Interference II

Lecture 7

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CS-210: Concurrency

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What did we do in the last session?



- Modelling concurrency: shared resources.
- Action synchronisation.
- Action hiding.
- Started to look at Java code for an ornamental garden.



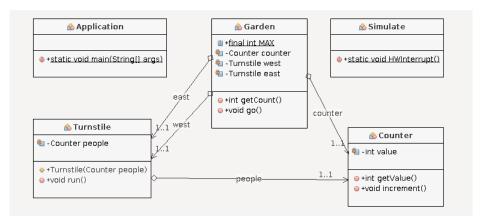
Learning outcomes.

- To compose concurrent processes from a given scenario and code.
- To evaluate the code and create models in FSP.
- To test the model and ensure that the model does what was intended.

Outline.

- 1 Java code for a concurrent scenario.
- ② Create a FSP model.
- Testing the model for correct behaviour.





Application is the main controller program and the Simulate class mimics a hardware interrupt. Full code is in github repository.



```
public class Counter {
    private int value;
    public int getValue(){
        return value;
    public void increment(){
        int temp = value; //read value
        Simulate. HWInterrupt();
        value = temp + 1; // set value
public class Simulate {
    public static void HWInterrupt(){
        if (Math.random()>0.5)
            Thread. yield();
```

Thread.yield() allows the thread to releases control and lets other threads to run.

increment in Counter class reads value, but before incrementing lets the control go so that some other thread can run with 50% probability. It will come back and set the value.



```
@Override
public void run() {
    double randomFactor:
    int waitingTime;
    for (int i=0; i< Garden.MAX; i++){</pre>
        randomFactor = Math.random():
        waitingTime = (int) Math.ceil(randomFactor * 10);
        try {
            Thread.sleep(waitingTime);
        } catch (InterruptedException ex) {
        people.increment();
```

Turnstile class implements the Runnable interface and provides an implementation of run method. After waiting for arbitrary amount of time, it increments its own Counter called people.



```
public void qo() throws InterruptedException{
    counter = new Counter();
    // create Turnstiles
    west = new Turnstile(counter);
    east = new Turnstile(counter);
    // create threads
    Thread westThread = new Thread(west, "west");
    Thread eastThread = new Thread(east, "east"):
    // start threads
    westThread.start();
    eastThread.start();
    // wait for threads to die
    westThread.join();
    eastThread.join();
```

Garden class has a go method that creates a new instance of its Counter, threads for its west and east turnstiles, and starts the threads. Finally, it exits when both threads have finished what they were doing.



```
public static void main(String[] args) {
    Garden garden;
    Scanner scan = new Scanner(System.in);
    while(true){
        garden = new Garden(); // create a garden
        // printout initial count
        System.out.println(garden.getCount());
        try {
            garden.go(); // run the garden
        } catch (InterruptedException ex) {
        // count at the end
        System.out.println(garden.getCount());
        // ask to continue
        System.out.println("Have another go?");
        int choice = scan.nextInt();
        if (choice == 0){
            break:
```

main method displays the count before and after the go method in Garden class is invoked.



Please note: The Simulate class and the user interactions (scan.nextInt() within a while loop) are only here for demonstration purposes. We used these to clarify understanding of how the sequence of events unfold. In a real application, we may not use these unless necessary. A similar argument can be made for the use of sleep within a Thread.



```
run:
0
198
Have another go?
1
0
199
Have another go?
1
0
195
```

Running the program generates arbitrary results for the count at the end, but sometimes it does show the count to be 200. This is a standard example of race condition, and also shows that it may not be easy to identify such errors through formal unit testing.

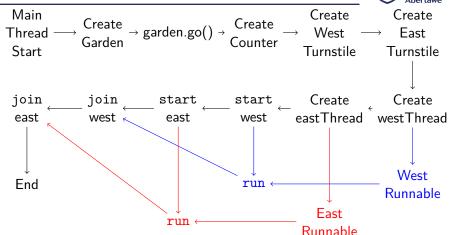


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```

How many threads do we have in the ornamental garden?

Please go to www.menti.com and use the code 11 00 50 4.





Event flow: The main thread creates threads for west and east turnstiles, and invokes respective run methods. The run method in threads are executed in parallel. We will see if shared counter lead to any issues.





What is happening here? The turnstiles do not know that the other exists, and their job is to look at this scoreboard and add 1. Imagine that west looks first and sees the number 4, it goes away to add 1. Now, east looks and sees the number 4, and it goes away to add 1. West has now come back and made the scoreboard 5. East comes back afterwards and makes it 5 too (instead of 6, because it does not know that west has updated it already). No synchronisation between their increment actions.



Shared Objects

The simplest way for processes to interact is to share objects, and the state of the object can be changed using its methods. Such sharing of resources can lead to interference.

Interference or race condition

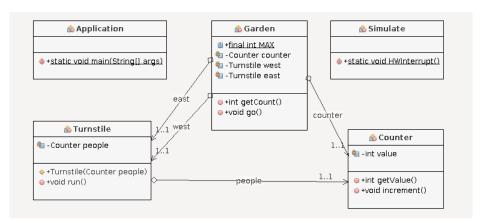
The execution of the instructions from a set of threads can be interleaved in an arbitrary fashion. This interleaving may result in incorrect updates to the state of a shared object. This phenomenon is known as *interference* or *race condition*.

Any questions?









We will not model Application and Simulate. Note that Garden — a composite process — is really bringing together Turnstile and Counter classes to do something useful.



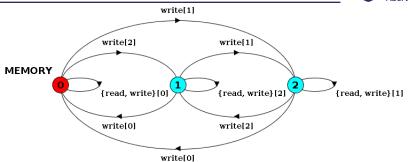
Actions

- arrive
- read
- write
- go
- reset
- return

Processes

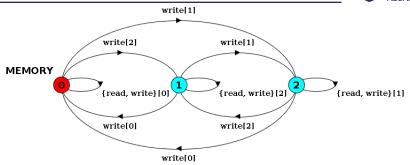
- Memory (Counter)
- Turnstile
- Garden





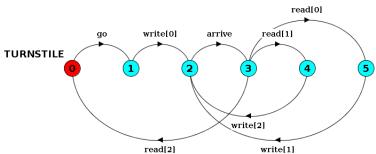
We can consider the Counter class we had to be a piece of Memory. You can read or write, and change the state. Here we consider a Memory model with three states.





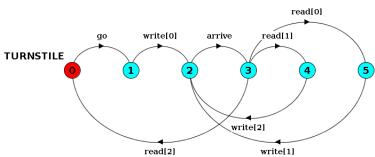
```
const N = 3
T = 0..N
MEMORY = MEMORY[0],
MEMORY[u:T] = (read[u] -> MEMORY[u] | write[v:T] ->
MEMORY[v]).
```





Thinking back the sequence of events in the go method once you called the method were: reset everything (like writing 0 in the memory), and then at every arrive event in a loop, we read a value and increment that value by adding 1 to it (inside the increment method of the Counter class).





```
TURNSTILE = (go -> write[0] -> RUN),
RUN = (arrive -> INCREMENT),
INCREMENT = (read[v:0..N-1] -> write[v+1] -> RUN |
read[N] -> TURNSTILE).
```



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How many states would this have?

So, TURNSTILE has 6 states, and we have two of those in a GARDEN, and then we also share a MEMORY with 3 states. Therefore, the total number of states would be at most: $6^2 \times 3 = 2^8$. Too large for us to spend time on.



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In these circumstances, we will write down the state changes that we expect to happen, and then verify with the animator and alphabets.



The GARDEN supposed to have east and west turnstiles, and a shared counter memory. So, the processes are: east:TURNSTILE, west:TURNSTILE and {east, west}::MEMORY.

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Presumably some of the actions are shared? How do we identify that?



We can look at the process alphabets for the components.

```
\label{eq:process} Process & Alphabet \\ east:TURNSTILE & east.\{\{arrive,go\}, \{read,write\}[0..2]\} \\ west:TURNSTILE & west.\{\{arrive,go\}, \{read,write\}[0..2]\} \\ \{east,west\}::MEMORY & \{east,west\}.\{read,write\}[0..2] \\ \end{cases}
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read/write independent actions, except write[0] and read[N] .

write[0] this is common as it resets the memory.

read[N] this means that we have written up to N and now ready to return.
```





```
||GARDEN = (east:TURNSTILE
|| west:TURNSTILE || {east,
west}::MEMORY/{reset/{east,
west}.write[0],
go/{east, west}.go,
east.return/east.read[N],
west.return/west.read[N].
```

Note that our model is slightly simplistic and partial version of the reality. In the real code east and west check for reaching their own counter up to *N* independently.

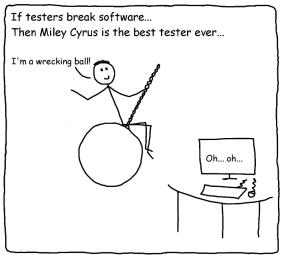
Any questions?





Testing for ERROR





cartoontester.blogspot.com @ 2014

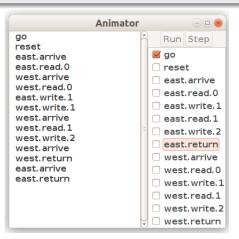
code or FSP we have, it is difficult to say whether race condition is present in it or not. One way we could find out is through animation. But that can be hectic too. Alternative is to write

If we analysed the

test code in FSP.



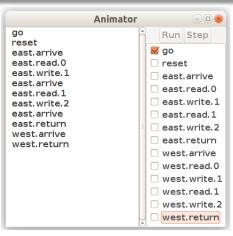
Let's see how we can identify a problem through animation...



Two parallel read[0] followed by two parallel write[1].



How would you test this?

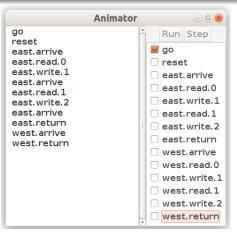


Please go to www.menti.com and use the code $82\ 83\ 79\ 2$

Testing for ERROR



How would you test this?



Please go to www.menti.com and use the code 82 83 79 2

Any questions?







- Arbitrary interleaving on which we have no control can cause interference.
- We model the behaviours and identify where the problem is through FSP and testing.