

SVEUČILIŠTE U ZAGREBU
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Predmet: Teorija informacije (34315)
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Zadatak
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Zadatak /zi03/:

Komunikacijskim kanalom prenose se četiri poruke generirane iz skupa od četiri simbola $\mathbf{X} = \{x_1, \dots, x_4\}$. Vjerojatnosti pojavljivanja simbola su sljedeće: $\mathbf{p}_x = [p/2, p/2, (1-p)/2, (1-p)/2]$, slijedno gledano ($p \in (0, 1)$). Matrica uvjetnih vjerojatnosti prijelaza u kanalu je:

$$[p(y_j | x_i)] = \begin{bmatrix} 1-f & f & 0 & 0 \\ f & 1-f & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}, \text{ uz } 0 \leq f \leq 1.$$

Odredite općeniti izraz za varijablu p koji osigurava maksimalnu količinu informacije po simbolu koja se u prosjeku može prenijeti danim kanalom. (**Napomena:** $H(f) = f \cdot \log(1/f) + (1-f) \cdot \log(1/(1-f))$)

$$C = \max I(\mathbf{X}; \mathbf{Y})$$

$$I(\mathbf{X}; \mathbf{Y}) = H(\mathbf{Y}) - H(\mathbf{Y} | \mathbf{X})$$

$$[p(y_j, x_i)] = [p(x_i)] * [p(y_j | x_i)]$$

$$[p(y_j, x_i)] = \begin{bmatrix} p/2 \\ p/2 \\ (1-p)/2 \\ (1-p)/2 \end{bmatrix} \begin{bmatrix} 1-f & f & 0 & 0 \\ f & 1-f & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} = \begin{bmatrix} (1-f)p/2 & fp/2 & 0 & 0 \\ fp/2 & (1-f)p/2 & 0 & 0 \\ 0 & 0 & (1-p)/2 & 0 \\ 0 & 0 & 0 & (1-p)/2 \end{bmatrix}$$

$$[p(y_j)] = [p/2 \quad p/2 \quad (1-p)/2 \quad (1-p)/2]$$

$$H(\mathbf{Y}) = \sum_{j=1}^4 p(y_j) * \log_2 \frac{1}{p(y_j)}$$

$$\begin{aligned} H(\mathbf{Y}) &= p/2 * \log(2/p) + p/2 * \log(2/p) + (1-p)/2 * \log(2/(1-p)) + (1-p)/2 * \log(2/(1-p)) = \\ &= p * \log(2/p) + (1-p) * \log(2/(1-p)) \end{aligned}$$

$$H(\mathbf{Y} | \mathbf{X}) = \sum_{j=1}^4 \sum_{i=1}^4 p(y_j, x_i) * \log_2 \frac{1}{p(y_j | x_i)}$$

$$\begin{aligned} H(\mathbf{Y} | \mathbf{X}) &= (1-f) * p/2 * \log(1/(1-f)) + f * p/2 * \log(1/f) + f * p/2 * \log(1/f) + (1-f) * p/2 * \log(1/(1-f)) + \\ &\quad + (1-p)/2 * \log 1 + (1-p)/2 * \log 1 = \\ &= (1-f) * p * \log(1/(1-f)) + f * p * \log(1/f) \end{aligned}$$

$$I(\mathbf{X}; \mathbf{Y}) = H(\mathbf{Y}) - H(\mathbf{Y} | \mathbf{X}) =$$

$$\begin{aligned}
&= p \cdot \log(2/p) + (1-p) \cdot \log(2/(1-p)) - (1-f) \cdot p \cdot \log(1/(1-f)) - f \cdot p \cdot \log(1/f) = \\
&= H(X) - (1-f) \cdot p \cdot \log(1/(1-f)) - f \cdot p \cdot \log(1/f) ?? \leftarrow [\text{ovaj zadnji red bi zanemario}]
\end{aligned}$$

C tj količina informacija je maksimalna kad je derivacija $I(X;Y)$ po p jednaka 0.

$$\begin{aligned}
dI(X;Y)/dp &= (p \cdot \log(2/p))' + ((1-p) \cdot \log(2/(1-p)))' - ((1-f) \cdot p \cdot \log(1/(1-f)) - f \cdot p \cdot \log(1/f))' = \\
&= (\log(2/p) - 1/\ln 2) + (\log(2/(1-p)) + 1/\ln 2) - (1-f) \cdot \log(1/(1-f)) - f \cdot \log(1/f) = \\
&= \log(2/p) + \log(2/(1-p)) - (1-f) \cdot \log(1/(1-f)) - f \cdot \log(1/f) = \\
&= \log((1-p)/p) - (1-f) \cdot \log(1/(1-f)) - f \cdot \log(1/f) = 0 \\
\log((1-p)/p) &= (1-f) \cdot \log(1/(1-f)) + f \cdot \log(1/f) \\
(1-p)/p &= 2^{(1-f) \cdot \log(1/(1-f)) + f \cdot \log(1/f)} \\
p &= 1 / (1 + 2^{(1-f) \cdot \log(1/(1-f)) + f \cdot \log(1/f)}) = 1 / (1 + 2^{H(f)})
\end{aligned}$$