

# Concurrent Execution II

## Lecture 5

Alma Rahat

CS-210: Concurrency

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Questions and comments: [shorturl.at/knJM5](https://shorturl.at/knJM5)

# What did we do in the last session?

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- We can model a thread's lifecycle using FSP.
- We can use indices in FSP: `store[i:0..3]`.
- We can use guarded actions: `(when B x- > | y -> Q)`.
- We model processes as logically parallel, but in reality they may be interleaved.

## Learning outcomes.

- 1 To explain how processes are interleaved during execution.
- 2 To construct models of parallel processes and shared actions between them.

## Outline.

- 1 Revisit multi-process execution and interleaving.
- 2 Modelling Concurrency.
  - Parallel composition.
  - Labelling.
  - Sharing between processes
    - Actions.

# Process Scheduling (Interleaving) Simulation

If we have one process running on one processor, we have no issues: it runs as you would expect. When we have multiple processes – running in parallel, and potentially each have idle times in-built (e.g. input/output – we need to decide how best to utilise the processors. This is where process scheduling comes to play. There are various algorithms, e.g. first-come-first-serve and round-robin.

## Scheduling Criteria

- CPU utilisation: Is the CPU busy at least most of the time?
- Throughput: What is the number of processes completed per unit time?
- Turnaround time: How much does a particular process take to complete?
- Waiting time: How long does a process queue waiting for its turn?
- Response time: How long does it take to handle interactions?

Simulator: [ess.cs.tu-dortmund.de/Software/AnimOS/CPU-Scheduling/](http://ess.cs.tu-dortmund.de/Software/AnimOS/CPU-Scheduling/)

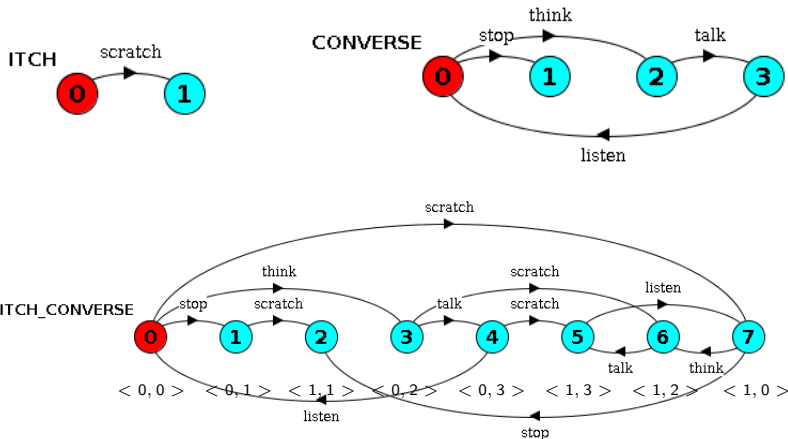
As discussed before, we will not worry about the interleaving: we will write models and codes to scenarios with multiple processes that work in a concurrent manner and correctly irrespective to scheduling.

If  $P$  and  $Q$  are processes then  $(P \parallel Q)$  represents the concurrent execution of  $P$  and  $Q$ . The operator  $\parallel$  is the parallel composition operator.

Consider itching while having a conversation!

```
ITCH = (scratch -> STOP).\\process for itching.  
CONVERSE = (think -> talk -> listen -> CONVERSE |  
stop -> STOP).\\process for conversing.  
||ITCH_CONVERSE = (ITCH || CONVERSE).\\composition of  
two processes.
```

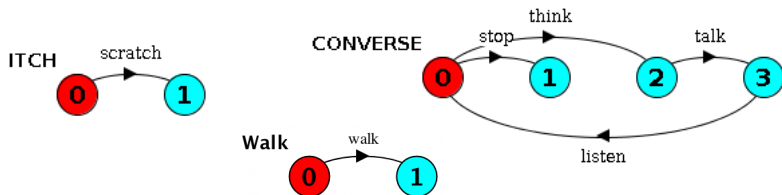
# How many combined states?



Cartesian product of two sets of states:  $\{ \langle i, j \rangle \mid \forall i \in P \wedge \forall j \in Q \wedge P \parallel Q \}$ .  
Therefore, total number of states:  $|P| \times |Q|$ .

ITCH has two states, and CONVERSE has four states: so ITCH || CONVERSE should  $2 \times 4 = 8$  states (all possible combinations of actions.)

# How many combined states?



Exercise: How many states do you think the parallel composition of *ITCH*, *CONVERSE* and *Walk* would have?

Please go to [www.menti.com](http://www.menti.com) and use the code 13 58 67 0.



Commutative  $(P \parallel Q) = (Q \parallel P).$

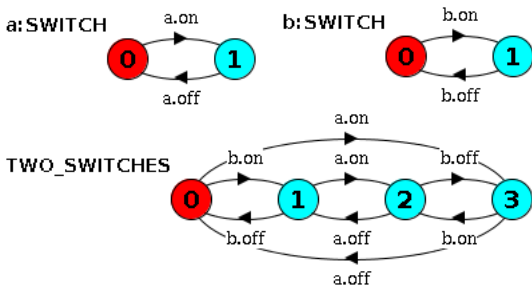
Associative  $((P \parallel Q) \parallel R)$   
 $= (P \parallel (Q \parallel R))$   
 $= (P \parallel Q \parallel R).$

# Process Instances and Labelling

$a:P$  creates an instance of process  $P$  and prefixes each action label in the alphabet of  $P$  with  $a$ .

Two instances of SWITCH process.

```
SWITCH = (on -> off -> SWITCH).  
||TWO_SWITCH = (a:SWITCH || b:SWITCH).
```



What happens if there are no labels?

```
SWITCH = (on -> off -> SWITCH).  
||SWITCHES(N = 3) = (forall[i:1..N] s[i]:SWITCH).
```

```
SWITCH = (on -> off -> SWITCH).  
||SWITCHES(N = 3) = (s[i:1..N]:SWITCH).
```

# Any questions?

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Simple rules to understand the nature of sharing:

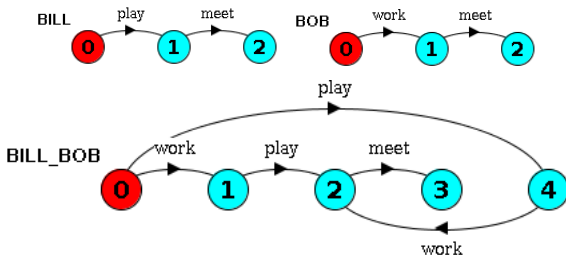
**Action** Common **behaviour** (what the process does; method) across processes, e.g. two users can **play** chess with their computers.

**Remember** You cannot perform a common action until all pre-requisite actions by all sharer have been performed.

**Resource** Common **property** (has-a, uses-a or controls-a type relationship; process or object) across processes, e.g. multiple users at a household have a **toaster**.

So far, we have seen composition of processes with disjoint alphabets: no actions are common between processes, and we can arbitrarily interleave them. In case of shared (or common) actions, we must complete all pre-requisite actions (in any order) before performing a common action.

Scenario: Bill plays tennis and then meets Bob before stopping, while Bob works and then meets Bill before stopping.

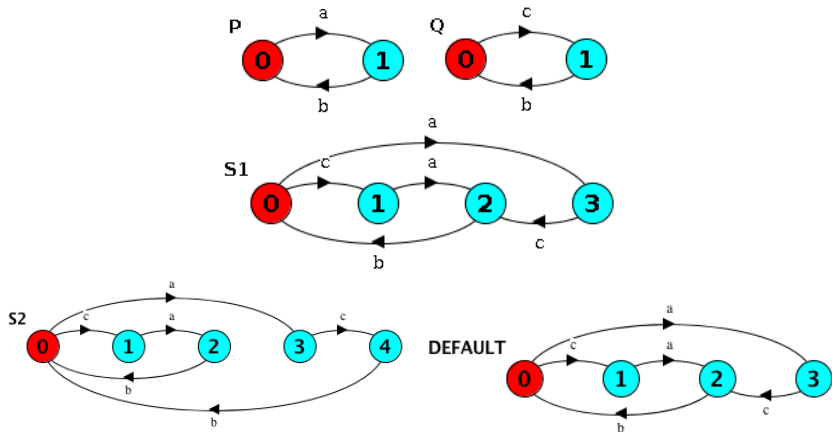


Do you think S1 and S2 describe the same behaviour?

$$P = (a \rightarrow b \rightarrow P).$$
$$Q = (c \rightarrow b \rightarrow Q).$$
$$||S1 = (P || Q).$$
$$S2 = (a \rightarrow c \rightarrow b \rightarrow S2 \mid c \rightarrow a \rightarrow b \rightarrow S2).$$

Hint: Draw the LTS of S1, and see if it is equivalent to S2.

Please go to [www.menti.com](https://www.menti.com) and use the code 10 83 21 7.



$$S2 \equiv \text{DEFAULT}$$

Use Build > Minimise to produce DEFAULT from S2 on the LTSA tool.



- We use  $(P \parallel Q)$  for parallel composition of processes  $P$  and  $Q$ .
- We can model shared actions.