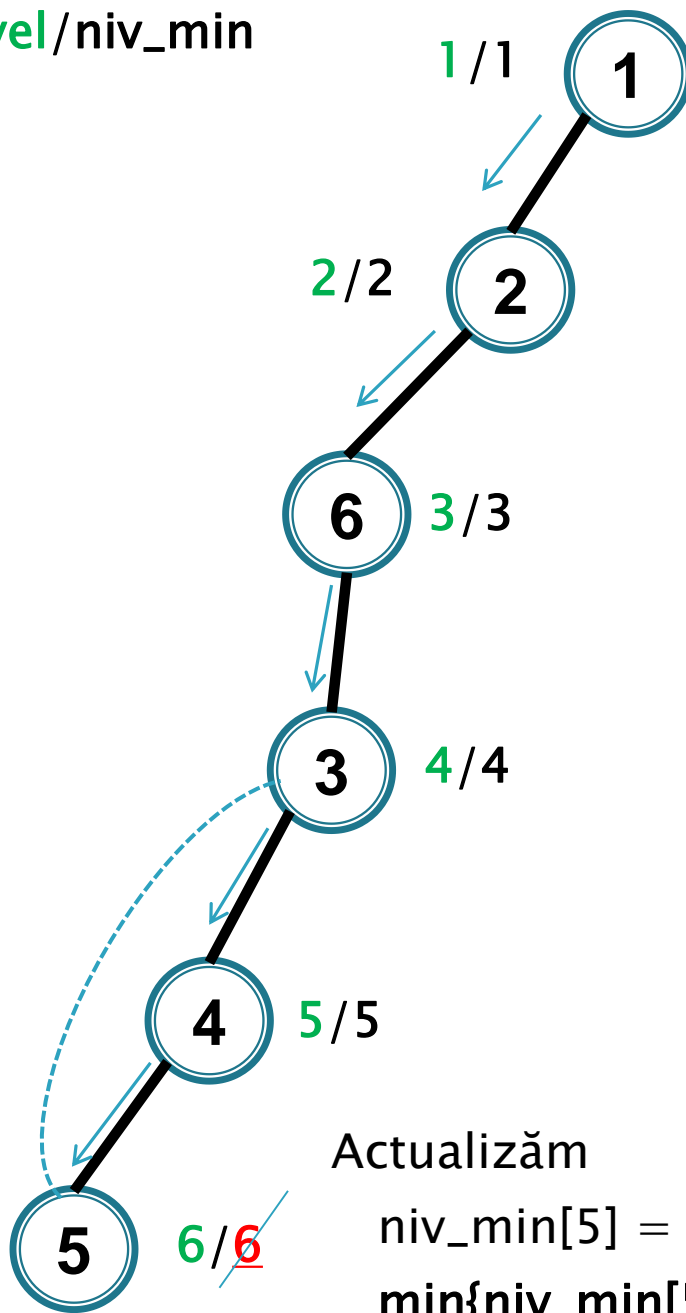


nivel/niv\_min





nivel/niv\_min

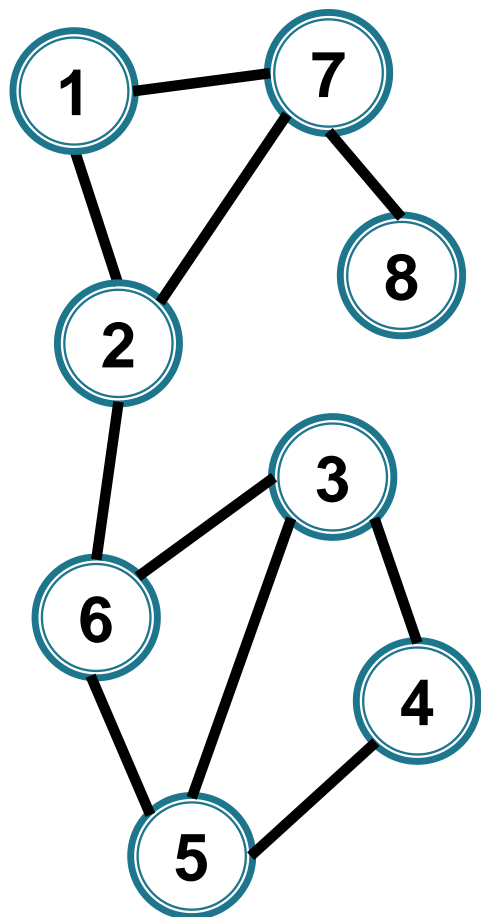


Actualizăm

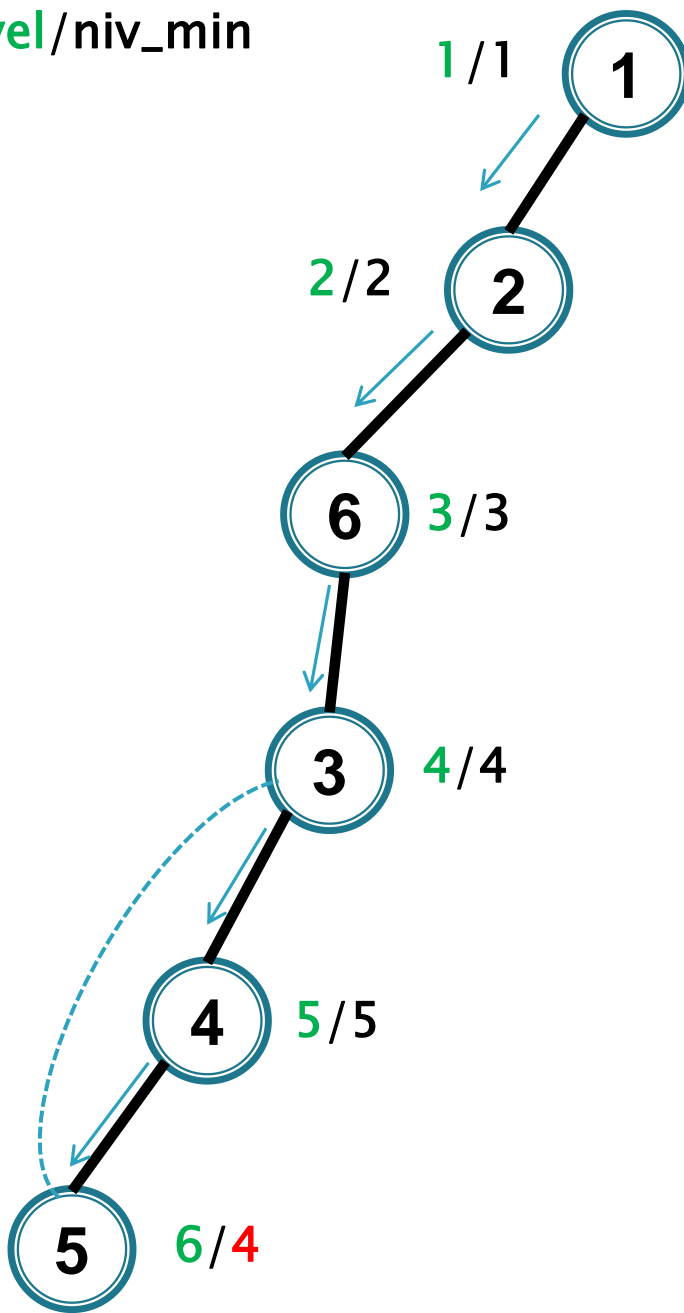
$\text{niv\_min}[5] =$

$\min\{\text{niv\_min}[5], \text{nivel}[3]\}$

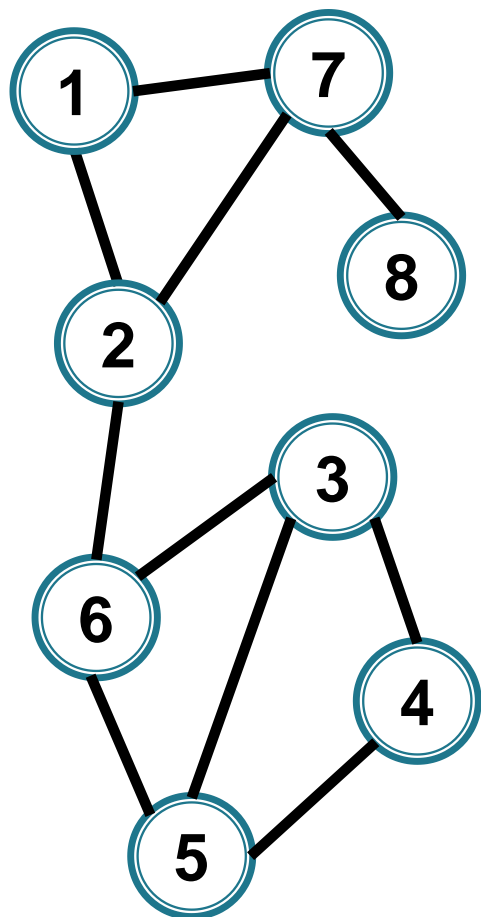
(cazul A)



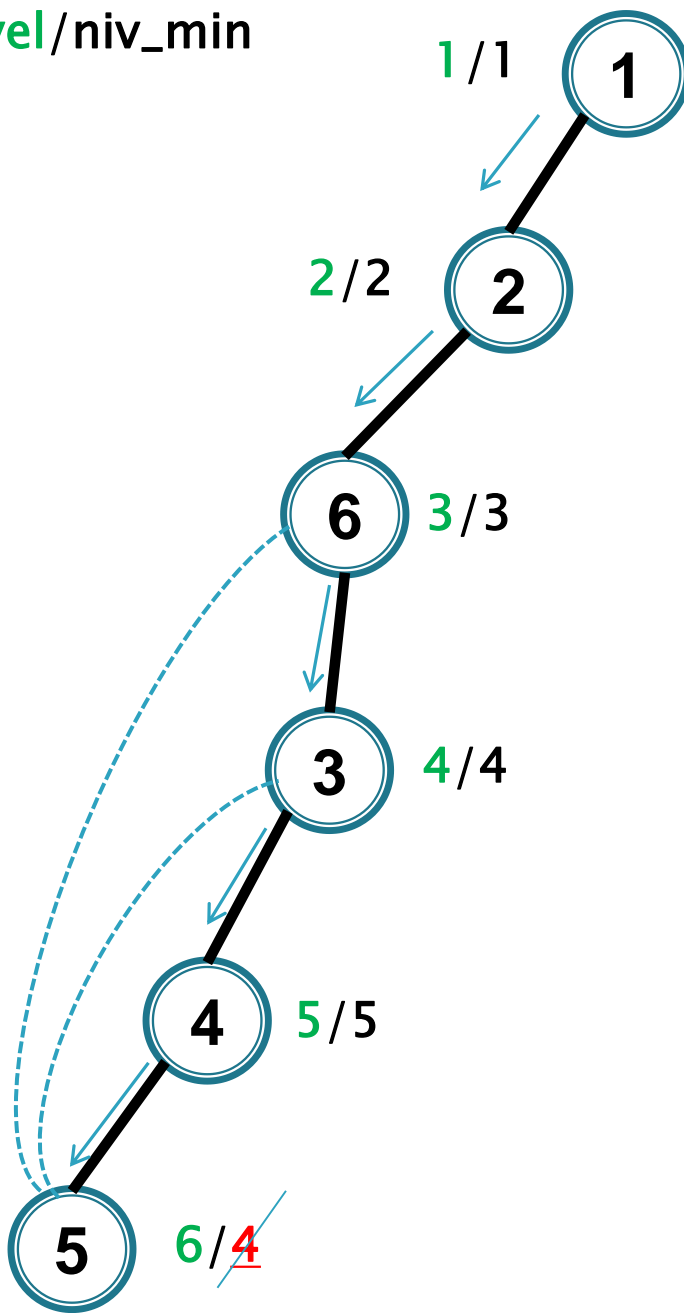
nivel/niv\_min

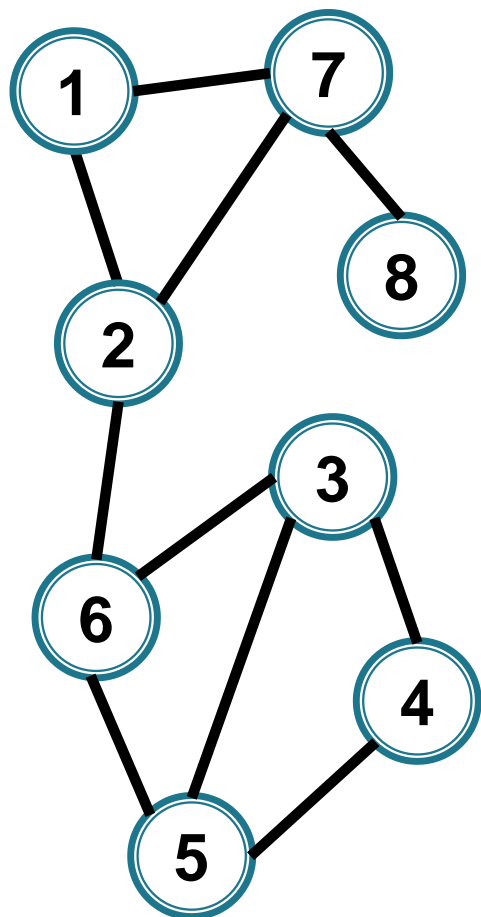




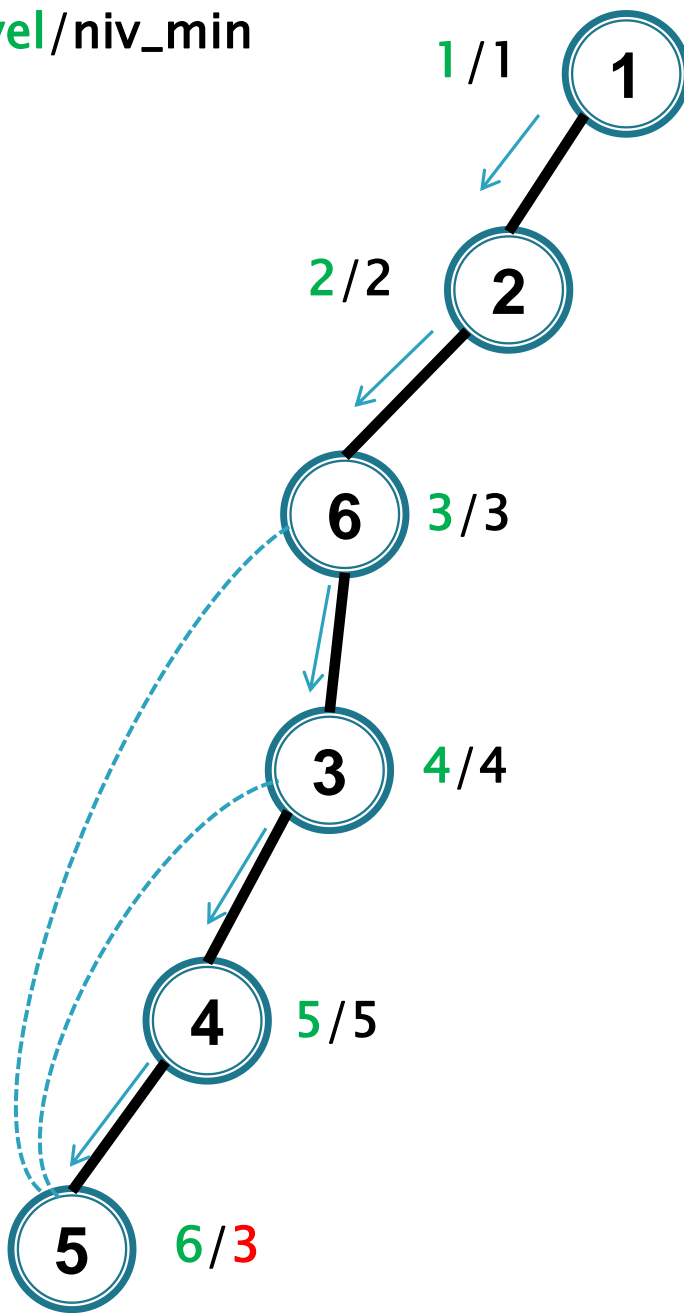


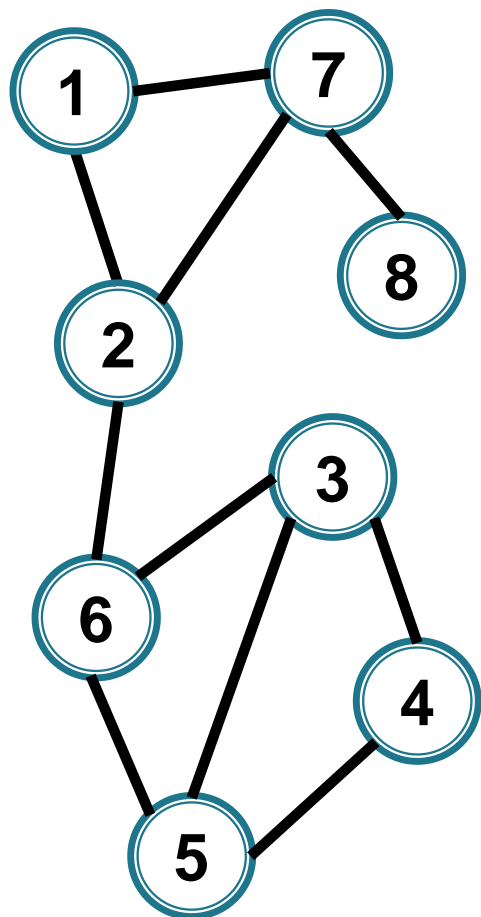
nivel/niv\_min



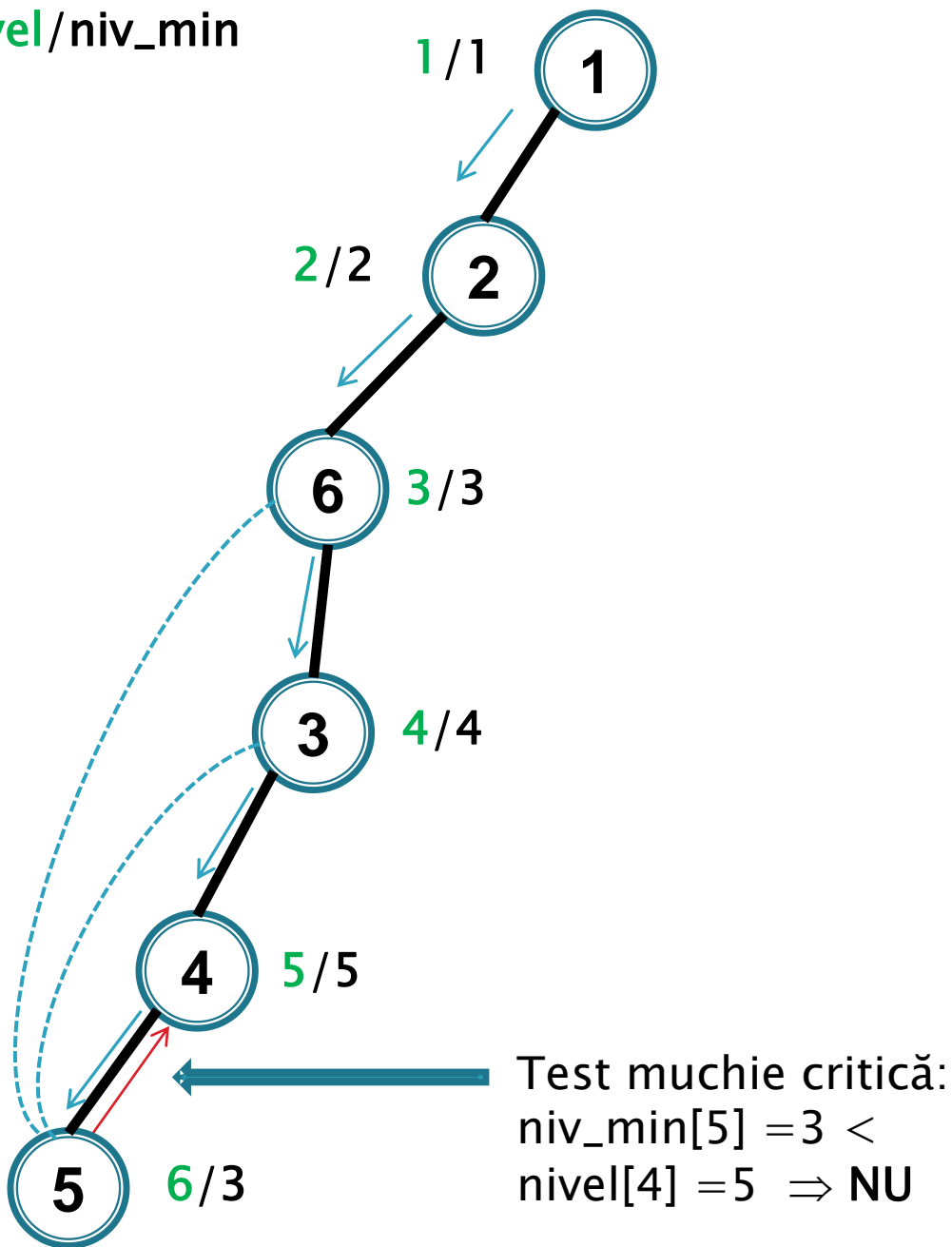


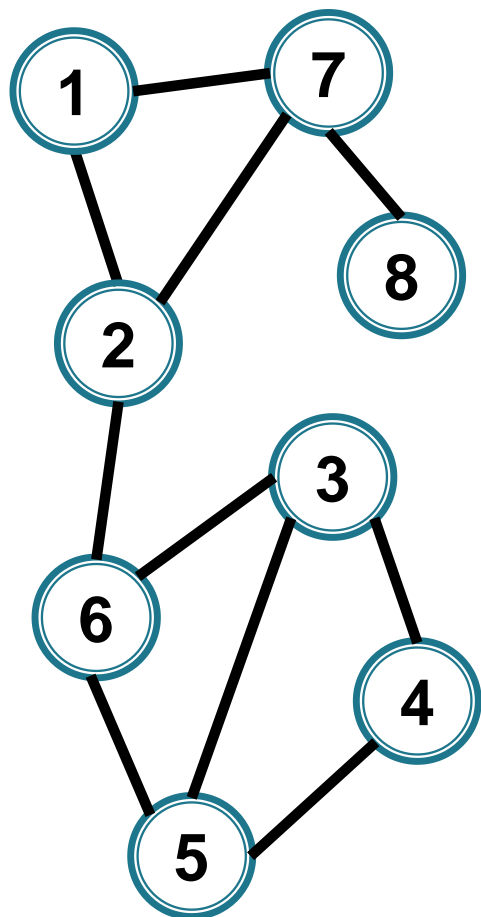
nivel/niv\_min



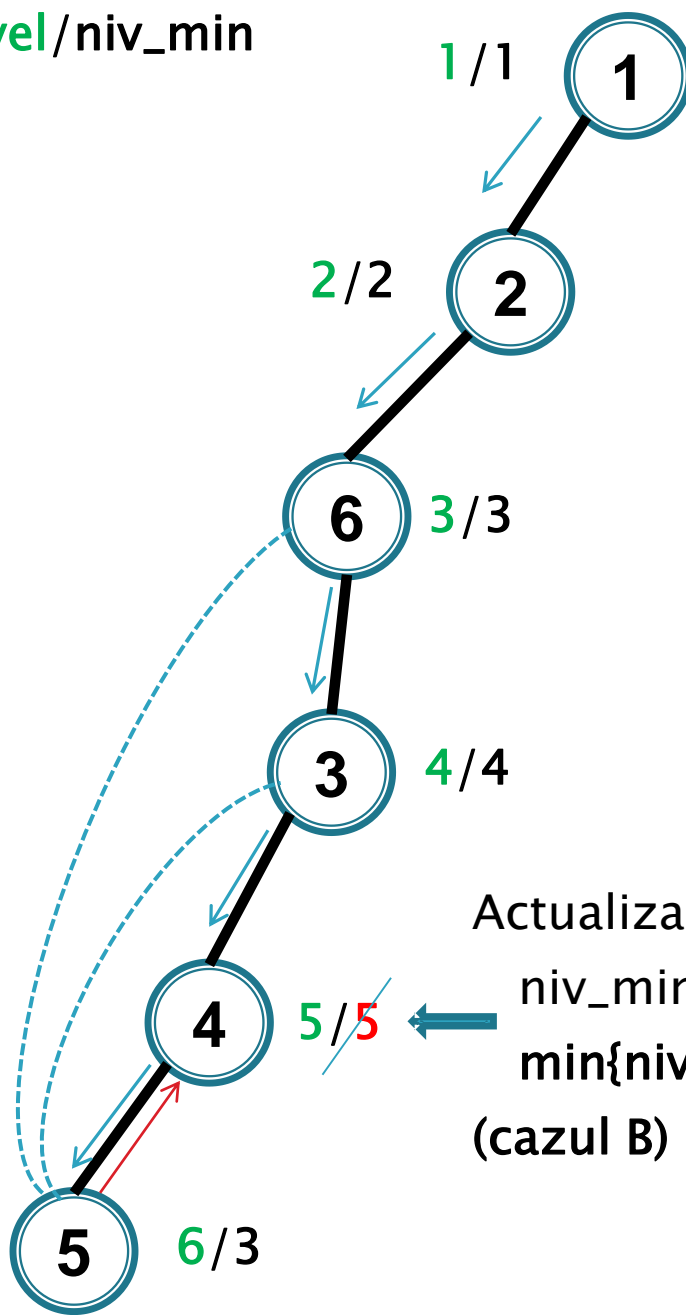


nivel/niv\_min





nivel/niv\_min

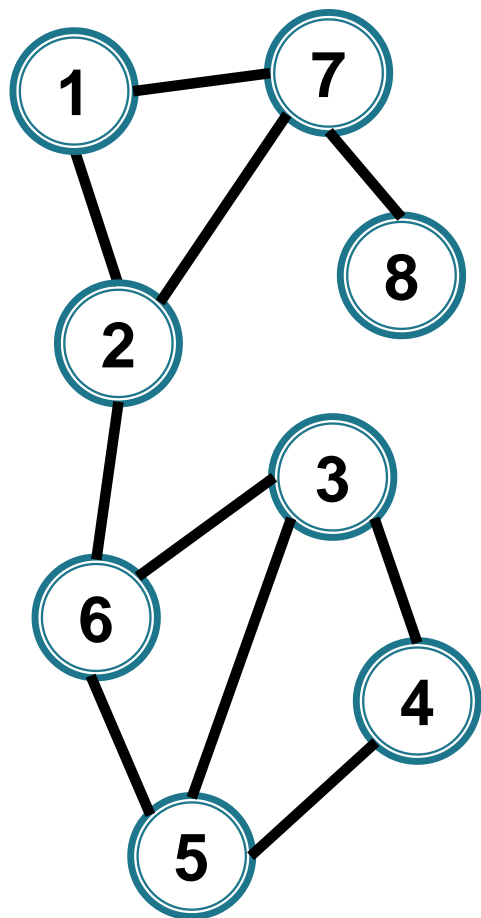


Actualizam

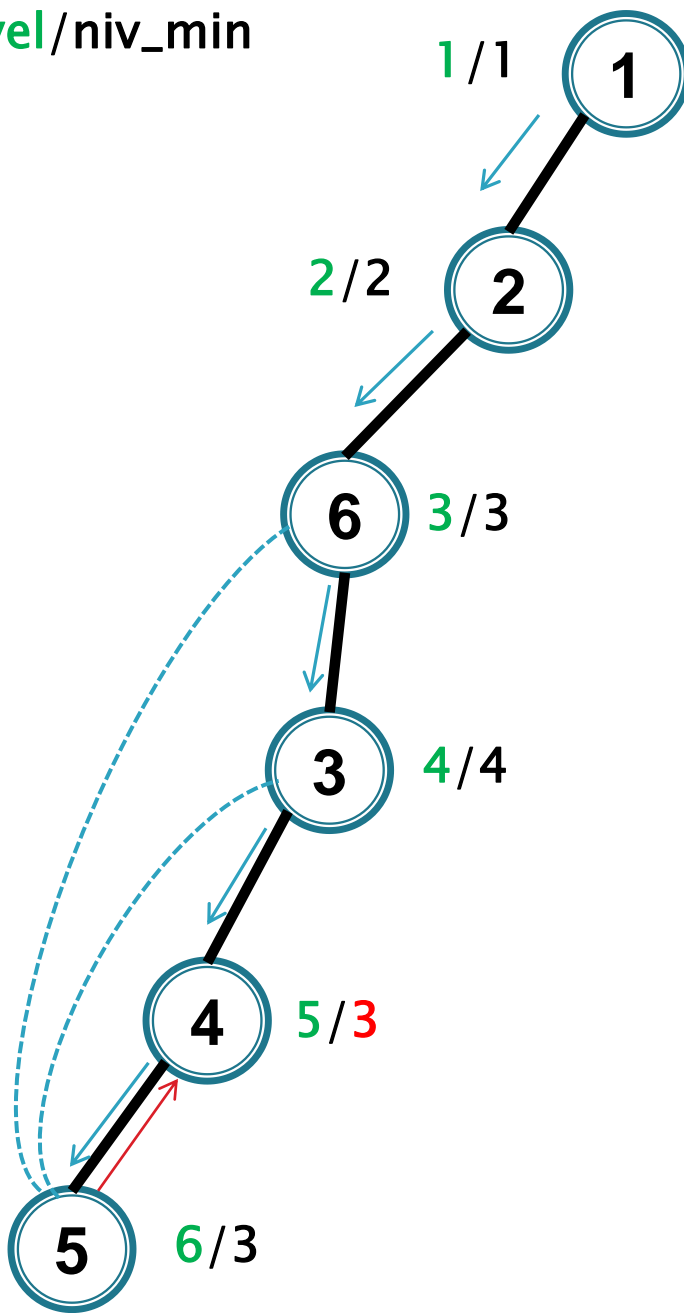
$\text{niv\_min}[4] =$

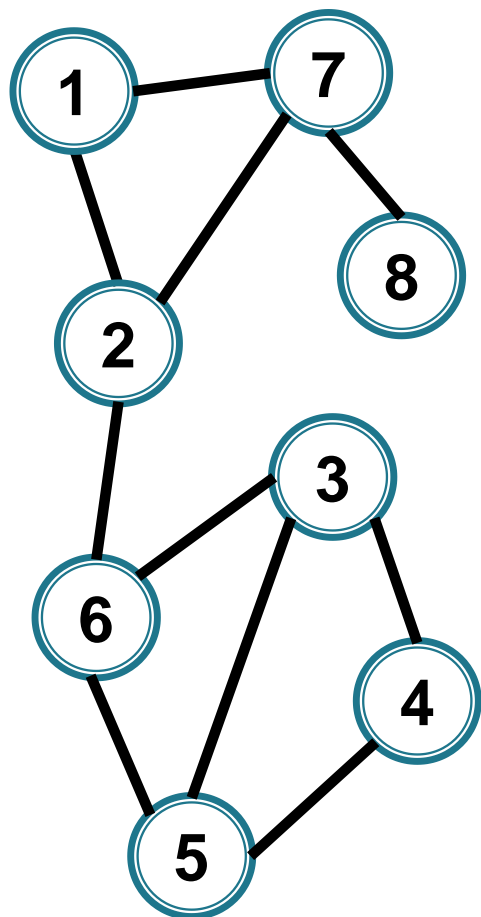
$\min\{\text{niv\_min}[4], \text{niv\_min}[5]\}$

(cazul B)

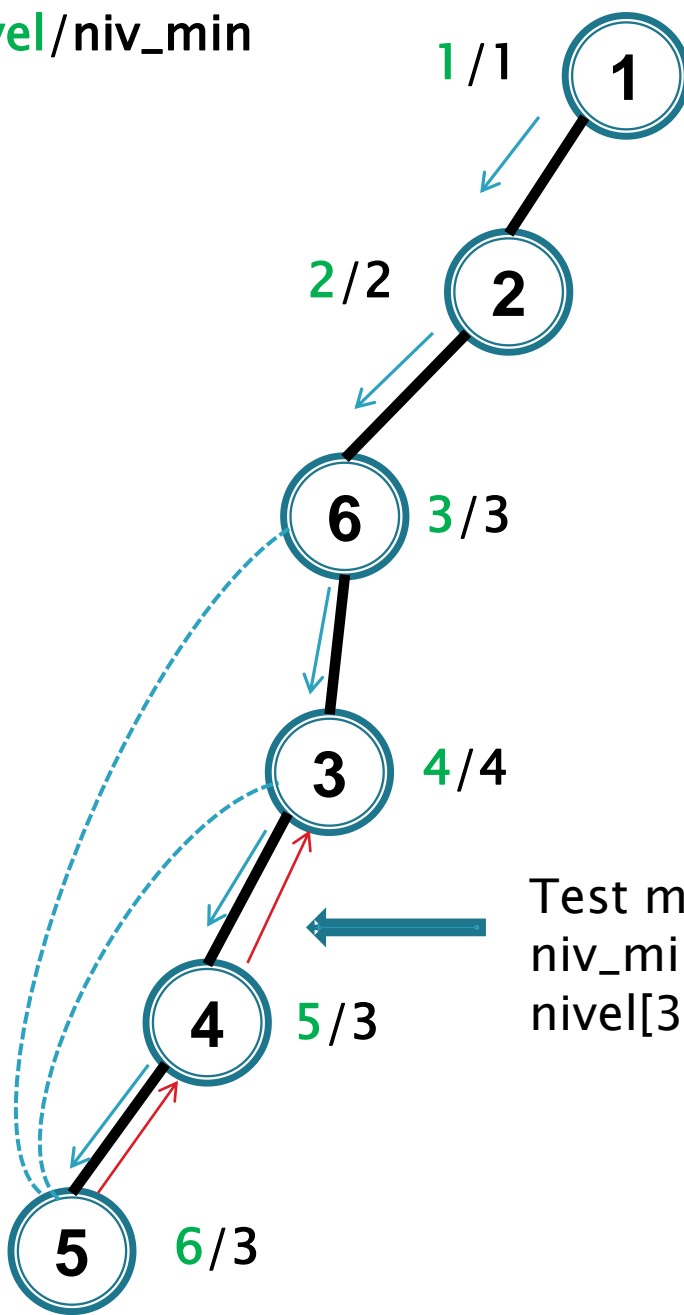


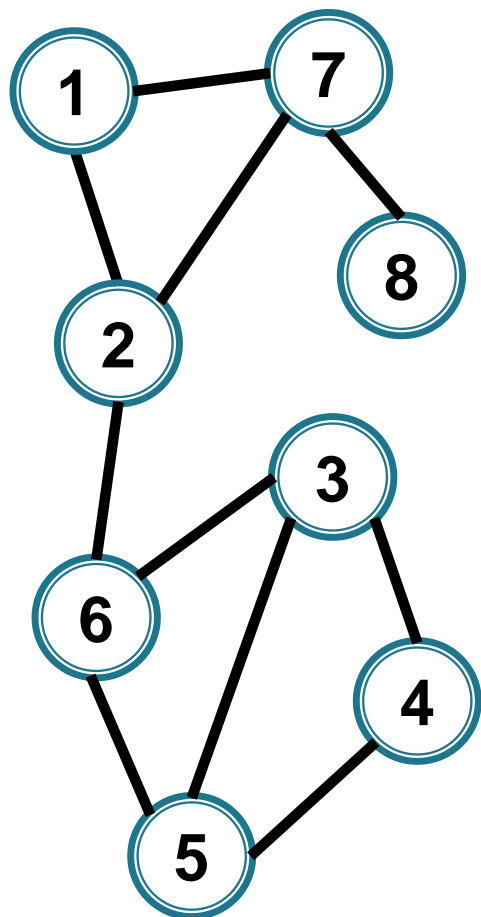
nivel/niv\_min



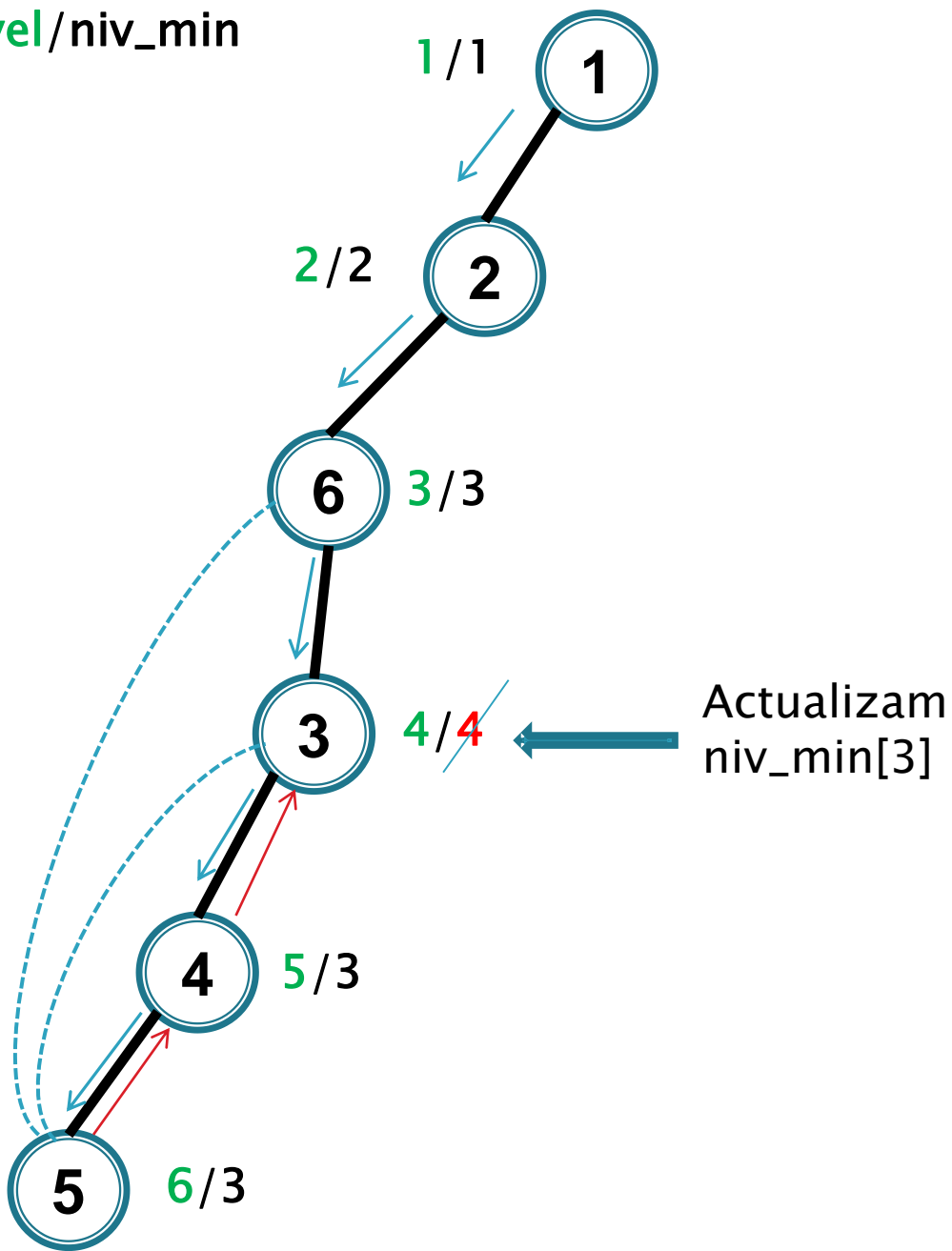


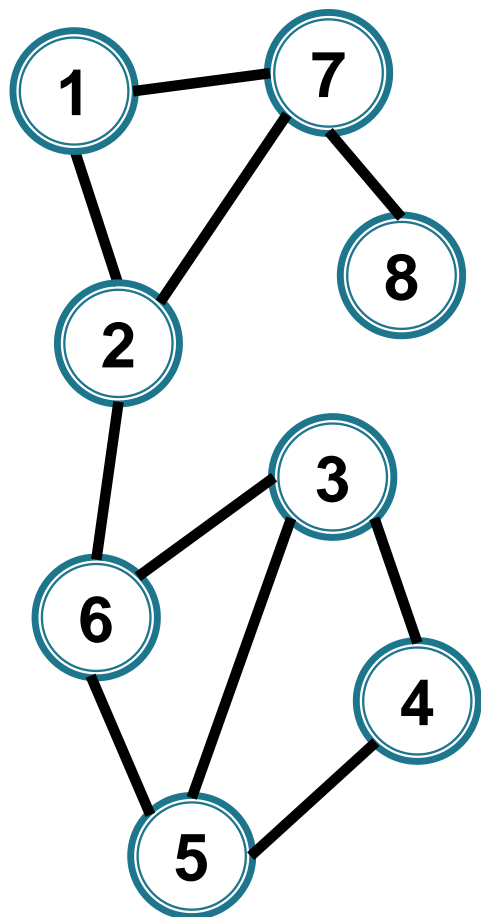
nivel/niv\_min



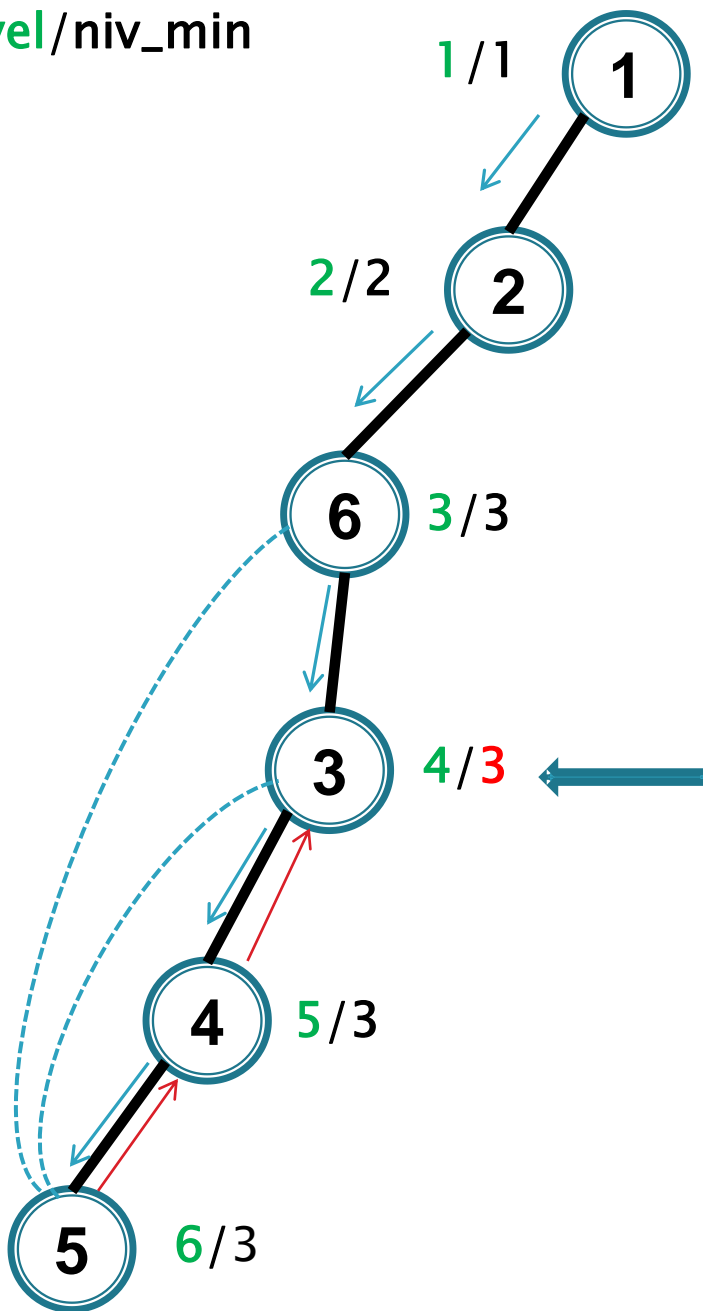


nivel/niv\_min

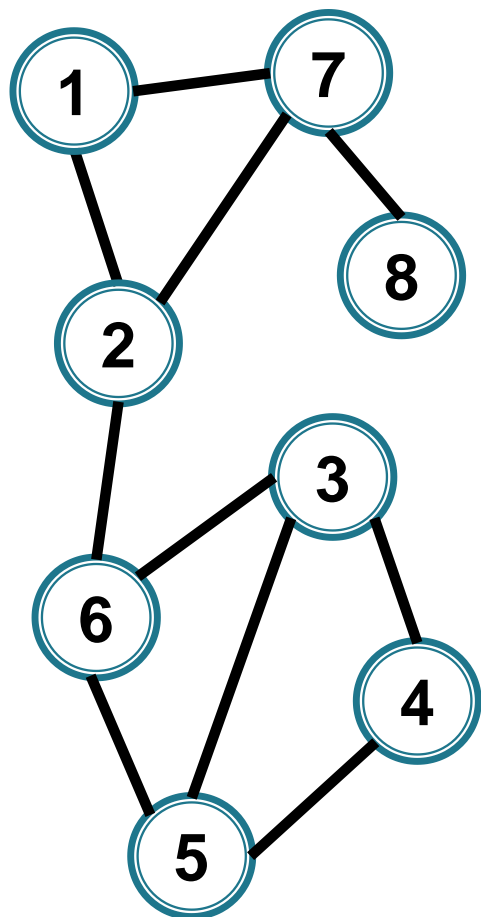




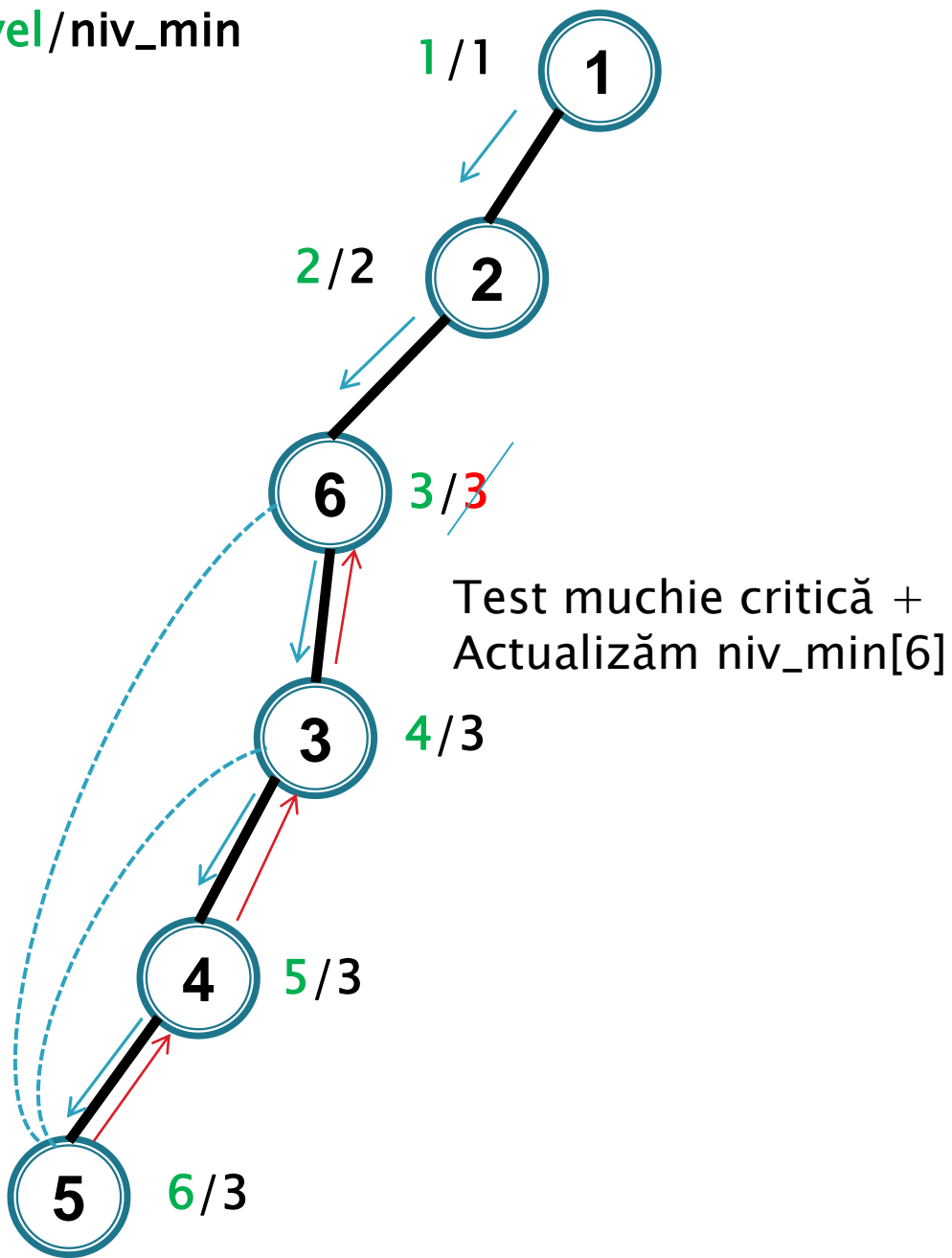
nivel/niv\_min

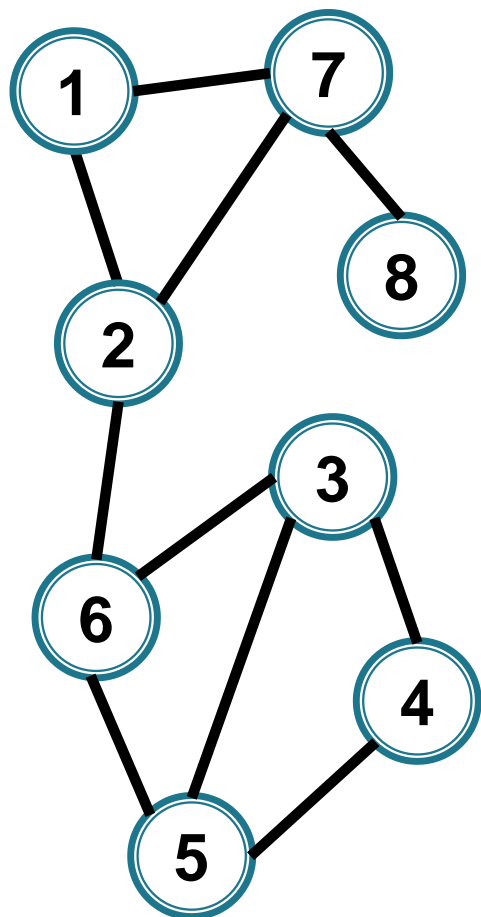




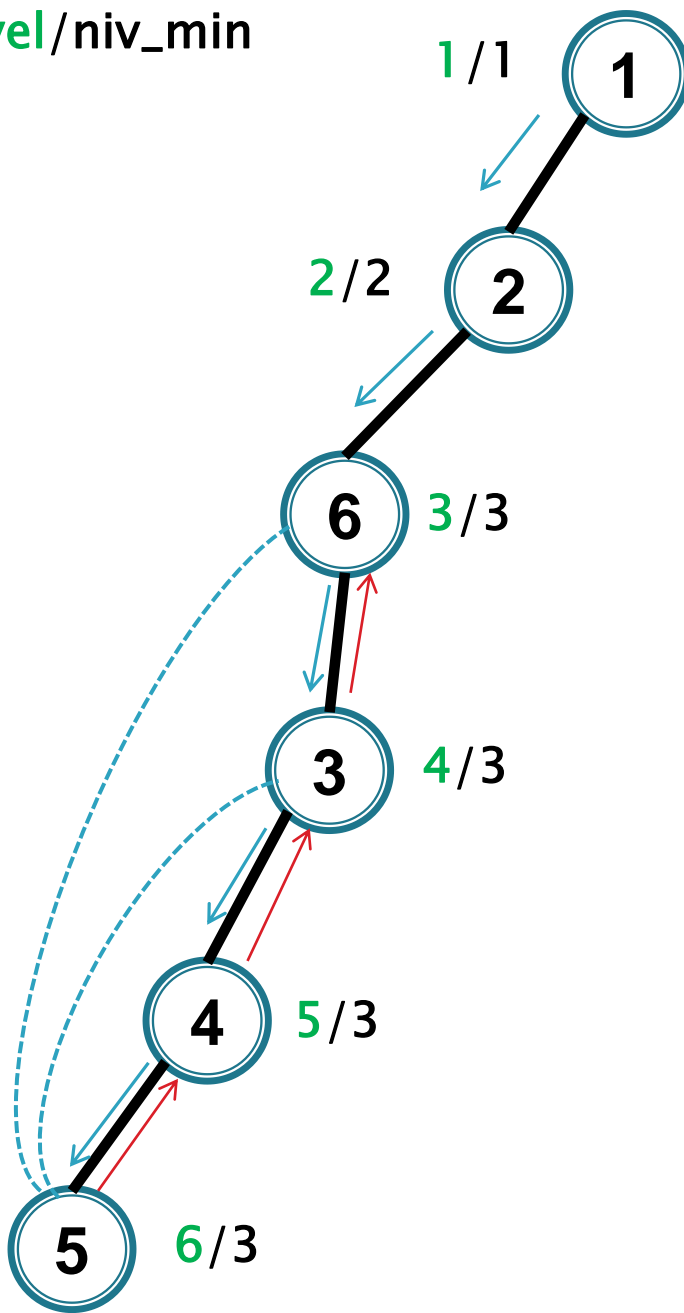


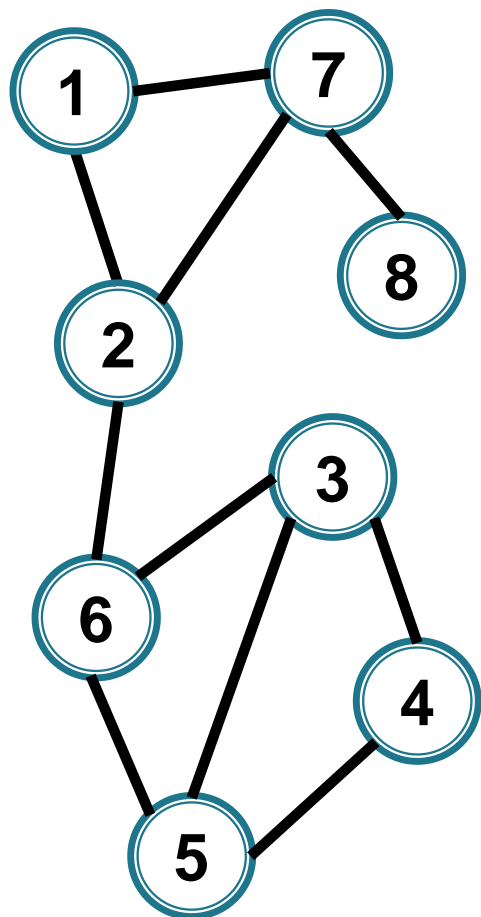
nivel/niv\_min



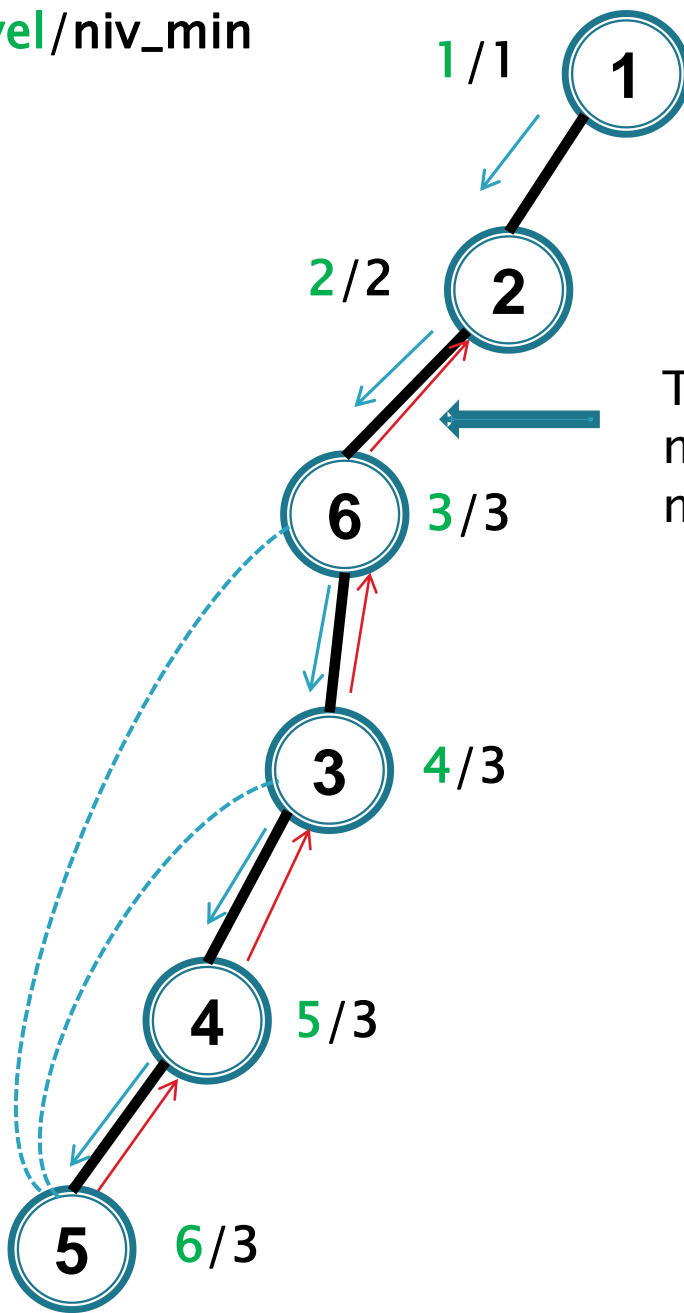


nivel/niv\_min

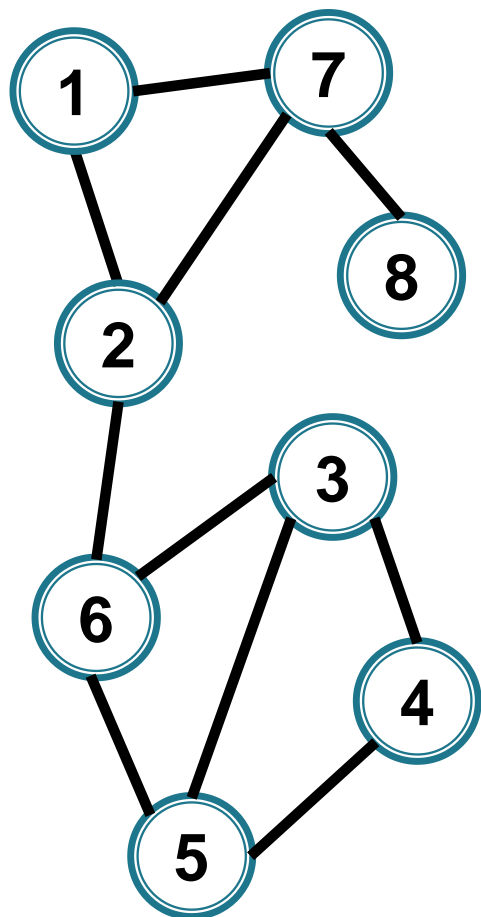




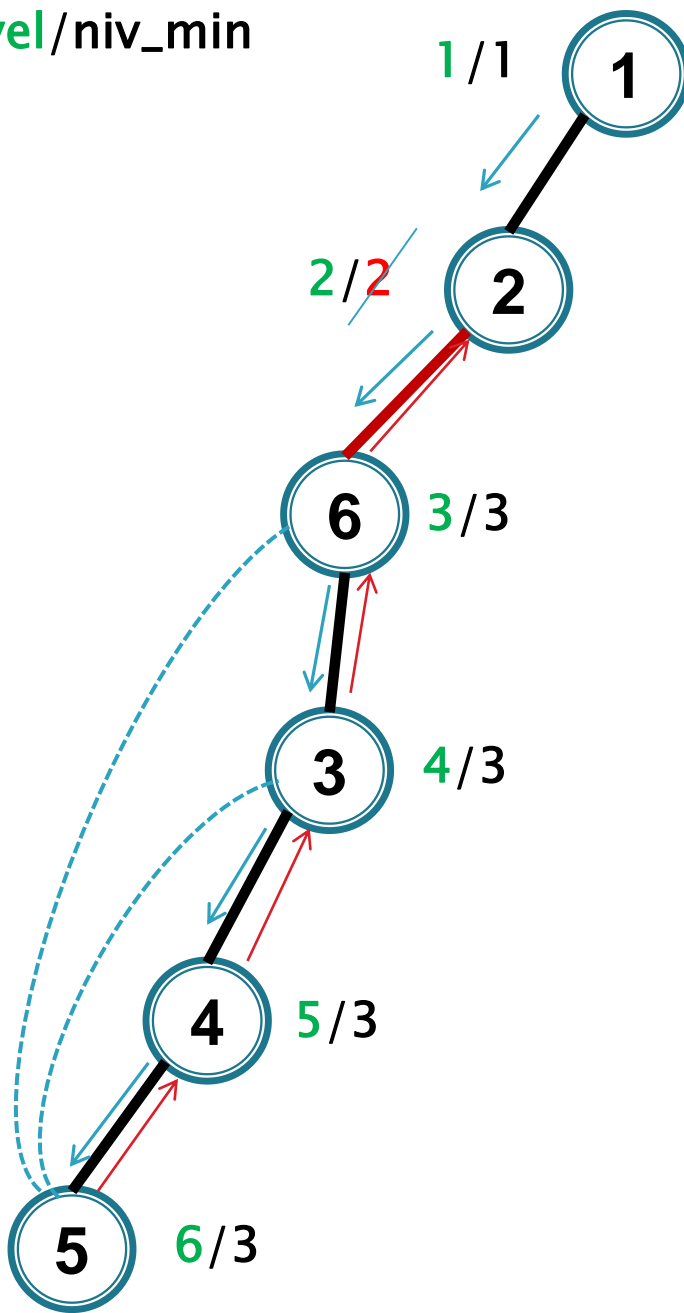
nivel/niv\_min

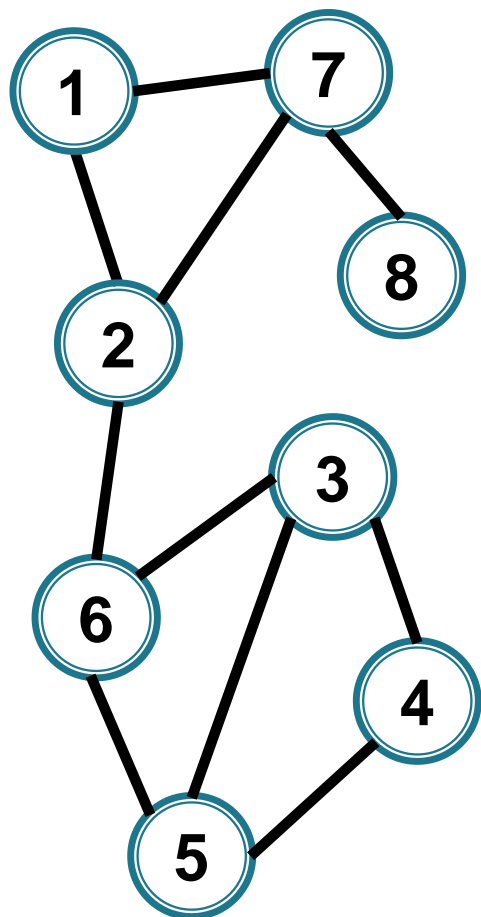


Test muchie critică:  
 $\text{niv\_min}[6] = 3 >$   
 $\text{nivel}[2] = 2 \Rightarrow \text{DA}$

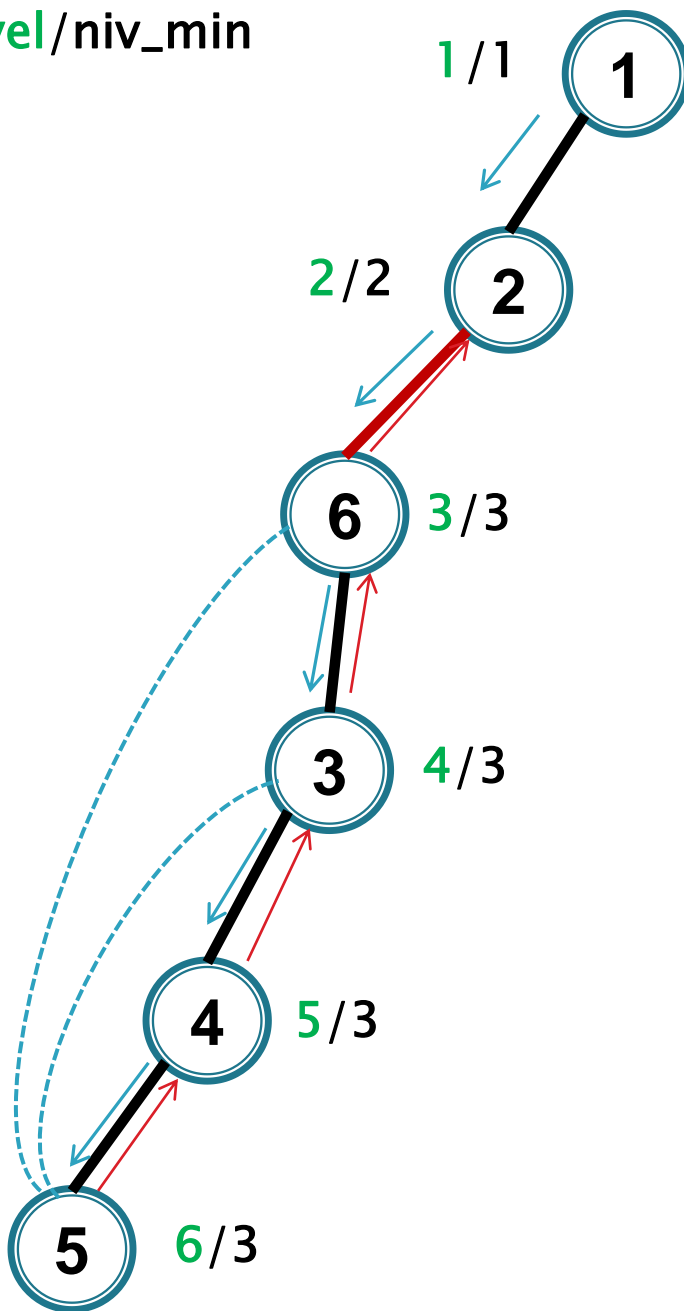


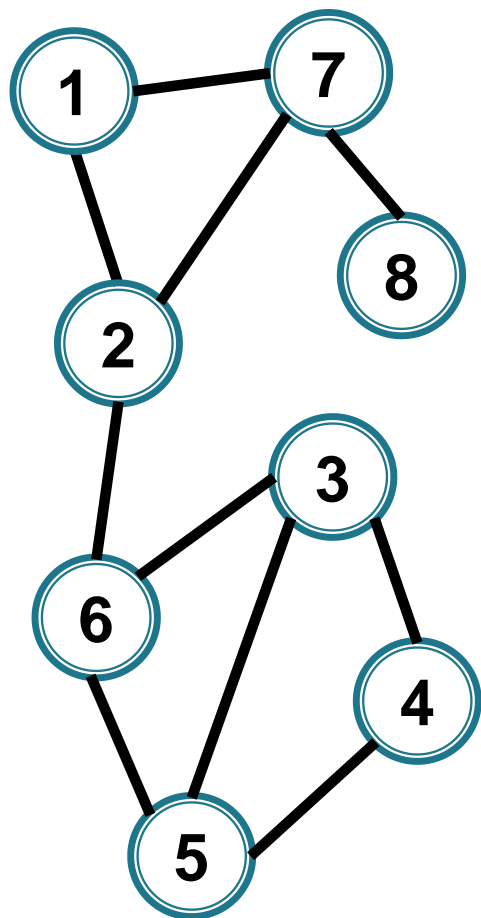
nivel/niv\_min



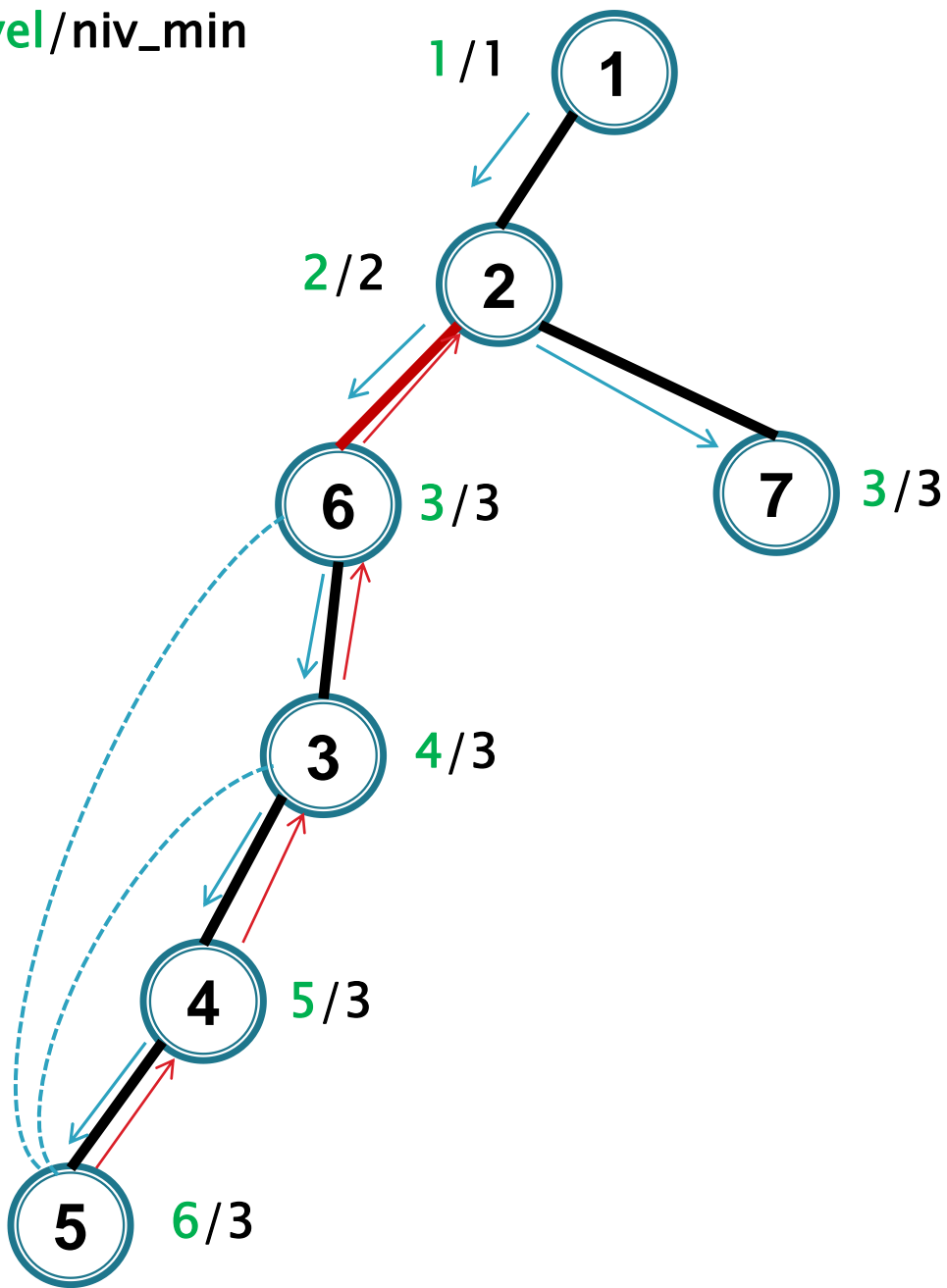


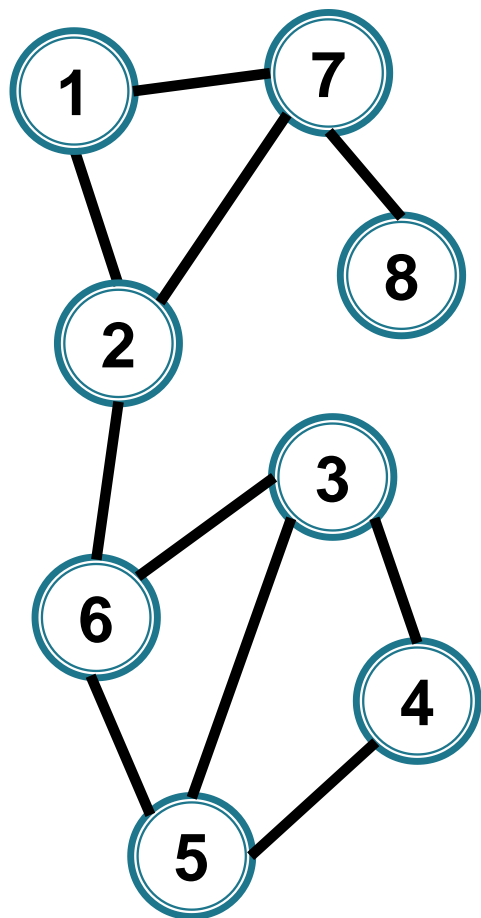
nivel/niv\_min



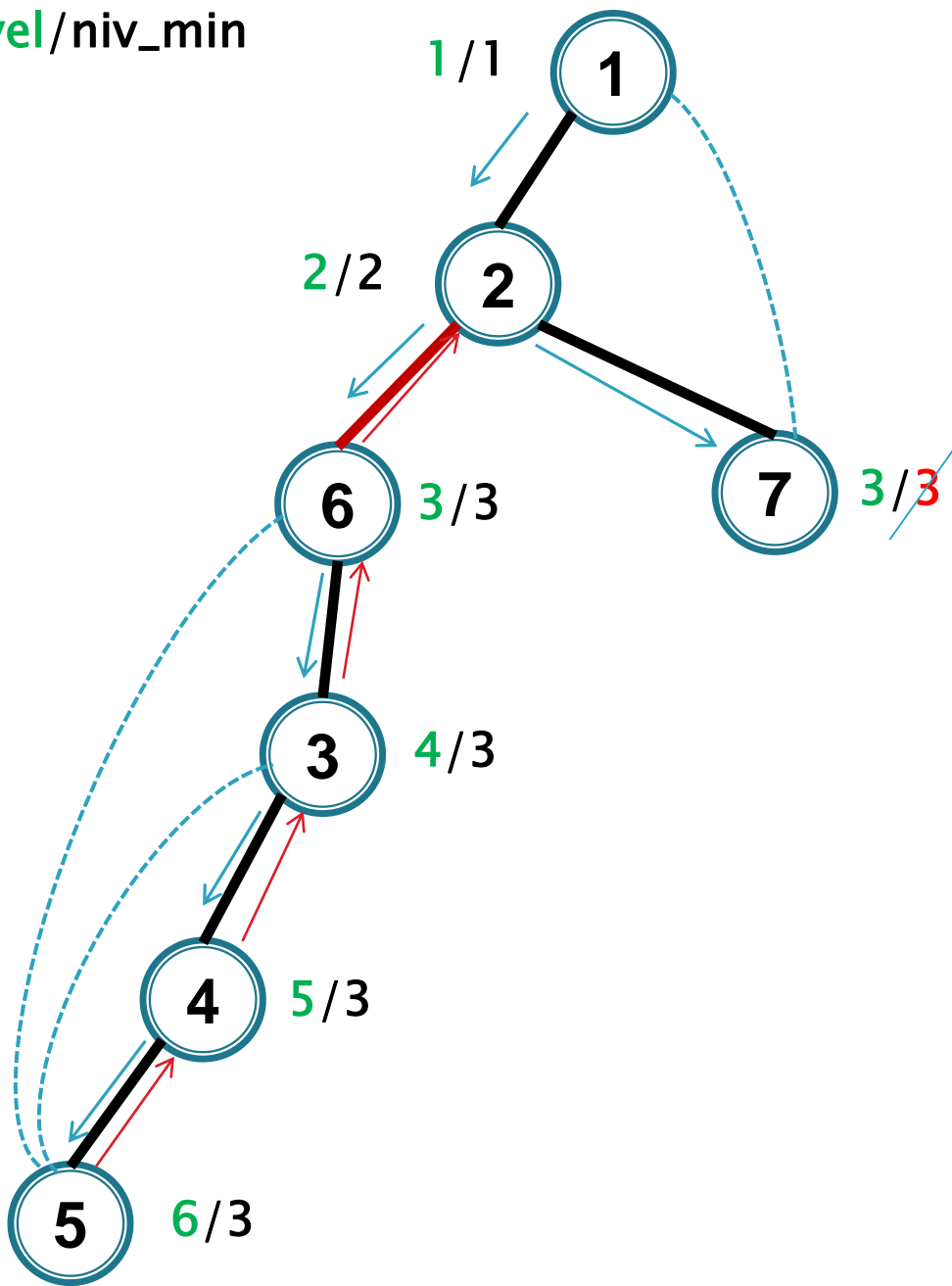


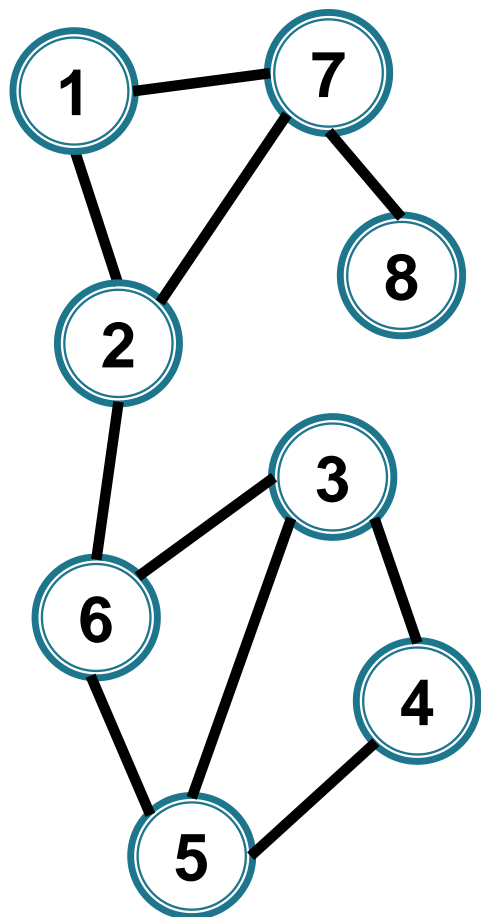
nivel/niv\_min



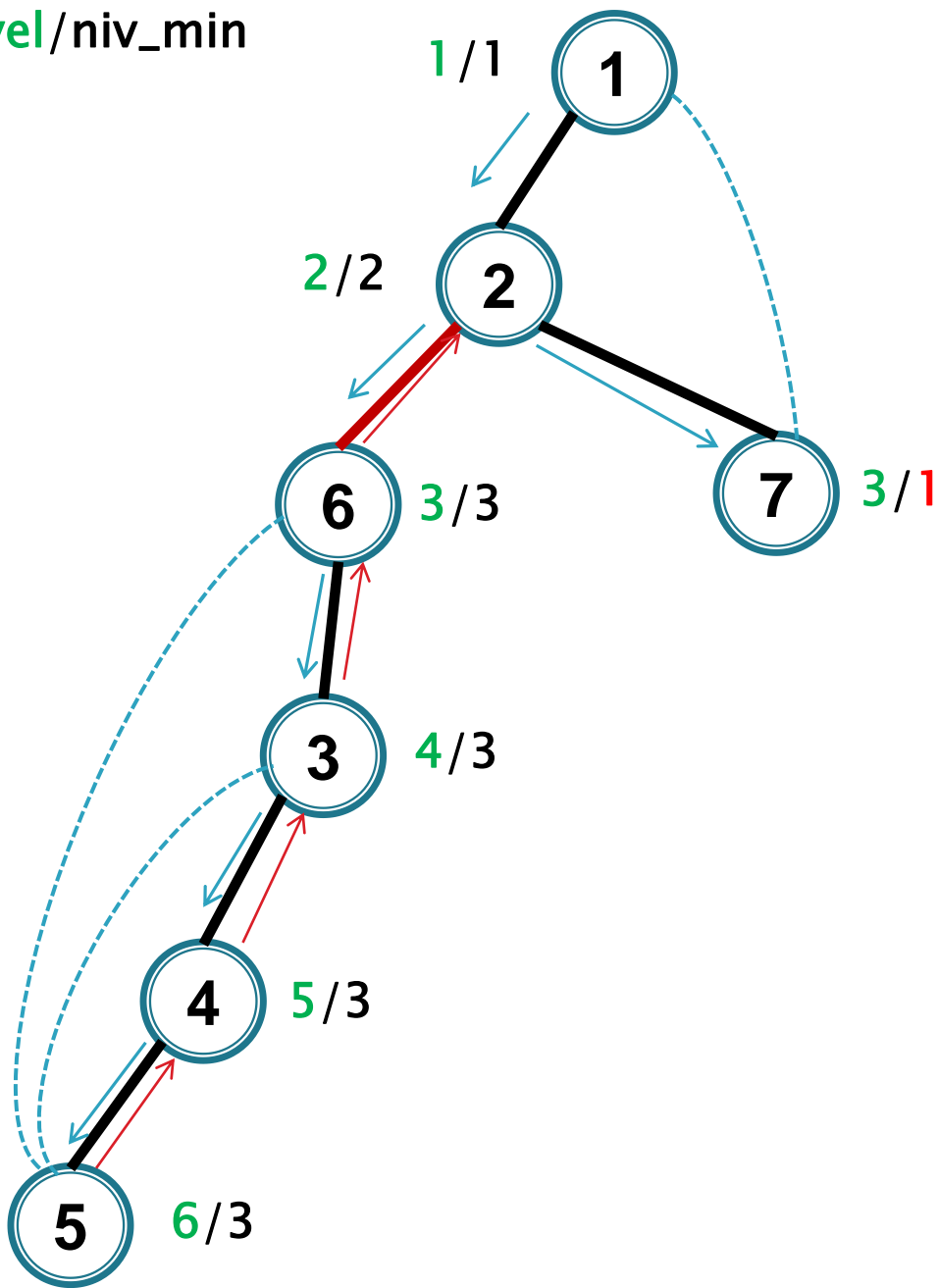


nivel/niv\_min

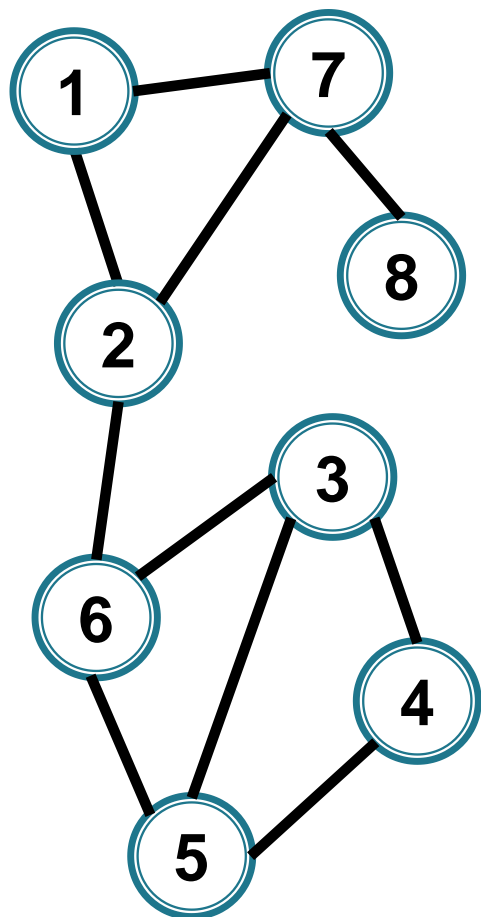




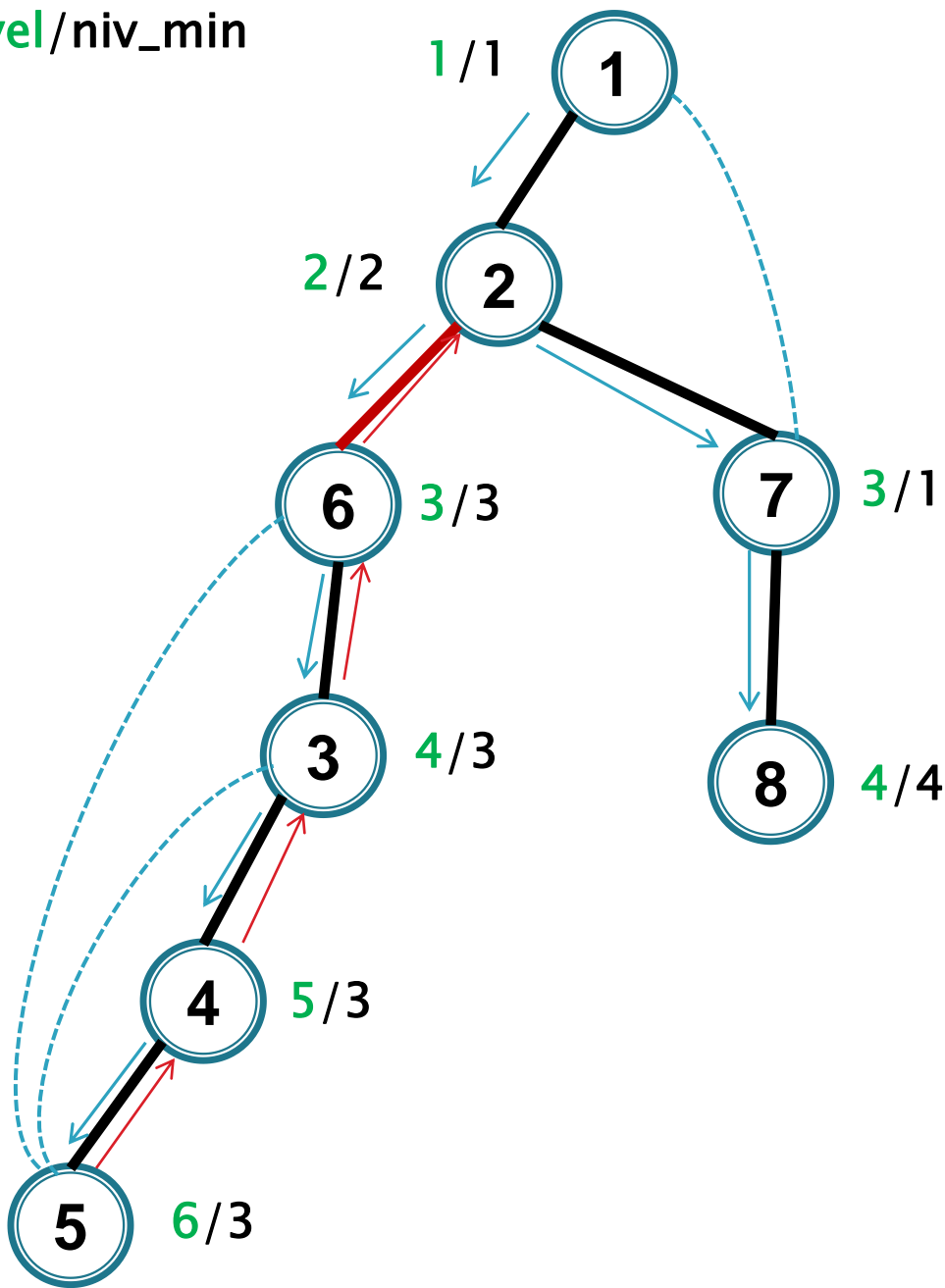
nivel/niv\_min

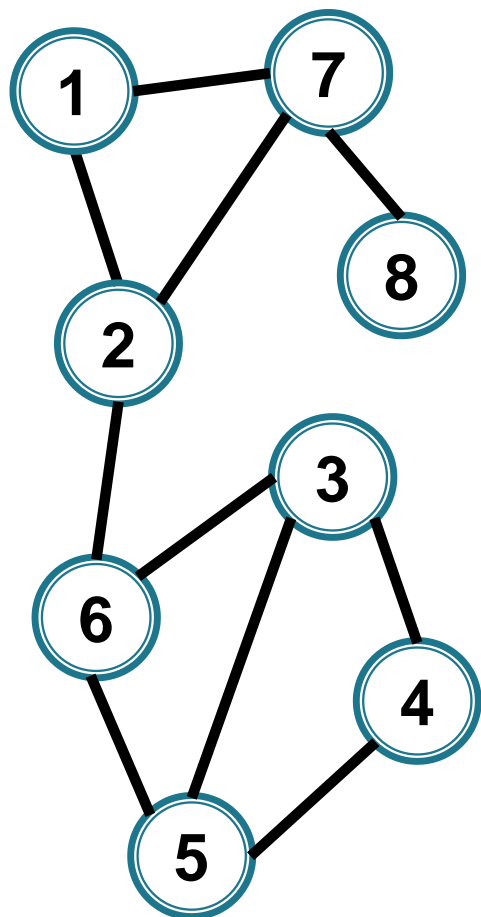




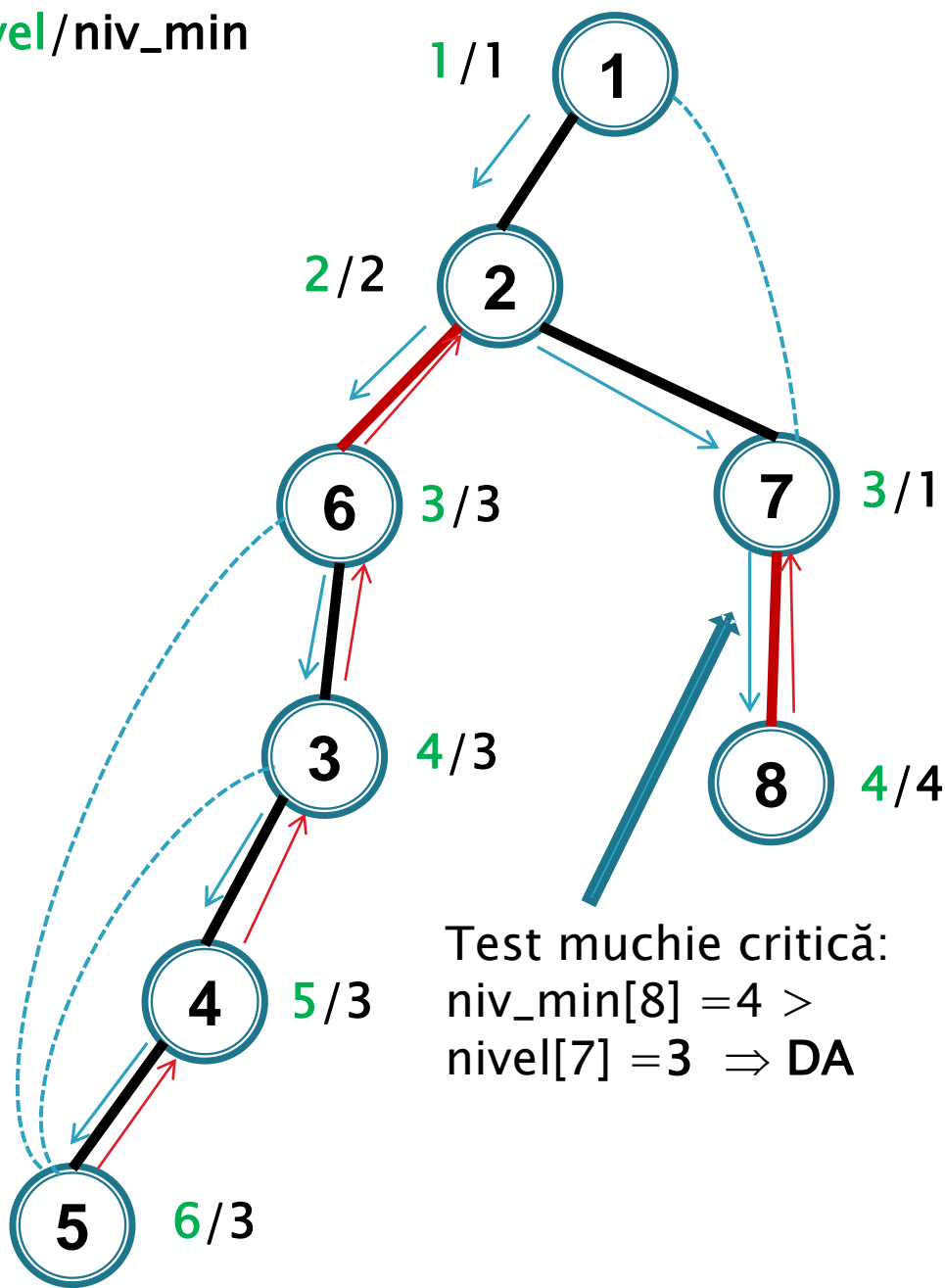


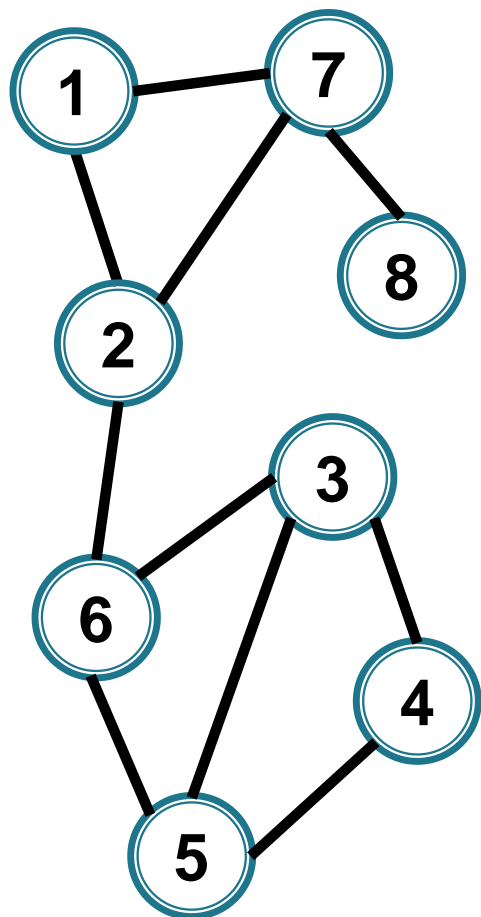
nivel/niv\_min



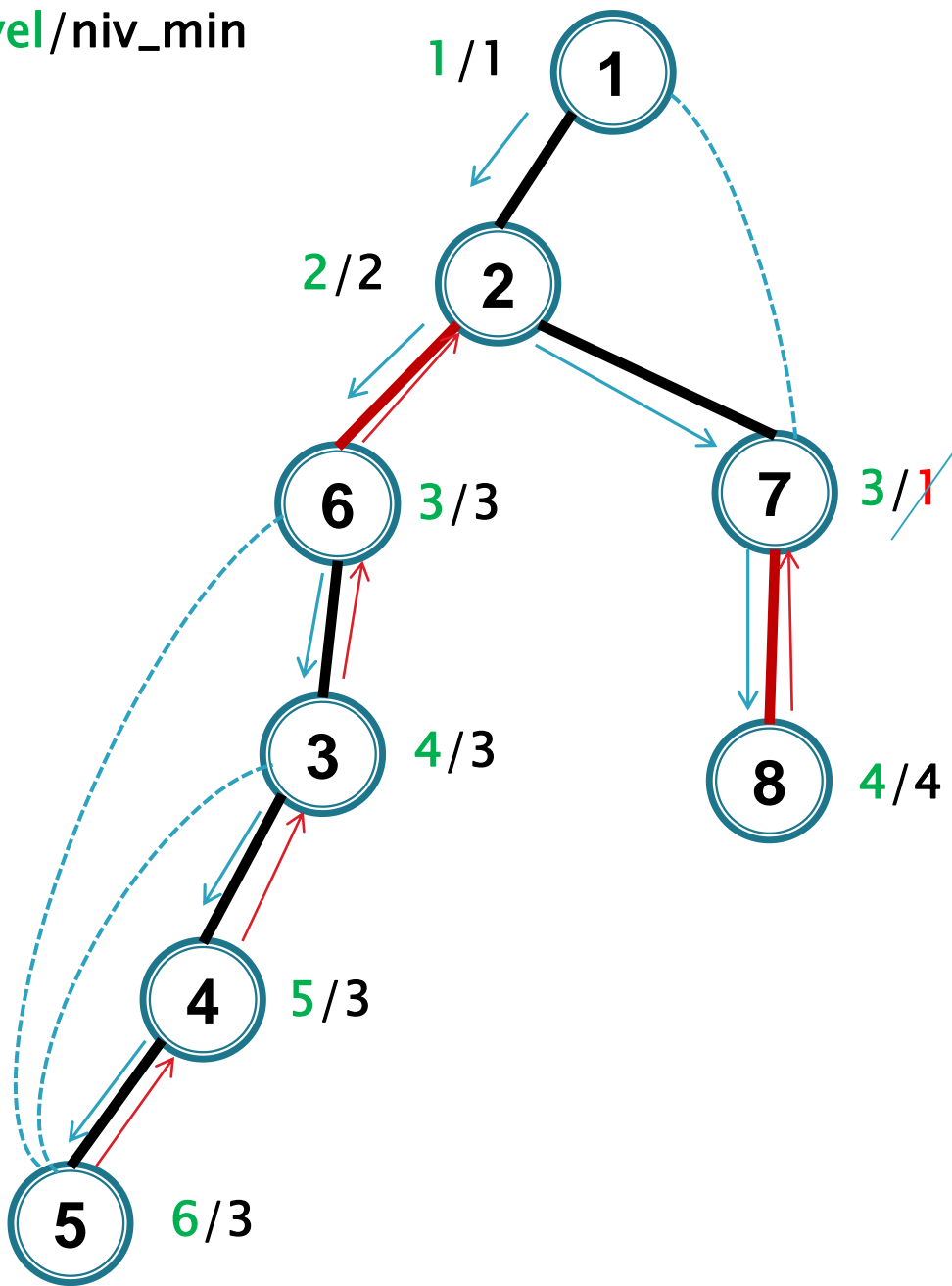


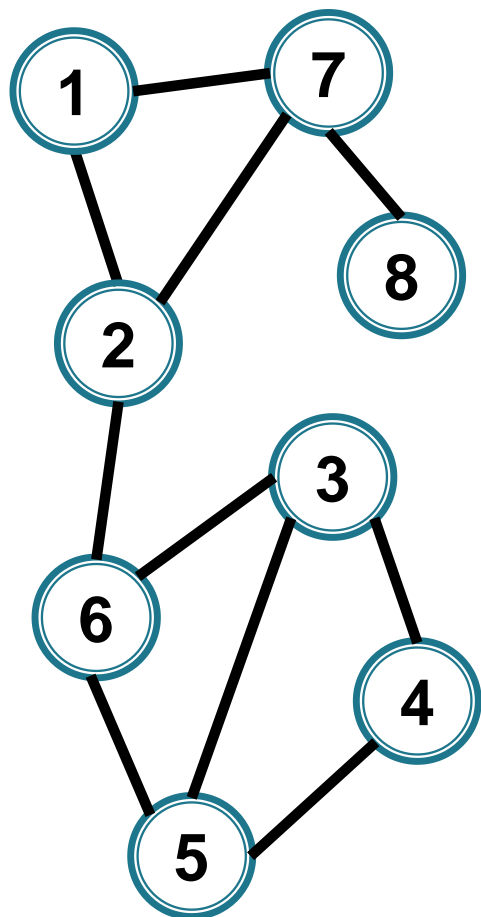
nivel/niv\_min



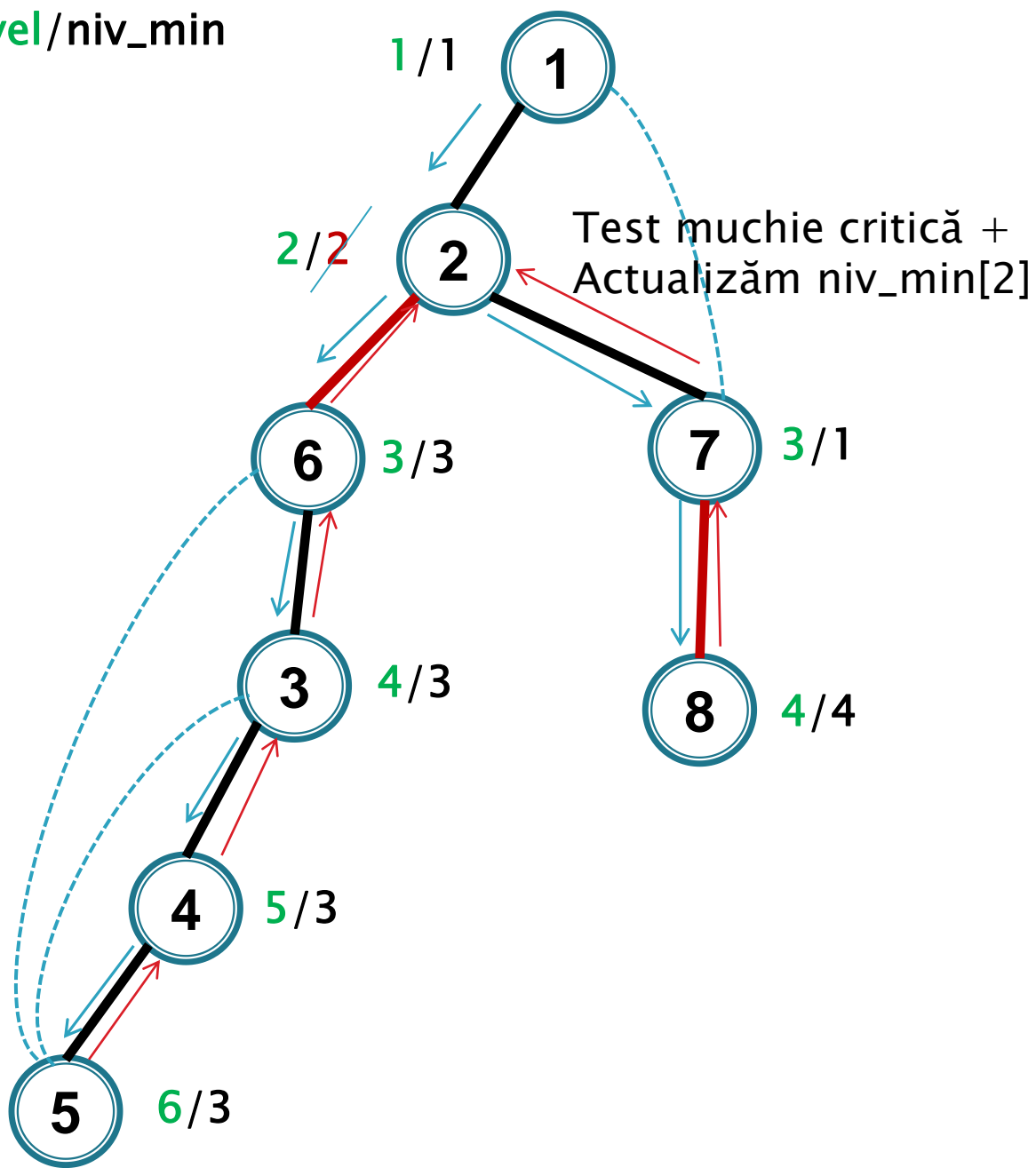


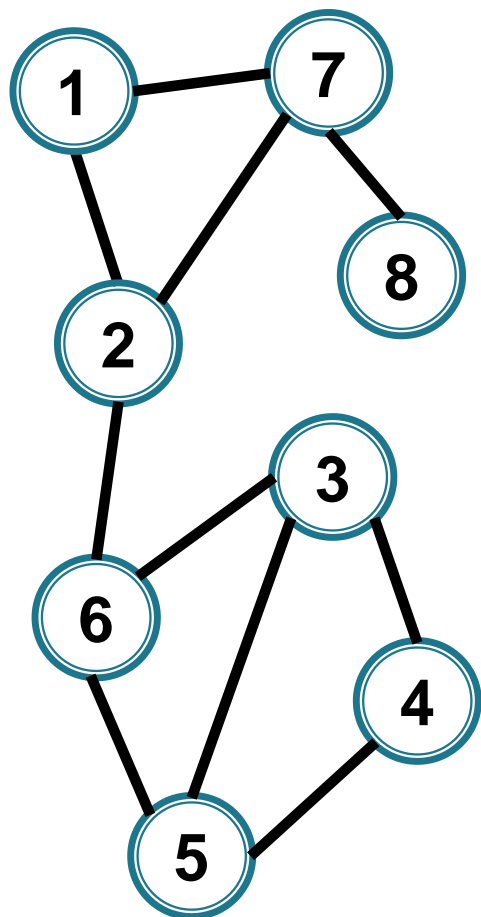
nivel/niv\_min



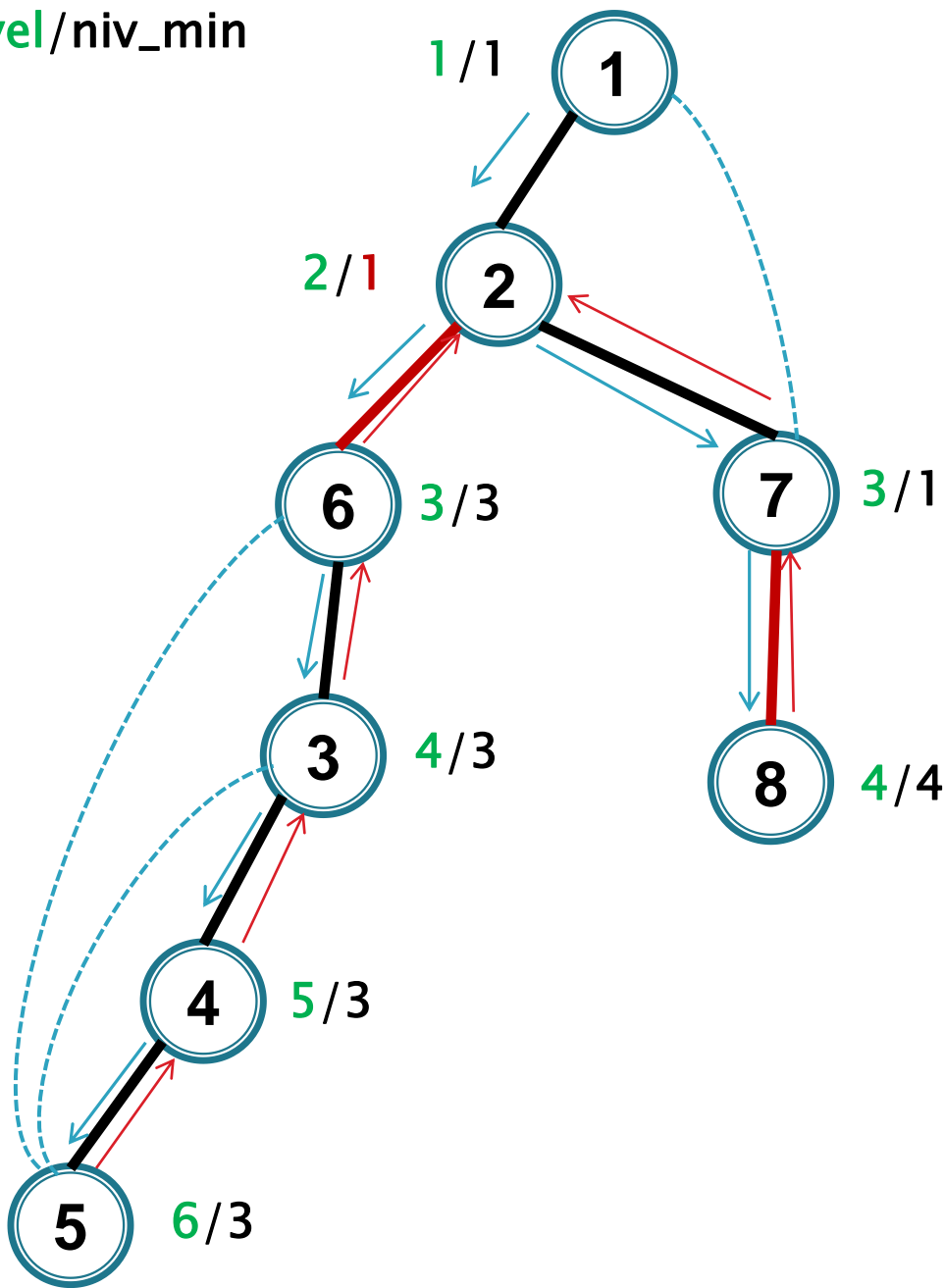


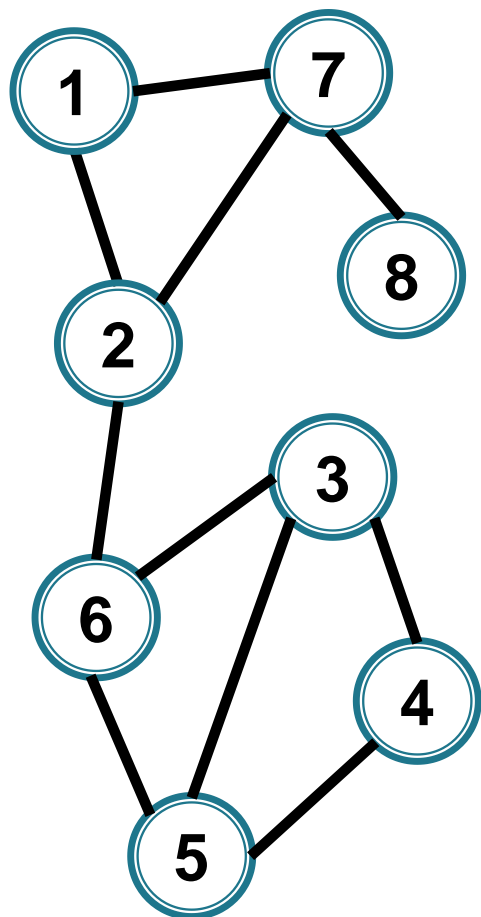
nivel/niv\_min





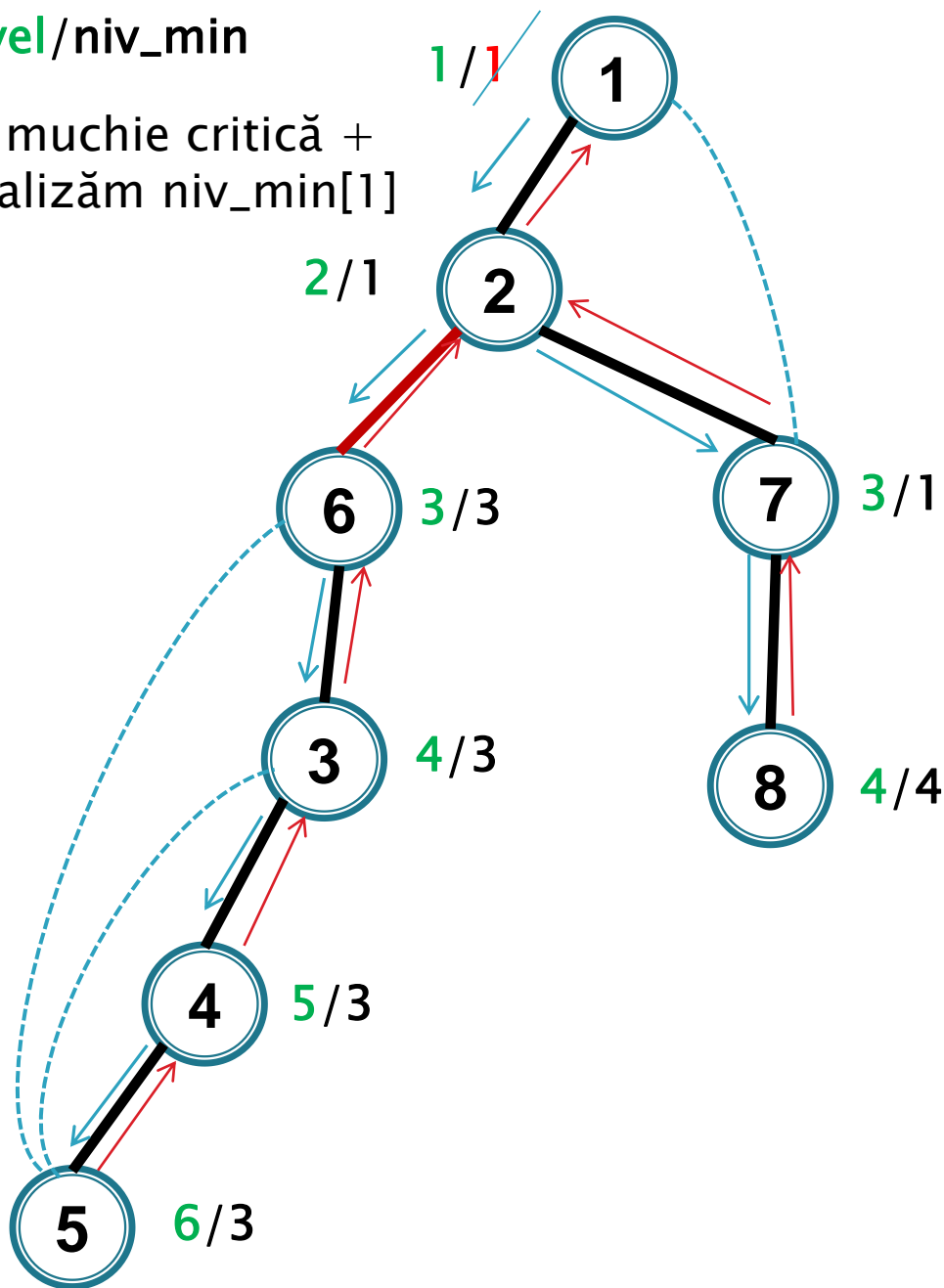
nivel/niv\_min

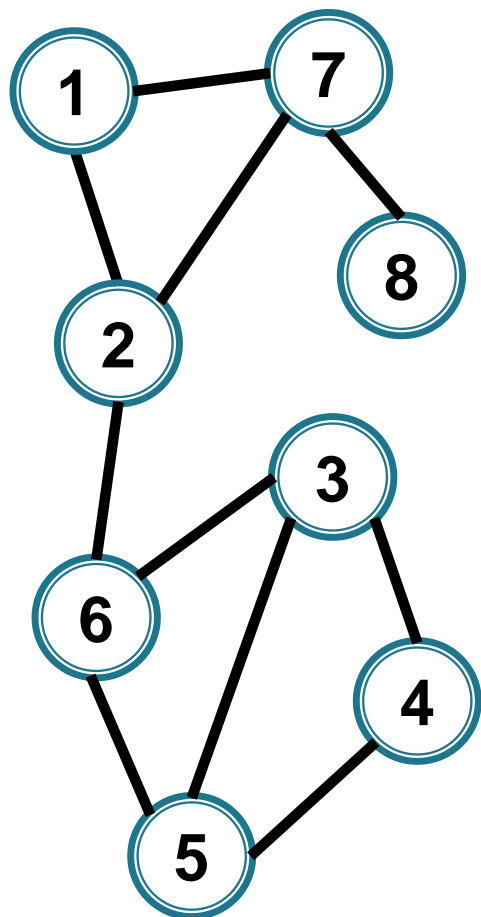




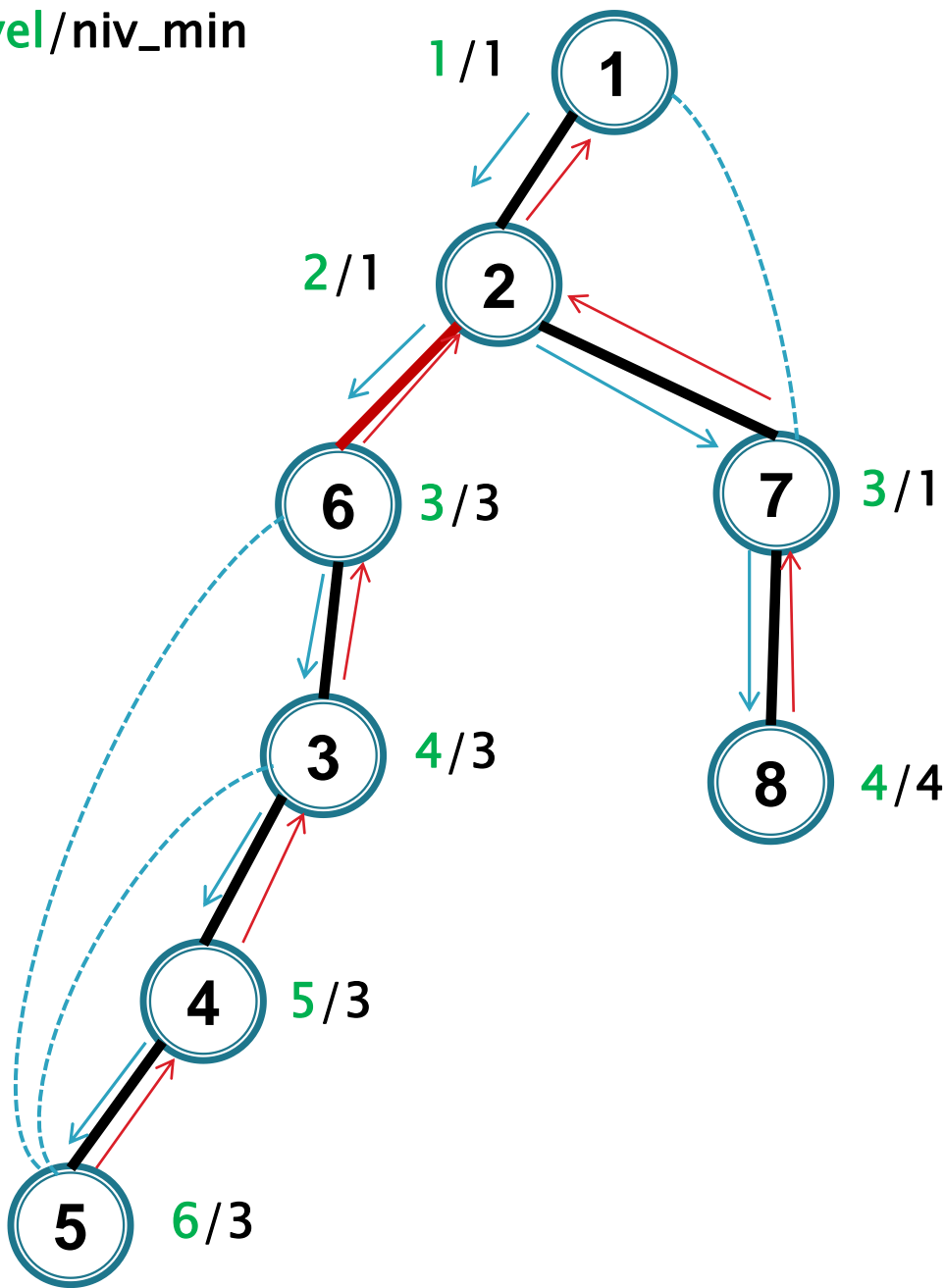
nivel/niv\_min

Test muchie critică +  
Actualizăm niv\_min[1]





nivel/niv\_min



# Indicații implementare

```
void df(int i){
    viz[i] = 1;
    niv_min[i] = nivel[i];
    for(j vecin al lui i)
        if(viz[j]==0){ //ij muchie de avansare
            nivel[j] = nivel[i]+1;
            df(j);

            //actualizare niv_min[i]- formula B
            ...

            //test ij este muchie critica
            ...
        }
    else
        if(nivel[j]<nivel[i]-1) //ij muchie de intoarcere
            //actualizare niv_min[i]- formula A
            ...
}
```



# Indicații implementare

```
void df(int i){
    viz[i] = 1;
    niv_min[i] = nivel[i];
    for(j vecin al lui i)
        if(viz[j]==0){ //ij muchie de avansare
            nivel[j] = nivel[i]+1;
            df(j);

            //actualizare niv_min[i]- formula B
            niv_min[i] = min{niv_min[i], niv_min[j] }

            //test ij este muchie critica
            ...
        }
    else
        if(nivel[j]<nivel[i]-1) //ij muchie de intoarcere
            //actualizare niv_min[i]- formula A
            ...
}
```

# Indicații implementare

```
void df(int i){
    viz[i] = 1;
    niv_min[i] = nivel[i];
    for(j vecin al lui i)
        if(viz[j]==0){ //ij muchie de avansare
            nivel[j] = nivel[i]+1;
            df(j);

            //actualizare niv_min[i]- formula B
            niv_min[i] = min{niv_min[i], niv_min[j] }

            //test ij este muchie critica
            if (niv_min[j]>nivel[i]) scrie muchia ij
        }
    else
        if(nivel[j]<nivel[i]-1) //ij muchie de intoarcere
            //actualizare niv_min[i]- formula A
            ...
}
```

# Indicații implementare

```
void df(int i){
    viz[i] = 1;
    niv_min[i] = nivel[i];
    for(j vecin al lui i)
        if(viz[j]==0){ //ij muchie de avansare
            nivel[j] = nivel[i]+1;
            df(j);

            //actualizare niv_min[i]- formula B
            niv_min[i] = min{niv_min[i], niv_min[j] }

            //test ij este muchie critica
            if (niv_min[j]>nivel[i]) scrie muchia ij
        }
    else
        if(nivel[j]<nivel[i]-1) //ij muchie de intoarcere
            //actualizare niv_min[i]- formula A
            niv_min[i] = min{niv_min[i], nivel[j] }
}
```

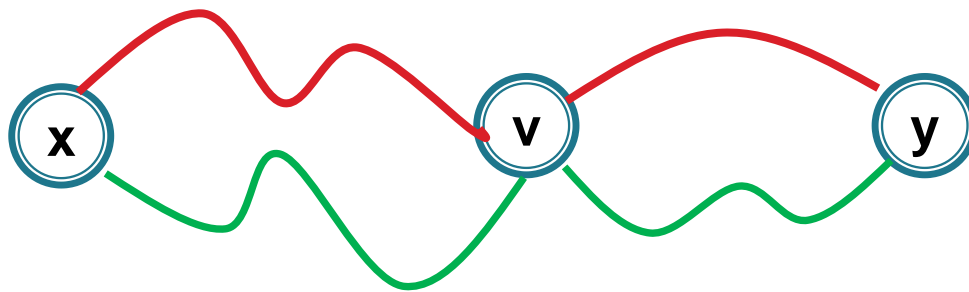
# Puncte critice

# Puncte critice

- ▶ Un vârf  $v$  este punct critic  $\Leftrightarrow$

există două vârfuri  $x, y \neq v$  astfel

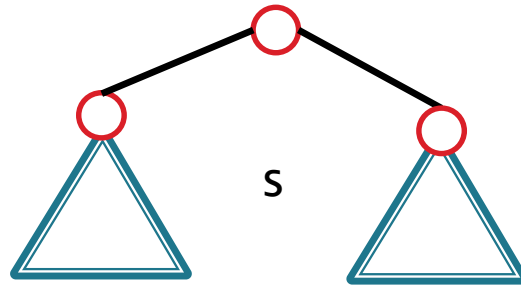
încât  $v$  aparține oricărui  $x, y$ -lanț



# Puncte critice

- ▶ Arborele DF

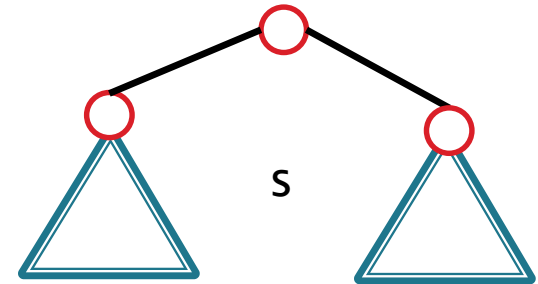
- rădăcina  $s$  este punct critic  $\Leftrightarrow$



# Puncte critice

## ▶ Arborele DF

- rădăcina s este punct critic  $\Leftrightarrow$

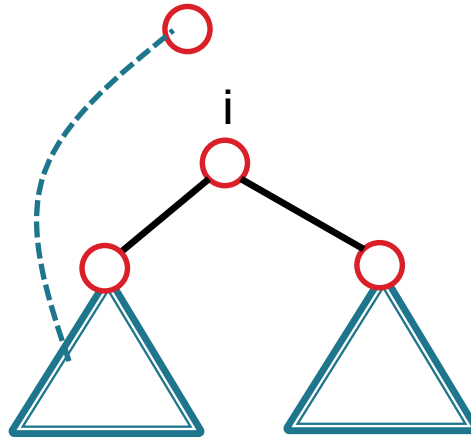


are cel puțin 2 fii în arborele DF

# Puncte critice

## ▶ Arborele DF

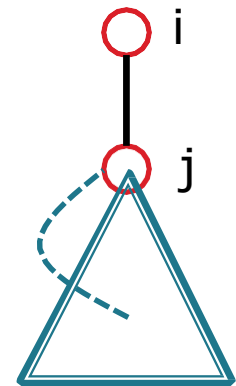
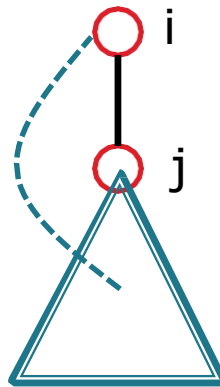
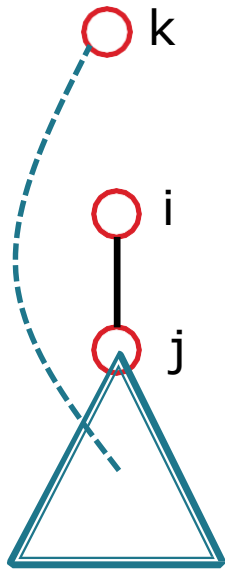
- un alt vârf  $i$  din arbore este critic  $\Leftrightarrow$





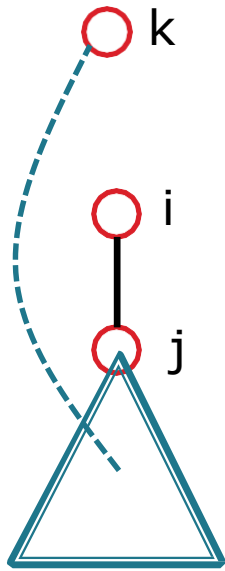
# Puncte critice

Pentru  $i \neq s$



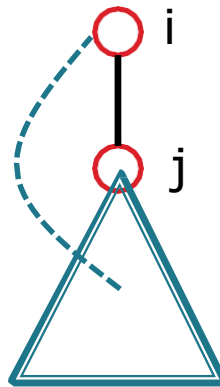
# Puncte critice

Pentru  $i \neq s$



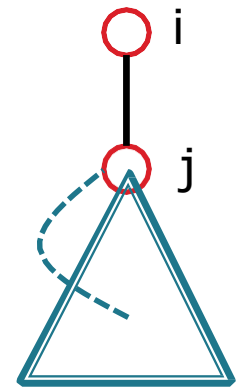
$i$  NU este critic

$niv\_min[j] < nivel[i]$



$i$  ESTE critic

$niv\_min[j] = nivel[i]$



$i$  ESTE critic

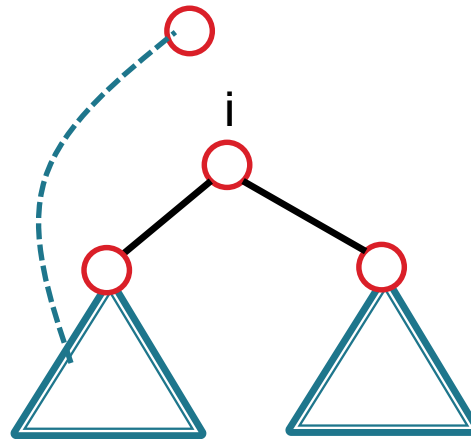
$niv\_min[j] > nivel[i]$

# Puncte critice

## ▶ Arborele DF

- un alt vârf  $i$  din arbore este critic  $\Leftrightarrow$

are cel puțin un fiu  $j$  cu  
 $niv\_min[j] \geq nivel[i]$



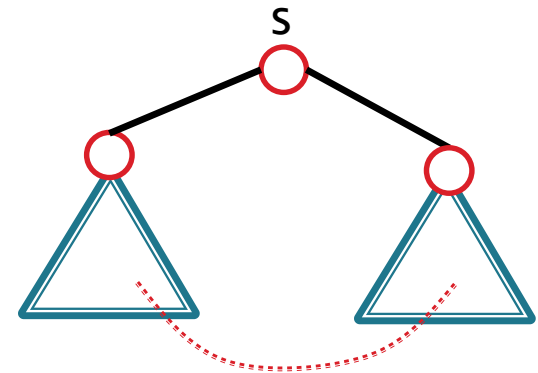
# Puncte critice

- ▶ Un vârf  $v$  este punct critic  $\Leftrightarrow$  există două vârfuri  $x, y \neq v$  astfel încât  $v$  aparține oricărui  $x, y$ -lanț

- ▶ Arborele DF

- rădăcina  $s$  este punct critic  $\Leftrightarrow$

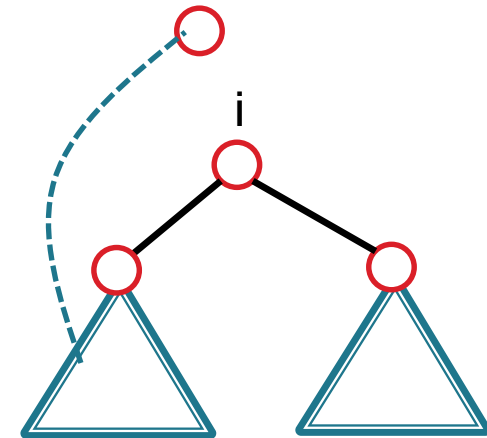
are cel puțin 2 fii în arborele DF



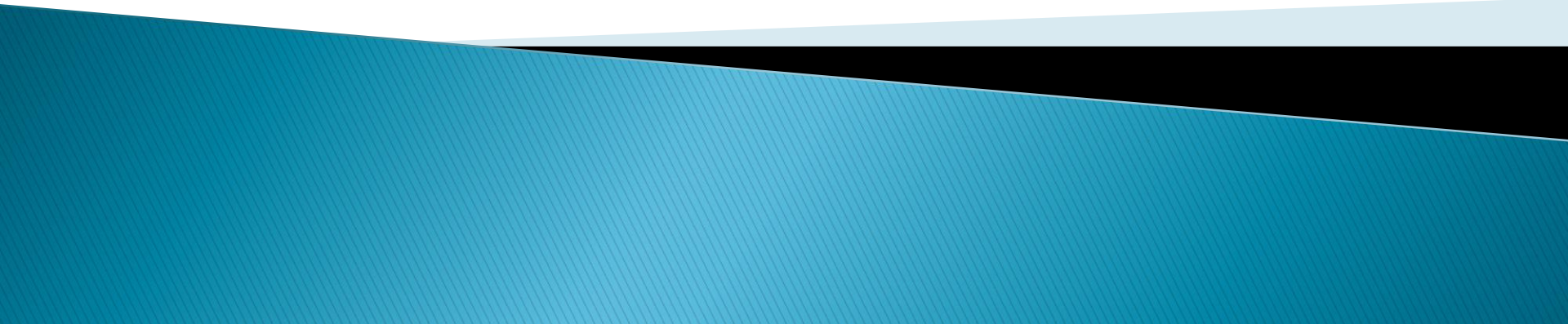
nu există muchii între subarbori  
(de traversare)

- un alt vârf  $i$  din arbore este critic  $\Leftrightarrow$

are cel puțin un fiu  $j$  cu  
 $niv\_min[j] \geq nivel[i]$

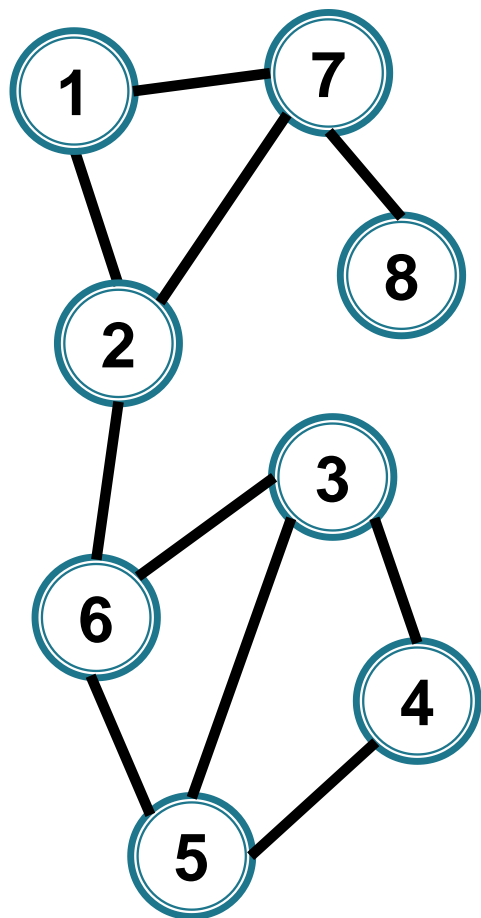


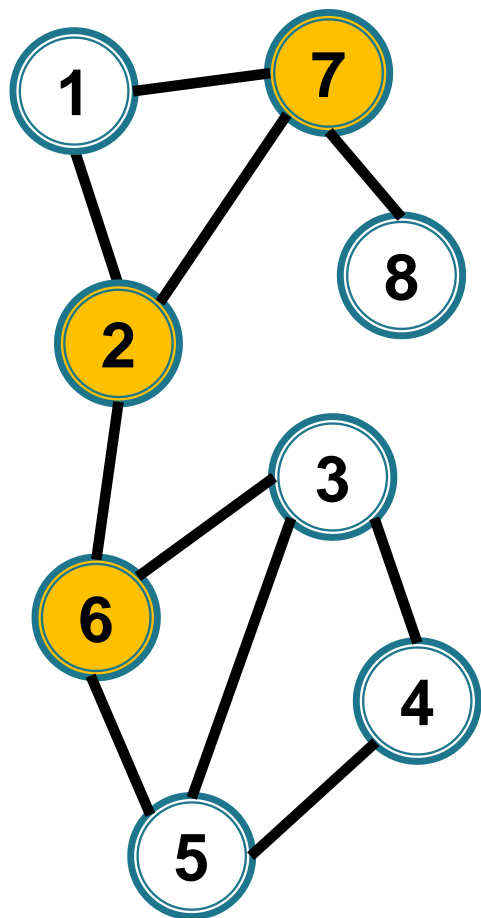
# Componente biconexe



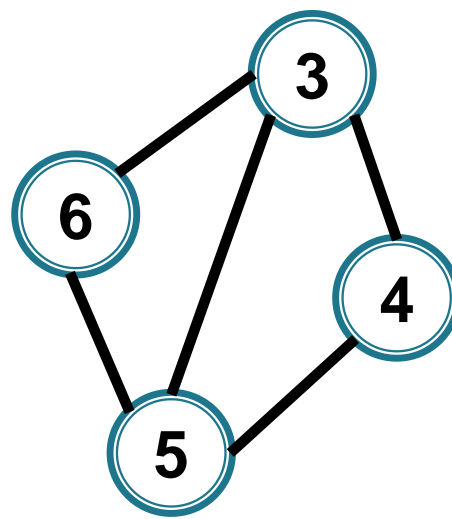
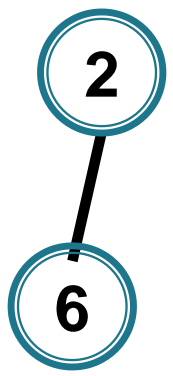
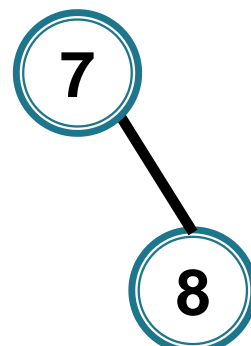
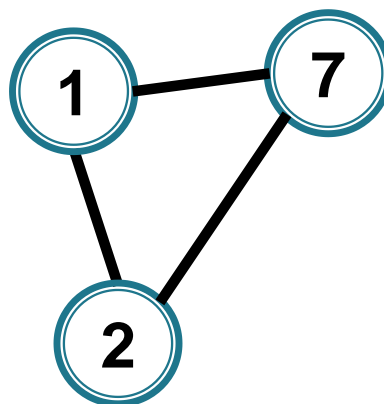
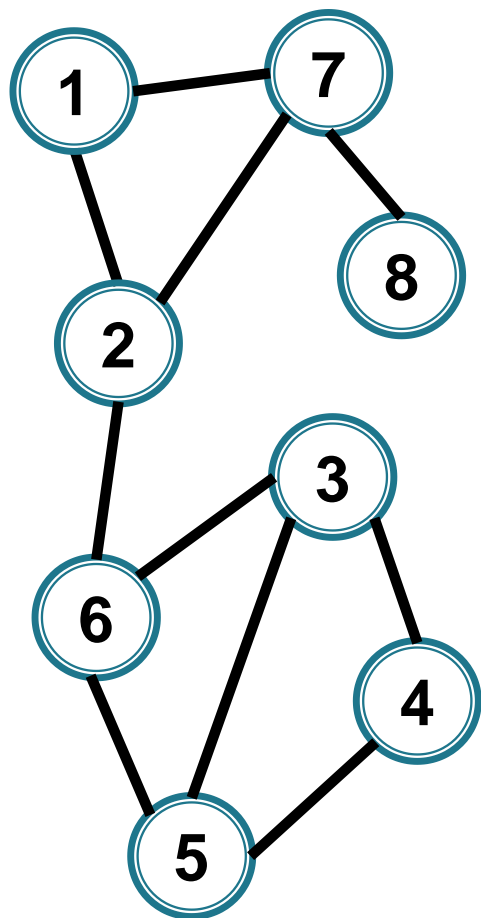
# Componente biconexe

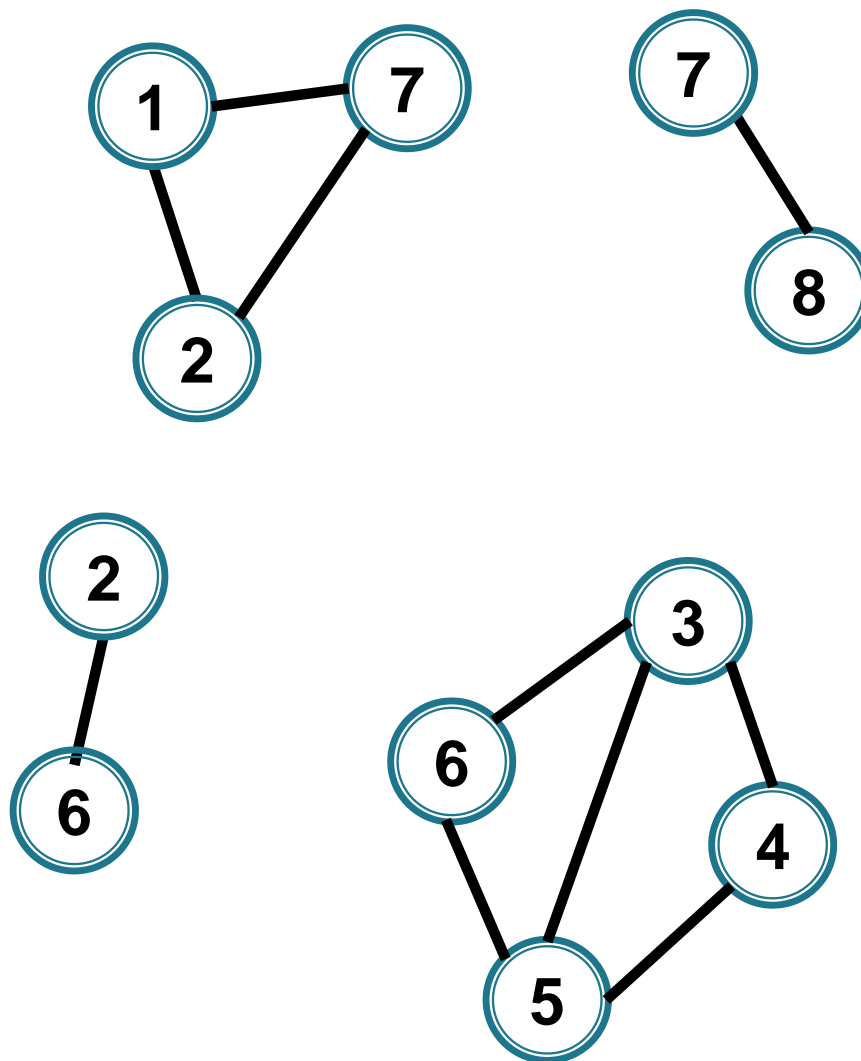
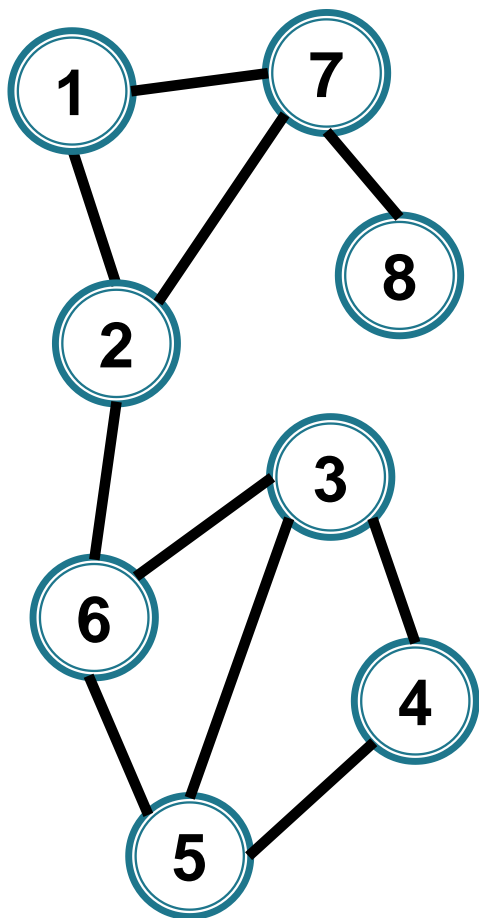
- ▶  $G$  – graf neorientat
- ▶  $G=(V,M)$  biconex = nu are puncte de articulație.
- ▶ Componentă biconexă (bloc) a lui  $G$  = subgraf biconex maximal.









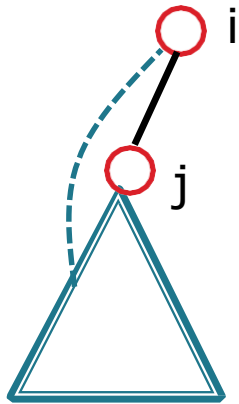


Muchie-disjuncte, nu vârf-disjuncte

# Componente biconexe

## ► Algoritm

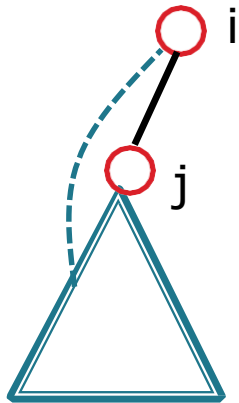
Intuitiv: când un fiu  $j$  semnalează că vârful curent  $i$  este critic, se poate afișa componenta care conține  $ij$  (care “se rupe” din vârful  $j$ )



# Componente biconexe

## ► Algoritm

Intuitiv: când un fiu  $j$  semnalează că vârful curent  $i$  este critic, se poate afișa componenta care conține  $ij$  (care “se rupe” din vârful  $j$ )

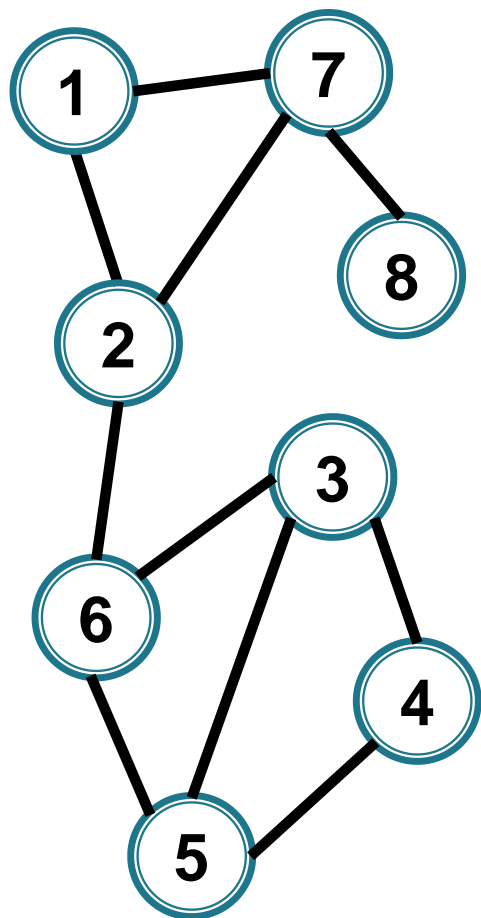


Memorăm muchiile într-o **stivă**

Când detectăm o componentă biconexă –  
scoatem muchiile din stivă până la muchia  $ij$   
 $\Rightarrow$  acestea formează o componentă

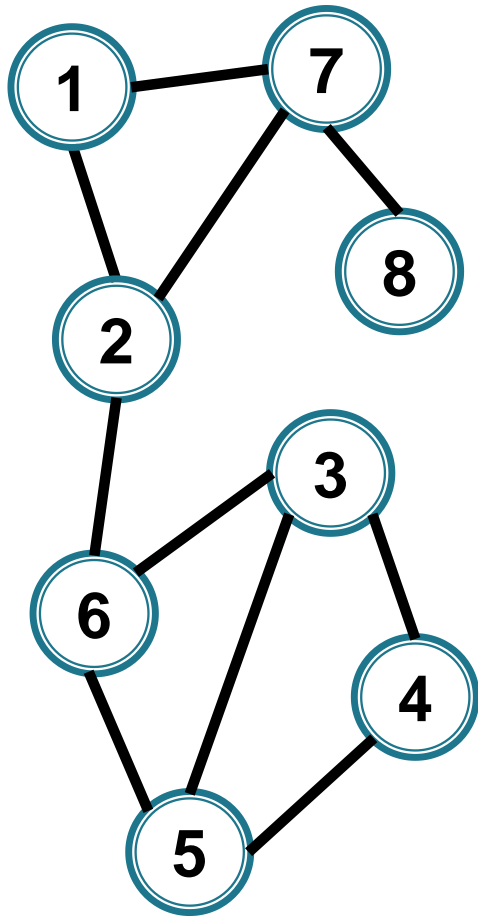
# Indicații implementare

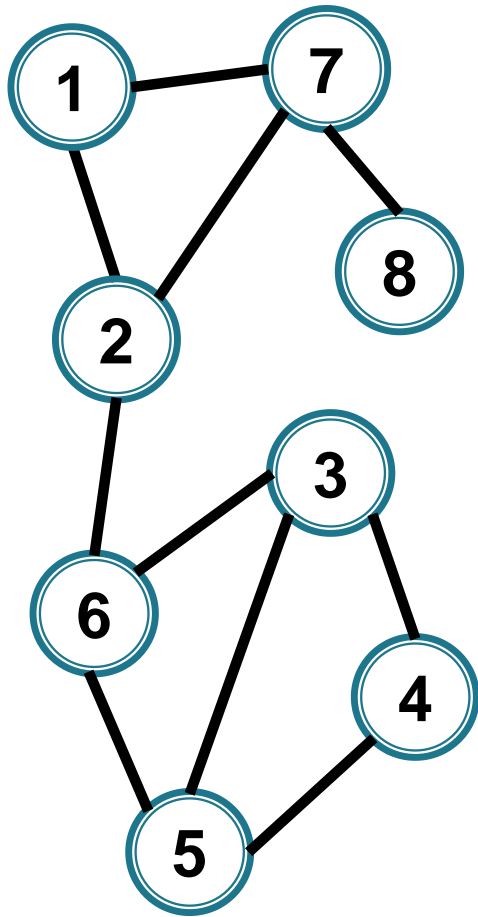
```
void df(int i){
    viz[i] = 1;
    niv_min[i] = nivel[i];
    for(j vecin al lui i)
        if(viz[j]==0){ //ij muchie de avansare
            nivel[j] = nivel[i]+1;
            adauga(S,ij)
            df(j);
            niv_min[i] = min{niv_min[i], niv_min[j] }
            if (niv_min[j] >= nivel[i])
                elimina din S toate muchiile pana la ij
        }
    else
        if(nivel[j]<nivel[i]-1) //ij muchie de intoarcere
            //actualizare niv_min[i]- formula A
            niv_min[i] = min{niv_min[i], nivel[j]}
            adauga(S,ij)
        }
}
```



nivel/niv\_min

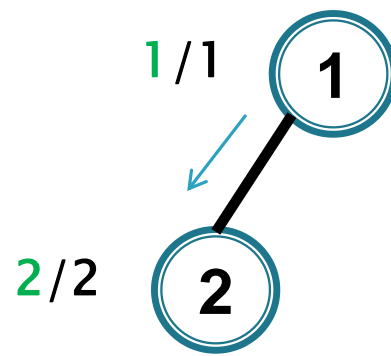
1/1



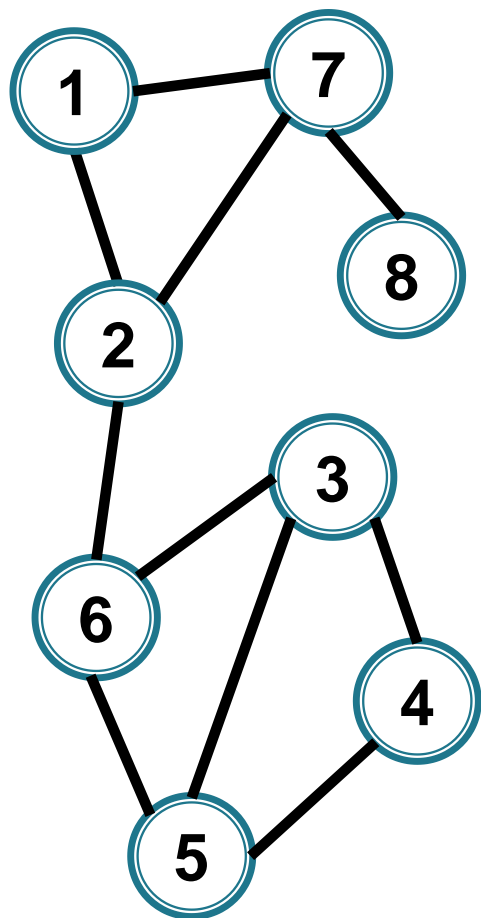


S:  
1 2

nivel/niv\_min

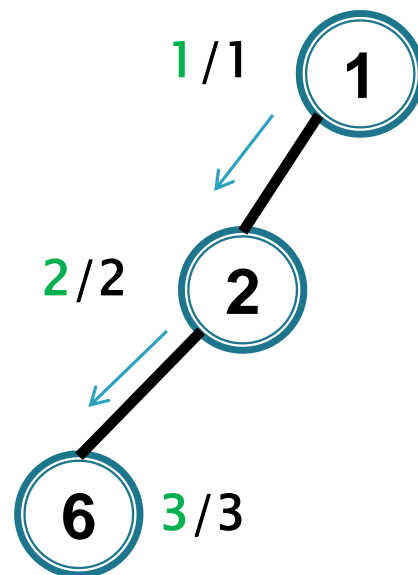


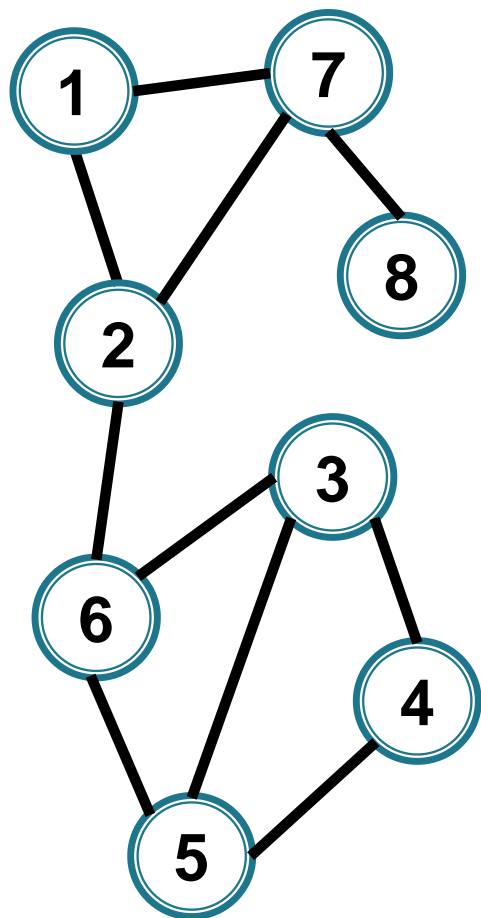




S:  
1 2  
2 6

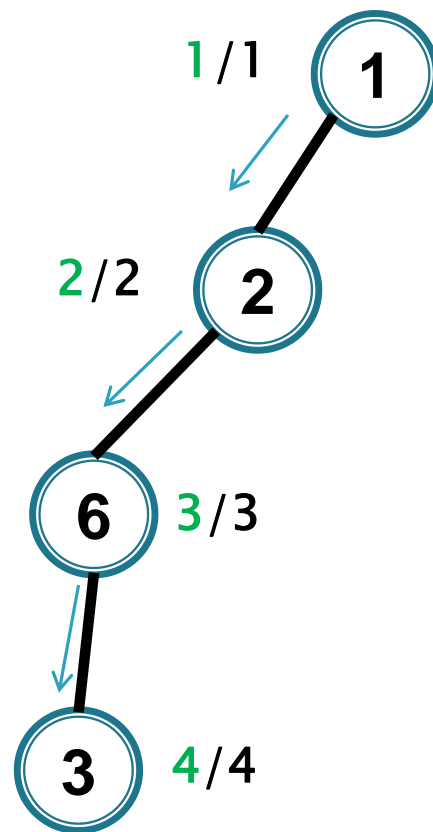
nivel/niv\_min

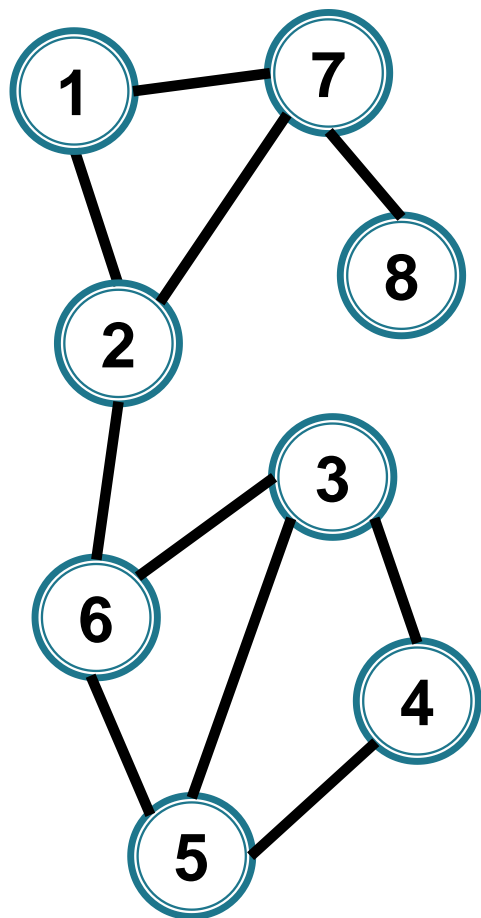




S:  
 1 2  
 2 6  
 6 3

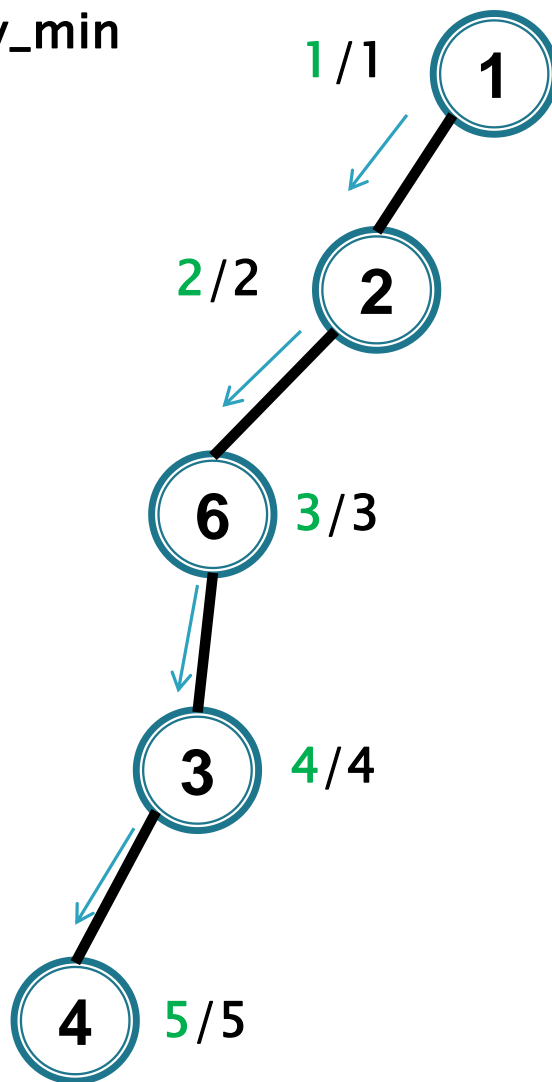
nivel/niv\_min

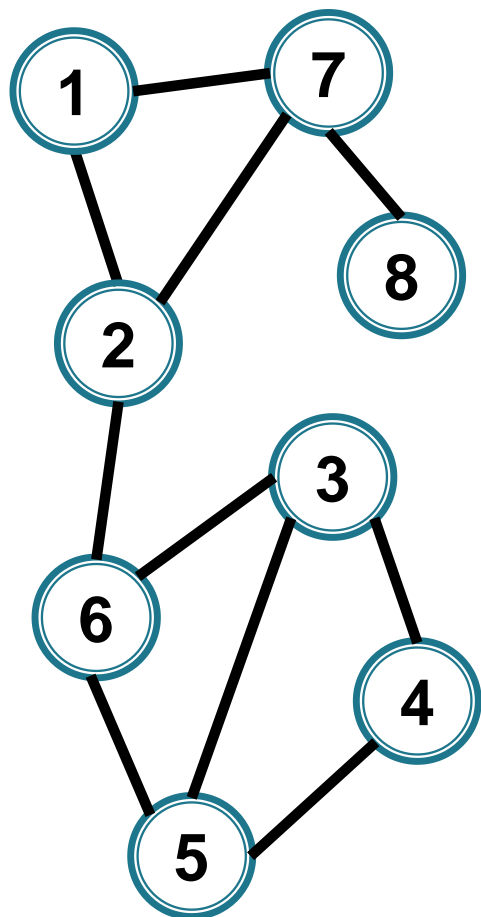




S:  
 1 2  
 2 6  
 6 3  
 3 4

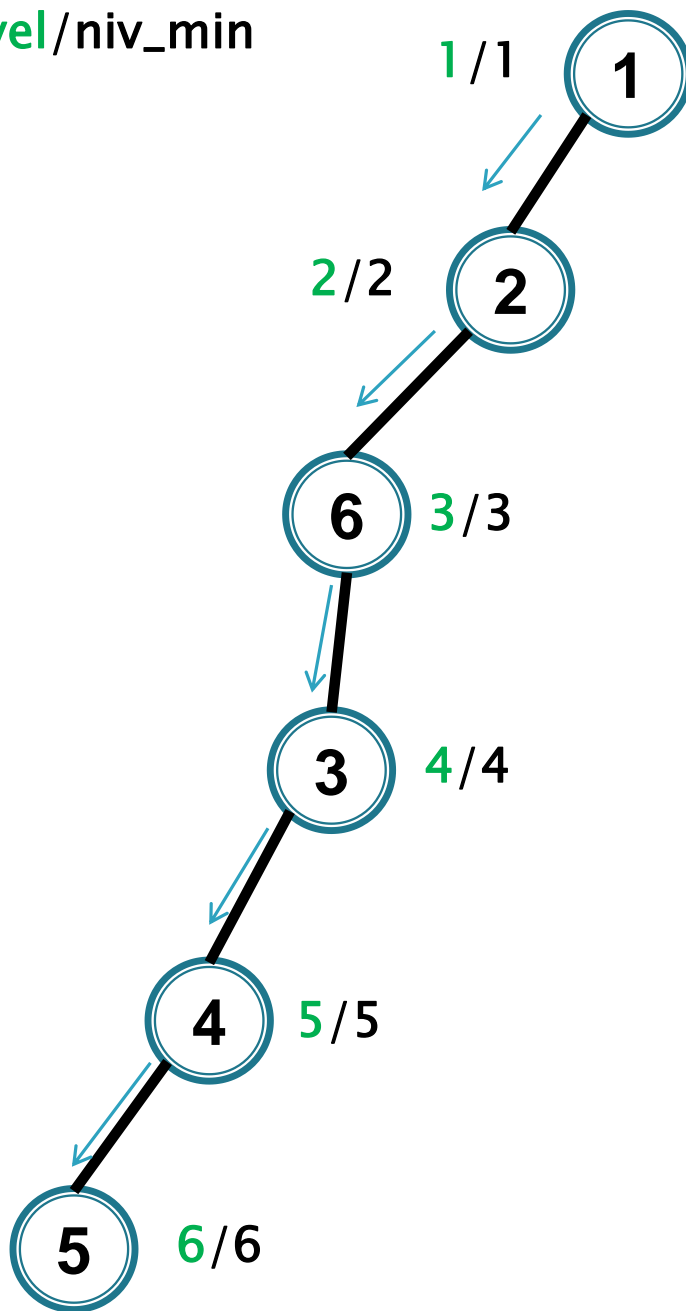
nivel/niv\_min

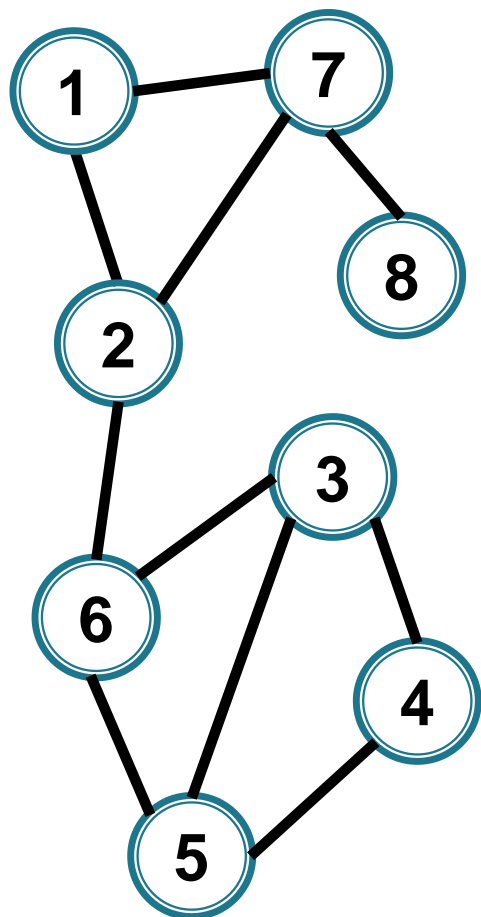




S:  
 1 2  
 2 6  
 6 3  
 3 4  
 4 5

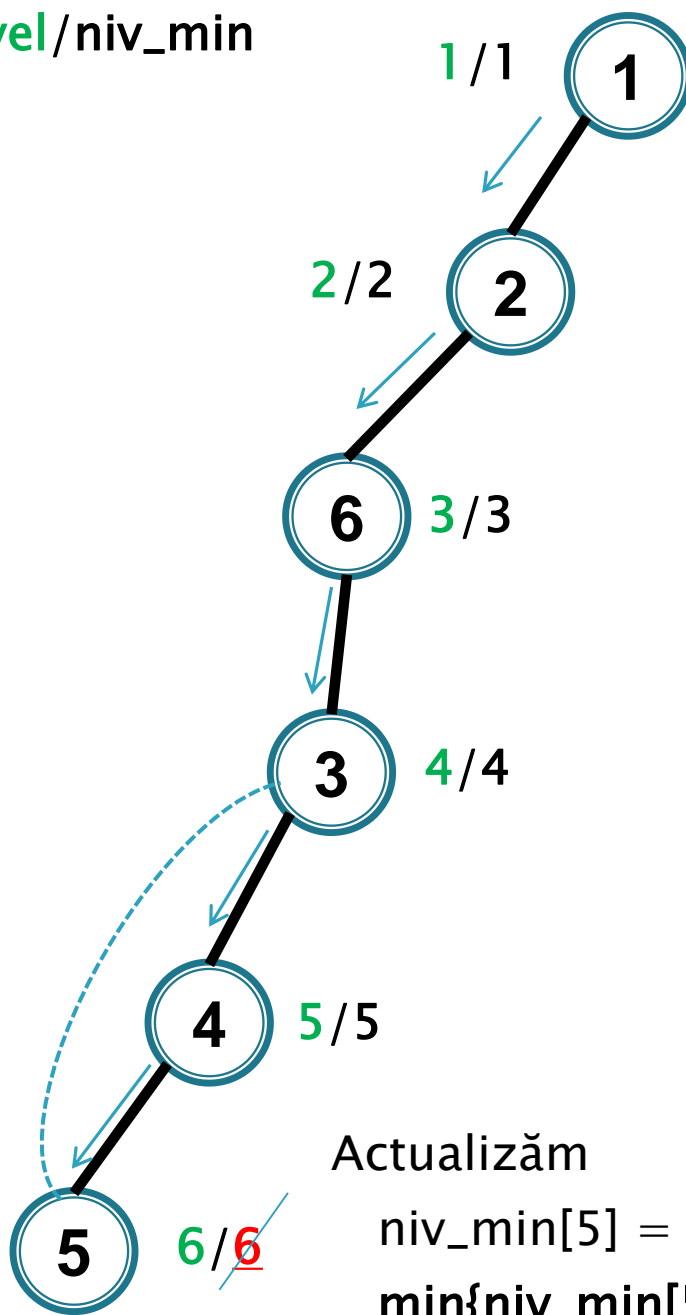
nivel/niv\_min





S:  
 1 2  
 2 6  
 6 3  
 3 4  
 4 5  
 5 3

nivel/niv\_min

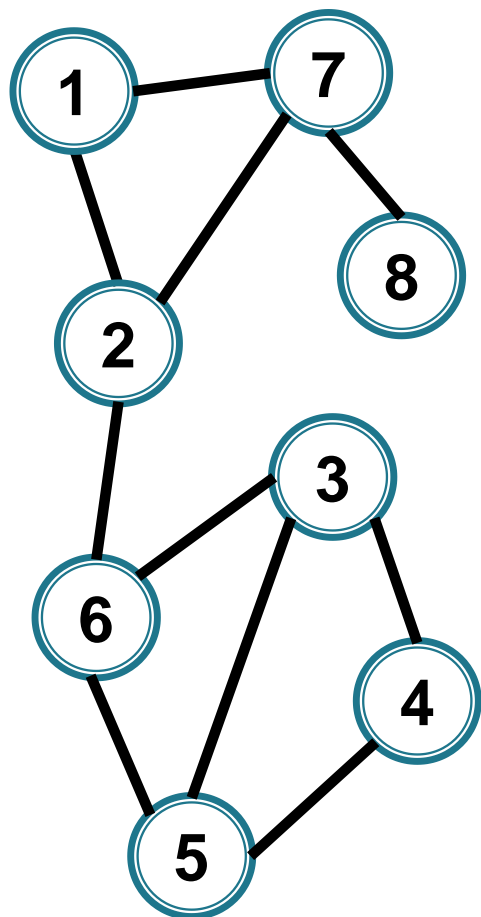


Actualizăm

$\text{niv\_min}[5] =$

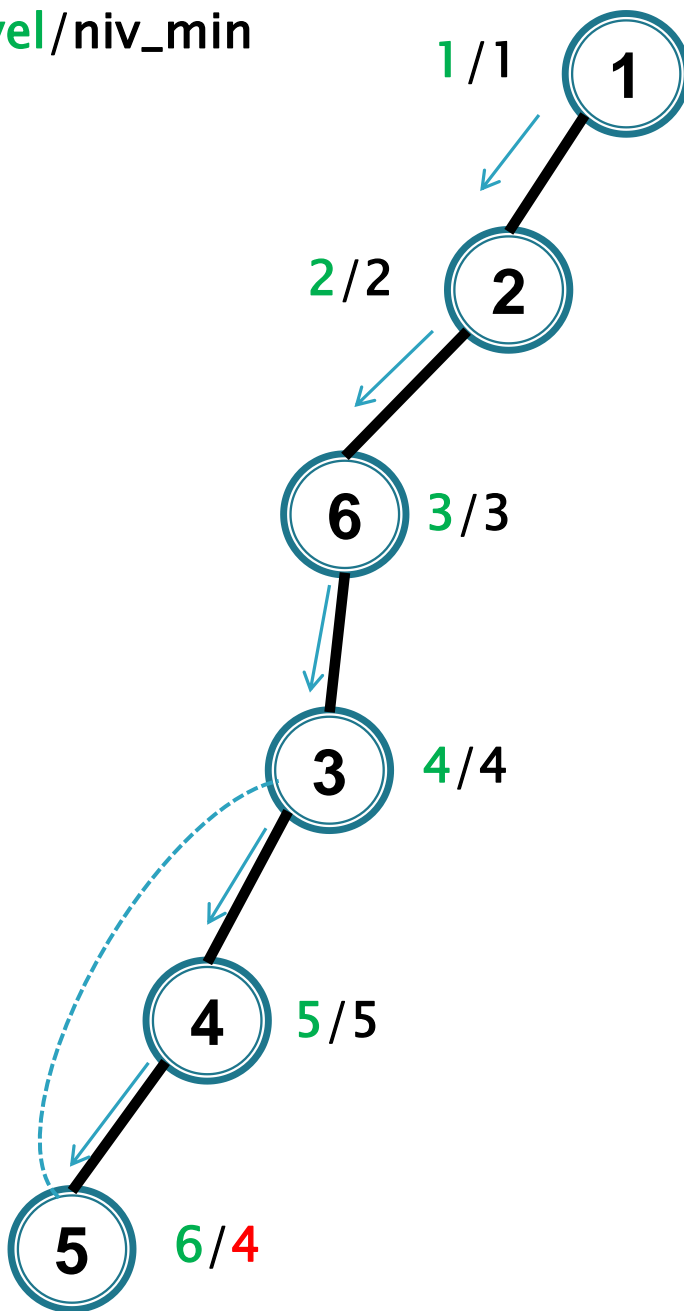
$\min\{\text{niv\_min}[5], \text{nivel}[3]\}$

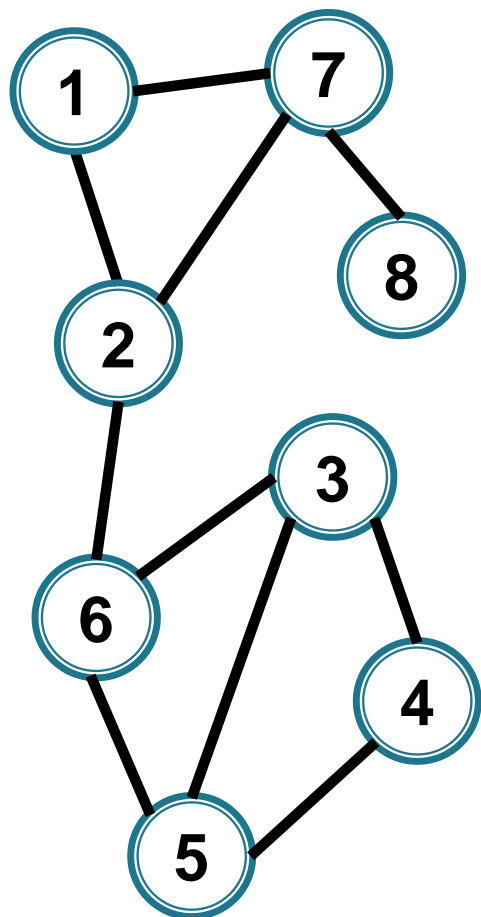
(cazul A)



S:  
 1 2  
 2 6  
 6 3  
 3 4  
 4 5  
 5 3

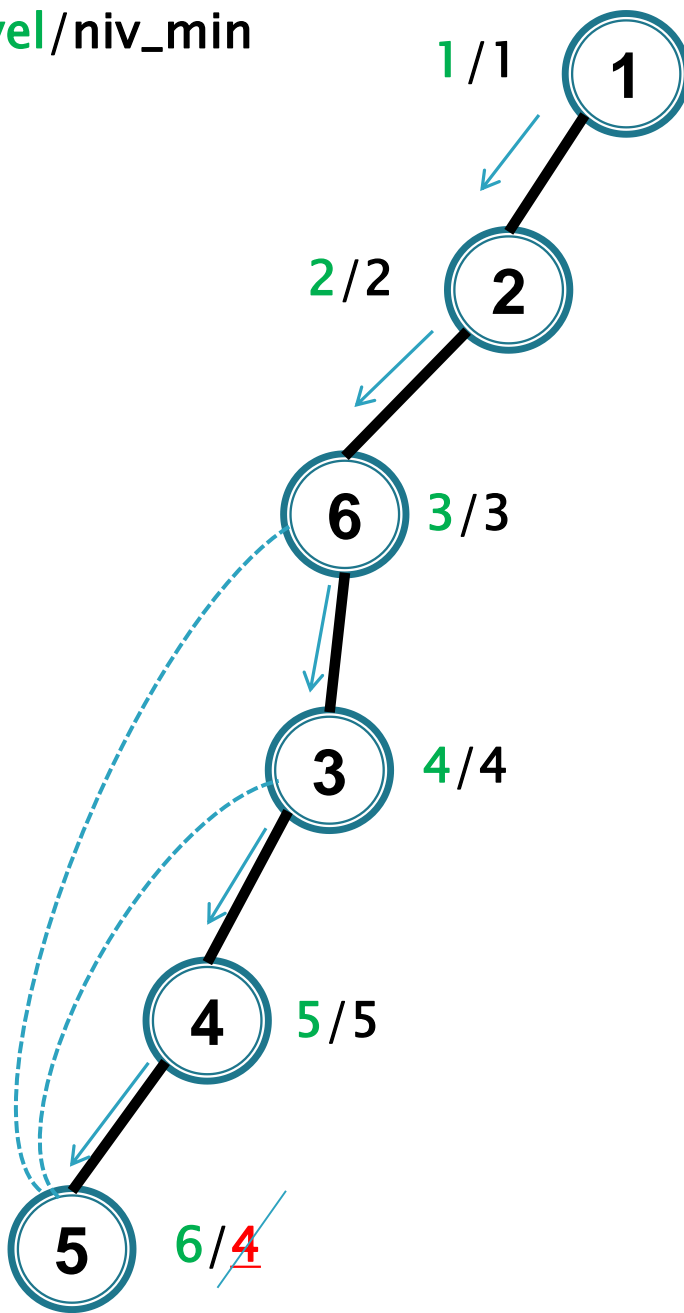
nivel/niv\_min

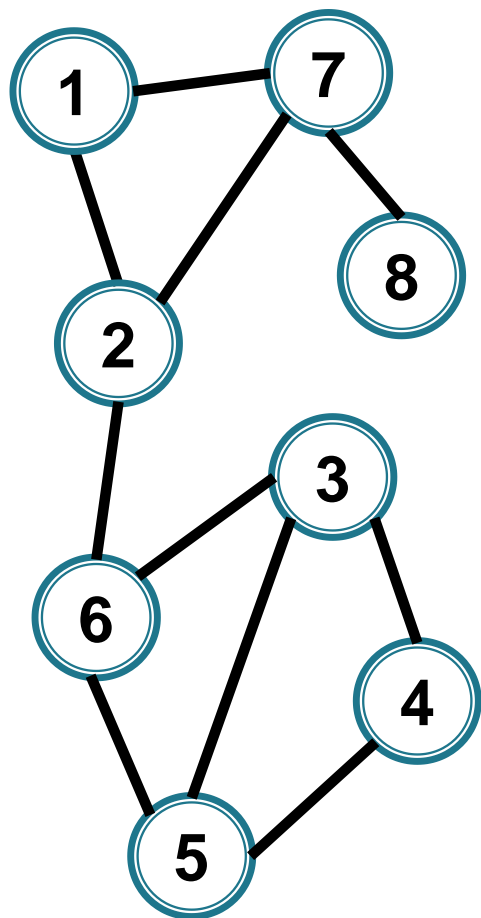




S:  
 1 2  
 2 6  
 6 3  
 3 4  
 4 5  
 5 3  
 5 6

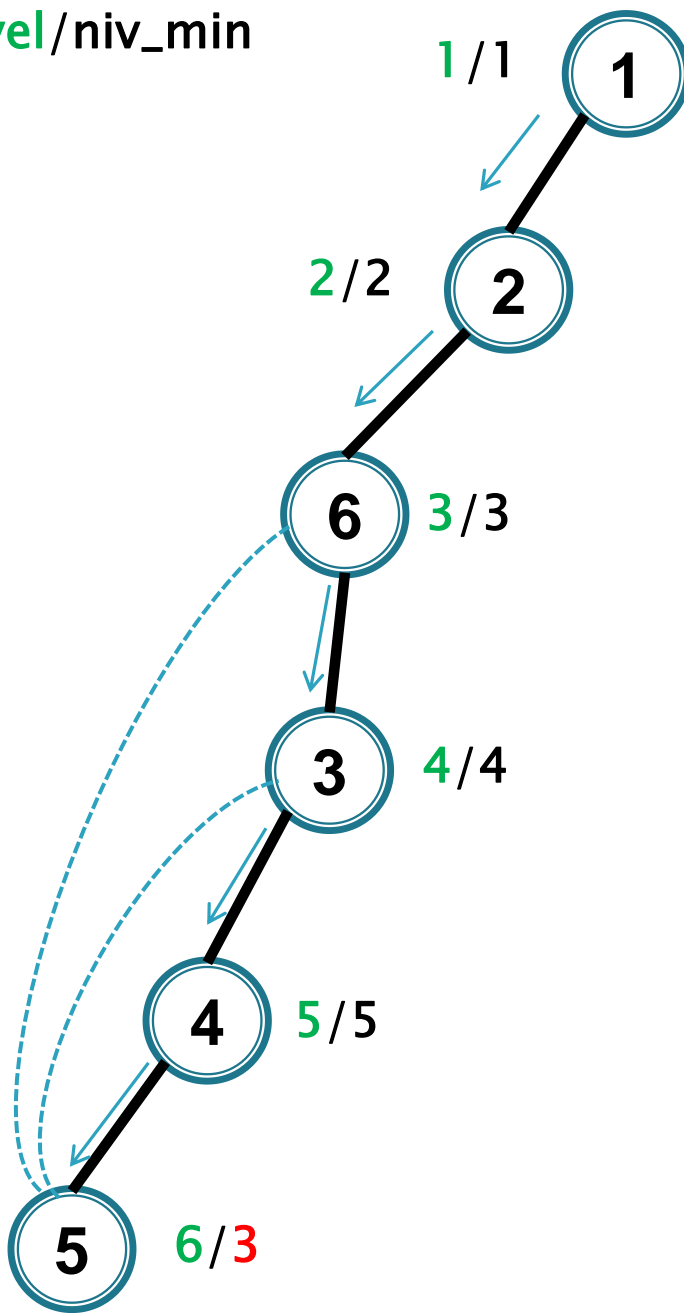
nivel/niv\_min



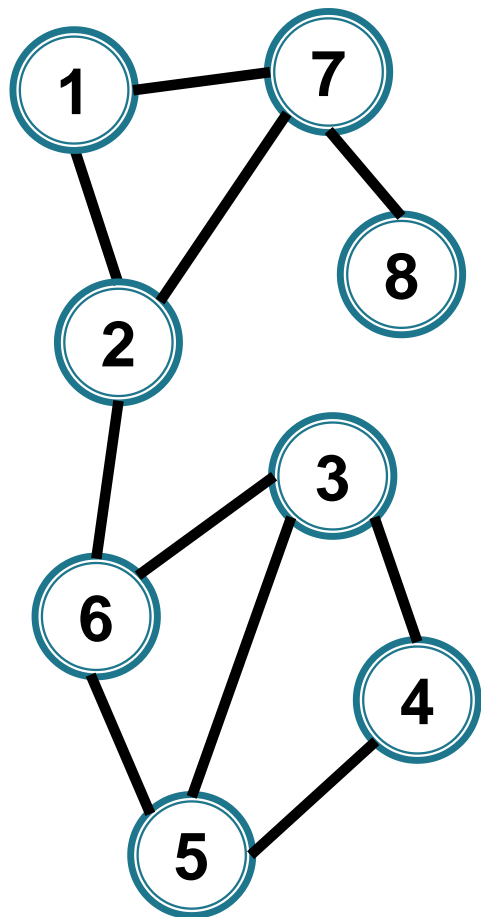


S:  
 1 2  
 2 6  
 6 3  
 3 4  
 4 5  
 5 3  
 5 6

nivel/niv\_min

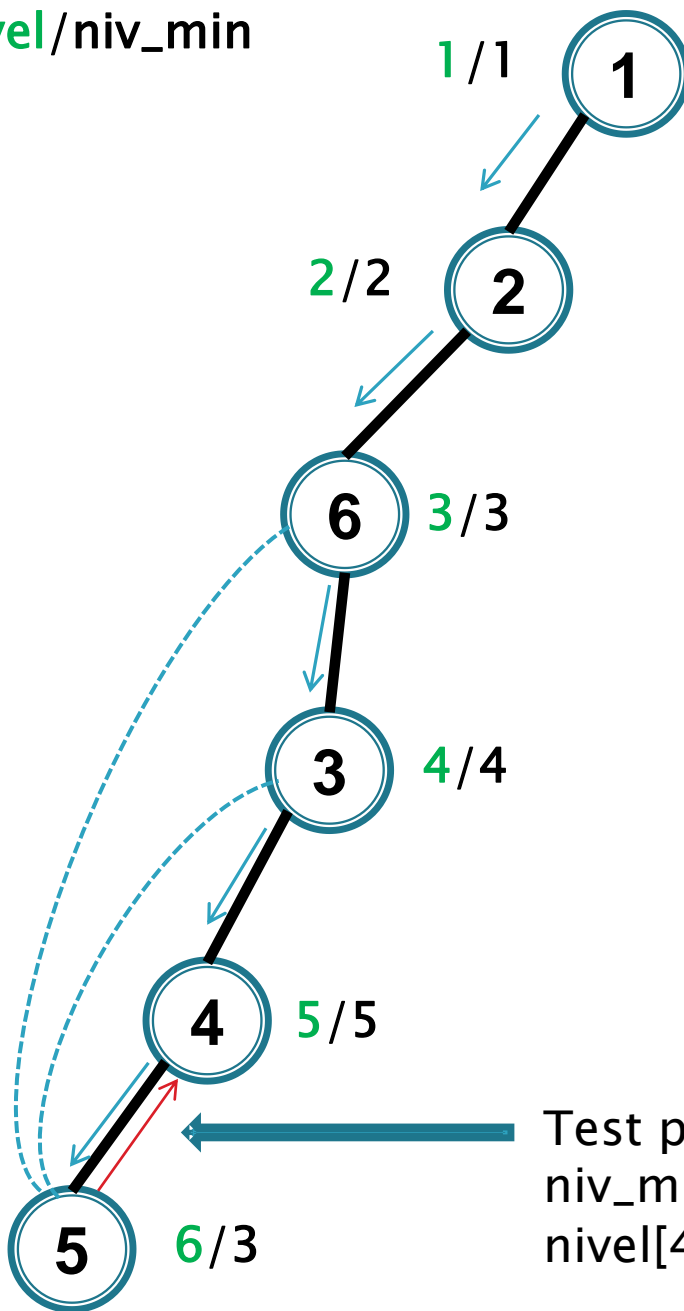




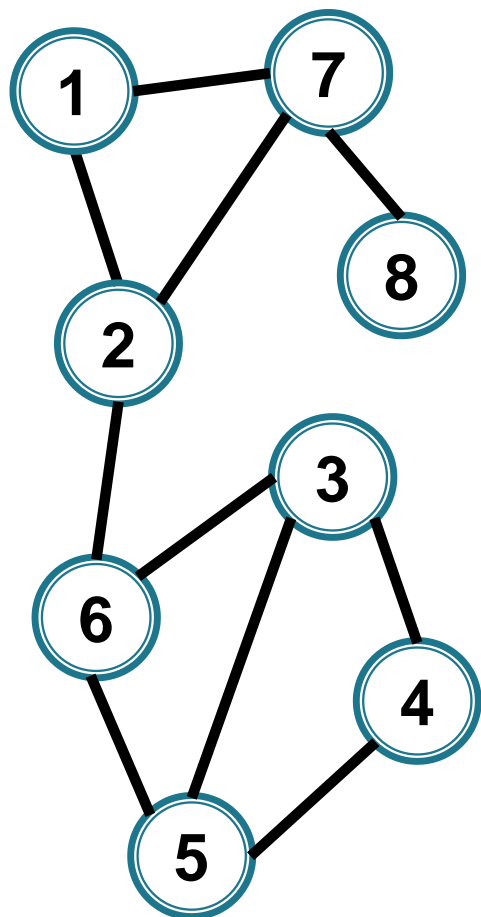


S:  
 1 2  
 2 6  
 6 3  
 3 4  
 4 5  
 5 3  
 5 6

nivel/niv\_min

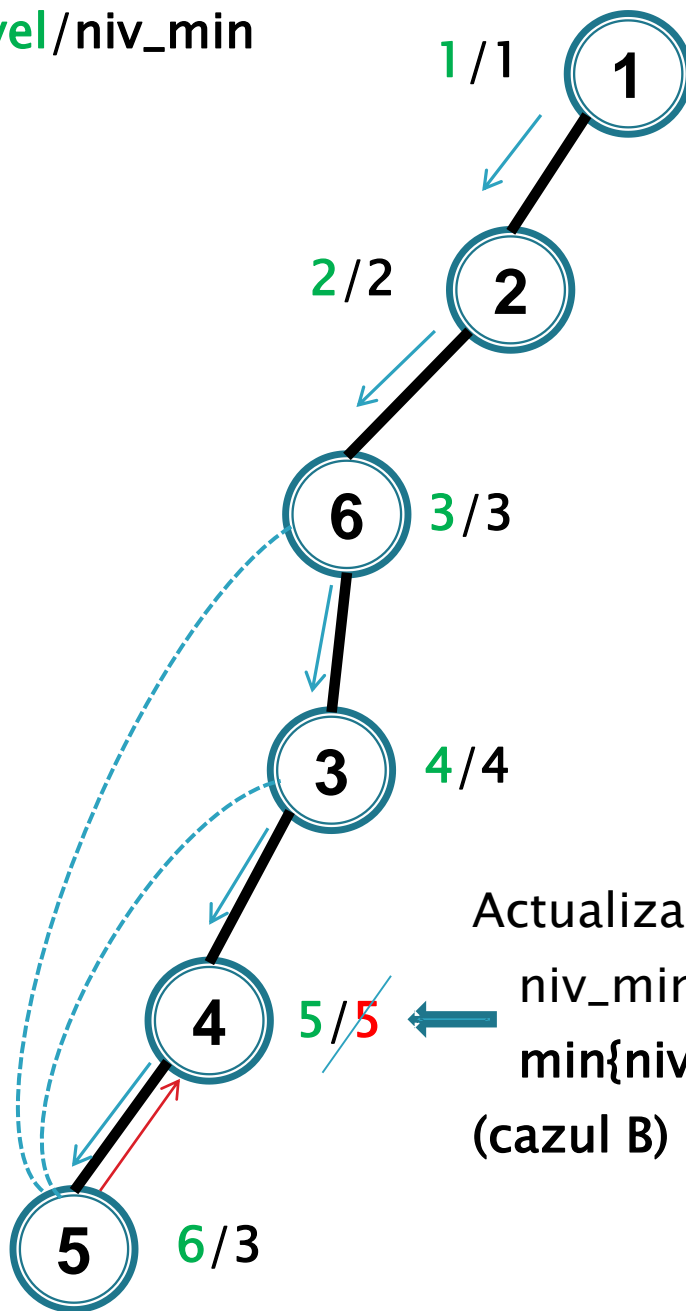


Test punct critic:  
 $niv\_min[5] = 3 < nivel[4] = 5 \Rightarrow \text{NU}$



S:  
 1 2  
 2 6  
 6 3  
 3 4  
 4 5  
 5 3  
 5 6

nivel/niv\_min

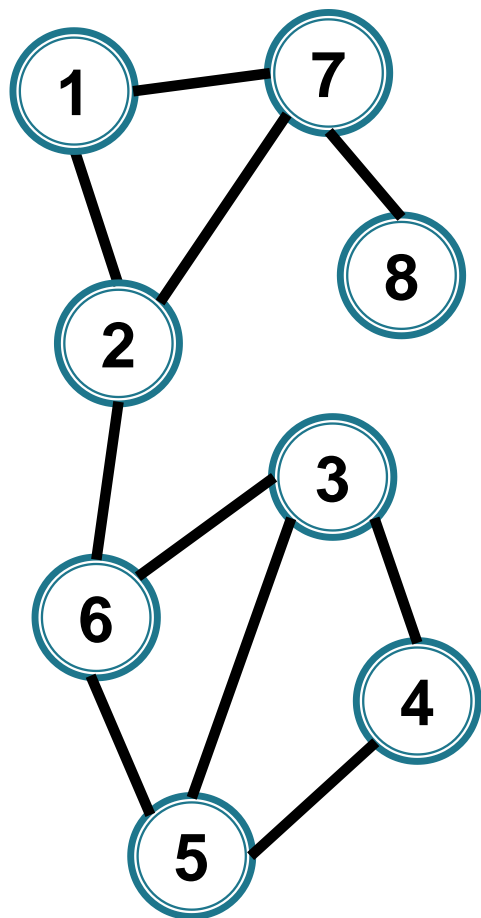


Actualizam

$\text{niv\_min}[4] =$

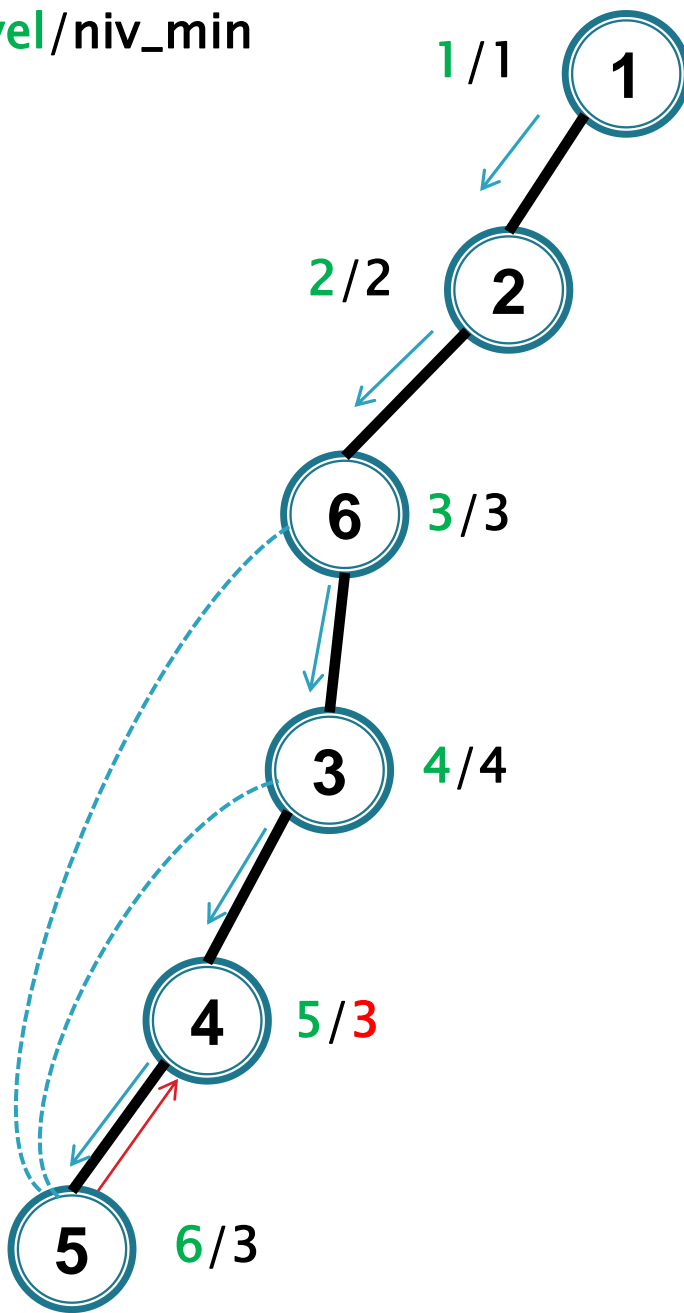
$\min\{\text{niv\_min}[4], \text{niv\_min}[5]\}$

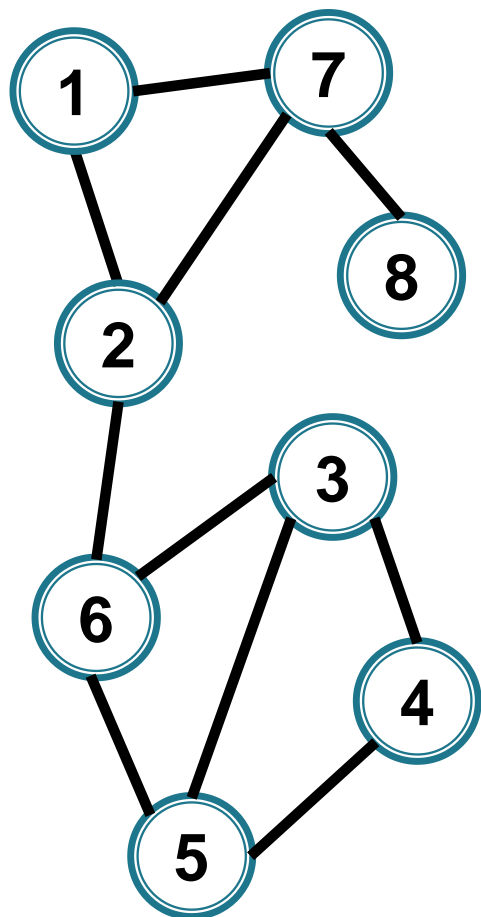
(cazul B)



S:  
 1 2  
 2 6  
 6 3  
 3 4  
 4 5  
 5 3  
 5 6

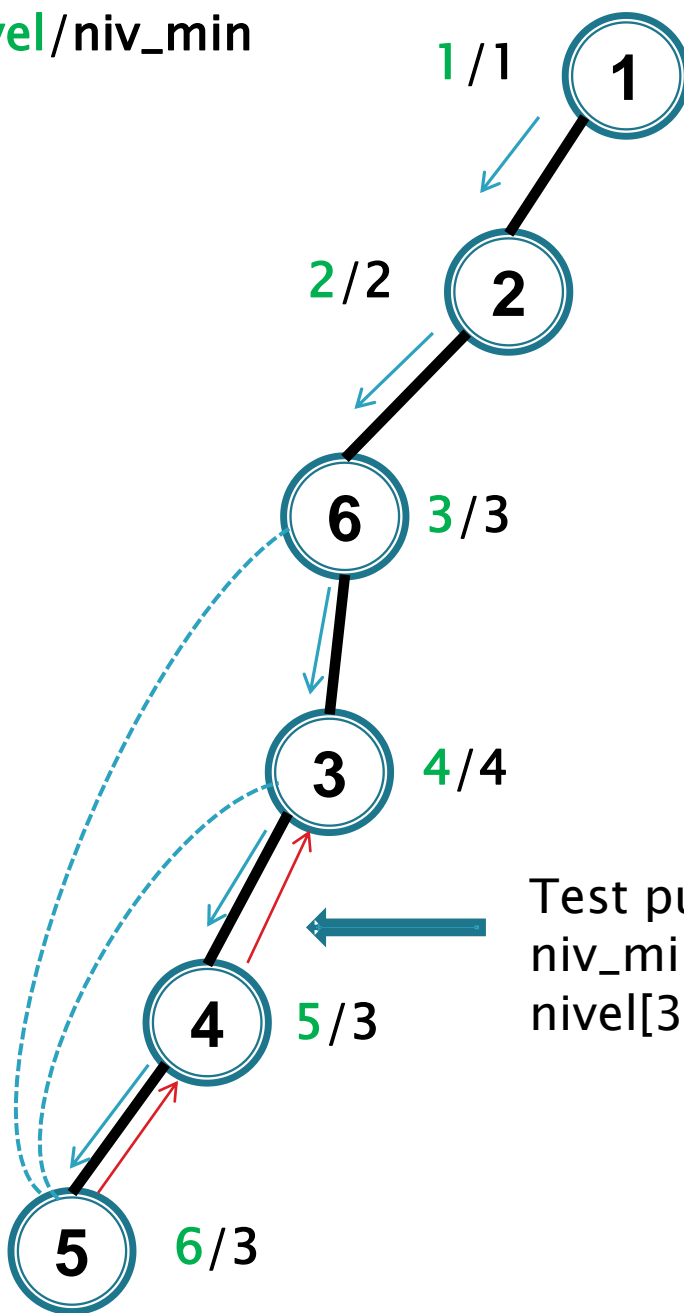
nivel/niv\_min



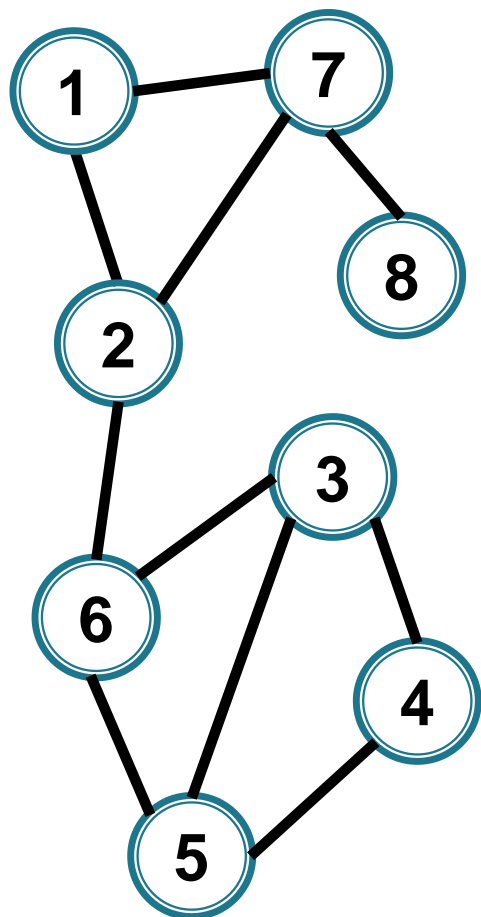


S:  
 1 2  
 2 6  
 6 3  
 3 4  
 4 5  
 5 3  
 5 6

nivel/niv\_min

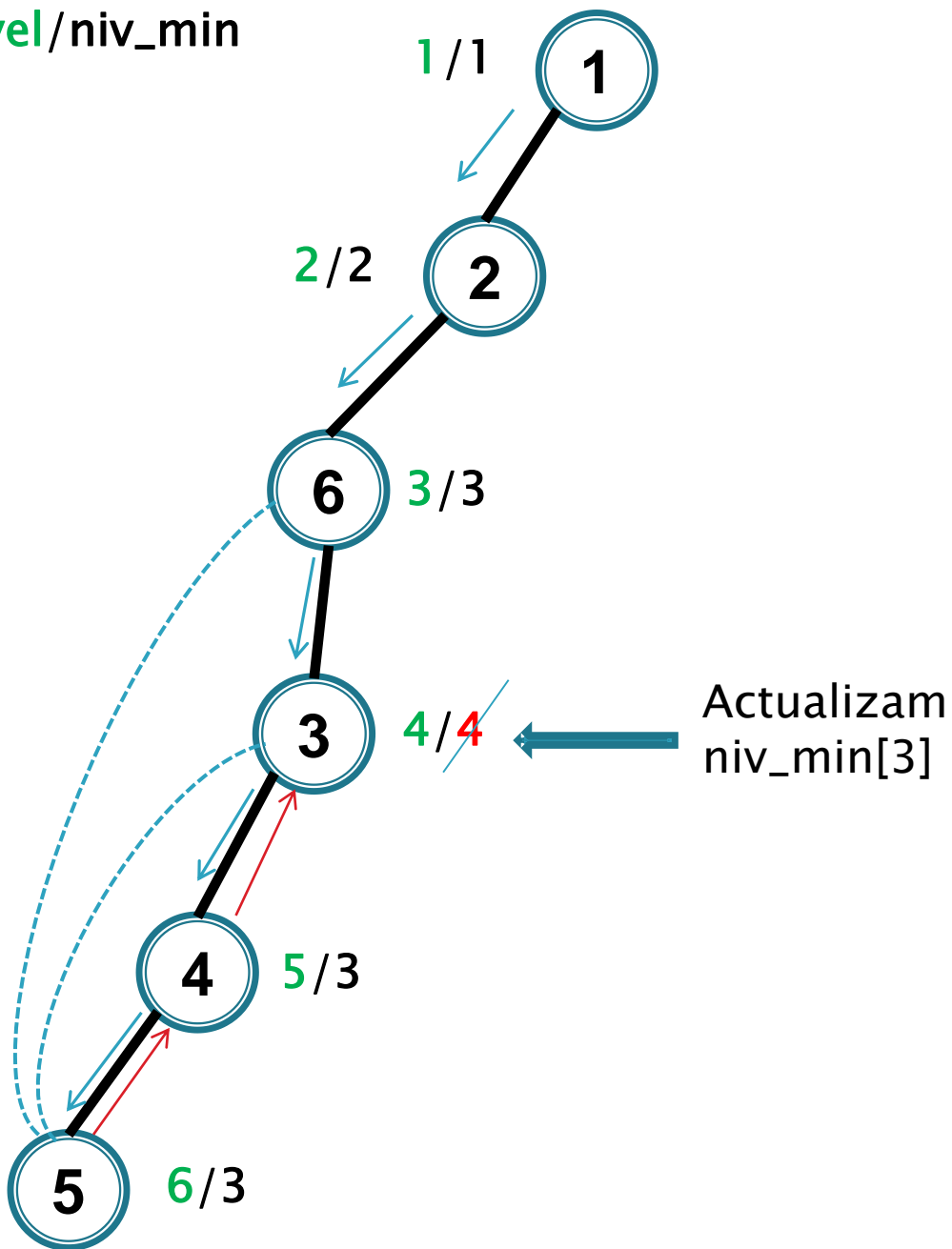


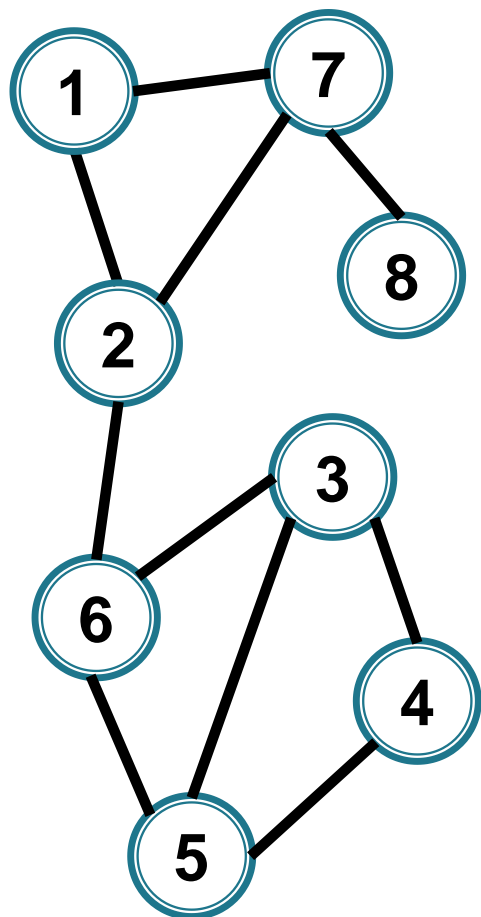
Test punct critic:  
 $\text{niv\_min}[4] = 3 < \text{nivel}[3] = 4 \Rightarrow \text{NU}$



S:  
 1 2  
 2 6  
 6 3  
 3 4  
 4 5  
 5 3  
 5 6

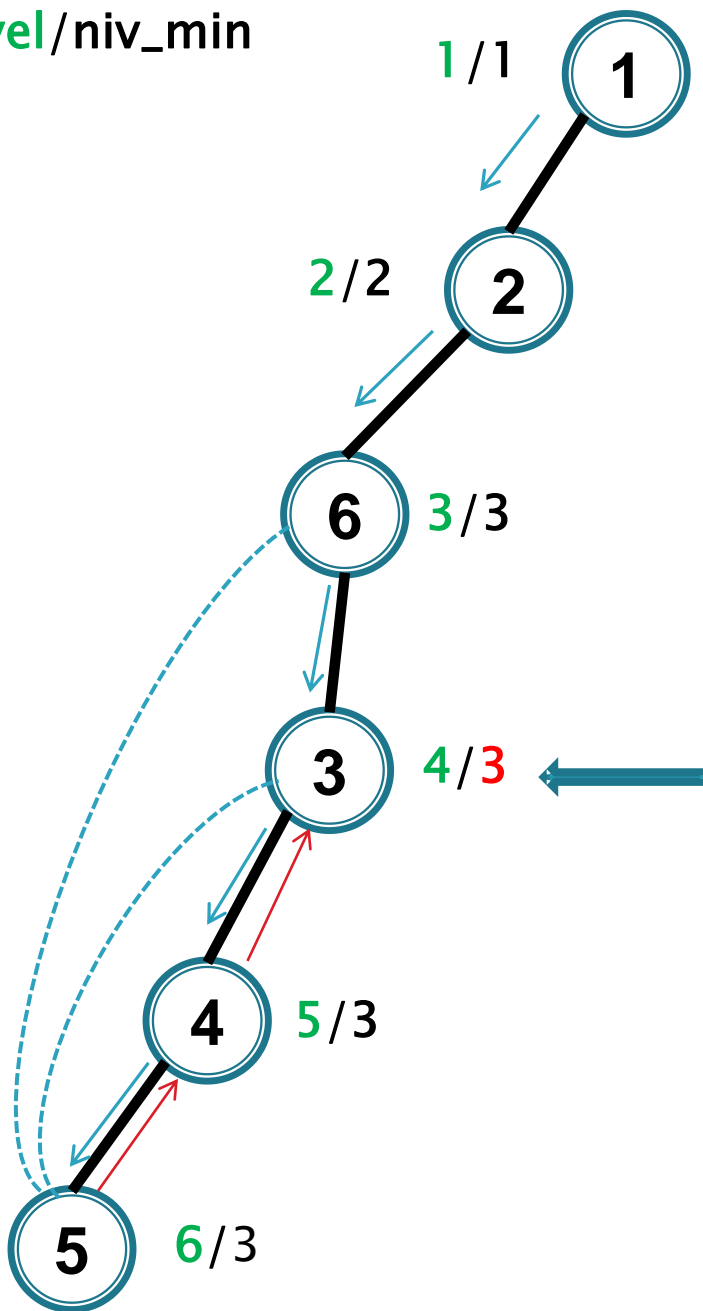
nivel/niv\_min

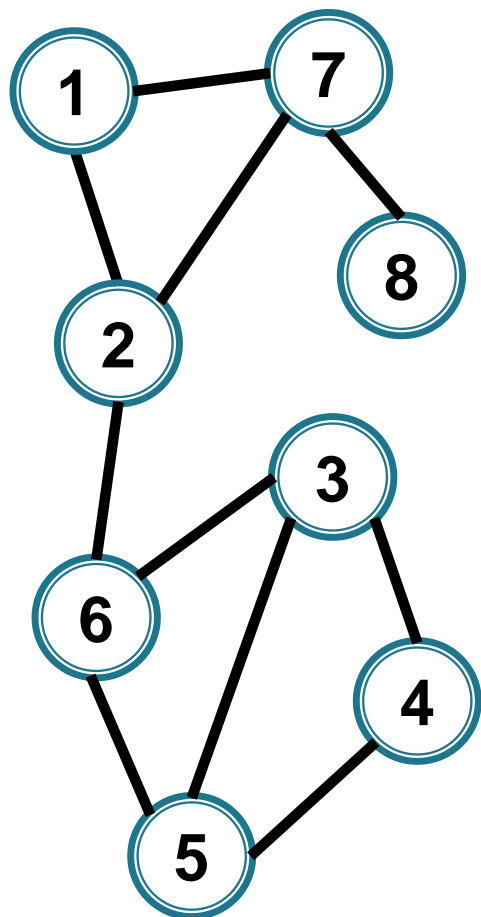




S:  
 1 2  
 2 6  
 6 3  
 3 4  
 4 5  
 5 3  
 5 6

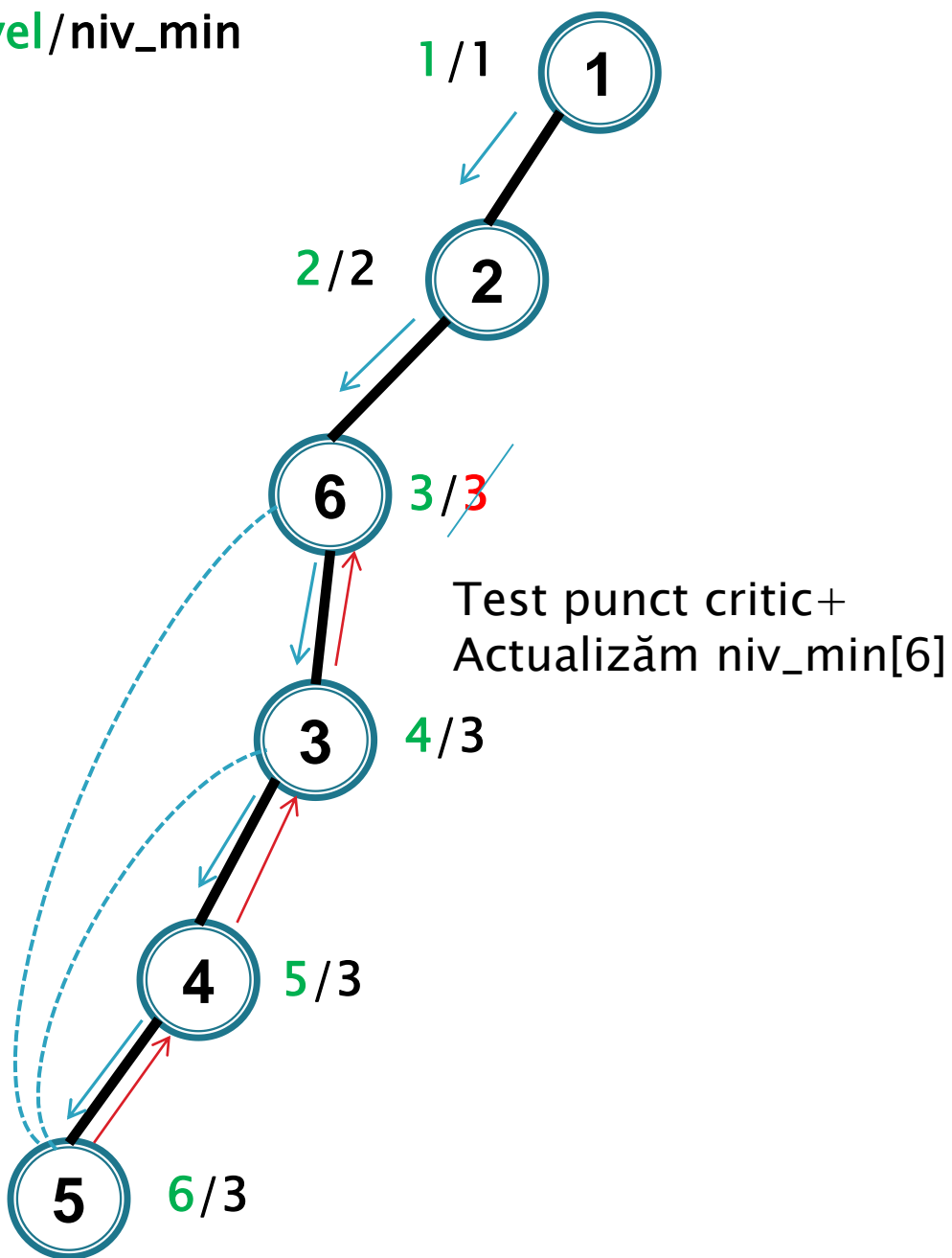
nivel/niv\_min

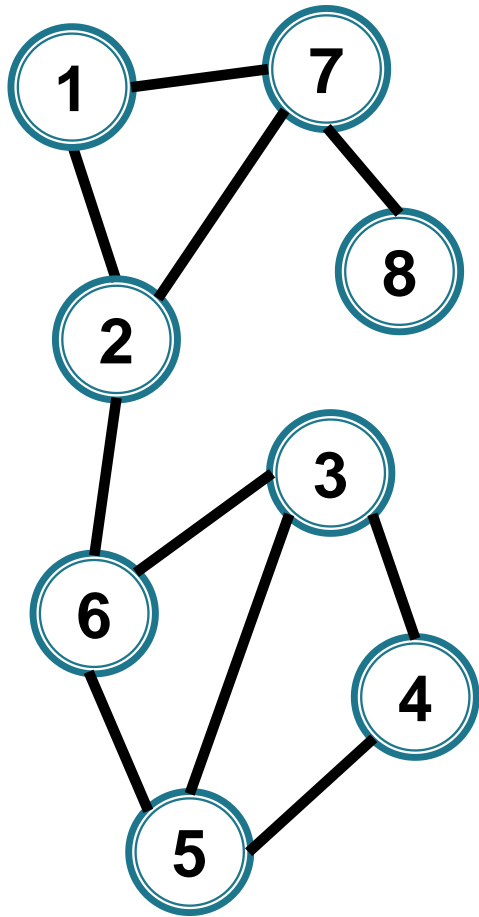




S:  
 1 2  
 2 6  
 6 3  
 3 4  
 4 5  
 5 3  
 5 6

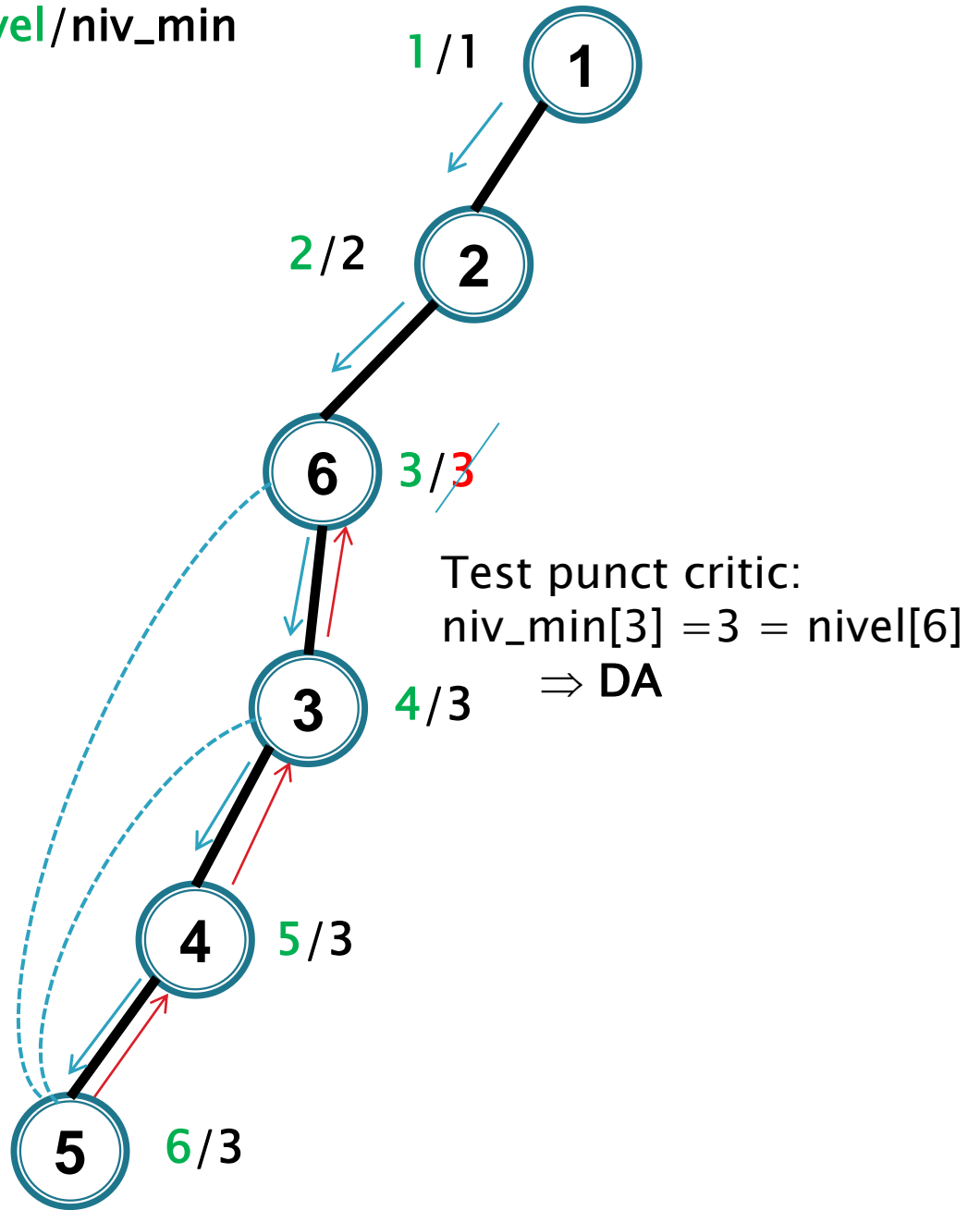
nivel/niv\_min



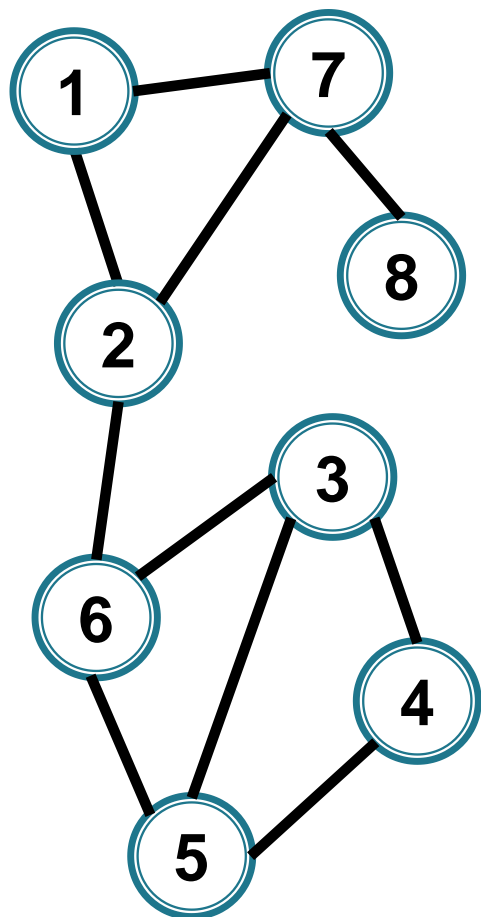


S:  
 1 2  
 2 6  
 6 3  
 3 4  
 4 5  
 5 3  
 5 6

nivel/niv\_min

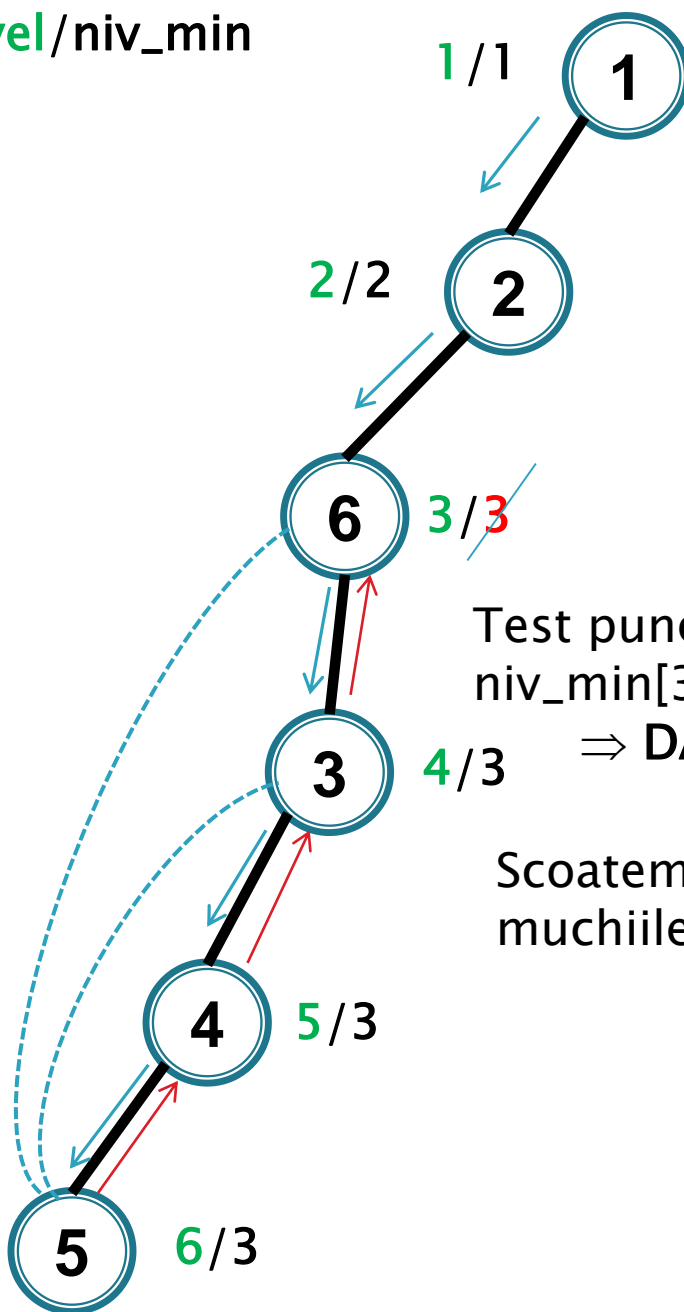






S:  
 1 2  
 2 6  
 6 3  
 3 4  
 4 5  
 5 3  
 5 6

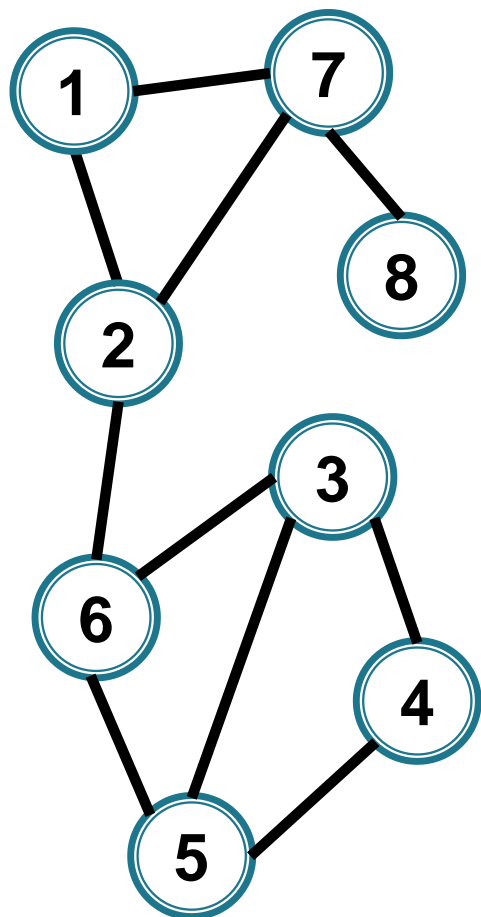
nivel/niv\_min



Test punct critic:  
 $\text{niv\_min}[3] = 3 = \text{nivel}[6]$

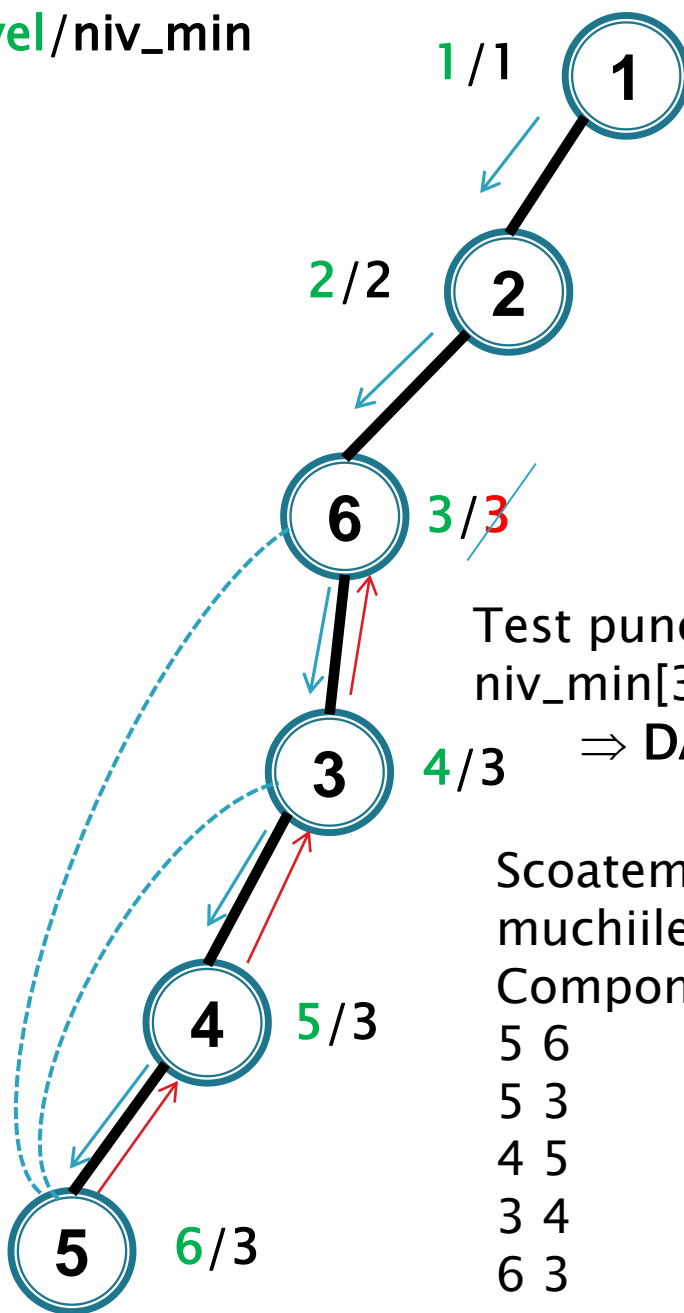
$\Rightarrow \text{DA}$

Scoatem din stiva toate  
 muchiile pana la 3 6  $\Rightarrow$



S:  
1 2  
2 6

nivel/niv\_min



Test punct critic:  
 $\text{niv\_min}[3] = 3 = \text{nivel}[6]$

$\Rightarrow$  DA

Scoatem din stiva toate  
muchiiile pana la 3 6  $\Rightarrow$   
Componenta cu muchiiile

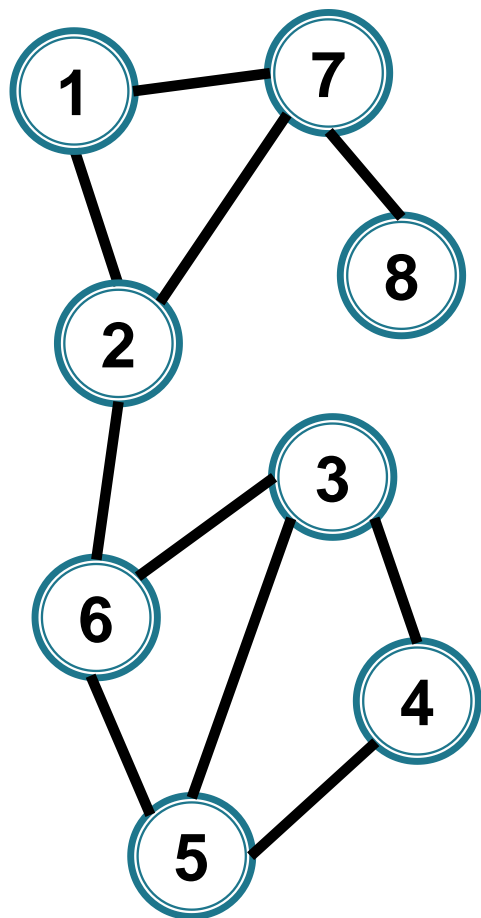
5 6

5 3

4 5

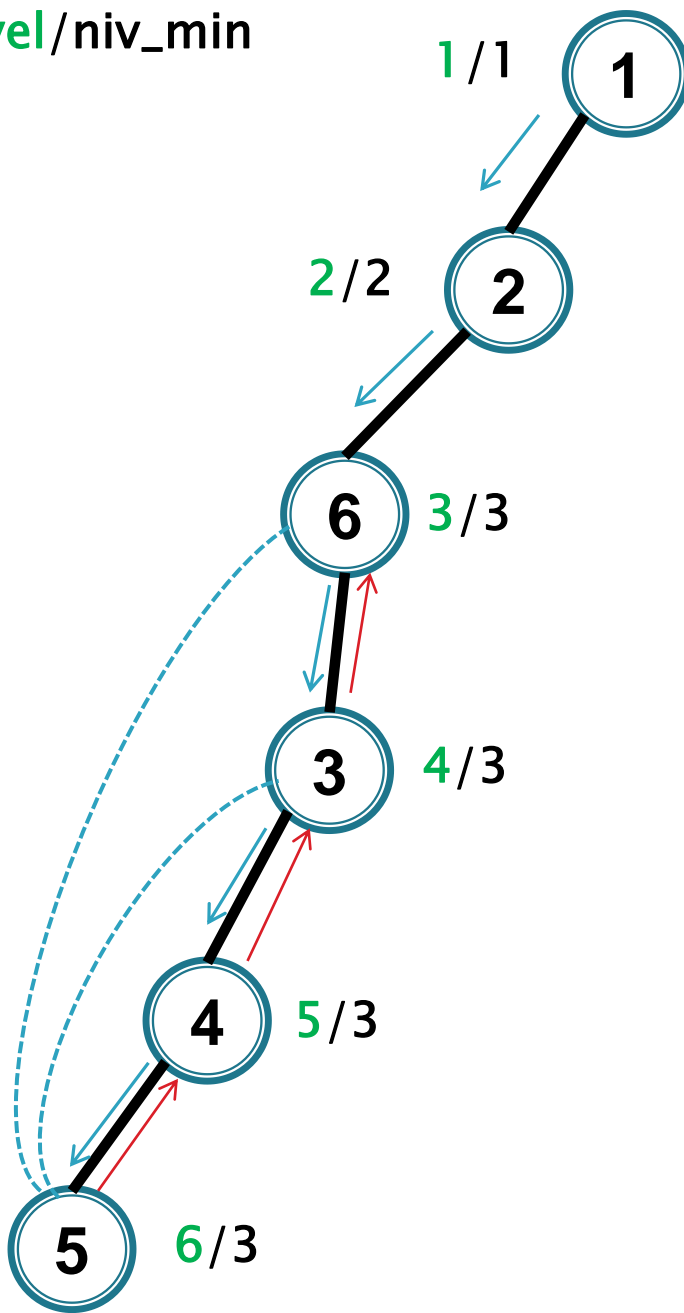
3 4

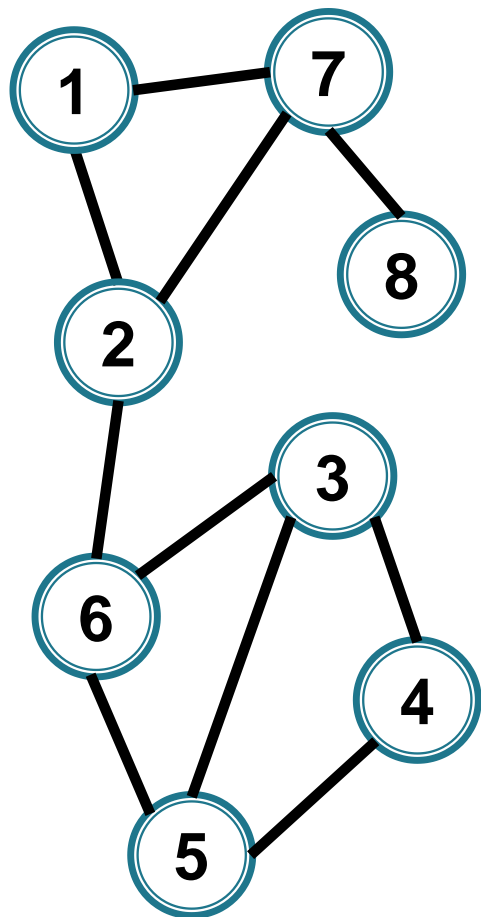
6 3



S:  
1 2  
2 6

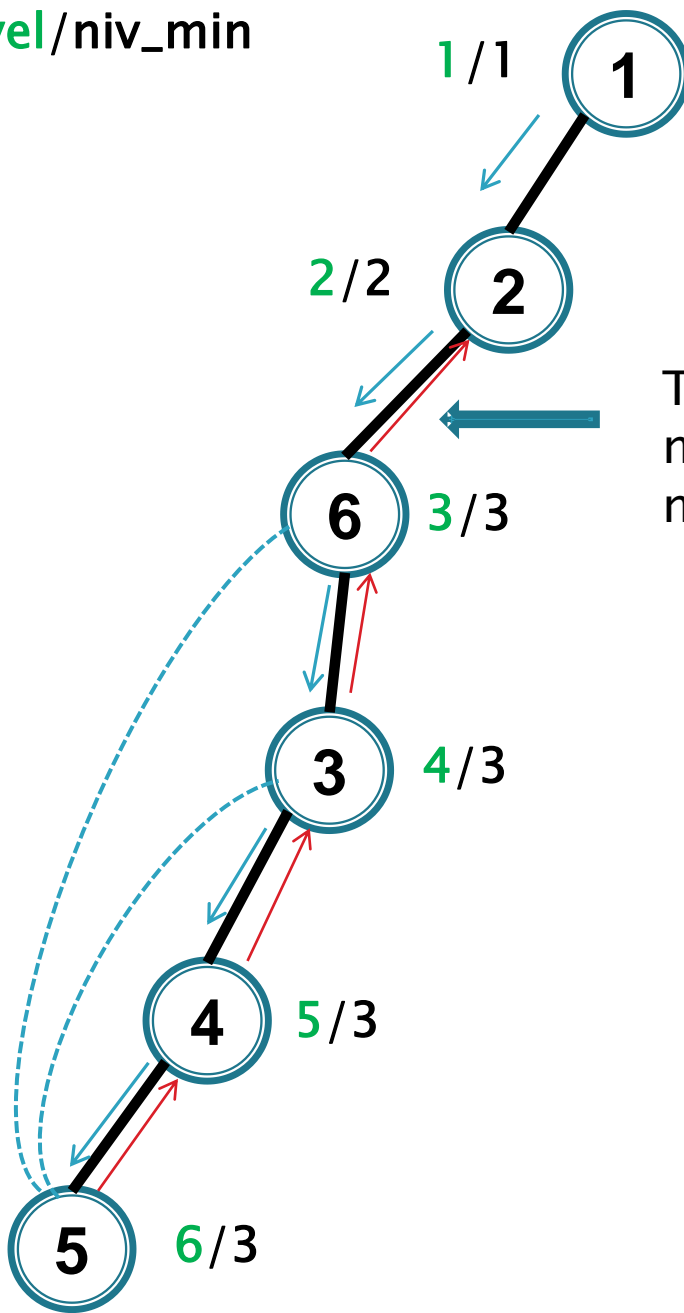
nivel/niv\_min



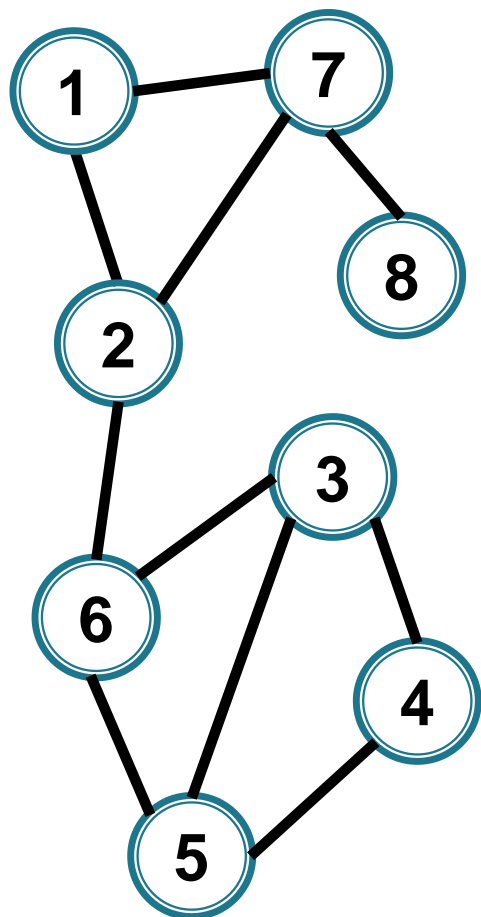


S:  
1 2  
2 6

nivel/niv\_min

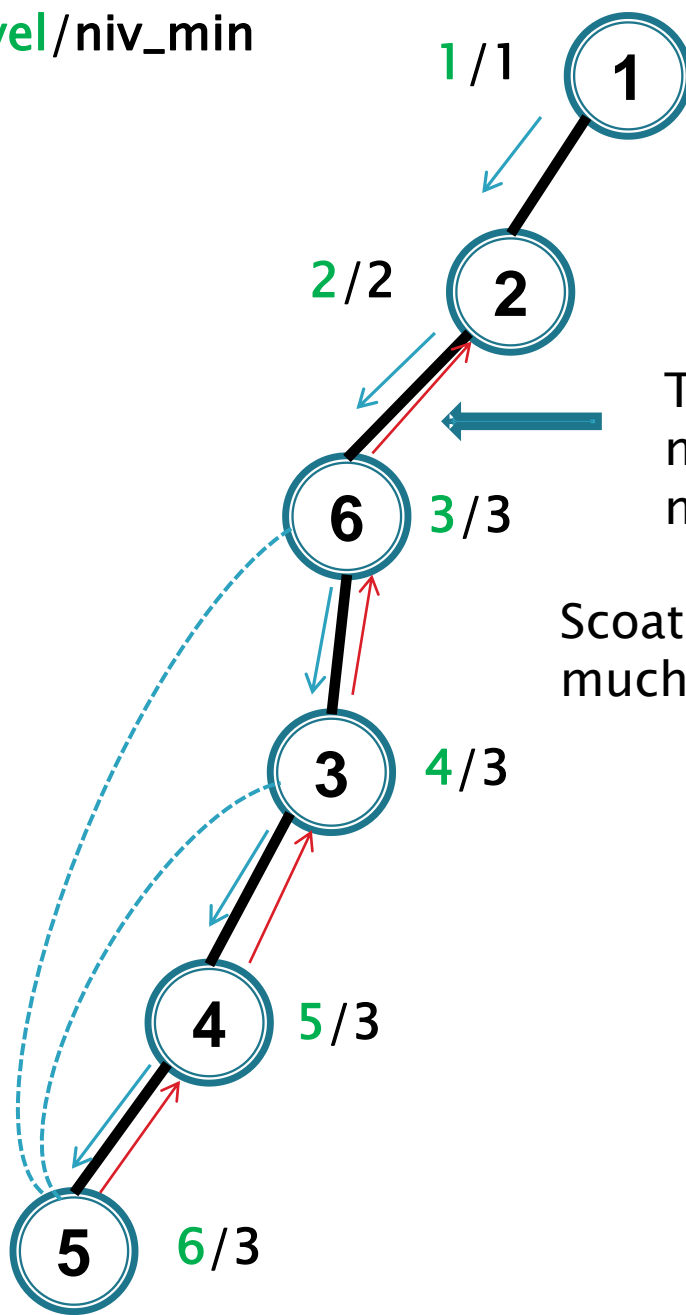


Test punct critic:  
 $\text{niv\_min}[6] = 3 >$   
 $\text{nivel}[2] = 2 \Rightarrow \text{DA}$



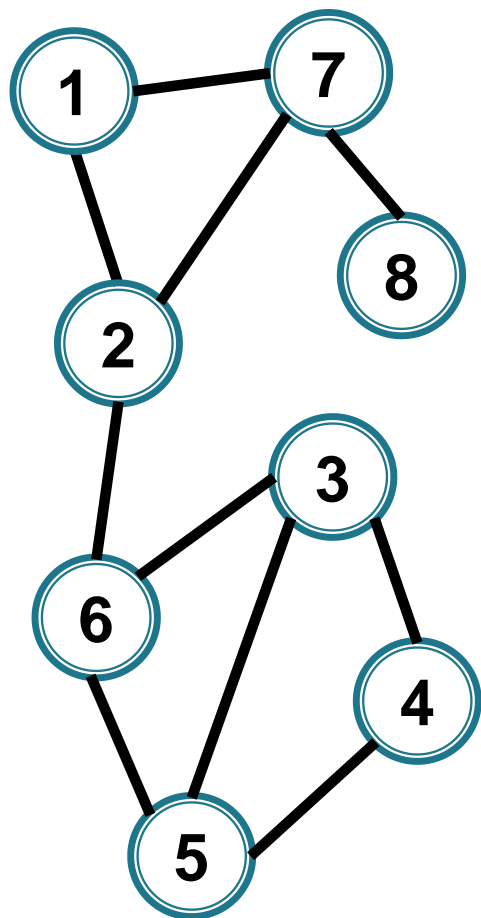
S:  
1 2  
2 6

nivel/niv\_min



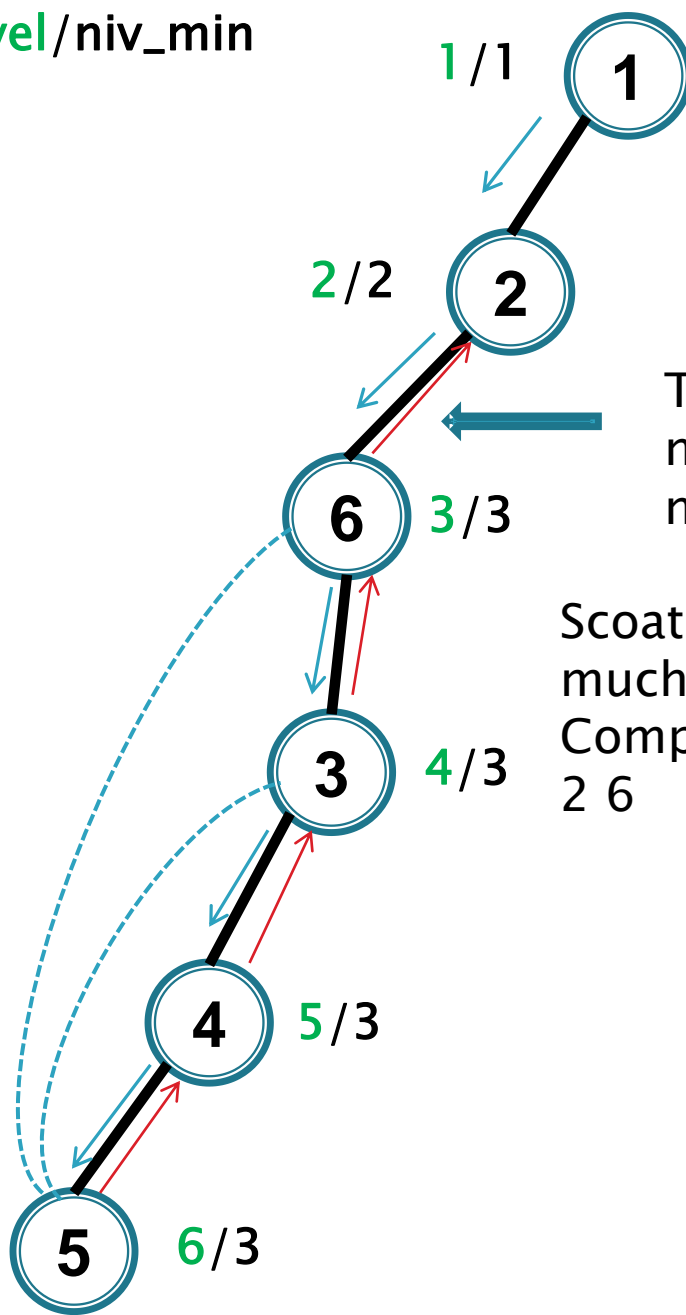
Test punct critic:  
niv\_min[6] = 3 >  
nivel[2] = 2 ⇒ **DA**

Scoatem din stiva toate  
muchiiile pana la 2 6 =>



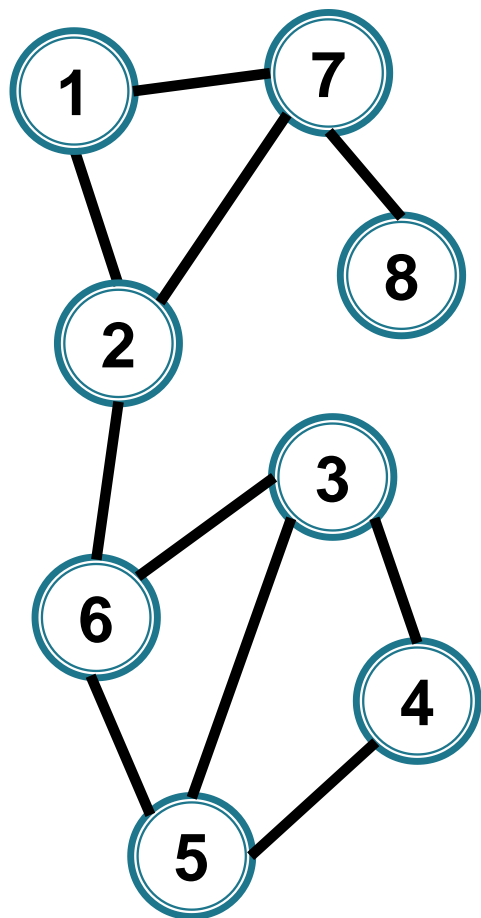
S:  
1 2

nivel/niv\_min



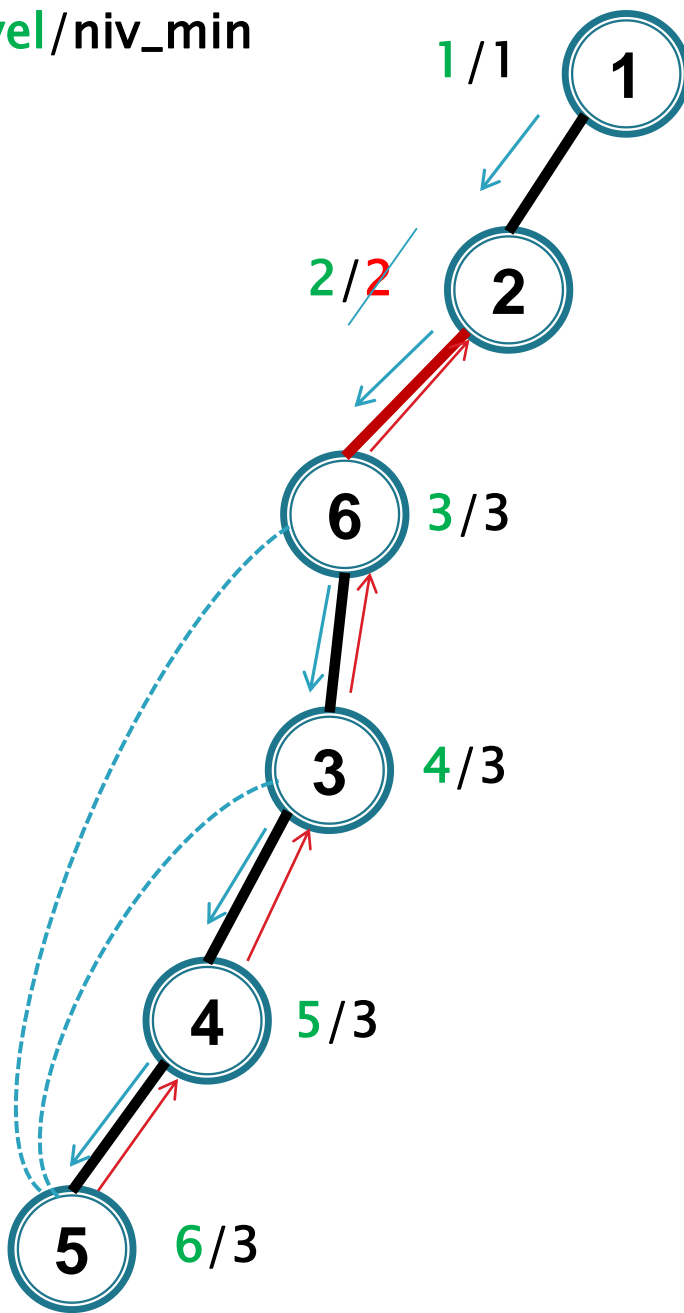
Test punct critic:  
 $\text{niv\_min}[6] = 3 >$   
 $\text{nivel}[2] = 2 \Rightarrow \text{DA}$

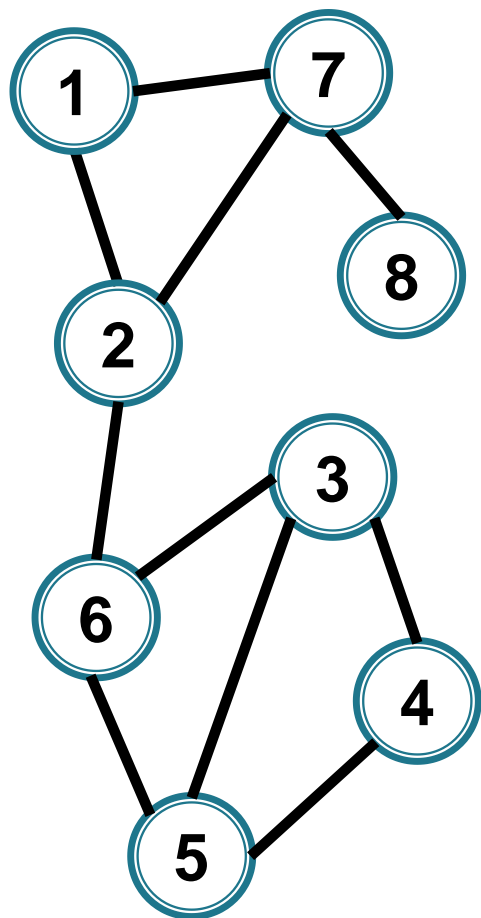
Scoatem din stiva toate  
 muchiile pana la 2 6  $\Rightarrow$   
 Componenta cu muchiile  
 2 6



S:  
1 2

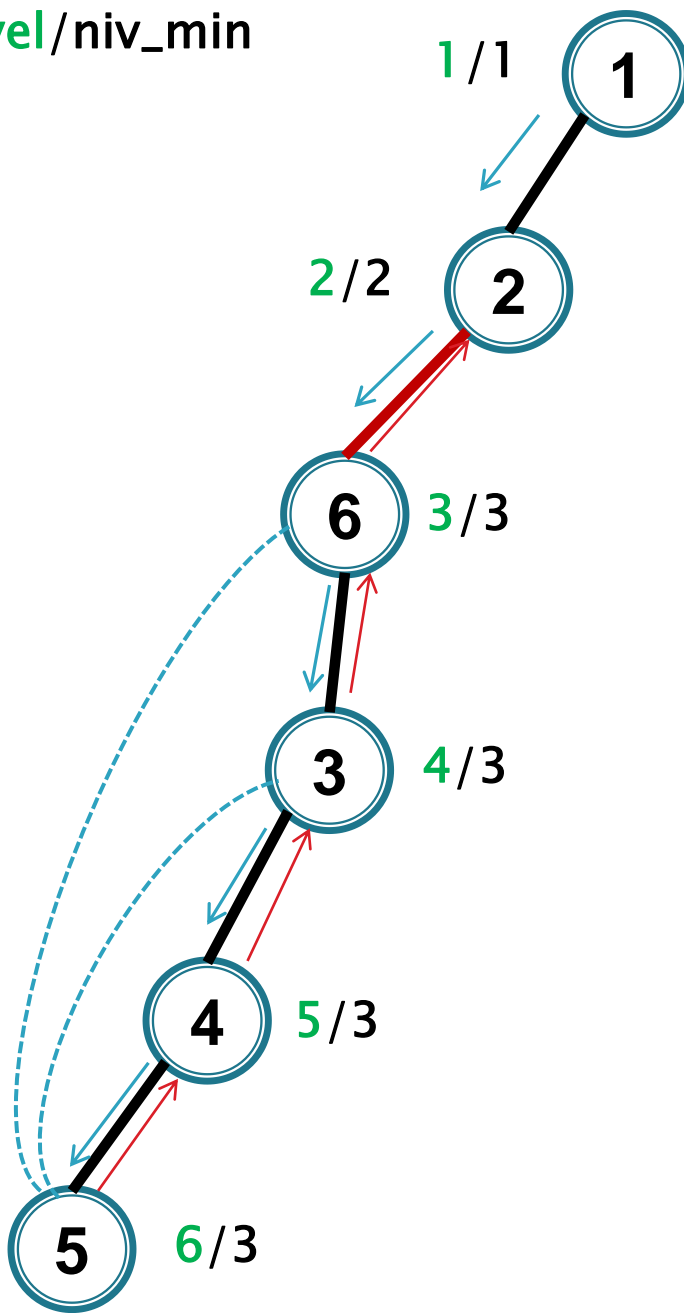
nivel/niv\_min



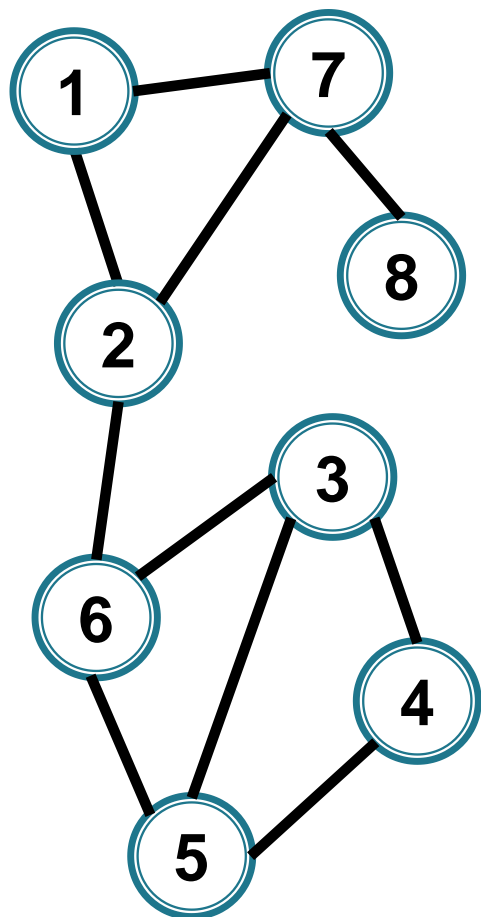


S:  
1 2

nivel/niv\_min

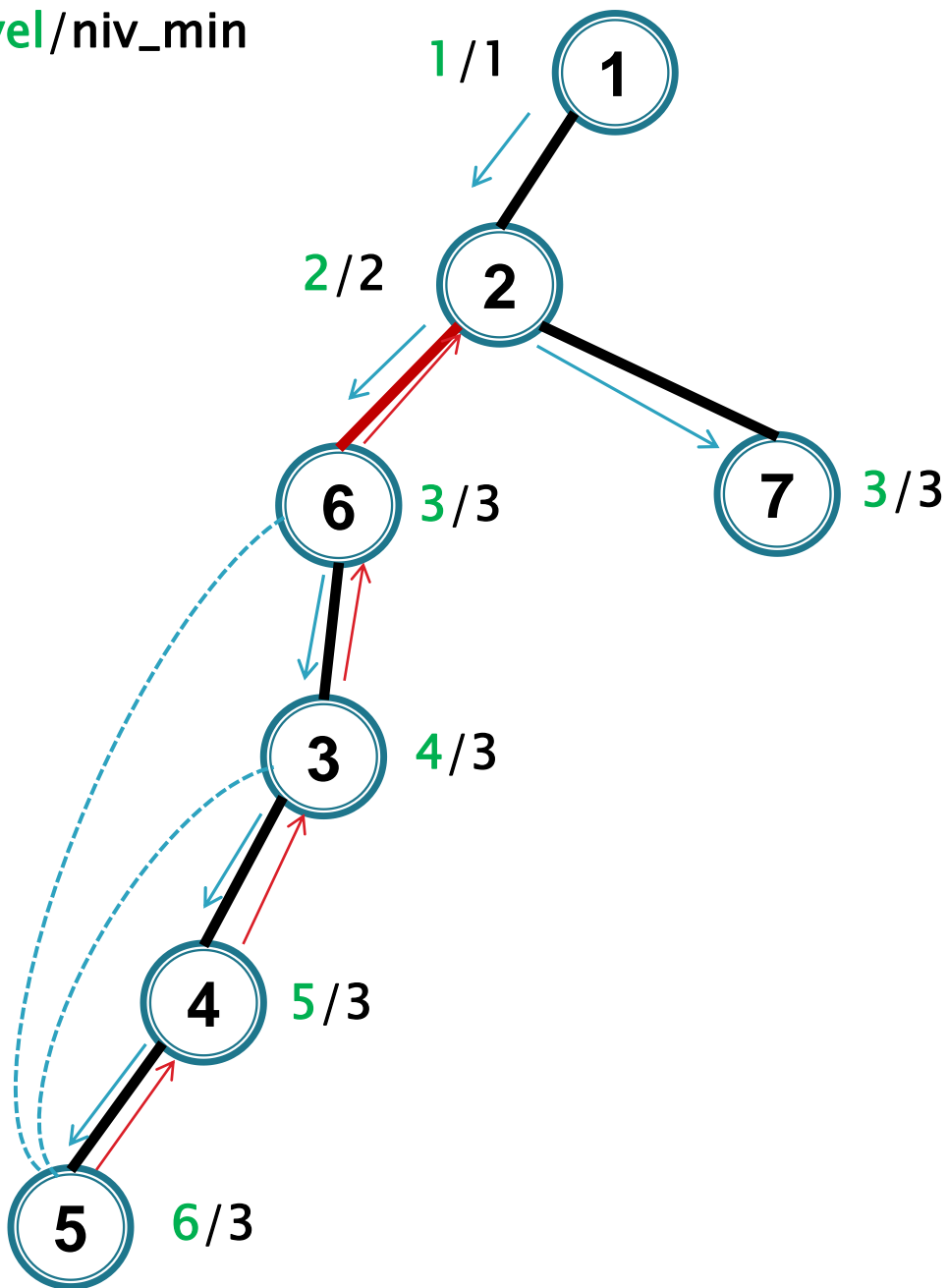


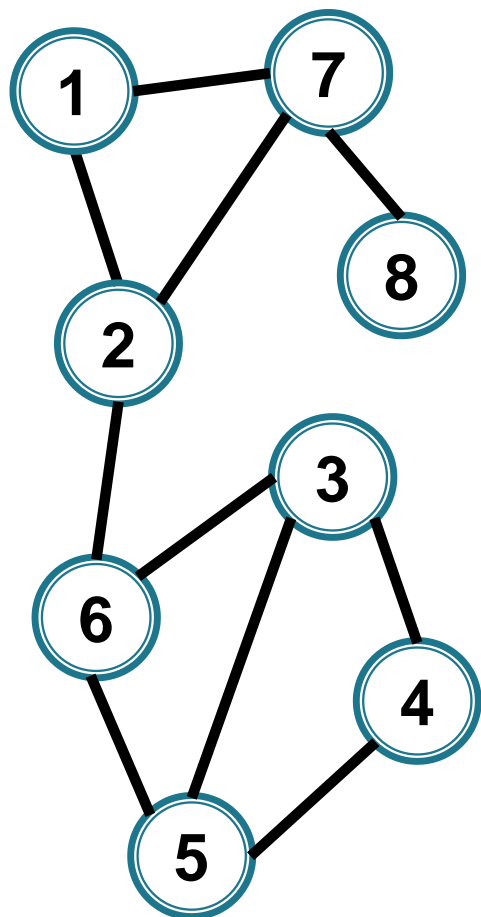




S:  
1 2  
2 7

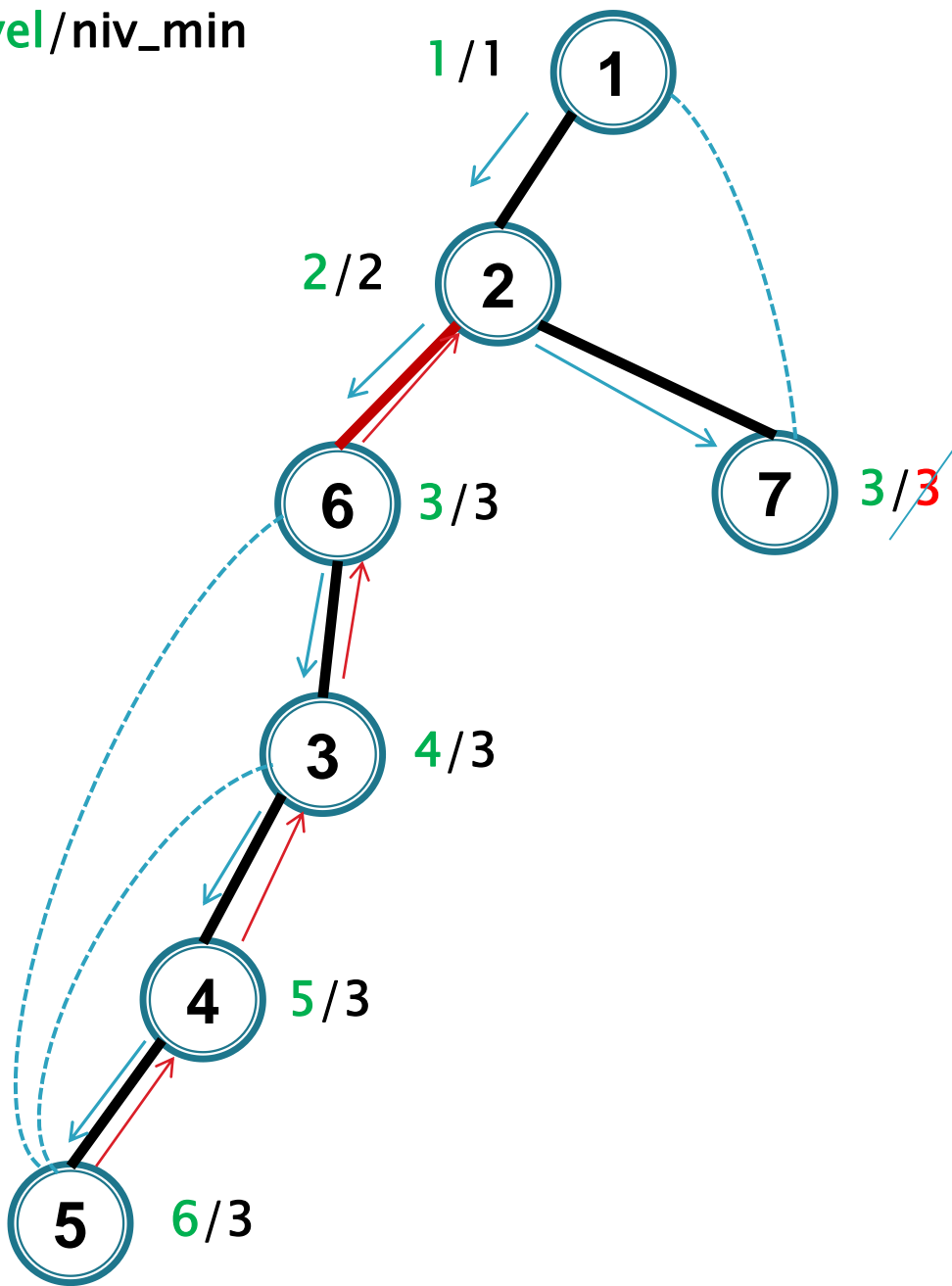
nivel/niv\_min

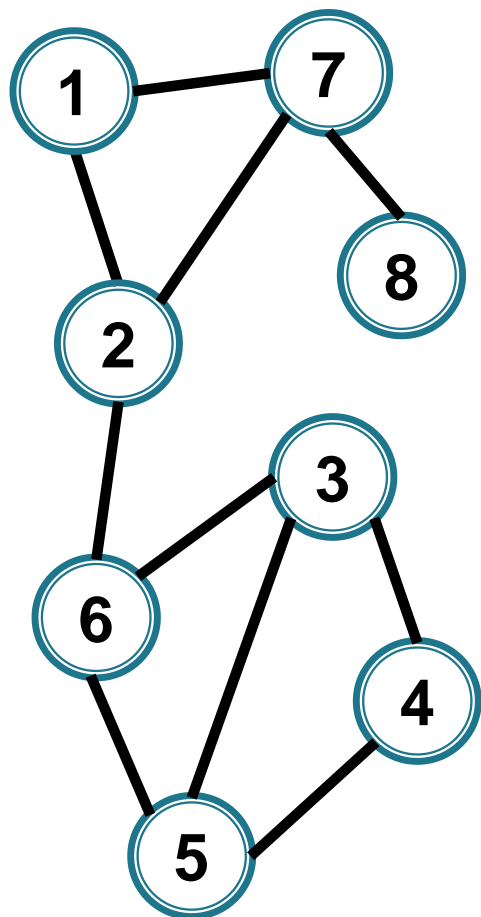




S:  
 1 2  
 2 7  
 1 7

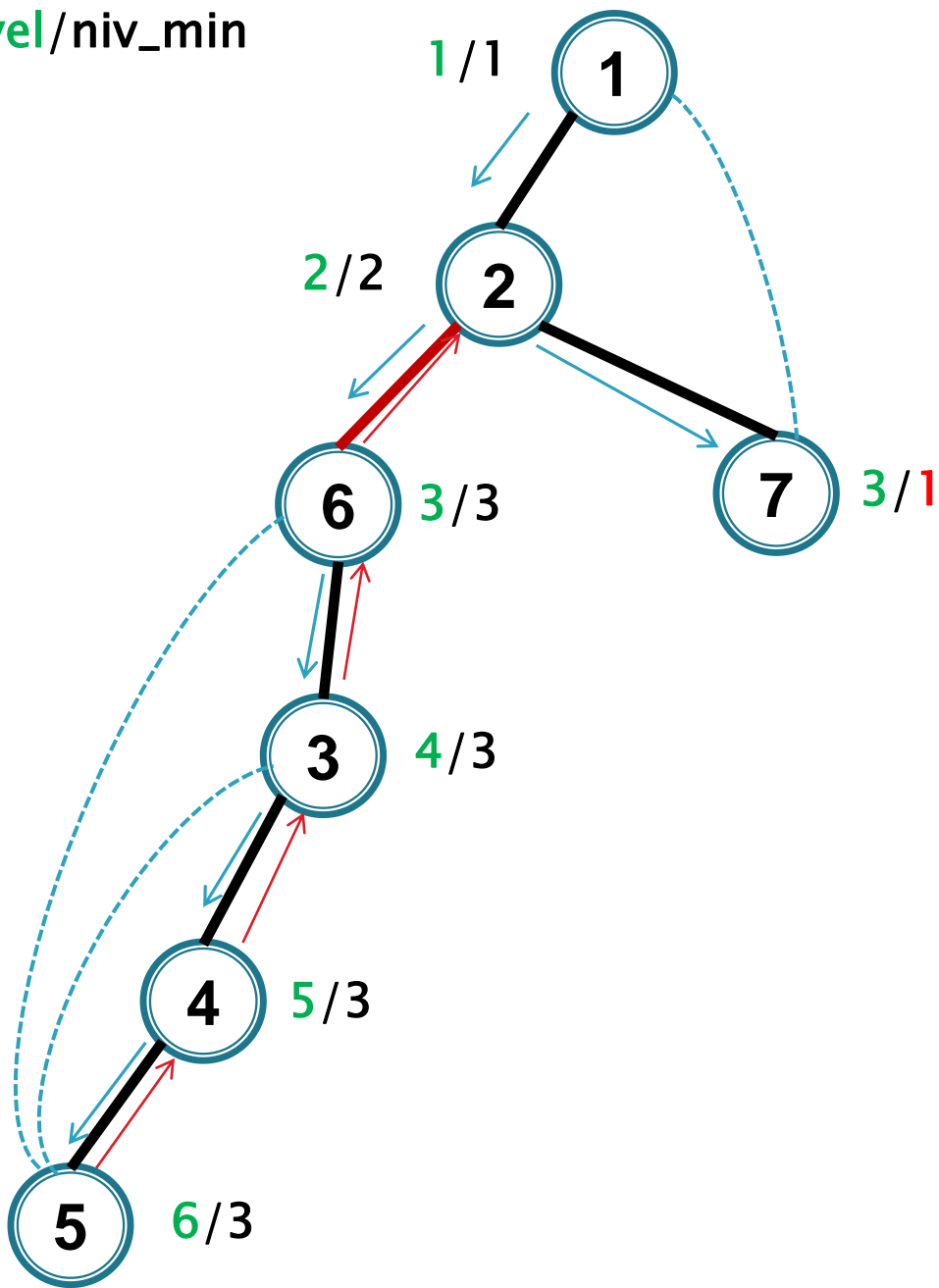
nivel/niv\_min

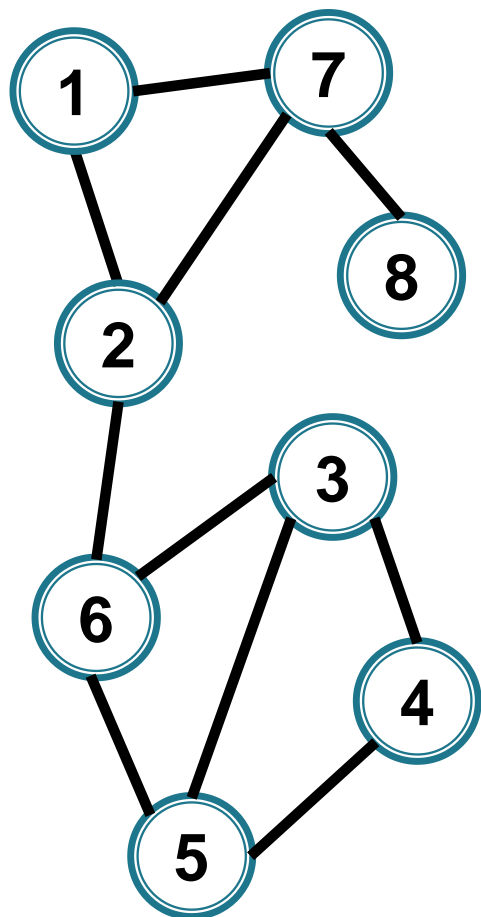




S:  
1 2  
2 7  
1 7

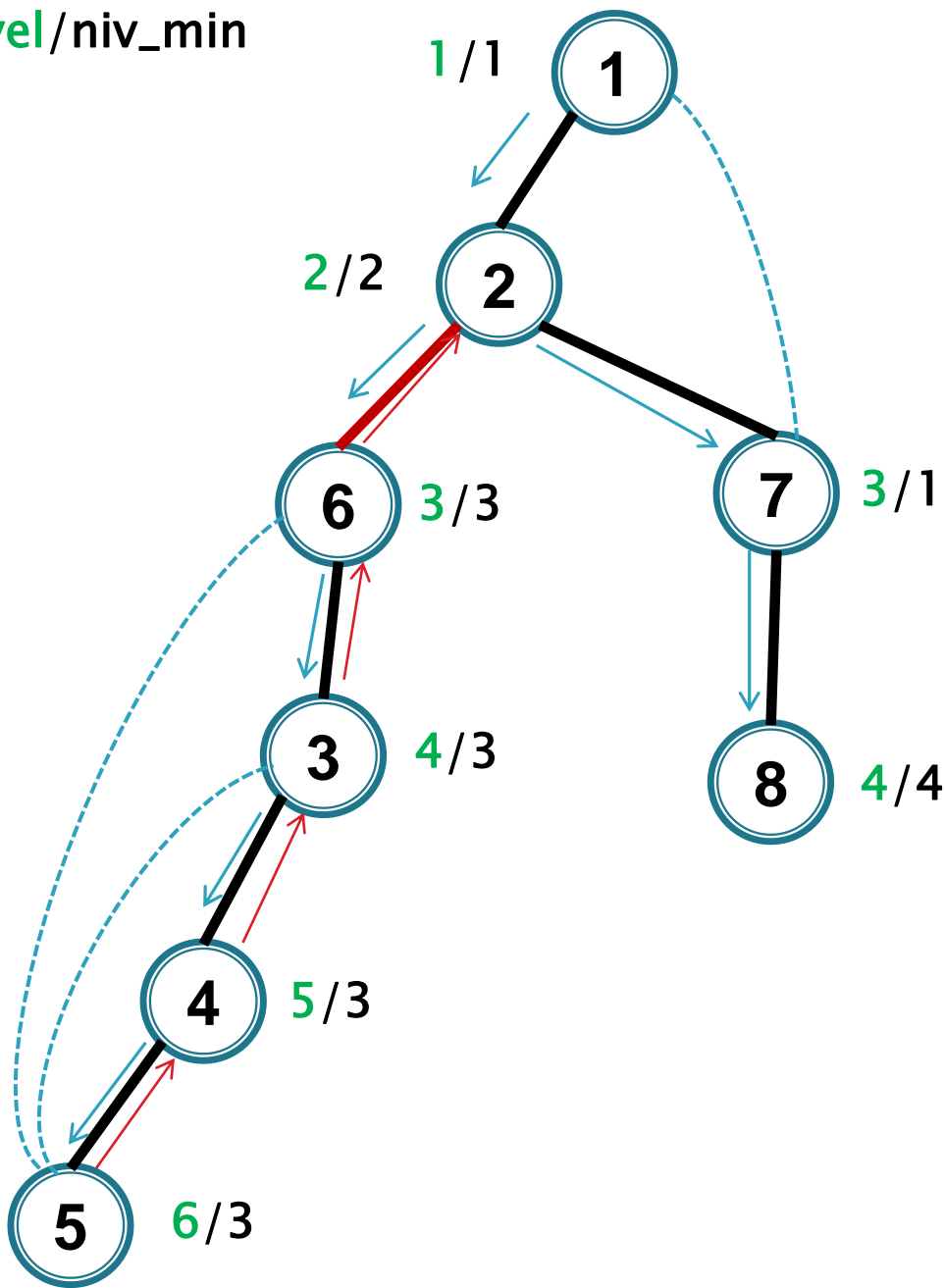
nivel/niv\_min

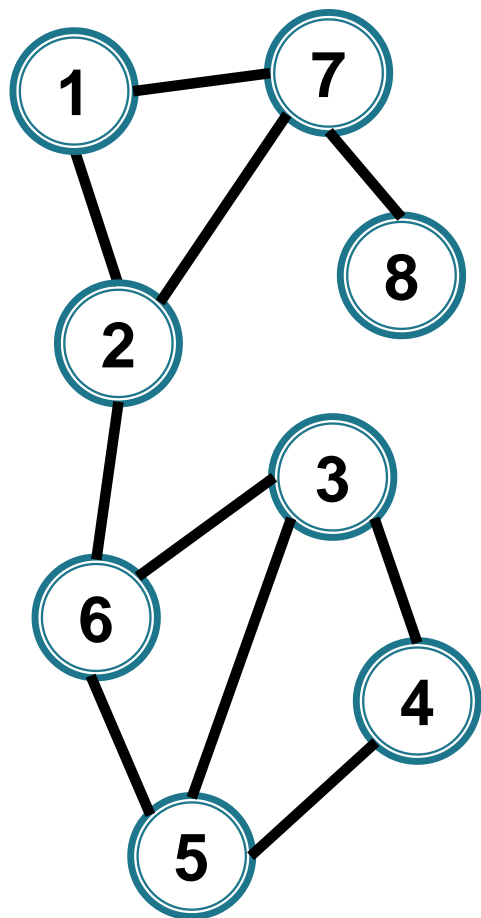




S:  
 1 2  
 2 7  
 1 7  
 7 8

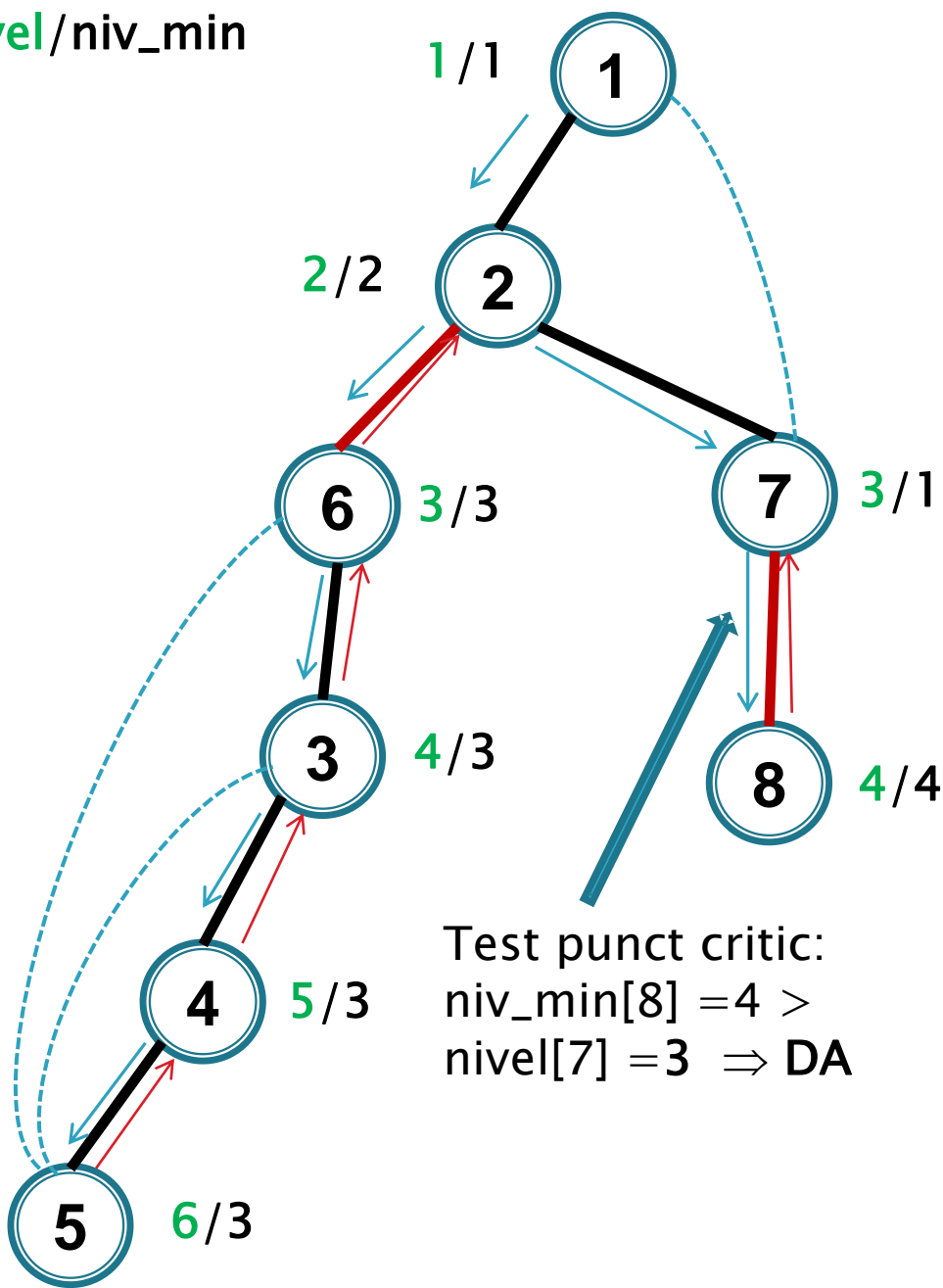
nivel/niv\_min



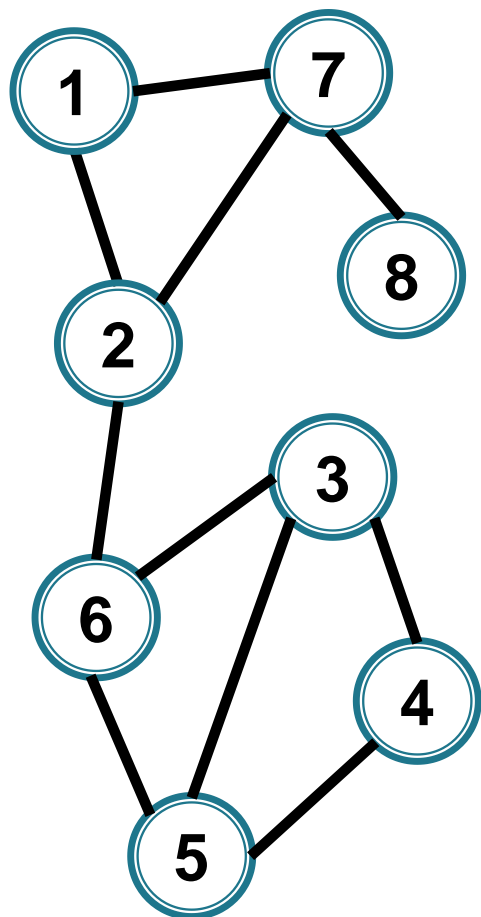


S:  
 1 2  
 2 7  
 1 7  
 7 8

nivel/niv\_min

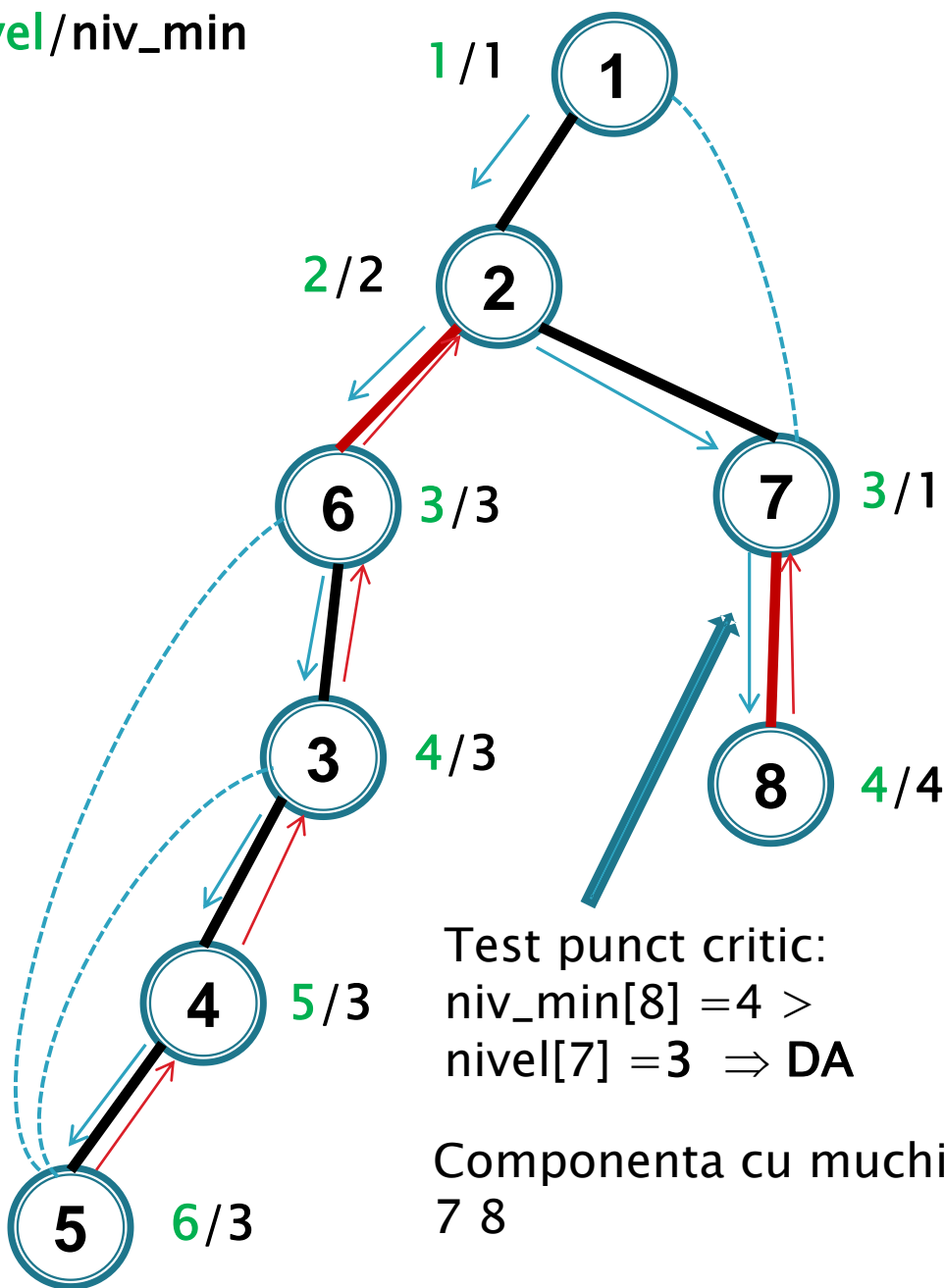


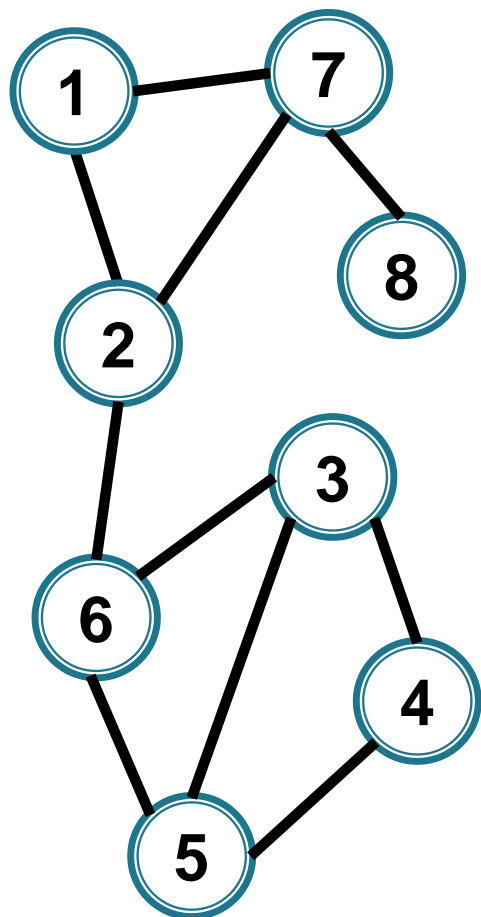
Test punct critic:  
 $\text{niv\_min}[8] = 4 >$   
 $\text{nivel}[7] = 3 \Rightarrow \text{DA}$



S:  
1 2  
2 7  
1 7

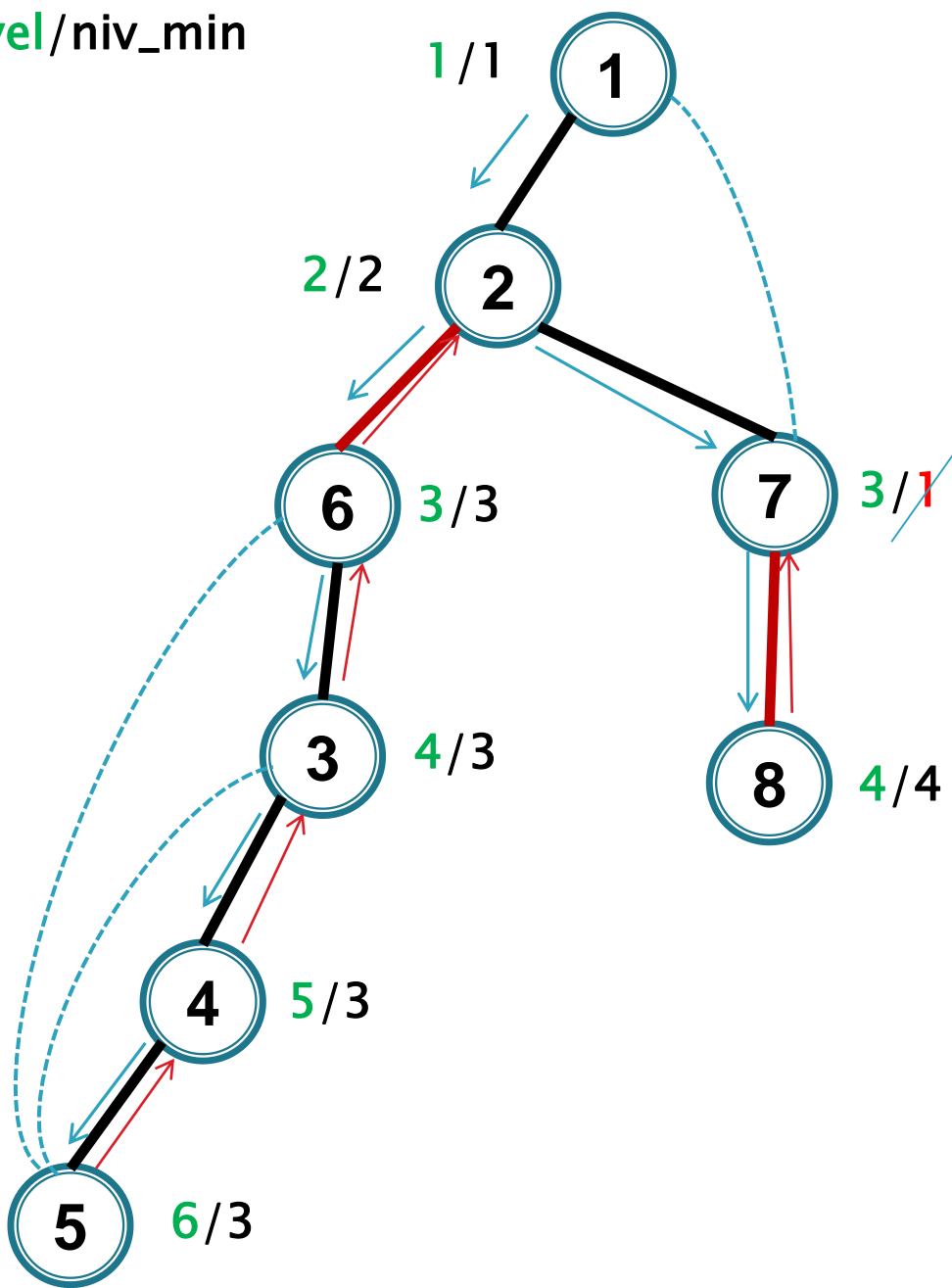
nivel/niv\_min

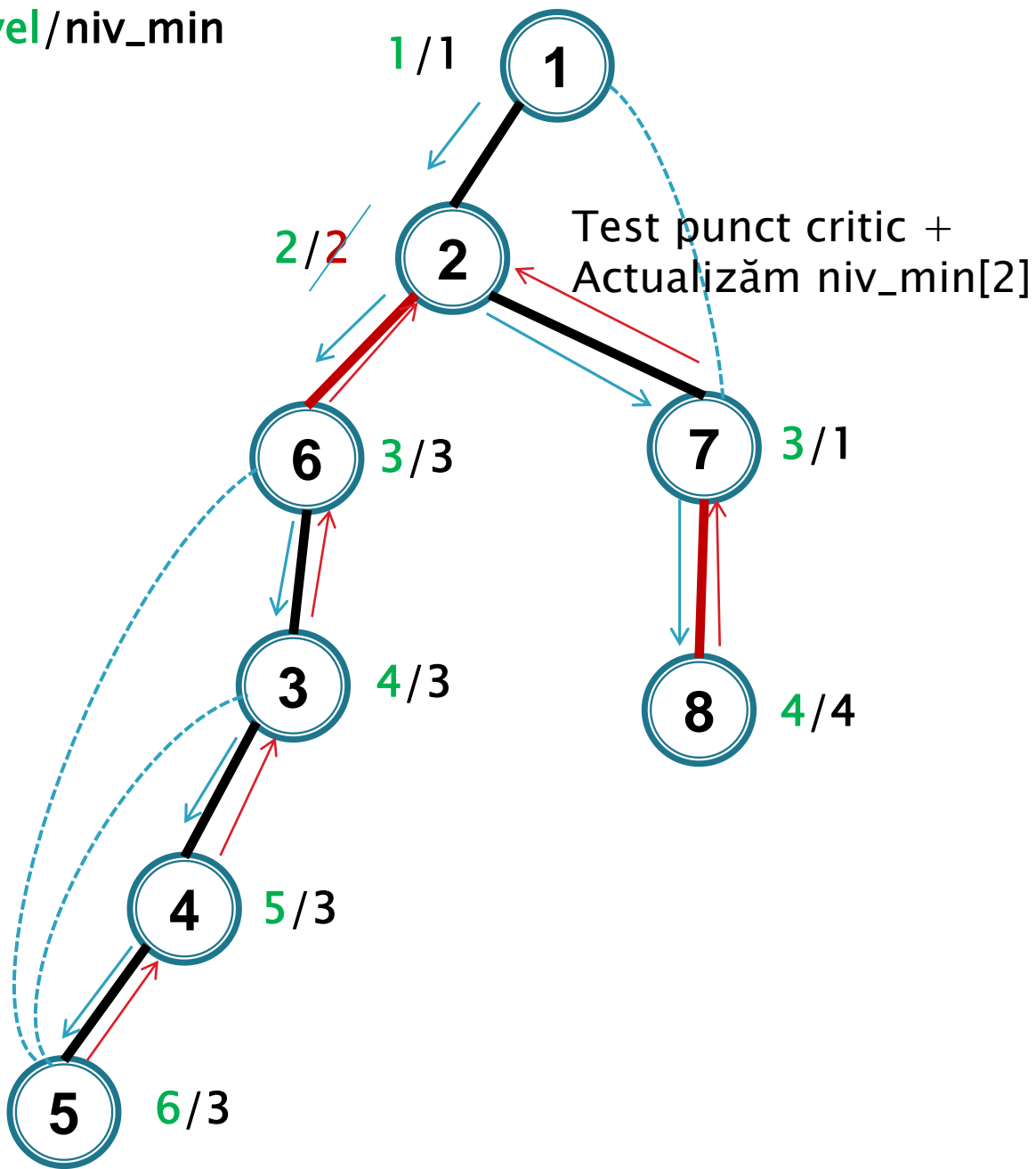
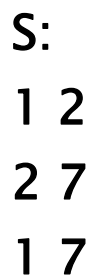




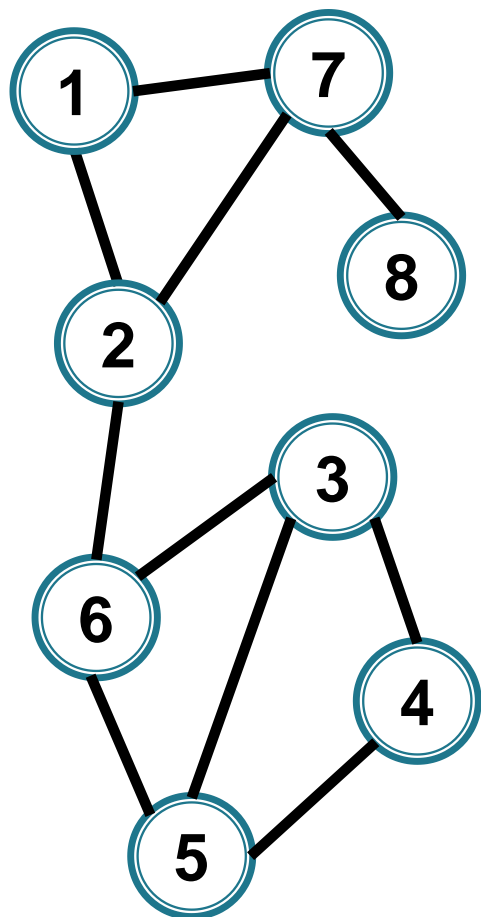
S:  
1 2  
2 7  
1 7

nivel/niv\_min



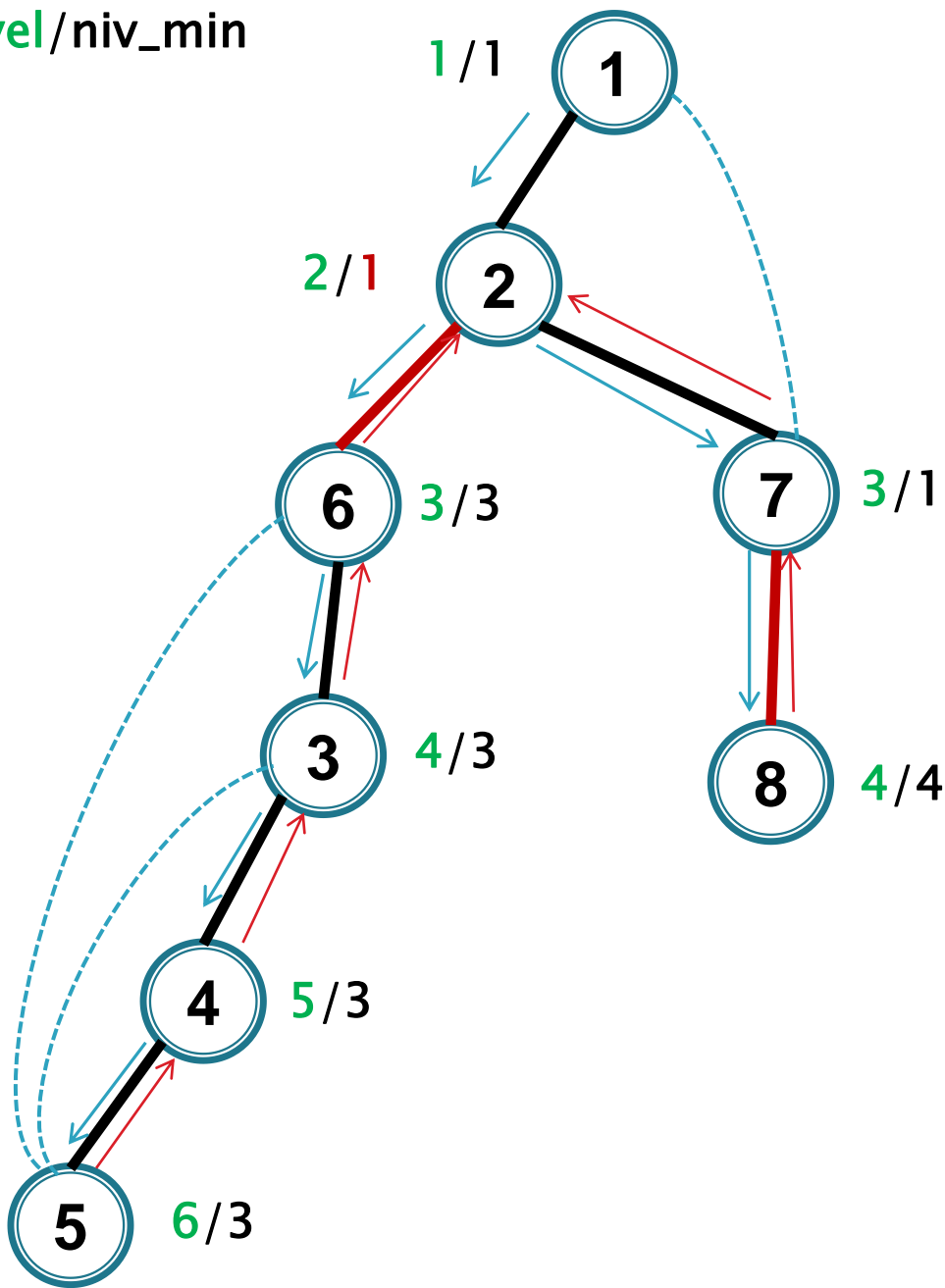


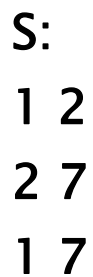




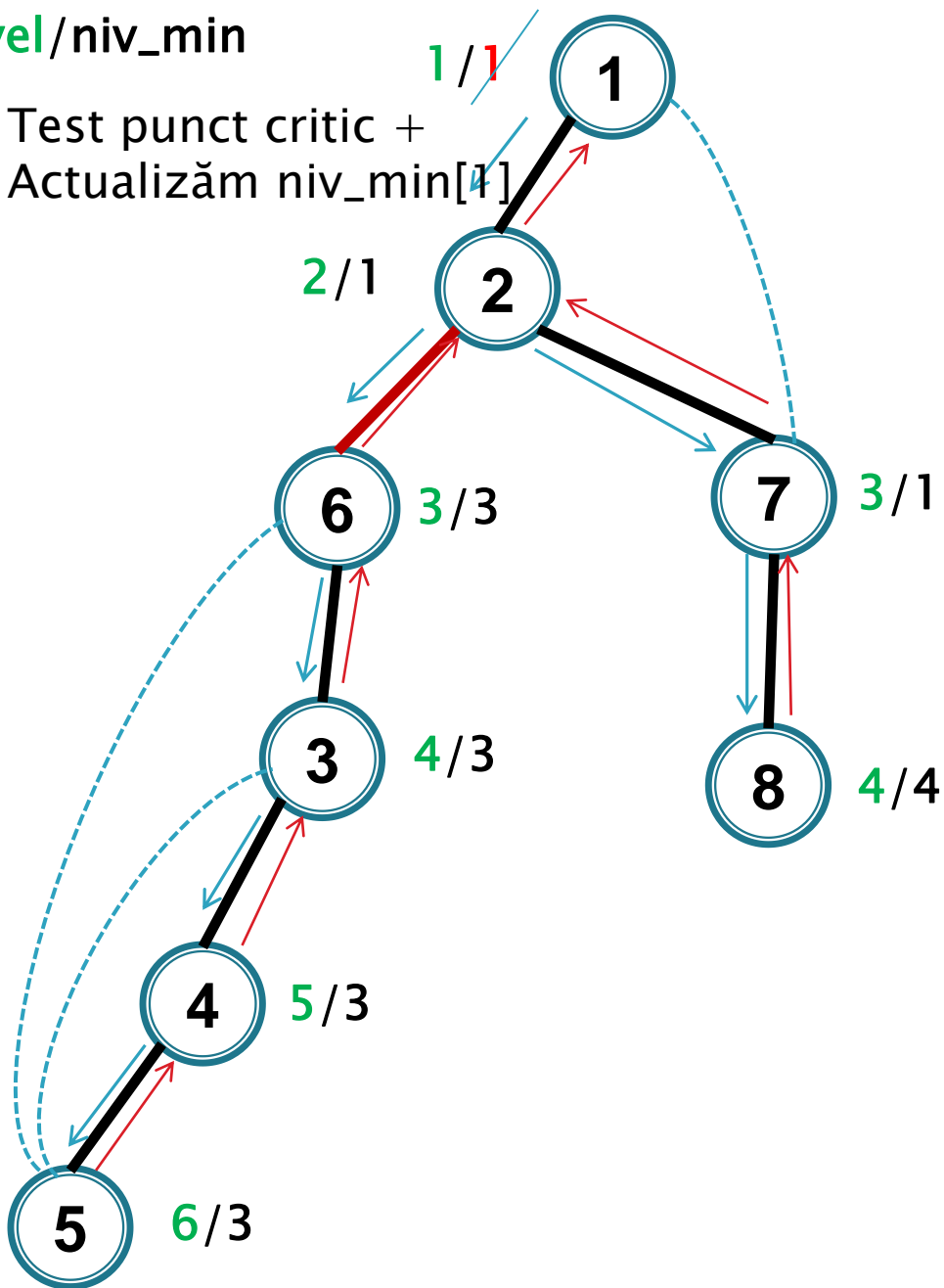
S:  
1 2  
2 7  
1 7

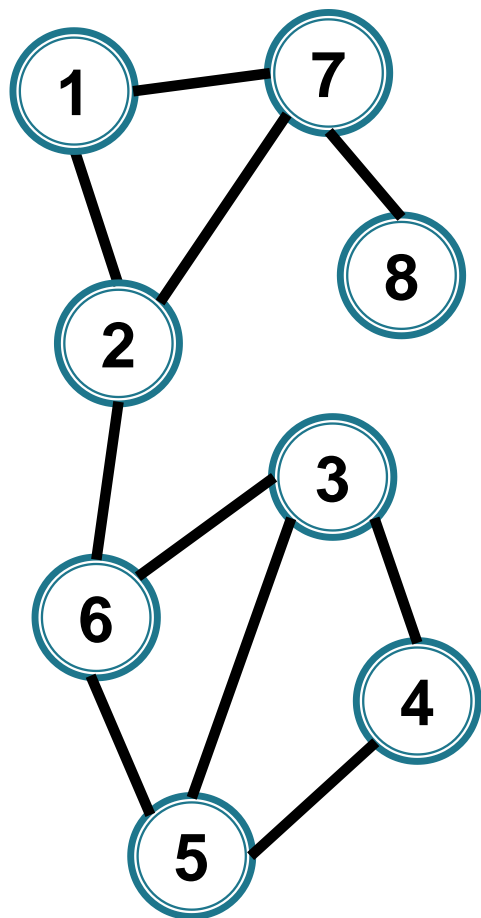
nivel/niv\_min





Test punct critic +  
Actualizăm niv\_min[1]

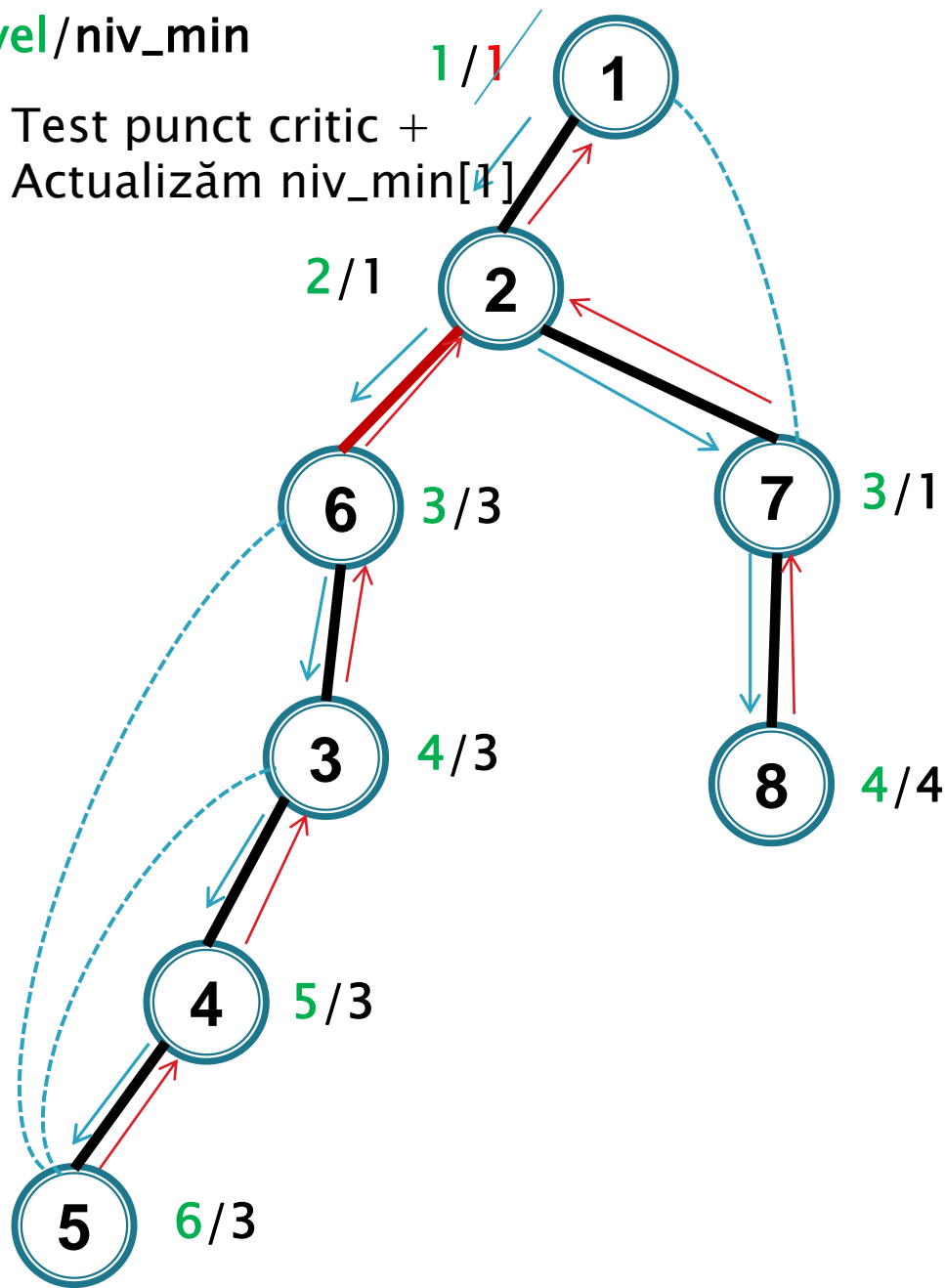




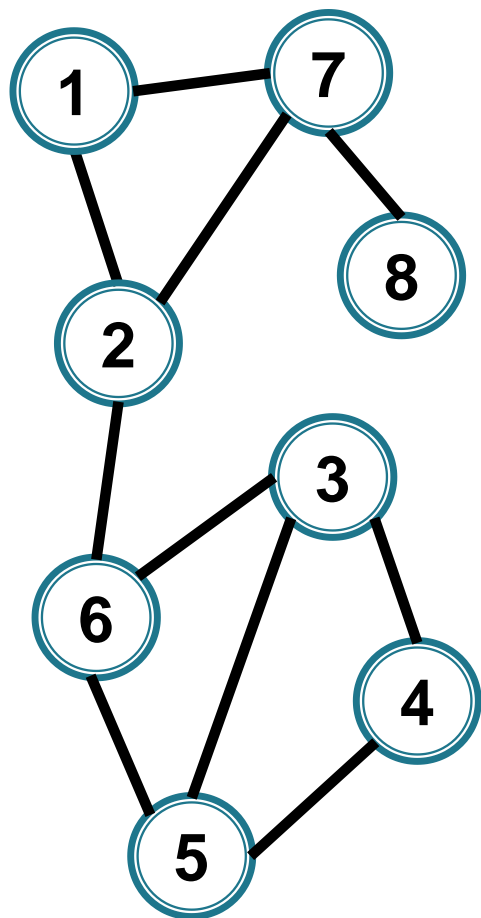
S:  
1 2  
2 7  
1 7

nivel/niv\_min

Test punct critic +  
Actualizăm niv\_min[4]



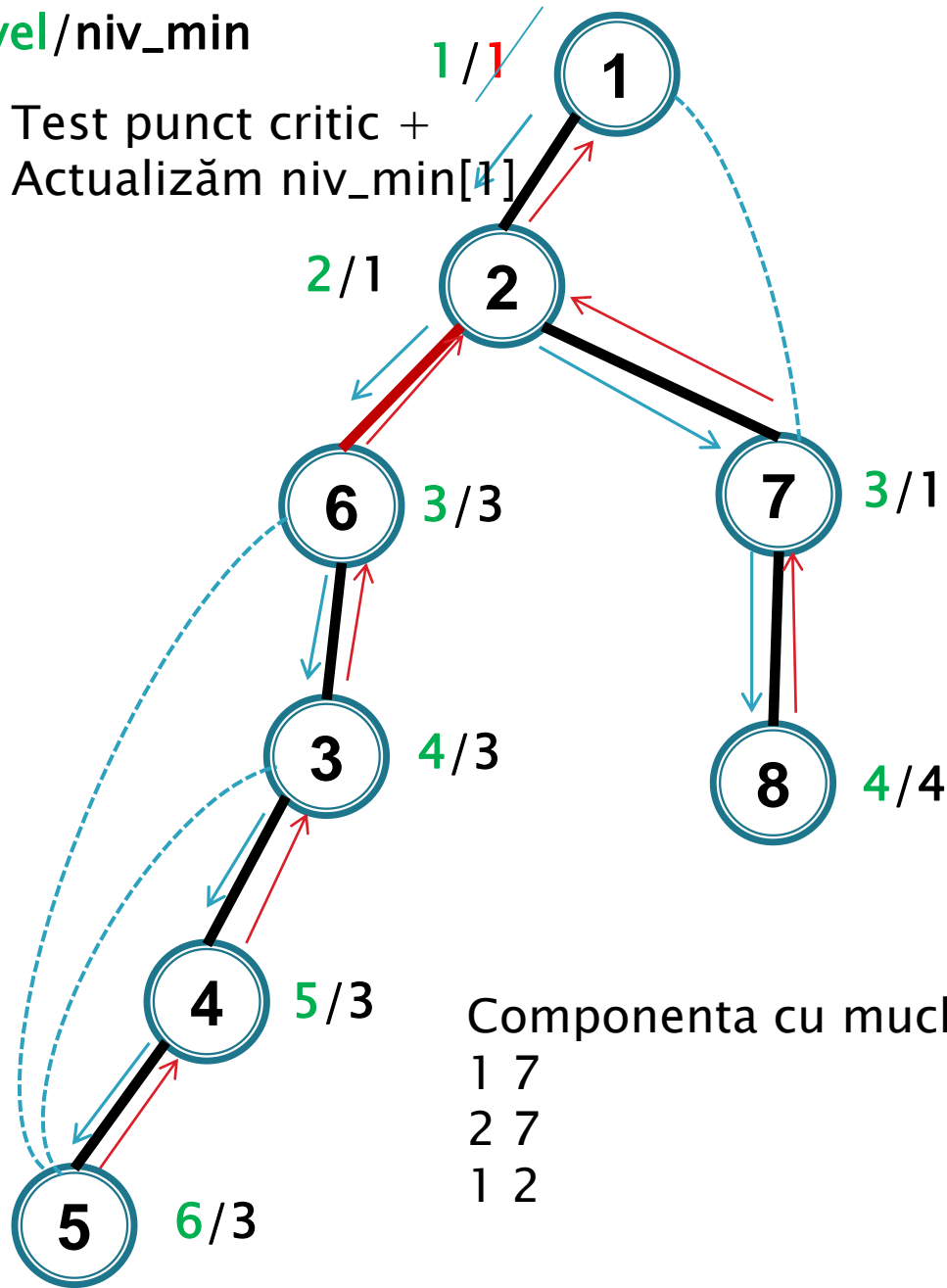
Rădăcina caz particular de  
test de punct critic – aici nu  
mai este necesar, se “rupe”  
o componenta biconexă  
chiar dacă are un singur fiu



S:  
1 2  
2 7  
1 7

nivel/niv\_min

Test punct critic +  
Actualizăm niv\_min[4]



Rădăcina caz particular de  
test de punct critic – aici nu  
mai este necesar, se “rupe”  
o componenta biconexă  
chiar dacă are un singur fiu

Componenta cu muchiile  
1 7  
2 7  
1 2

# Exerciții

# Exerciții

1. Se dă un graf neorientat conex  $G = (V, E)$ . Se consideră parcurgerile DFS și BFS care pornesc din nodul 1, iar vecinii unui nod sunt considerați în ordine crescătoare.

Ce proprietăți trebuie să respecte  $G$  pentru ca cele două parcurgeri să obțină același arbore parțial?

# Exerciții

2. Fie  $G_1$  și  $G_2$  două grafuri cu proprietatea că ordinea în care sunt parcurse vârfurile în BF este aceeași pentru ambele grafuri, la fel și în DF (vecinii unui vârf sunt parcurși în ordine crescătoare).

Sunt  $G_1$  și  $G_2$  egale?

# Exerciții

3. Fie  $G_1$  și  $G_2$  două grafuri cu proprietatea că arborii BF și DF ai celor două grafuri sunt egali (vecinii unui vârf sunt parcurși în ordine crescătoare).

Sunt  $G_1$  și  $G_2$  egale?



# Exerciții

4. Se dă un graf neorientat conex  $G = (V, E)$  și un vârf  $s$ . Care este numărul minim de muchii ale unui graf parțial al lui  $G$  care conservă distanțele de la  $s$  la celelalte vârfuri

# Exerciții

5. Fie  $T$  un graf neorientat cu  $n$  vârfuri.

Arătați că  $T$  este arbore  $\Leftrightarrow T$  este conex și are  $n-1$  muchii

# Exerciții

6. Arătați că în orice grup de  $n$  persoane între care există relații de prietenie reciprocă există două persoane cu același număr de prieteni