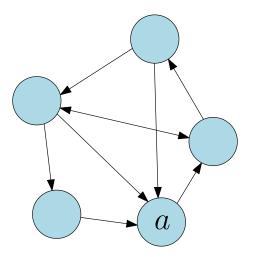
#### Complexity and Data Structures

**Daniel Archambault** 



Graphs connect the world!

CS-115: Complexity

• What differentiates a graph from a tree?

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- What are two graph data structures that implement the ADT?
- What are advantages/disadvantages?

Complexity of algorithms? What about data structures?

# **Complexity and Data Structures**

#### Complexity, Data Structures, and Sorting

- Complexity tries to quantify how fast or big something is
  - given an input size of n how long could this take?
  - given an input size of n how much memory/space do we need?
- We use *n* as the input size because it's variable
- With big n, sometimes a lot of time and space is required!
- In this lecture, we look at worst case (how bad it can get) and average case complexity.

CS-115: Complexity

#### **Rules of Complexity**

- In complexity, we worry only about the largest exponent
- We ignore all constant values and lower exponents
  - $ightharpoonup 2n^3 + 32n^2 + 4 \text{ is } O(n^3)$
- Why can we do this? As n becomes really large, other terms don't matter so much.

#### **Space Complexity**

- You have already done time complexity? Think sorting...
- Well, you can use the same tool to talk about how much space something takes in memory.
  - ▶ Given an input size of n, it's possible for it to take  $O(n^2)$  space
- Important for data structures
  - time complexity for the operations on the data
  - space complexity for the data itself

### Warning!

- In this lecture we estimate complexity
- We are providing estimates not proofs
- Complexity theory is hard and involves proofs
  - particular algorithm has performance...
  - any algorithm for a problem will have performance...
- Our arguments are not proofs, but estimates
- In second year, you will do rigourous proofs.

```
int[] a = new int [n];
```

• Insertion complexity?

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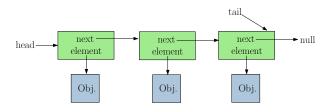
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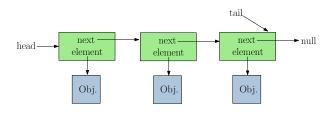
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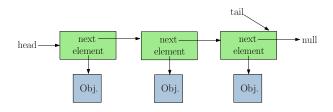
- Insertion complexity? O(n)
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- Space complexity? O(n)... usually



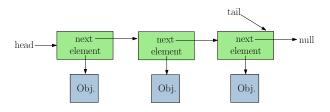
• Insertion complexity (if location known)?



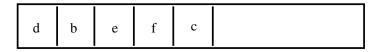
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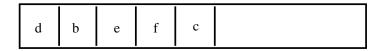
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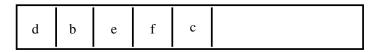
• isEmpty complexity?



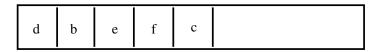
- isEmpty complexity? O(1)
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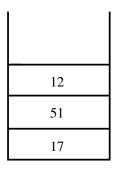
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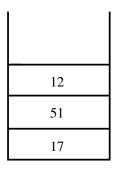
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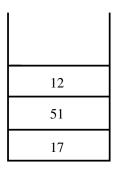
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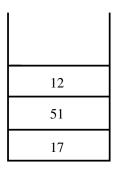
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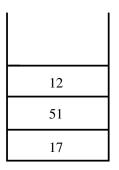


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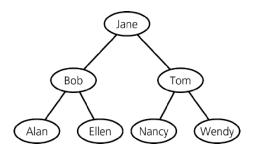
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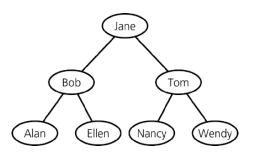


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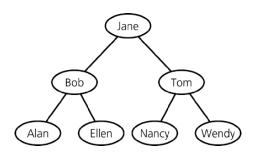




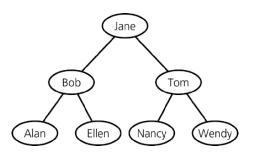
• Find a element?



- Find a element? can be O(n) but usually  $O(\lg(n))$ 
  - a balanced binary tree has maximum height of O(lg(n))
  - ▶  $O(\lg(n))$  log base 2 of n
- Print out in alphabetical order?

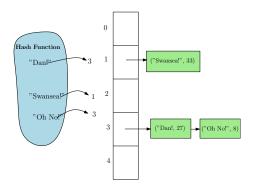


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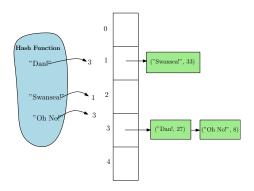
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# Hash Maps



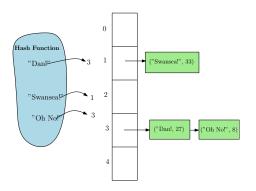
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# Hash Maps



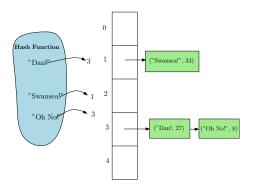
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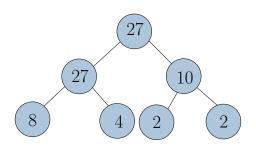


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- Space complexity?

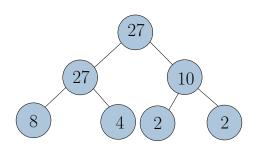
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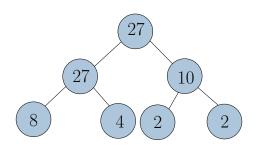
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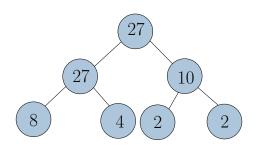
• Look at the top?



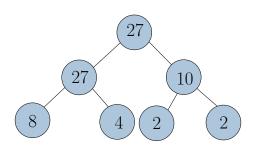
- Look at the top? O(1)
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- Look at the top? O(1)
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Edge[][] matrix; //nxn matrix specified by constructor
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- Space complexity?  $O(n^2)$

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Node[] nodes;
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public class Node {
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### Complexity of Comparison-Based Sorting

- All sorting algorithms currently are comparison based
  - general sorting algorithms that involves knowing nothing about data
- We know that there are algorithms that take  $O(n \lg n)$  comparisons
- But, surely, there are faster algorithms, Dan?

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- But, surely, there are faster algorithms, Dan?
- For comparison-based sorts, no.
  - ▶ mathematical proof says you need at least  $\Omega(n \lg n)$  comparisons
  - proof is hard, but we can do it for fun sometime...
- However, if you know something about the data, you can do better!

### Bucket Sort: *O*(*n*)

- If you know that your data can be easily divided into b buckets, you can use bucket sort which is O(n).
- The algorithm:
  - Create an array of b elements. Each stores an array or linked list.
  - Traverse all n unsorted elements, throwing each into their correct bucket.
  - Assuming each bucket has a small number of elements, sort by any means.
- If any particular bucket ends up with n elements it is  $O(n^2)$
- Bucket sort is really useful, especially with hashmaps
  - I've personally used this strategy many times.

CS-115: Complexity

# Radix Sort: O(n)

- If you know that each element of your data has a key of b digits, you can use radix sort which is O(bn).
- Algorithm:
  - start with the least significant digit and apply bucket sort
  - loop through all digits up to the greatest digit
- Assumption is good bucket sorts and small number lengths.
- What to do with non-integer keys...
- Still much debate about its complexity