

Lecture 22

Last time

- ▷ Radamacher complexity
- ▷ Polynomial discrimination

Today

- ▷ Vapnik-K-Chervonenkis (VC) Theory.

VC Theory

We will develop another method to certify polynomial discrimination of $\{0, 1\}$ -valued functions \mathcal{F} .

Def: We say that

$$x = \{x_1, \dots, x_n\}$$

is shattered by \mathcal{F} if

$$\#\mathcal{F}(x) = 2^{\#x}$$

The VC dimension is

$$vc(\mathcal{F}) = \sup \{n \in \mathbb{N} \mid \exists x \in X^n \text{ shattered by } \mathcal{F}\}.$$

We will use the following notation, if S is a class of sets and

$$\mathcal{F} = \{\mathbb{1}_A \mid A \in S\},$$

then,

$$S(X) = \mathcal{F}(X) \text{ and } \text{vc}(S) = \text{vc}(\mathcal{F}).$$

Example: Recall the last example from last class

$$S_{\text{one}} = \{(-\infty, t] \mid t \in \mathbb{R}\}$$

then, we have that $\text{vc}(S_{\text{one}}) = 1$.

Similarly,

$$S_{\text{two}} = \{(a, b] \mid a, b \in \mathbb{R}, a < b\}.$$

With $n=2$

| | | |
|-----|---|----|
| | | |
| (] | → | 00 |
| (] | → | 10 |
| (] | → | 01 |
| (] | → | 11 |

But with $n=3$ we cannot form
101.

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Lemma (Sauer and Shelah): Suppose that $d = \text{vc}(S)$. Then for any $X = \{x_1, \dots, x_n\}$,

$$\# S(X) \stackrel{(1)}{\leq} \sum_{i=0}^d \binom{n}{i} \stackrel{(2)}{\leq} (n+1)^{\text{vc}(S)}.$$

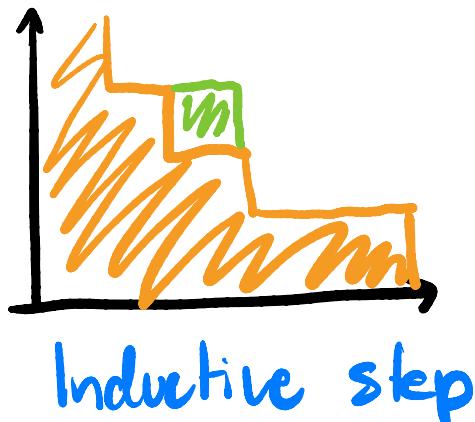
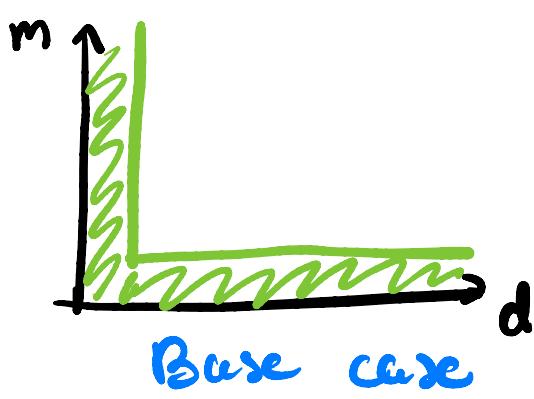
$$\underbrace{\Phi_d(n)}_{\text{Oskn else.}} \quad \binom{n}{k} = \begin{cases} \frac{n!}{k!(n-k)!} & \text{Oskn} \\ 0 & \text{else.} \end{cases}$$

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Thus, $\tilde{\mathcal{F}}$ has polynomial discrimination of order $vc(s)$ and by the results from last lecture we can bound

$$R_n(\mathcal{F}) \leq \sqrt{\frac{2vc(s) \log(n+1)}{n}}.$$

Proof: We prove (1) and leave (2) as an exercise. We will prove (1) using induction on $d+n$



Base case: Assume d arbitrary and $n=0$,

$$\#\tilde{\mathcal{F}}(\emptyset) = 0 \leq 1 = \sum_{i=0}^d \binom{0}{i} = \Phi_d(0)$$

$\Phi_d(\emptyset) = 1$

Assume n arbitrary $d=0$.

$$\#\tilde{\mathcal{F}}(x) = 1 \leq 1 = \binom{n}{0} = \Phi_0(n).$$

Inductive step: Fix d and m and suppose the inequality holds for any d', m' s.t. $d' + m' \leq d + m$. Let X be such that

$$X = \underset{|X|=m}{\operatorname{argmax}} |\mathcal{F}(X)|.$$

Take any $a \in X$, and define

$$X^- = X \setminus \{a\}.$$

Label $H = \mathcal{F}(X) \subseteq \{0, 1\}^{\#S}$, and define

$$H_0 = \{ h|_{X^-} \mid h \in H \text{ and } h(a) = 0 \},$$

$$H_1 = \{ h|_{X^-} \mid h \in H \text{ and } h(a) = 1 \}.$$

Define $H_n = H_0 \cap H_1$ these are binary strings that can be extended to X with both labels.

(Example \downarrow
 $H = \{ \begin{array}{c} 010 \\ 000 \\ 011 \end{array} \} \Rightarrow H_0 = \{ 00, 0 \}, \quad H_1 = \{ 01 \}$, and $H_n = \{ 01 \} \).$

Therefore,

$$\# H = \#(H_0 \cup H_1) + \# H_n.$$

Convince yourself of this!

$$\begin{aligned} &= \# \mathcal{F}(X^-) + \# H_n \quad (\star) \\ \text{Inductive hypothesis} \rightarrow &\leq \Phi_d(n-1) + \# H_n. \end{aligned}$$

Consider now

$$\mathcal{F} = \left\{ g: X^- \rightarrow \{0, 1\} \mid \exists f_1, f_2 \in \mathcal{F} \text{ s.t. } f_1(a) \neq f_2(a), \text{ and } \forall x \in X^-, g(x) = f_1(x) = f(x). \right\}$$

Then

$$H_n \subseteq \mathcal{F}_a(X^-).$$

We claim that $\text{vc}(\mathcal{F}_a) \leq d-1$. Suppose that it is not $\Rightarrow \mathcal{F}_a$ can shatter some $Y \subseteq X^-$ and using the two extensions at a we would have

$$\# \mathcal{F}(Y \cup \{a\}) = 2^{d+1} \quad \downarrow$$

Thus, we obtain

$$\# H_n \leq \# \mathcal{F}_a(X^-) \stackrel{\text{Inductive hypothesis}}{\leq} \Phi_{d-1}(n-1). \quad (\beta)$$

Fact (Pascal's Identity):

$$\binom{n}{i} = \binom{n-1}{i} + \binom{n-1}{i-1} \quad (\because)$$

Number of \uparrow subsets of $[n]$ of size i

$$= \#\{S \subseteq [n-1] \mid \#S = i\}$$

$$+ \#\{S \cup \{n\} \mid S \subseteq [n-1], \#S = i-1\}$$

Then, (★), (◑) and (◐)

$$\begin{aligned} \#\mathcal{F}(X) &\leq \sum_{i=0}^d \binom{n-1}{i} + \sum_{i=0}^{d-1} \binom{n-1}{i} \\ &= \sum_{i=0}^d \binom{n-1}{i} + \binom{n-1}{i-1} \\ &= \sum_{i=0}^d \binom{n}{i}. \end{aligned}$$

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