

## Knowledge Representation

2023/2024

### Exercise Sheet 1 – Classical and Description Logics

6th November 2023

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**Exercise 1.1** Consider an interpretation  $I$  with

$$I(a) = \text{false} \quad I(b) = \text{false} \quad I(c) = \text{false} \quad I(d) = \text{false}$$

Which of the following propositional formulas are satisfied by this interpretation:

- |  |  |
|--|--|
| (a) $(a \wedge b) \vee \neg c \vee \neg d$ | (c) $(a \rightarrow \neg b) \vee (\neg c \rightarrow d)$   |
| (b) $(a \wedge b) \vee (\neg c \wedge d)$  | (d) $(\neg a \rightarrow b) \wedge (c \rightarrow \neg d)$ |

**Exercise 1.2** Consider the two formulas

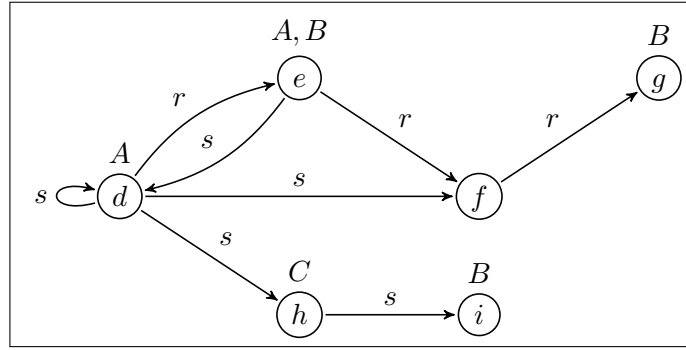
$$F = \neg p \wedge \neg q \quad G = p \rightarrow q$$

- (a) Which truth assignments to  $p$  and  $q$  give models of  $F$  and  $G$ ?
- (b) Does  $F$  entail  $G$ ?
- (c) Does  $G$  entail  $F$ ?

**Exercise 1.3** Assume that we use propositional logic for knowledge representation and that the reasoning problem we have to solve is the following: given a formula  $F$ , decide whether  $F$  is satisfiable. Which of the following algorithms is sound/complete/terminating?

- (a) Always return “yes”.
- (b) Always return “no”.
- (c) Enter an infinite loop, never return.
- (d) Go through all truth assignments for the variables in  $F$  one after the other. For each truth assignment, check whether it gives a model for  $F$ . If a satisfying truth assignment is found, return “yes”. Otherwise return “no”.

**Exercise 1.4** Consider the interpretation  $\mathcal{I}$  represented as the following graph:



For each of the following concepts  $D$ , write down the elements in their interpretation  $D^{\mathcal{I}}$ :

- |                       |                              |                                  |
|-----------------------|------------------------------|----------------------------------|
| (a) $\neg A$          | (c) $\exists r.(A \sqcup B)$ | (e) $\exists s.\exists s.\neg A$ |
| (b) $A \sqcap \neg B$ | (d) $\forall r.(A \sqcup B)$ | (f) $\exists r.\forall r.\neg A$ |

**Exercise 1.5**

- (a) Express the following phrases using  $\mathcal{ALC}$  concepts:
- (i) "persons that do not have a friendly neighbour,"
  - (ii) "persons that have a neighbour that is not friendly."
- (b) Construct an interpretation where an element satisfies the concept for (i), but not the concept for (ii).
- (c) Construct an interpretation where an element satisfies the concept for (ii), but not the concept for (i).

**Exercise 1.6** Which of the following  $\mathcal{ALC}$  axioms corresponds to the following English statement:

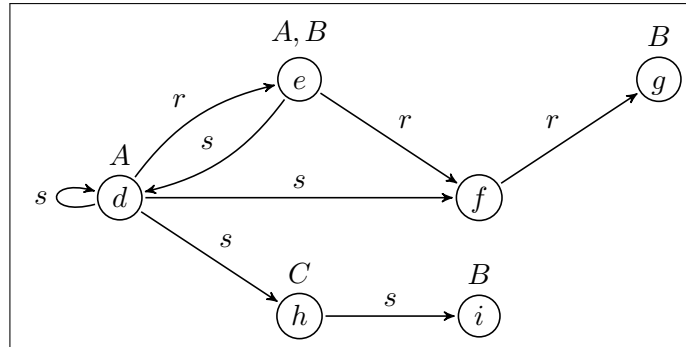
- "A Mule is an animal that has a horse and a donkey as a parent."

- (a)  $\text{Mule} \equiv \text{Animal} \sqcap \exists \text{hasParent} . (\text{Horse} \sqcap \text{Donkey})$
- (b)  $\text{Mule} \equiv \text{Animal} \sqcap \exists \text{hasParent} . \text{Horse} \sqcap \exists \text{hasParent} . \text{Donkey}$
- (c)  $\text{Mule} \equiv \text{Animal} \sqcap (\exists \text{hasParent} . \text{Horse} \sqcup \exists \text{hasParent} . \text{Donkey})$

**Exercise 1.7** Translate the following axiom into English:

$$\text{KRTeacher} \equiv \exists \text{teaches} . (\text{Course} \sqcap \forall \text{hasTopic} . (\text{DL} \sqcup \text{Arg} \sqcup \text{PGM}))$$

**Exercise 1.8** We consider the same graph from Exercise 1.4, but with a different reading: This time, we see it as an *ABox*  $\mathcal{A}$ , where the nodes are individual names, i.e.  $d, e, f, g, h, i \in \mathbf{I}$ , and the labels correspond to concept and role assertions, e.g.  $e : A$ ,  $g : B$ ,  $(d, e) : r$ , etc.



We also consider again the following set of concepts:

$$S = \{ \neg A, A \sqcap \neg B, \exists r.(A \sqcup B), \forall r.(A \sqcup B), \exists s.\exists s.\neg A, \exists r.\forall r.\neg A \}.$$

- (a) For each concept  $D \in S$ , what are the instances of  $D$  w.r.t.  $\mathcal{A}$ , i.e. for which individual names  $x \in \{d, e, f, g, h, i\}$  does  $\mathcal{A} \models x : D$  hold?
- (b) For each concept  $D \in S$ , what are the instances of  $D$  w.r.t. the ontology  $\mathcal{O} = \mathcal{A} \cup \mathcal{T}$ , where  $\mathcal{T} = \{C \sqsubseteq \neg A, B \sqsubseteq \forall r.C\}$ ?