Introduction to Game Theory 5 Iterated Games

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Reading

Recommended

Shoham and Leyton-Brown: Chapter 5, sections 5.1

Optional

- William Spaniel: Game Theory 101: The Complete Textbook (paperback): Very accessible and clear, teaching through examples. Accompanying YouTube channel.
- N. Nisan, T. Roughgarden, E. Tardos, V. Vazirani: Algorithmic Game Theory. Cambridge UP. Solid, mathematical. Advanced.
- A. Dixit, B. Nalebuff: Thinking Strategically. Norton. Lots of context and background. Interesting and non-technical.

Overview

Bargaining as example of sequential game

Repeated Games

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Bargaining as example of sequential game

Repeated Games

The Ultimatum Game

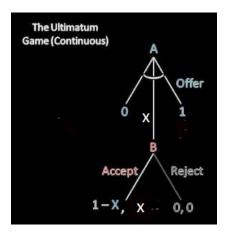
Ultimatum Game (UG): baseline (simplest non-trivial) model for **bargaining:** *take-it-or-leave-it!*

Assumptions

- **Surplus** can be divided continuously $(0 \le x \le 1)$
- Two agents:
 - A proposes split x versus 1 x, (proposal power)
 - B accepts or rejects;
- No deal (conflict deal) is considered worst outcome;
- Both agents aim to maximize their utility;

One round ultimatum game

• A makes single offer, B either accepts or rejects!

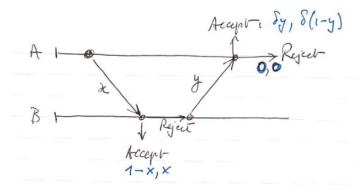


Two rounds ultimatum game with impatient players

Power of counter-offer:

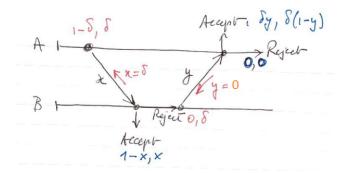
- A makes offer;
- B either accepts, or makes counter-offer.
- However, in each round the total is reduced by factor $\delta < 1$ (the icecream is melting!)

A makes offer, but B can make counter-offer!



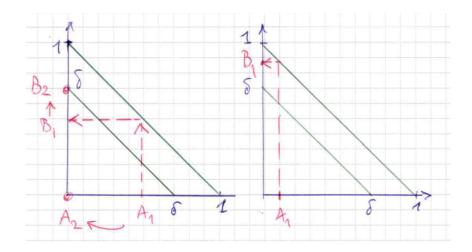
Two rounds ultimatum game: BI solution

- A makes offer, but B can make counter-offer!
- Use backward induction to find optimal solution.

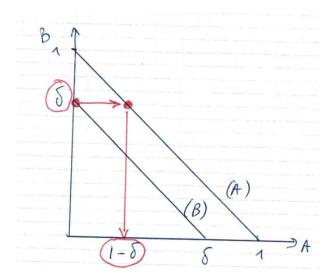


Conclusion: A offers split $(1 - \delta, \delta)$ which B accepts.

Alternative interpretation

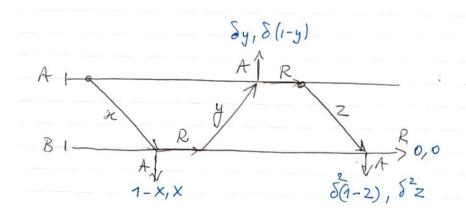


Two rounds ultimatum game: "Pareto Shuffle"

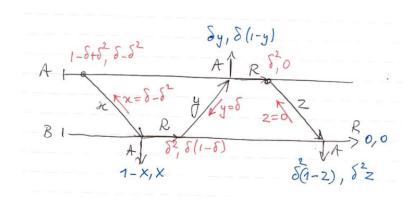


Three rounds ultimatum game

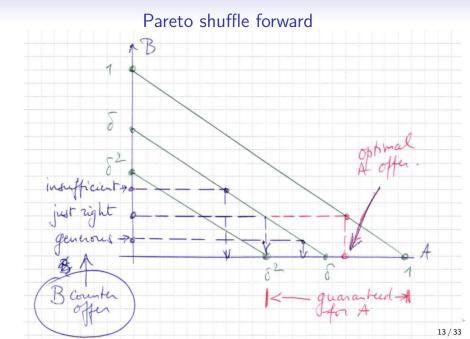
Room for two counter-offers!



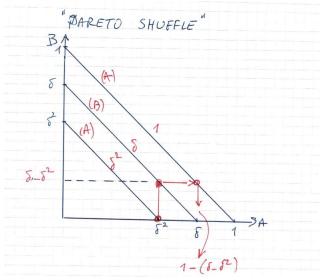
Three rounds ultimatum game: Backwards induction



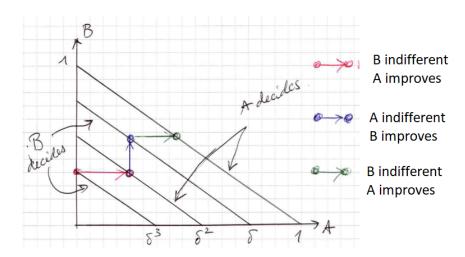
Conclusion: A proposes split $(1 - \delta + \delta^2, \delta - \delta^2)$ which B accepts.



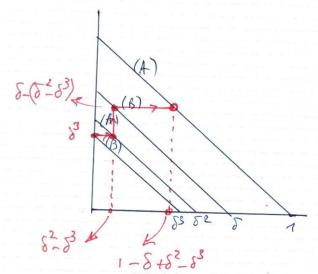
Three rounds ultimatum game: "Pareto Shuffle"



Four rounds ultimatum game: backward induction



Four rounds ultimatum game: "Pareto Shuffle"

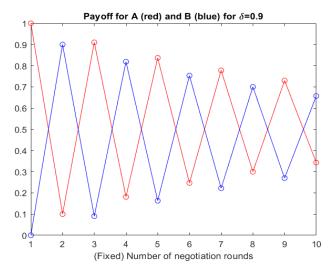


Alternating offers bargaining

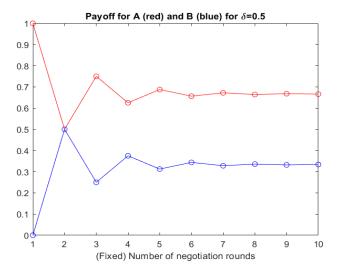
- Generalizing to negotiation with n rounds;
- When number of negotiation rounds is fixed at n, the equilibrium split is given in table below:

payoff for				
<i>n</i> rounds	n = 1	n=2	n = 3	n = 4
for A	1	$1-\delta$	$1 - \delta + \delta^2$	$1 - \delta + \delta^2 - \delta^3$
for B	0	δ	$\delta - \delta^2$	$\delta - \delta^2 + \delta^3$

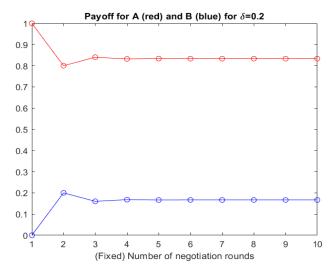
Alternating offers bargaining (mildly impatient, $\delta = 0.9$)



Alternating offers bargaining (seriously impatient, $\delta = 0.5$)



Alternating offers bargaining (very impatient, $\delta = 0.2$)



Mathematical aside

$$\sum_{k=0}^{\infty} x^k = 1 + x + x^2 + x^3 + \dots = \frac{1}{1 - x} \qquad \text{(for } |x| < 1\text{)}$$

$$\sum_{k=1}^{\infty} x^k = x + x^2 + x^3 + x^4 \dots = \frac{x}{1 - x} \qquad \text{(for } |x| < 1\text{)}$$

$$1 + x + x^2 + \dots + x^n = \frac{1}{1 - x} - x^{n+1} \frac{1}{1 - x}$$

$$= \frac{1 - x^{n+1}}{1 - x}$$

General conclusions

Limit behaviour for pay-offs:

$$A(n) = 1 - \delta + \delta^{2} - \dots \pm \delta^{n-1} = \frac{1 - (-\delta)^{n}}{1 - (-\delta)}$$

$$\implies \lim_{n \to \infty} A(n) = \frac{1}{1 + \delta}$$

$$B(n) = 1 - A(n) \implies \lim_{n \to \infty} B(n) = \frac{\delta}{1 + \delta}$$

- First offer advantage: $\lim A(n) \ge \lim B(n)$
- First offer advantage disappears for very patient negotiators:

$$\delta \longrightarrow 1 \implies A(n) \setminus 1/2, B(n) \nearrow 1/2;$$

Rubinstein's Model: Infinite Horizon Bargaining

- 2 agents, infinite horizon (# rounds not fixed in advance);
- **Time is valuable:** discount factors for both agents (δ_1, δ_2) ;
- Optimal split:

$$u_1=rac{1-\delta_2}{1-\delta_1\delta_2}$$
 and $u_2=rac{\delta_2(1-\delta_1)}{1-\delta_1\delta_2}$

- Tragedy of bargaining: The more time matters, the lower the share! E.g. if $\delta_2 \approx 0$, then $u_2 \approx 0$, etc.
- Notice **first mover's advantage** for $\delta_1 = \delta_2 = \delta$:

$$u_1=rac{1-\delta}{1-\delta^2}=rac{1}{1+\delta}$$
 and $u_2=rac{\delta(1-\delta)}{1-\delta^2}=rac{\delta}{1+\delta}$

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Bargaining as example of sequential game

Repeated Games

Repeated Games

- Repeat the same one-shot (stage) game (special case of sequential game);
- At each stage, information about preceding games is known;
- **Finite number** *n* (known!) of repetitions;
 - Maximize total reward:

$$R_n = \sum_{t=1}^n r_t$$

- Infinite (unlimited) number of repetitions
 - Maximize discounted total reward:

$$R = \sum_{t=1}^{\infty} \delta^{t-1} r_t$$
 discount factor $0 < \delta < 1$

equivalently: ending after random number of repetitions

Mathematical aside

Recall:

$$\sum_{k=0}^{\infty} x^k = 1 + x + x^2 + x^3 + \dots = \frac{1}{1-x} \quad \text{(for } |x| < 1\text{)}$$

$$\sum_{k=1}^{\infty} x^k = x + x^2 + x^3 + x^4 \dots = \frac{x}{1-x} \qquad \text{(for } |x| < 1\text{)}$$

Interpretation of discount factor

Repeated games: interpretation of discount factor

- After every stage-game there is a
 - probability 1δ that game will be ended;
 - ullet probability δ that game will proceed to next round
- Reward R now becomes random variable:

$$R_N = \sum_{t=1}^N r_t$$
 where $P(N = n) = \delta^{n-1}(1 - \delta)$

(Assuming at least one stage game is played, hence $N \geq 1$).

Then:

$$E(R_N) := E\left(\sum_{t=1}^N r_t\right) = r_1 + \delta r_2 + \delta^2 r_3 + \ldots = \sum_{t=1}^\infty \delta^{t-1} r_t.$$

Interpretation of discount factor

$$E(R_{N}) = E\left(\sum_{t=1}^{N} z_{t}\right)$$

$$= \sum_{n=1}^{\infty} E\left(\sum_{t=1}^{N} z_{t} \mid N=n\right) P(N=n)$$

$$= \sum_{n=1}^{\infty} \left(\sum_{t=1}^{N} z_{t}\right) \frac{P(N=n)}{S^{N-1}(1-\delta)}$$

$$= (1-\delta) \sum_{n=1}^{\infty} \left(\sum_{t=1}^{N} z_{t}\right) S^{N-1}$$

Interpretation of discount factor

$$E(\mathcal{R}_{N}) = (1-\delta) \sum_{N=1}^{\infty} \left(\sum_{k=1}^{N} z_{k} \right) \delta^{N-1}$$

$$= {}^{t}z_{1} + \delta z_{2} + \delta z_{3} + \delta^{2}z_{1} + \delta^{2}z_{2} + \delta z_{3}$$

$$= (1-\delta) \left[\frac{1}{1-\delta} z_{1} + \delta \left(\frac{1}{1-\delta} \right) z_{2} + \cdots \right]$$

$$= z_{1} + \delta z_{2} + \delta^{2}z_{3} + \cdots = \sum_{k=1}^{\infty} \delta^{k-1} z_{k}.$$

Example: Repeated Prisoner's dilemma

Pay-off matrix for stage game:

	C(oop)	D(efect)
С	3,3	1,4
D	4, 1	2, 2

- For finite number of repetitions:
 - Single Nash eq.: Play D-D for each repetition;
 - Can be proved using backward induction
- For **infinite number** of repetitions:
 - More interesting strategies, including cooperation
 - Examples: Tit-for-Tat, Grim Trigger

Grim Trigger Strategy for Repeated Prisoner's Dilemma

- Grim Trigger and Tit-for-Tat strategy Start by cooperating . . .
 - GT: continue cooperating, until someone defects; from then onwards, always defect!
 - TfT: from then onwards, copy last move of opponent.
- If both parties play GT, is it rational to defect?
 Consider player 1:
 - Utility for continued cooperation:

$$u_1(C) = 3 + 3\delta + 3\delta^2 + \ldots = \frac{3}{1 - \delta}$$

Utility for defection:

$$u_1(D) = 4 + 2\delta + 2\delta^2 + \ldots = 4 + 2\frac{\delta}{1 - \delta}$$

• Continued **cooperation is rational** if $u_1(C) > u_1(D)$.

Grim Trigger Strategy for Repeated Prisoner's Dilemma

• Continued **cooperation is rational** if $u_1(C) > u_1(D)$:

$$u_1(C) > u_1(D)$$
 \iff $\frac{3}{1-\delta} = 4 + 2\frac{\delta}{1-\delta}$ \iff $\delta > 1/2.$

 Interpretation: If the players are sufficiently patient (i.e. future rewards are sufficiently valuable) then it's rational to cooperate.