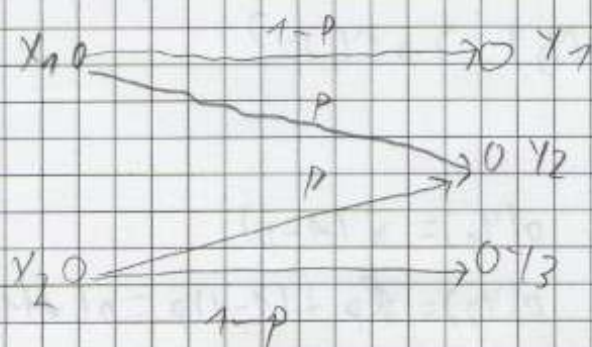


① $p(x_1) = x$
 $p(x_2) = 1-x$

$$P = p(y_j | x_i) = \begin{bmatrix} 1-p & p & 0 \\ 0 & p & 1-p \end{bmatrix}$$



a) $H(Y|X) = ?$

$$H(Y|X) = - \sum_{i=1}^2 \sum_{j=1}^2 p(x_i, y_j) \log p(y_j | x_i) = - \sum_{i=1}^2 p(x_i, y_j) \log p(y_j | x_i)$$

$$p(x_i, y_j) = p(y_j | x_i) \cdot p(x_i)$$

$$p(x_i, y_j) = \begin{bmatrix} x(1-p) & xp & 0 \\ 0 & (1-x)p & (1-x)(1-p) \end{bmatrix}$$

$$H(Y|X) = -x(1-p) \log(1-p) - xp \log p - (1-x)p \log p - (1-x)(1-p) \log(1-p)$$

$$H(Y|X) = -\log(1-p) [x(1-p) + (1-x)(1-p)] - \log p [xp + (1-x)p] =$$

$$= -\log(1-p) [(1-p)(x+1-x)] - \log p [p(x+1-x)] =$$

$$\boxed{H(Y|X) = -(1-p) \log(1-p) - p \log p}$$

$$b) I(X;Y) = ?$$

$$p(y_i, x_i) = \begin{bmatrix} x(1-p) & xp & 0 \\ 0 & (1-x)p & (1-x)(1-p) \end{bmatrix}$$

$$I(X;Y) = H(Y) - H(Y|X) \Rightarrow H(Y) = ?$$

$$H(Y) = - \sum_{j=1}^3 p(Y_j) \log p(Y_j) \Rightarrow p(Y_j) = ?$$

$$p(Y_j) = \sum_{i=1}^2 p(X_i, Y_j) \Rightarrow p(Y_1) = x(1-p)$$

$$p(Y_2) = xp + (1-x)p = p(x+1-x) = p$$

$$p(Y_3) = (1-x)(1-p)$$

$$\begin{aligned} H(Y) &= -x(1-p) \log[x(1-p)] - p \log p - (1-x)(1-p) \log[(1-x)(1-p)] = \\ &= -x(1-p) \log x - x(1-p) \log(1-p) - p \log p - (1-x)(1-p) \log(1-x) - (1-x)(1-p) \log(1-p) = \\ &= -x(1-p) \log x - \log(1-p) [x(1-p) + (1-x)(1-p)] - p \log p - (1-x)(1-p) \log(1-x) = \\ &= -x(1-p) \log x - \log(1-p) [(1-p)(x+1-x)] - p \log p - (1-x)(1-p) \log(1-x) = \\ H(Y) &= -x(1-p) \log x - (1-p) \log(1-p) - p \log p - (1-x)(1-p) \log(1-x) \end{aligned}$$

$$H(Y|X) = -(1-p) \log(1-p) - p \log p$$

$$I(X;Y) = -x(1-p) \log x - (1-p) \log(1-p) - p \log p - (1-x)(1-p) \log(1-x) + (1-p) \log(1-p) + p \log p$$

$$I(X;Y) = -x(1-p) \log x - (1-x)(1-p) \log(1-x)$$

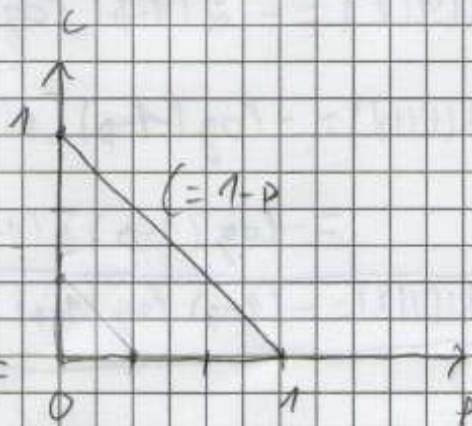
c)

$$C = \max_{\sum p(X_i)} I(X;Y) \Rightarrow p(X_1) = p(X_2)$$

$$x = 1-x$$

$$2x = 1$$

$$x = \frac{1}{2}$$



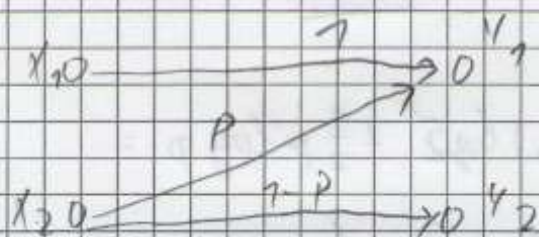
$$C = -\frac{1}{2}(1-p) \log 2^{-1} - \frac{1}{2}(1-p) \log 2^{-1} =$$

$$= -(1-p) \log 2^{-1} = -(1-p) \cdot (-1) = 1-p$$

$$C = 1-p$$

② $p(X_1) = 1-x$
 $p(X_2) = x$

$$P = p(Y; X_i) = \begin{bmatrix} 1 & 0 \\ p & 1-p \end{bmatrix}$$



a) $H(Y) = ?$

$$H(Y) = - \sum_{j=1}^2 p(Y_j) \log p(Y_j) \Rightarrow p(Y_j) = ?$$

$$p(Y_j) = \sum_{i=1}^2 p(X_i, Y_j) \Rightarrow p(X_i, Y_j) = ?$$

$$p(X_i, Y_j) = p(Y_j | X_i) \cdot p(X_i) \Rightarrow p(Y_i, Y_j) = \begin{bmatrix} 1-x & 0 \\ xp & x(1-p) \end{bmatrix}$$

$$p(Y_1) = 1-x + xp, \quad p(Y_2) = x(1-p)$$

$$H(Y) = -(1-x+xp) \log(1-x+xp) - x(1-p) \log x(1-p)$$

$$H(Y) = -(1-x+xp) \log(1-x+xp) - x(1-p) \log x - x(1-p) \log(1-p)$$

b) $I(X; Y) = ?$

$$I(X; Y) = H(Y) - H(Y|X) \Rightarrow H(Y|X) = ?$$

$$H(Y|X) = - \sum_{i=1}^2 \sum_{j=1}^2 p(X_i, Y_j) \log p(Y_j | X_i)$$

$$H(Y|X) = -xp \log p - x(1-p) \log(1-p)$$

$$I(X; Y) = -(1-x+xp) \log(1-x+xp) - x(1-p) \log x - x(1-p) \log(1-p) + xp \log p + x(1-p) \log(1-p)$$

$$I(X; Y) = -(1-x+xp) \log(1-x+xp) - x(1-p) \log x + xp \log p$$

$$C) C = ?$$

$$C = \max_{\{p(x,y)\}} L(y, y) \Rightarrow p(x_1) = p(x_2) \Rightarrow \underline{p = \frac{1}{2}}$$

$$C = -\left(1 - \frac{1}{2} + \frac{1}{2}p\right) \log\left(1 - \frac{1}{2} + \frac{1}{2}p\right) - \frac{1}{2}(1-p) \log 2^{-1} + \frac{1}{2}p \log p =$$

$$= -\frac{1}{2}(1+p) \log\left[\frac{1}{2}(1+p)\right] + \frac{1}{2}(1-p) + \frac{1}{2}p \log p =$$

$$= -\frac{1}{2}(1+p) \log 2^{-1} - \frac{1}{2}(1+p) \log(1+p) + \frac{1}{2}(1-p) + \frac{1}{2}p \log p =$$

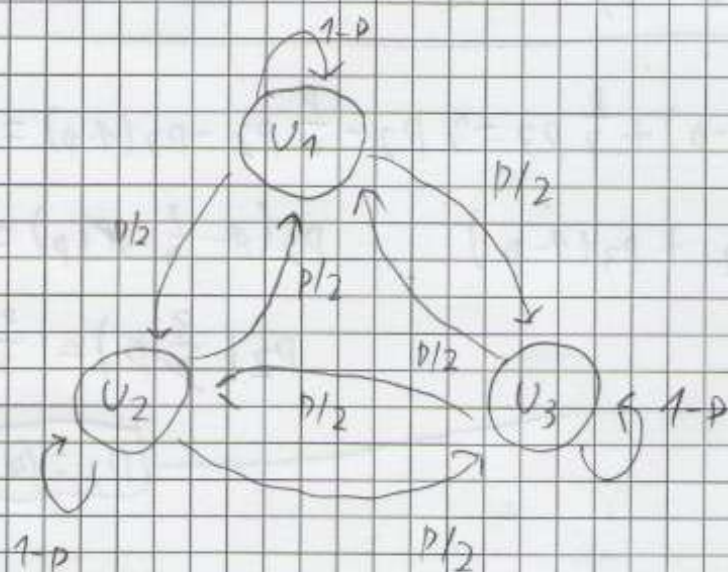
$$= \frac{1}{2}(1+p) - \frac{1}{2}(1+p) \log(1+p) + \frac{1}{2}(1-p) + \frac{1}{2}p \log p =$$

$$C = 1 - \frac{1}{2}(1+p) \log(1+p) + \frac{1}{2}p \log p =$$

(3) $V = \{U_1, U_2, U_3\}$

$p(U_j | U_i) = p/2, i \neq j$

(2)



b) $p(U_1) = ? = p_1$
 $p(U_2) = ? = p_2$
 $p(U_3) = ? = p_3$

$p(U_j | U_i) = \begin{bmatrix} 1-p & p/2 & p/2 \\ p/2 & 1-p & p/2 \\ p/2 & p/2 & 1-p \end{bmatrix}$

$p(U_i) = \sum_{j=1}^3 p(U_i, U_j) = ? \quad p(U_i, U_j) = ?$

$p(U_i, U_j) = p(U_j | U_i) p(U_i) = ? \quad p(U_j | U_i) = ?$

$p(U_i, U_j) = \begin{bmatrix} p_1(1-p) & \frac{p}{2} p_2 & \frac{p}{2} p_3 \\ \frac{p}{2} p_1 & p_2(1-p) & \frac{p}{2} p_3 \\ \frac{p}{2} p_1 & \frac{p}{2} p_2 & (1-p)p_3 \end{bmatrix}$

$p_1 + p_2 + p_3 = 1$

$$p_1 = p_1(1-p) + \frac{p}{2} p_2 + \frac{p}{2} p_3 = 1 \quad p_1 p = \frac{p(p_2 + p_3)}{2}$$

$$p_2 = \frac{p}{2} p_1 + p_2(1-p) + \frac{p}{2} p_3$$

$$p_3 = \frac{p}{2} p_1 + \frac{p}{2} p_2 + p_3(1-p)$$

$$p_2 = \frac{p(p_2 + p_3)}{4} + p_2(1-p) + \frac{p}{2} p_3 = 1 \quad p_2 - \frac{p}{4} p_2 - p_2(1-p) = \frac{p}{4} p_3 + \frac{p}{2} p_3$$

$$p_3 = \frac{p(p_2 + p_3)}{4} + \frac{p}{2} p_2 + p_3(1-p)$$

$$p_2(1 - \frac{p}{4} - 1 + p) = \frac{3p}{4} p_3$$

$$p_2(\frac{3}{4}p) = \frac{3}{4}p p_3$$

$$p_1 = \frac{p_2 + p_3}{2}$$

$$p_2 = p_3$$

$$p_1 = \frac{2p_2}{2} = p_2$$

$$p_1 + p_2 + p_3 = 1$$

$$p_2 + p_2 + p_2 = 1$$

$$3p_2 = 1 \Rightarrow$$

$$p_2 = p_1 = p_3 = \frac{1}{3}$$

$$\Rightarrow p(U_i, U_j) =$$

$$\begin{bmatrix} \frac{1-p}{3} & \frac{p}{6} & \frac{p}{6} \\ \frac{p}{6} & \frac{1-p}{3} & \frac{p}{6} \\ \frac{p}{6} & \frac{p}{6} & \frac{1-p}{3} \end{bmatrix}$$

$$C) H(U_j | U_i) = ?$$

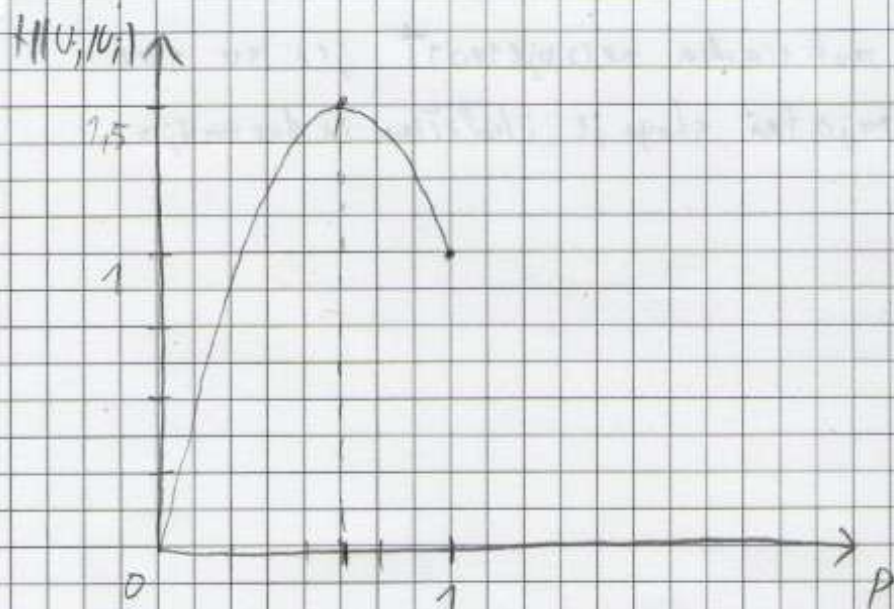
$$H(U_j | U_i) = - \sum_{i=1}^3 \sum_{j=1}^3 p(U_j | U_i) \log_2 p(U_j | U_i) =$$

$$= \left(-\left(\frac{1-p}{3}\right) \log_2(1-p) - \frac{p}{6} \log_2 \frac{p}{6} - \frac{p}{6} \log_2 \frac{p}{6} \right) \cdot 3 =$$

$$= \left(-\frac{1-p}{3} \log_2(1-p) - \frac{p}{3} \log_2 \frac{p}{6} \right) \cdot 3 =$$

$$= -(1-p) \log_2(1-p) - p \log_2 \frac{p}{6} = -(1-p) \log_2(1-p) - p \log_2 p + p \log_2 6 =$$

$$H(U_j | U_i) = p - p \log_2 p - (1-p) \log_2(1-p)$$



d) $p = ?$ (max $H(U_j|U_i)$)

$$H(U_j|U_i) = p - p \log p - (1-p) \log (1-p)$$

$$\begin{aligned} (H(U_j|U_i))' &= 1 - \log p - p \frac{1}{p \ln 2} + \log(1-p) + (1-p) \frac{1}{(1-p) \ln 2} = \\ &= 1 - \log p + \log(1-p) \end{aligned}$$

$$(H(U_j|U_i))' = 0 \Rightarrow 1 - \log p + \log(1-p) = 0$$

$$\log 2 - \log p + \log(1-p) = \log 1$$

$$\log \left[\frac{2}{p} (1-p) \right] = \log 1$$

$$\frac{2}{p} (1-p) = 1$$

$$2 - 2p = p$$

$$3p = 2$$

$$\boxed{p = \frac{2}{3}}$$

e) Za $p=0$ svaki simbol može jedino prići sam u sebi stoga ne postoji nikakva neizvesnost pa je količina informacije 0 bit/symbol.

Za $p=1$ svaki simbol može jedino prići u jedan od prostih dva simola sa jednakom vjerojatnošću, u tom slučaju je količina informacije 1 bit/symbol.

>>>

Za $p = \frac{2}{3}$ ostvaruje se maksimalna nezavisnost jer su svi
prijelazi jednako vjerojatni stoga je i količina informacija
maksimalna.

4.

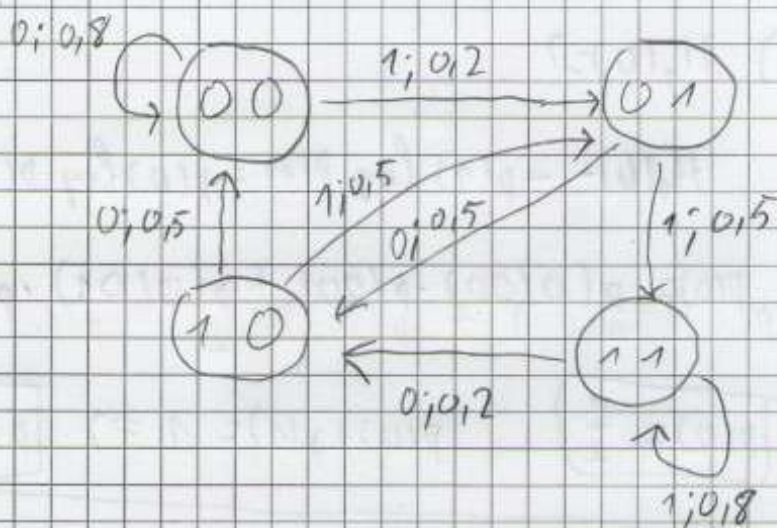
$$U = \{0, 1\}$$

$$p(0/00) = 0,8$$

$$p(0/01) = 0,5$$

$$p(0/10) = 0,5$$

$$p(0/11) = 0,2$$



a) $H_3(\text{tri}) = ?$

$$H_3(U) = ?$$

$$p(00) = p(11) = \frac{5}{14}$$

$$p(01) = p(10) = \frac{1}{7}$$

$$p(000) = p(00) \cdot p(0/00) = \frac{2}{7}$$

$$p(001) = p(00) \cdot p(1/00) = \frac{1}{14}$$

$$p(010) = p(01) \cdot p(0/01) = \frac{1}{14}$$

$$p(011) = p(01) \cdot p(1/01) = \frac{1}{14}$$

$$p(100) = p(10) \cdot p(0/10) = \frac{1}{14}$$

$$p(101) = p(10) \cdot p(1/10) = \frac{1}{14}$$

$$p(110) = p(11) \cdot p(0/11) = \frac{1}{14}$$

$$p(111) = p(11) \cdot p(1/11) = \frac{2}{7}$$

$$H_3(\text{tri}) = -2 \cdot \frac{2}{7} \log \frac{2}{7} - 6 \cdot \frac{1}{14} \log \frac{1}{14} = 2,664 \text{ bit/trigram}$$

$$H_3(U) = \frac{H_3(\text{tri})}{3} = 0,888 \text{ bit/symbol}$$

b) $H_2(\text{bi}) = ?$

$$H_2(U) = ?$$

$$H_2(\text{bi}) = -2 \cdot \frac{5}{14} \log \frac{5}{14} - 2 \cdot \frac{1}{7} \log \frac{1}{7}$$

$$H_2(\text{bi}) = 1,863 \text{ bit/bigram}$$

$$H_2(U) = \frac{H_2(\text{bi})}{2} = 0,932 \text{ bit/symbol}$$

c) $H_1(U) = ?$

$$H_1(U) = -p(1) \log p(1) - p(0) \log p(0) \Rightarrow p(0) = ?, p(1) = ?$$

$$p(0) = p(0/00) \cdot p(00) + p(0/01) \cdot p(01) + p(0/10) \cdot p(10) + p(0/11) \cdot p(11)$$

$$p(0) = \frac{1}{2} \quad p(0) + p(1) = 1 \Rightarrow p(1) = \frac{1}{2}$$

$$H_1(U) = -2 \cdot \frac{1}{2} \log \frac{1}{2} = 1 \text{ bit/symbol}$$

d) $F_1(U), F_2(U), F_3(U) = ?$

$$F_N(U) = N H_N(U) - (N-1) H_{N-1}(U)$$

$$F_1(U) = H_1(U) = 1 \text{ bit/symbol}$$

$$F_2(U) = 2 H_2(U) - H_1(U) = 0,863 \text{ bit/symbol}$$

$$F_3(U) = 3 H_3(U) - 2 H_2(U) = 0,801 \text{ bit/symbol}$$

e) $F_1(U) < F_2(U) < F_3(U)$

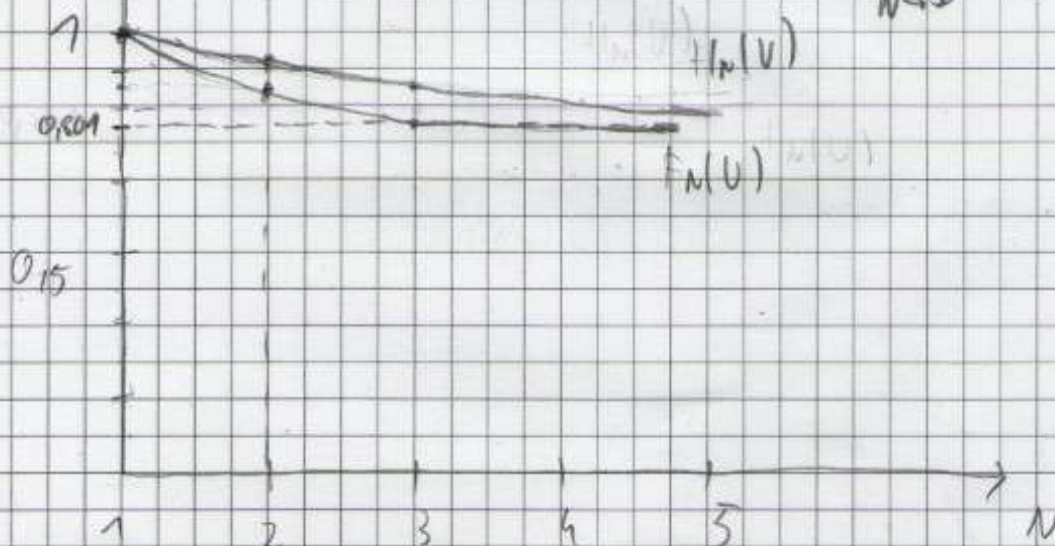
Za fiksni N , $F_N(U)$ će biti određen vjerojatnostima prijela za čija vrijednost ovisi o broju prethodnih simbola. Povećanjem broja prethodnih simbola razmjernost se smanjuje stoga će se i $F_N(U)$ talasito smanjiti.

f) $F_4(U)$ će biti jednako $F_3(U)$ jer je markov lanac drugog reda, stoga samo dva prethodna simbola utječu na vjerojatnosti prijela. Zato za $N \geq 3$ se vrijednost $F_N(U)$ više ne mijenja.

$H_4(U)$ će biti manji od $H_2(U)$ jer prva dva simbola i dulje nula imaju utjecaj u količini informacije poruke.

g) $f_N(u), f_N(u)$

$$\lim_{N \rightarrow \infty} f_N(u) = f_3(u)$$



5. $L \subseteq W$

posloup = a b b b b b b a b a b a a a b

$D[1] = a$ $D[2] = b$

| | | | | | | | | | | | | | | | | |
|---|---|---|---|---|---|---|---|---|----|----|----|----|----|----|----|----|
| 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 | 15 | 16 | 17 |
| a | b | b | b | b | b | b | a | b | a | b | a | b | a | a | a | b |

| konk | mjesto | sadržaj zjčnku | kolik je kodu |
|------|--------|----------------|---------------|
| 1. | 1 | (1) a | (1) |
| 2. | 2 | (2) b | (2) |
| 3. | 3 | (3) a b | (3) |
| 4. | 5 | (4) b b | (5) |
| 5. | 8 | (5) b b b | (3) |
| 6. | 10 | (6) b b b a | (7) |
| 7. | 13 | (7) a b a | (2) |
| 8. | 14 | (8) a b a b | (1) |
| 9. | 15 | (9) b a | (10) |
| 10. | 17 | (10) a a | (2) |
| | | (11) a a b | |