

Lecture 6

Last time

- ▷ Gordon's Theorem of alternatives
- ▷ Optimality conditions with functional constraints.

Fenchel conjugate

We will introduce a transformation that plays a key role in duality.

Question: Given a function $f: E \rightarrow \bar{\mathbb{R}}$ and a point $z \in E$, how to find \bar{x} s.t. $z \in \partial f(x)$?

Recall that $g \in \partial f(x)$ iff

$$f(x) - \langle z, \bar{x} \rangle \leq f(x) - \langle z, x \rangle \quad \forall x \in E.$$

This is equivalent to

$$\bar{x} \in \operatorname{argmin}_x f(x) - \langle z, x \rangle$$

$$\bar{x} \in \underset{x}{\operatorname{argmax}} \downarrow \langle z, x \rangle - f(x).$$

This is the core idea behind the following definition.

Def: The Fenchel conjugate of a function $f: E \rightarrow \bar{\mathbb{R}}$ is

Today

- ▷ Fenchel conjugate
- ▷ Fenchel duality

$$f^*(z) = \sup_{x \in E} \langle z, x \rangle - f(x).$$

-1

Lemma: f^* is convex and closed.

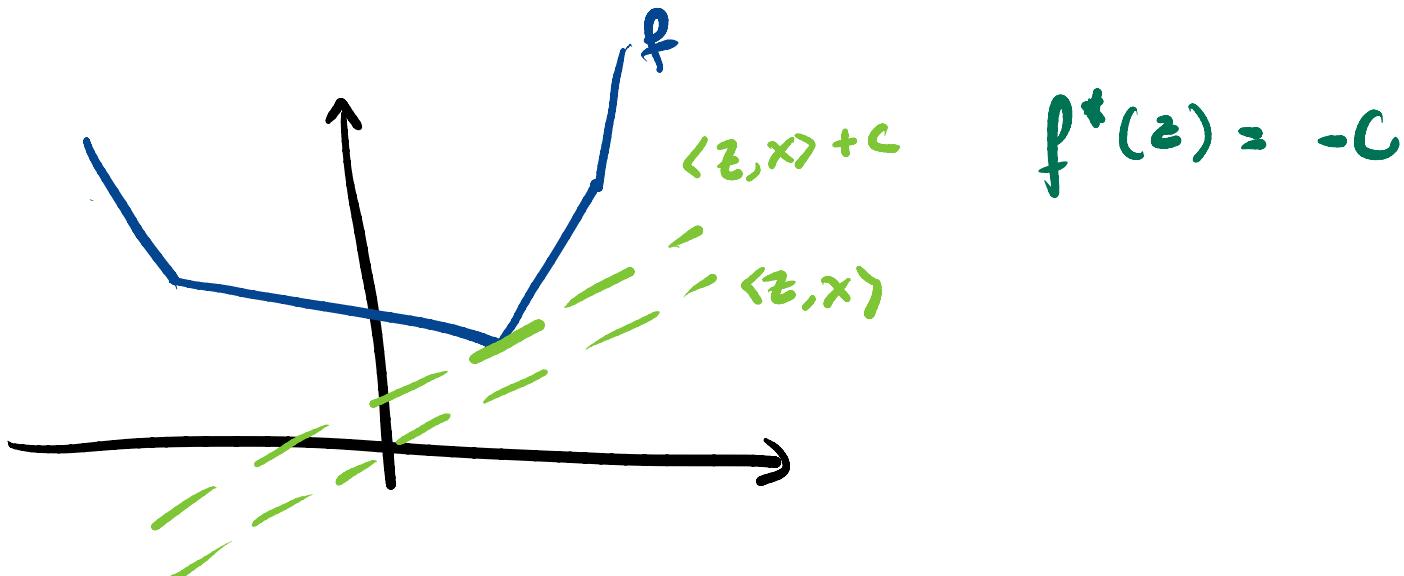
Proof: Note that

$$t \geq f^*(z) \Leftrightarrow t \geq \langle z, x \rangle - f(x) \quad \forall x \in E.$$

Then, $\text{epi } f^*$ is an intersection of halfspaces. \square

Another interpretation

By our derivation, the sup in $f^*(z)$ is attained whenever $z \in \partial f(x)$. Thus, we could compute f^* geometrically by raising the graph of $x \mapsto \langle z, x \rangle + c$ until it is tangent to f .



Proposition: For any $f: E \rightarrow \bar{\mathbb{R}}$ we have that

$$f(x) + f^*(z) \geq \langle z, x \rangle$$

If $f(x)$ is finite, then equality holds iff $z \in \partial f(x)$. →

Proof: Follows easily by our discussion above. □

Examples

▷ Linear functions Suppose $f(x) = \langle c, x \rangle$

Then,

$$\begin{aligned} f^*(z) &= \sup_x \langle z, x \rangle - \langle c, x \rangle \\ &= \sup_x \langle z, -c, x \rangle \\ &= \begin{cases} 0 & \text{if } z = -c \\ +\infty & \text{otherwise.} \end{cases} \\ &= \tau_{fcy}(z). \end{aligned}$$

▷ Indicator of a point $f(x) = \mathbb{1}_{\{c\}}(x)$

$$\begin{aligned} f^*(z) &= \sup_x \langle z, x \rangle - \mathbb{1}_{\{c\}}(x) \\ &= \sup_{x \in \{c\}} \langle z, x \rangle \\ &= \langle c, z \rangle. \end{aligned}$$

We got $f = f^{**}$. We will see this is a recurrent pattern for convex closed functions.

► Indicator of $[-1, 1]$ suppose $f(x) = \mathbb{1}_{[-1, 1]}(x)$

$$\begin{aligned} f^*(z) &= \sup_x zx - \mathbb{1}_{[-1, 1]}(x) \\ &= \sup_{x \in [-1, 1]} zx \\ &= \text{sign}(z) z \\ &= |z|. \end{aligned}$$

► Absolute value $f(x) = |x|$

$$f^*(z) = \sup_x zx - |x|$$

Let's consider two cases if $|z| \leq 1$

$$\Rightarrow \sup_x \underbrace{x(z - \text{sign}(x))}_{\text{always have opposite signs}} \leq 0 \quad \text{achieved with } x=0$$

if $|z| > 1$, then taking any $x = \lambda z$
for any $\lambda > 0$ yields

$$f^*(z) \geq \lambda z^2 - \lambda |z| > \lambda(z^2 - z) > 0$$

Taking $\lambda \uparrow \infty$ gives $f^*(z) = \infty$.

Therefore, $f^*(z) = \mathbb{1}_{[-1, 1]}(x)$.

Once more the two coincide.

► p norms Suppose that $f(x) = \frac{|x|^p}{p}$ $p \geq 1$.

Then

$$f^*(z) = \frac{|z|^q}{q} \quad \text{with} \quad \frac{1}{p} + \frac{1}{q} = 1.$$

(Exercise)

Fenchel duality

Recall that the adjoint A^* of a linear map $A: E \rightarrow Y$ is defined via
Euclidean.

$$\langle Ax, y \rangle = \langle x, A^*y \rangle \quad \forall x \in E, y \in Y.$$

Theorem (Fenchel Duality): For any functions $f: E \rightarrow \bar{\mathbb{R}}$ and $g: Y \rightarrow \bar{\mathbb{R}}$, and linear map $A: E \rightarrow Y$. Define the primal problem

$$p^* = \inf_{x \in E} f(x) + g(Ax)$$

and the dual problem

$$d^* = \sup_{z \in Y} -f^*(A^*z) - g^*(-z).$$

Then, without extra assumptions,

$$p^* \geq d^* \quad (\text{Weak duality}).$$

If further f and g are never $-\infty$ and convex, and

$$0 \in \text{int}(\text{dom } g - A\text{dom } f) := \text{int}\{u - Av \mid u \in \text{dom } g, v \in \text{dom } f\}$$

then, we have

(Constraint qualification)

$$p^* = d^* \quad (\text{Strong duality}).$$

In this case d^* is finite, and it is always attained. Moreover, \bar{x} and \bar{y} are optimal if, and only if,

$$A^*y \in \partial f(x) \text{ and } -\bar{y} \in \partial g(A\bar{x})$$

(Complementary slackness)

Moreover, constraint qualification holds if either

1. There is $\hat{x} \in \text{dom } f$ s.t. $A\hat{x} \in \text{int dom } g$.
2. There is $\hat{x} \in \text{int dom } f$ st. $A\hat{x} \in \text{dom } g$, and A is surjective.

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Proof: Weak duality follows from

Fenchel-Young since $f^*(x, y) \in E^* \times Y^*$

$$\begin{aligned} f(x) + f^*(A^*y) &\geq \langle x, A^*y \rangle \\ + g(Ax) + g^*(-y) &\geq \langle Ax, -y \rangle \end{aligned} \quad (\because)$$

$$f(x) + g(Ax) + f^*(A^*y) + g^*(-y) \geq 0$$

Taking an \inf yields $f^* \geq d^*$.

The key object to prove strong duality is the value function

$$V(z) := \inf_{x \in E} \{f(x) + g(Ax + z)\} \text{ for } z \in Y.$$

So $p^* = v(0)$. Note that if $p = -\infty$, we are strong duality holds trivially by weak duality. Assume that $z \in \text{dom } v$, thus the infimum is not $+\infty$. Therefore, $z \in \text{dom } v$ iff $x \in \text{dom } f(x)$ and $Ax + z \in \text{dom } g$.

Thus, $z \in \text{dom } g - A \text{dom } f$. Hence, constraint generalification is equivalent to $0 \in \text{int dom } v$.

Claim (Exercise): The function v is convex.

Then, since $0 \in \text{int dom } v$, we have $v(0)$ is finite and by HW1 P1c, v is never $-\infty$ in $\text{int dom } f$. Thus, there exists $-y \in \partial v(0)$. Moreover,

$$v(0) + v^*(-y) = \langle 0, -y \rangle = 0.$$

Therefore,

$$\begin{aligned} d^* \leq p^* &= v(0) \\ &= -v^*(-y) \\ &= -\sup_z -\langle y, z \rangle - v(z) \\ &= \inf_z v(z) + \langle y, z \rangle \\ &= \inf_z \inf_x \{f(x) + g(Ax + z)\} + \langle y, z \rangle \\ &= \inf_x \inf_z \{ \underbrace{\langle y, Ax + z \rangle}_{-g^*(-y)} + g(Ax + z) \} + \langle y, -Ax \rangle + f(x) \end{aligned}$$

$$\begin{aligned}
 &= -g^*(-y) + \underbrace{\inf_{x \in \mathbb{R}^n} \langle A^*y, x \rangle}_{-f^*(A^*y)} + f(x) \\
 &= -g^*(-y) - f^*(A^*y) \\
 &\leq d^*
 \end{aligned}$$

Thus, by weak duality $p^* = -g^*(y) - f^*(A^*y) = d^*$. Complementary slackness ensures that the inequalities in (2) become equalities.

Now, let's prove that constraint qualification holds given (1) or (2). If (1) holds

Then, $0 \in \text{dom } g - A \text{dom } f$ and since $A\hat{x} \in \text{int dom } g$ $\exists r > 0$ s.t. $\hat{x} + rB \subseteq \text{dom } g$
 $\subseteq \text{dom } g - A \text{dom } f$, thus $0 \in \text{int dom } g - A \text{dom } f$.

If (2) holds, then we can use the following.

Theorem (Open Mapping) Any surjective mapping $A: E \rightarrow Y$ maps open sets to open sets.

We have that $\exists r > 0$ s.t. $\hat{x} + rB \subseteq \text{dom } f$. Therefore $A\hat{x} + r\hat{A}B$ is an open map contained in $A \text{dom } f$. A similar argument as before, yields that constraint qualification holds. \square

Proof of the Open Mapping Theorem: A is
one-to-one

We start by assuming $\dim Y = \dim E$.

Notice it suffices to show $0 \in \text{int } AB$.

Seeking contradiction assume $0 \notin \text{int } AB$,

$\Rightarrow \exists (x_n) \subseteq B^c$ s.t. $(Ax_n)_n \subseteq (AB)^c$ with $Ax_n \rightarrow 0$. Therefore, $\underline{Ax_n} \rightarrow 0$.

WLOG assume $\frac{x_n}{\|x_n\|} \rightarrow x^*$. $\frac{\|x_n\|}{\text{This}}$ is a contradiction since $Ax^* \neq 0$.

If $\dim E > \dim Y$, we reduce by the previous case. Again assume $0 \notin \text{int } AB$. From Linear Algebra we know

$$A(\ker A)^\perp = \text{range } A \text{ and } A(B \cap (\ker A)^\perp) = AB.$$

We can run the same argument setting

$$E' = (\ker A)^\perp.$$

□