# Introduction to Game Theory 6 Shapley Value for Coalitional Games

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#### Reading

- Recommended
  - Shoham and Leyton-Brown: Chapter 12, sections 12.1-12.2

#### Overview

#### Coalitional Game Theory

- Basic modelling unit is group rather than indvidual agent.
- Transferable vs. non-transferable utility
- Coalitional game with transferable utility (N, v):
  - *N* finite set of players:
  - $v: 2^N \longrightarrow \mathbb{R}$  pay-off function  $(v(\emptyset) = 0)$
- Fundamental questions:
  - Which coalitions will form?
  - How should coalition divide its pay-off among its members?

#### Classes of coalitional games

Super-additive game ("synergy")
 Game (N, v) is super-additive iff

$$\forall S, T \subset N : S \cap T = \emptyset \Longrightarrow v(S \cup T) \ge v(S) + v(T).$$

In particular:  $v(S \cup i) \ge v(S) + v(i)$  for any  $S \subset N \setminus \{i\}$ .

 As a consequence, for super-additive game, the grand coalition has the highest pay-off of all coalitional structures:

$$v(N) = v(S \cup S^c) \ge v(S) + v(S^c) \ge v(S).$$

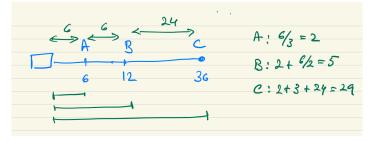
 Therefore focus on a fair redistribution of total pay-off among the members of the grand coalition.

## Motivating example: Ways to allocate common benefits

- Three friends sharing a taxi cab
- Splitting the profit from an investment

## Worked example: Fair division of taxi fare

- Alice, Bob and Charlize share a taxi to go home;
- The individual fares would be: A(6), B(12) and C(36);
- If they share the cab then they only have to pay the fare to the farthest destination (C = 36).
- What would be a fair way to share the fare?



**Sanity check:** 2+5+29=36, and everyone is better off!

#### Worked example: Fair division of taxi fare

- Consider a sequential version of the problem in which A, B and C arrive in random order, and pay whatever is lacking (i.e. their marginal contribution);
- Permutation ACB indicates that coalition grows as follows:

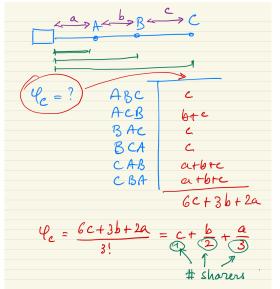
$$A \rightarrow AC \rightarrow ACB$$

- 1. When A joins, he pays the fare to his destination: 6
- 2. When C joins, he pays the remainder to get to C: 36-6=24;
- 3. Finally, when B joins, everything is already paid for.

## Worked example: Fair division of taxi fare

A B				
marginal Contribution	A	$\mathcal{B}$	С	
ABC	G	6,	24	4
ACB	6	0	30	5
BAC	0	12	24	4
BCA	0	12	24	4
CAB	0	0	36	
CBA	0 (	0 (	36	6
mean	12 (2)	30 5	174	29

#### Shapley: common sense vs. permutation solution



#### Shapley value: Alternative definition

- Marginal contribution only depends on what precedes a contributor;
- Marginal contribution of player *i* to subset *S*:

$$\delta_i(S) = v(S \cup i) - v(S)$$

• Shapley value of player *i*: (denoting #N = n, #S = s)

$$\varphi_i(N, v) := \frac{1}{n} \sum_{S \subset N \setminus i} {n-1 \choose s}^{-1} \delta_i(S)$$

Amplification: see next slides!

- We focus on Shapley value  $\varphi_i(N, v)$  for agent i;
- For any existing coalition S not including i, i.e.

$$S \subset N_i := N \setminus i$$

we consider the value increment due to *i* joining:

$$\delta_i(S) = v(S \cup i) - v(S)$$

• The size s := #S of the possible coalitions S that i joins, can range between  $0 \le s \le n-1$ .

• For fixed coalition size s there are

$$N_s := \binom{n-1}{s}$$

coalitions S of that size.

• Hence, the mean contribution  $\overline{\Delta}_i$  of i to existing coalitions S of size s is given by:

$$\overline{\Delta}_i(s) := \frac{1}{N_s} \sum_{S: \#S = s} \delta_i(S) = \binom{n-1}{s}^{-1} \sum_{S: \#S = s} \delta_i(S).$$

• Finally, since  $0 \le s \le n-1$  we compute the average over the n possible choices of s. This average is the Shapley value:

$$\varphi_{i} := \frac{1}{n} \sum_{s=0}^{n-1} \overline{\Delta}_{i}(s) = \frac{1}{n} \sum_{s=0}^{n-1} \binom{n-1}{s}^{-1} \sum_{S:\#S=s} \delta_{i}(S)$$

$$= \frac{1}{n} \sum_{s=0}^{n-1} \sum_{S:\#S=s} \binom{n-1}{s}^{-1} \delta_{i}(S)$$

$$= \frac{1}{n} \sum_{S=0}^{n-1} \binom{n-1}{s}^{-1} \delta_{i}(S)$$

• Double sum above is actually sum over all subsets  $S \subset N \setminus i$ .

Recall:

$$\binom{n}{k} = \frac{n!}{k!(n-k)!} \quad \text{where} \quad n! = n(n-1)(n-2)\dots 3\cdot 2\cdot 1.$$

Hence:

$$\varphi_{i} = \frac{1}{n} \sum_{S \subset N \setminus i} {n-1 \choose s}^{-1} \delta_{i}(S)$$

$$= \frac{1}{n} \sum_{S \subset N \setminus i} \frac{s!(n-1-s)!}{(n-1)!} \delta_{i}(S)$$

$$= \frac{1}{n!} \sum_{S \subset N \setminus i} s!(n-1-s)! \delta_{i}(S)$$

Claim: this corresponds to the permutation definition!

## Shapley: equivalence of definitions

Permutation formula:

$$\varphi_{i} = \frac{1}{n!} \sum_{S \in S} s! (n-1-s)! S_{i}(S) (s=\#,S)$$
 $\pi_{1} = 24315 \rightarrow S = \{2,4\}$ 
 $S = \{1,5\}$ 
 $\pi_{2} = 43215 \rightarrow S = \{4\}$ 
 $S = \{1,25\}$ 

#### Shapley: equivalence of definitions

**Recall:**  $\delta_i(S)$  depends on the **set** S, not on the permutation sequence of the elements in S!

- Hence,  $\delta_3(S)$  has the same value in  $\pi_1 = 24315$  and  $\pi_2 = 42351$  as in both cases  $S = \{2, 4\}$ .
- Consider an arbitrary permutation of 1, 2, ..., n and let's focus on i somewhere in the sequence, all numbers that appear to the left of i, constitute the set S. Likewise, all elements that appear to the right, are collected in the set S<sup>c</sup>:

$$\pi = \underbrace{* * * \dots * * *}_{S} i \underbrace{* * * \dots * * *}_{S^{c}}$$

• Any permutation of the *S*-elements in  $\pi$  yields the same value  $\delta_i(S)$  (see above), There are s! such permutations.

#### Shapley: equivalence of definitions

- Likewise,  $\delta_i(S)$  does **not depend** on the elements in  $S^c$ . Any permutation of these elements in  $S^c$  yields that same value  $\delta_i(S)$ . There are (n-1-s)! such permutations.
- Hence, we can conclude that of the n! permutations of  $\{1,2,\ldots,n\}$ , a total of s!(n-1-s)! give rise to the same value  $\delta_i(S)$ , which only depends on the **set** S.
- Averaging over all possible choices for  $S \subset N_i \equiv N \setminus i$  yields:

$$\frac{1}{n!} \sum_{S \subset N_i} s!(n-1-s)! \, \delta_i(S) = \phi_i(i)$$

## Shapley's Axioms: Some useful terminology

 Players i and j are interchangeable if their contributions to every coalition (subset) S is exactly the same:

$$\forall S \subset N \setminus \{i,j\} : \quad v(S \cup i) = v(S \cup j)$$

 A player i is a dummy player if the amount he contributes to any coalition is exactly the amount he's able to achieve alone:

$$\forall S \subset N \setminus \{i\} : v(S \cup i) = v(S) + v(i)$$

#### Shapley's Axioms

• **Symmetry:** If *i* and *j* are interchangeable then:

$$\psi_i(N, v) = \psi_j(N, v).$$

 Dummy Player: will only get what he can achieve on his own:

$$\psi_i(N,v)=v(i).$$

• Additivity: Consider two games  $G_1 = (N, v)$ ,  $G_2 = (N, w)$  and assume that we play  $G_1$  with probability p and  $G_2$  with prob q = 1 - p. Then

$$\psi_i(N, pv + qw) = p\psi_i(N, v) + q\psi_i(N, w)$$

#### Shapley's theorem

#### **Shapley (1951)**

Given a coalitional game (N, v), the Shapley values  $\varphi_i$ , i = 1, ..., n specifies the unique distribution of the total value v(N) that is both

- efficient, i.e.  $\sum_i \varphi_i = v(N)$
- satisfies Shapley's axioms,
   i.e. Symmetry, Dummy Player and Additivity.

#### Shapley value: worked example

An AI expert (E) developed a powerful new algorithm. However, in order to implement his ideas, he needs to create a startup and hire a programmer (P) for 2 years. An angel investor (A) provides funding. The value that each coalition of these three stakeholders (E, P, A) can generate satisfies the following rules:

- Without both investor and expert, no value can be generated.
- If he has no assistance from a programmer, the expert's value equals 3, but if he can delegate the programming and focus on R&D, his value rises to 10.
- The value created by the programmer is 5. This is in addition to the rise in value of the expert.

The startup is sold to a large software company for serious money. How to split this money fairly among the three stakeholder?

## Shapley value: Method 1

Shapley value computation: 
$$\#N=n=3$$
,  $\#S=s$ 

$$\begin{cases}
P_i = \frac{1}{n} \sum_{S \in \mathbb{N}} \binom{N-1}{s}^{-1} \delta_i(S) \\
Fxput (E)
\end{cases}$$

$$S=0 \longrightarrow S=\emptyset \Rightarrow \delta_E(S) = \overline{v}(E) - \overline{v}(\phi) = 0.$$

$$L \Rightarrow \binom{n-1}{s} = \binom{2}{0} = 1$$

$$S=1 \longrightarrow \delta_E(A) = \overline{v}(AE) - \overline{v}(A) = 3$$

$$\delta_E(P) = \overline{v}(EP) - \overline{v}(P) = 0$$

$$\begin{pmatrix} N-1 \\ s \end{pmatrix} = \binom{1}{1} = 2$$

$$S=2 \longrightarrow \delta_E(AP) = \overline{v}(APE) - \overline{v}(AP) = 15$$

$$L \Rightarrow \binom{n-1}{s} = \binom{1}{2} = 1$$

$$\mathcal{Q}_E = \frac{1}{3} \left[ \frac{1}{1} \cdot 0 + \frac{1}{2} \cdot 3 + \frac{1}{4} \cdot 15 \right] = \frac{11}{2}$$

## Shapley value: Method 2

Figure: Notice distribution is efficient: 
$$\varphi_E + \varphi_A + \varphi_P = 15 = v(N)$$