

KOLIČINA INFORMACIJE

$$I(x_i) = -\log_2 p(x_i) \left[\frac{\text{bit}}{\text{simbol}} \right]$$

$$I(x_1 x_2 \dots x_k) = -\log_2 [p(x_1) \cdot p(x_2) \cdot \dots \cdot p(x_k)] \left[\frac{\text{bit}}{\text{poruka}} \right]$$

VJEROJATNOSTI U KOMUNIKACIJSKOM SUSTAVU

$$[p(y_j|x_i)] = \begin{bmatrix} p(y_1|x_1) & p(y_2|x_1) & \dots & p(y_m|x_1) \\ p(y_1|x_2) & p(y_2|x_2) & \dots & p(y_m|x_2) \\ \dots & \dots & \dots & \dots \\ p(y_1|x_n) & p(y_2|x_n) & \dots & p(y_m|x_n) \end{bmatrix} \begin{matrix} \Sigma = 1 \\ \Sigma = 1 \\ \dots \\ \Sigma = 1 \end{matrix}$$

$$[p(y_j)] = [p(x_i)][p(y_j|x_i)] \quad [p(x_i)]^T = [p(x_i|y_j)][p(y_j)]^T$$

$$[p(x_i, y_j)] = [p(x_i)p(y_j|x_i)] = [p(x_i|y_j)p(y_j)]$$

$$p(x_i|y_j) = \frac{p(x_i, y_j)}{p(y_j)} = \frac{p(x_i, y_j)}{\sum_{i=1}^n p(x_i, y_j)} = \frac{p(x_i)p(y_j|x_i)}{\sum_{i=1}^n p(x_i)p(y_j|x_i)}$$

$$[p(x_i, y_j)] = \begin{bmatrix} p(x_1, y_1) & p(x_1, y_2) & \dots & p(x_1, y_m) \\ p(x_2, y_1) & p(x_2, y_2) & \dots & p(x_2, y_m) \\ \dots & \dots & \dots & \dots \\ p(x_n, y_1) & p(x_n, y_2) & \dots & p(x_n, y_m) \end{bmatrix} \begin{matrix} \Sigma = p(x_1) \\ \Sigma = p(x_2) \\ \dots \\ \Sigma = p(x_n) \end{matrix}$$

$$\Sigma = p(y_1) \quad \Sigma = p(y_2) \quad \dots \quad \Sigma = p(y_m)$$

ENTROPIJA

$$H(X) = -\sum_{i=1}^n p(x_i) \log_2 p(x_i) \left[\frac{\text{bit}}{\text{simbol}} \right]$$

$$H(Y) = -\sum_{j=1}^m p(y_j) \log_2 p(y_j) \left[\frac{\text{bit}}{\text{simbol}} \right]$$

$$H(X, Y) = -\sum_{i=1}^n \sum_{j=1}^m p(x_i, y_j) \log_2 p(x_i, y_j) \left[\frac{\text{bit}}{\text{simbol}} \right]$$

ENTROPIJA ŠUMA (IRELEVANTNOST)

$$H(Y|X) = -\sum_{i=1}^n \sum_{j=1}^m p(x_i, y_j) \log_2 p(y_j|x_i) \left[\frac{\text{bit}}{\text{simbol}} \right]$$

EKVIVOKACIJA (MNOGOZNAČNOST)

$$H(X|Y) = -\sum_{i=1}^n \sum_{j=1}^m p(x_i, y_j) \log_2 p(x_i|y_j) \left[\frac{\text{bit}}{\text{simbol}} \right]$$

SREDNJI SADRŽAJ INFORMACIJE (TRANSINFORMACIJA)

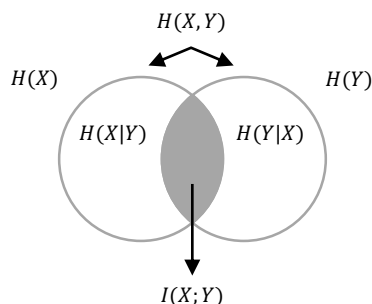
$$I(X; Y) = \sum_{i=1}^n \sum_{j=1}^m p(x_i, y_j) \log_2 \frac{p(x_i, y_j)}{p(x_i)p(y_j)} \left[\frac{\text{bit}}{\text{simbol}} \right]$$

$$I(X; Y) = H(X) - H(X|Y) \quad I(X; Y) = H(Y) - H(Y|X) \quad I(X; X) = H(X)$$

$$I(X; Y) = H(X) + H(Y) - H(X, Y)$$

$$H(X, Y) = H(X) + H(Y|X) \quad H(X, Y) = H(Y) + H(X|Y)$$

$$X \text{ i } Y \text{ nezavisni} \rightarrow H(X, Y) = H(X) + H(Y)$$



RELATIVNA ENTROPIJA

$$D(p||q) = \sum_{i=1}^n p(x_i) \log \frac{p(x_i)}{q(x_i)}$$

KAPACITET DISKRETNOG KOMUNIKACIJSKOG KANALA

$$C = \max_{\{p(x_i)\}} I(X; Y) = \max_{\{p(x_i)\}} (H(Y) - H(Y|X)) \left[\frac{\text{bit}}{\text{simbol}} \right]$$

INFORMACIJSKA BRZINA IZVORIŠTA

$$R = \frac{H(X)}{T_s} \left[\frac{\text{bit}}{\text{s}} \right] \quad (T_s - \text{prosječno trajanje simbola})$$

ENTROPIJSKO KODIRANJE

SREDNJA DULJINA KODNE RIJEČI

$$L(X) = \sum_{i=1}^n p(x_i) l_i \left[\frac{\text{bit}}{\text{simbol}} \right]$$

EFIKASNOST KODA

$$\varepsilon = \frac{H(X)}{L(X)}$$

KRAFTOVA NEJEDNAKOST (nužan i dovoljan uvjet za prefiksni kod)

$$\sum d^{-l_i} \leq 1$$

OPTIMALNOST KODA (nužan i dovoljan optimalnosti koda)

$$H(X) \leq L(X) < H(X) + 1$$

$$\min \left[L(X) = \sum_{i=1}^n p(x_i) l_i \right] \text{ uz uvjet } \sum_{i=1}^n d^{-l_i} \leq 1$$

SARDINAS-PATTERSONOV TEST

$C(y)$ se dodaje u skup S_{i+1} ako i samo ako:

$$\exists C(x) \in S_0 \text{ tako da } C(x)C(y) \in S_i$$

ili

$$\exists C(z) \in S_i \text{ tako da } C(z)C(y) \in S_0$$

■ Kod je JDK ako niti jedan S_i ($i \geq 1$) ne sadrži kodne riječi iz S_0 ■

HUFFMANNOVO KODIRANJE

$$N - \text{broj simbola}; \quad B - \text{baza kodiranja}; \quad k = \left\lfloor \frac{N-1}{B-1} \right\rfloor$$

$$N' = (B-1)k + 1; \quad N' \neq N \Rightarrow \text{dodaj } N' - N \text{ simbola s } p = 0$$

ARITMETIČKO KODIRANJE

$$D' = D + (G - D) \cdot D_s$$

$$G' = D + (G - D) \cdot G_s$$

$$\left\lceil \log_2 \left(\frac{1}{G' - D'} \right) \right\rceil + 1 \text{ znamenki} \rightarrow \text{kod koji se može JDK}$$

$$L_{a(d=10)} = \sum_{i=1}^N L_{a(d=2)}(i) \cdot 2^{-i}$$

SREDNJA DULJINA KODNE RIJEČI

$$L(X) = \frac{\sum_{i=1}^n p(x_i) \cdot l(x_i)}{\sum_{i=1}^n p(x_i) \cdot n(x_i)} \left[\frac{\text{bit}}{\text{simbol}} \right] \quad \begin{matrix} l(x_i) \rightarrow \text{duljina kodne riječi} \\ n(x_i) \rightarrow \text{broj simbola} \end{matrix}$$

VJEROJATNOST POJAVE POJEDINOG SIMBOLA

$$p(\text{pojava_simbola}) = \frac{p(x_i) \cdot n(x_i)}{L(X)}$$

$$n(x_i) \rightarrow \text{duljina simbola}$$

VJEROJATNOST ISPRAVNOG PRIJENOSA BINARNOG SIMETRIČNOG KANALA

$$P_{BSK} = \frac{1}{2} (1 + (1 - 2 \cdot p)^k) \quad p \rightarrow \text{vjerojatnost pogreške}$$

VJEROJATNOST POGREŠNOG PRIJENOSA BINARNOG SIMETRIČNOG KANALA

$$P_{BSK} = \frac{1}{2} (1 - (1 - 2 \cdot p)^k) \quad p \rightarrow \text{vjerojatnost pogreške}$$

OSTALO**VJEROJATNOST PRIJELAZA**

$$[p(y_j)] = [p(z_k)][p(y_j|z_k)]$$

$$[p(y_j)] = [p(x_i)][p(y_j|x_i)] \rightarrow [p(y_j|x_i)] = [p(z_k|x_i)][p(y_j|z_k)]$$

$$[p(z_k)] = [p(x_i)][p(z_k|x_i)]$$

HUFFMANOVO KODIRANJE (*m* simbola)

$$p(x_i) = \frac{1}{m}$$

$$2^n \leq m < 2^{n+1} \rightarrow k = m - 2^n$$

$$\text{kodnih riječi duljine } \mathbf{n} \text{ ima} \rightarrow 2^n - k$$

$$\text{kodnih riječi duljine } \mathbf{n + 1} \text{ ima} \rightarrow 2k$$

NAJMANJI BROJ ELEMENATA SKUPA

$$\min(\text{card}(X)) \rightarrow \max(H(X))$$

$$\text{card}(X) \geq \log_2(H(X))$$

NAJMANJA DULJINA KODNE RIJEČI (*Aritmetički algoritam*)

$$l(x) = \left\lceil \log_2 \left(\frac{1}{P(X)} \right) \right\rceil + 1 \text{ [bit]}$$

$$P_i(x) = \prod_{i=1}^n p(x_i)$$

$$\text{ako tražimo minimalnu duljinu} \rightarrow \text{tražimo } \max(P_i(x))$$