Distributed Hash Table

- primjer Chord

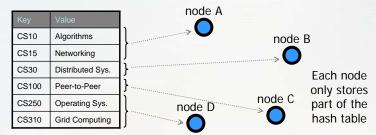
Preuzeto i dorađeno iz

Self-Organization in Structured Overlay Networks (DKS) http://dks.sics.se

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What is a Distributed Hash Table (DHT)? 1/2

 An infrastructure that enables the distribution of an ordinary hash table onto a set of cooperating nodes

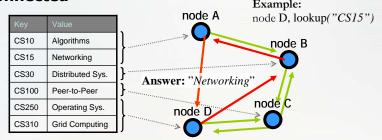


- The DHT provides a basic *lookup service*, which allows *any* node to find the value associated with a given key
- Example:

lookup("CS30"), at any node should return: "Distributed Sys."

What is a Distributed Hash Table (DHT)? 2/2

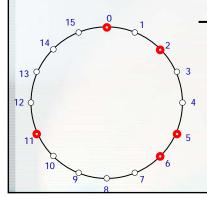
To provide the lookup service, the nodes must be interconnected



 Each node maintains a routing table with pointers to some other nodes such that lookup requests can be routed to the node storing the requested key/value-pair (a.k.a. item)

How do we distribute the hash table?

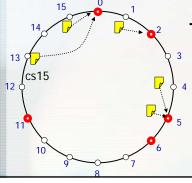
- Use a logical name space called the *identifier space*, consisting of identifiers $\{0,1,2,\ldots,N-1\}$
- Identifiers can be binary encoded, key length equals m
- ullet The identifier space can be perceived as a logical ring modulo N
- Every node is assigned an identifier using a function H_I .
- Note there are not necessarily N nodes in a network!



- Example: *N*=16 (*m*=4), nodes {**a,b,c,d,e**}
- Node a gets identifier 0 since H_I(a)=0, the other nodes b, c, d, e, get identifiers 2, 5, 6, 11 the same way (they are marked red in the figure)

How are items assigned to nodes?

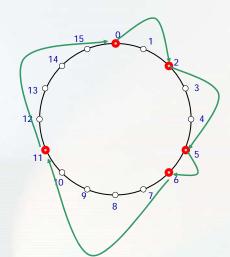
- Items are mapped to the same identifier space using another function H_2 , every node knows H_2
- Items are stored at nodes with the same identifier (if possible)
- As items may be mapped to identifiers of non-existing nodes, such items are stored at their successor nodes, i.e. the first node encountered moving in the clockwise direction



- Item ("cs15", "networking") is mapped to identifier 13 since H_2 ("cs15") = 13, other items are similarly mapped to 15, 2, 4, 5
- As there is no node with id=13 in the network, the item is stored at the successor node (node 0 in this case)

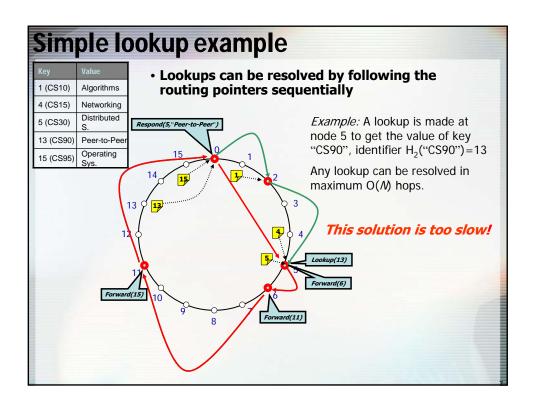
How do we interconnect the nodes?

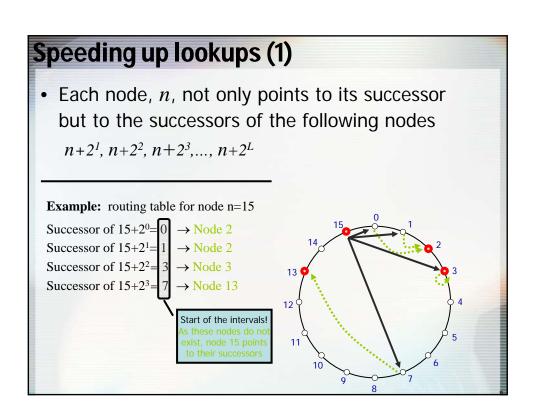
 Each node maintains a routing pointer to the successor in the ring (single routing pointer)



 The successor of a node n is the first node met going in clockwise direction starting at n +1

Successor of Node $0 \rightarrow$ Node 2 Successor of Node $2 \rightarrow$ Node 5 Successor of Node $5 \rightarrow$ Node 6 Successor of Node $6 \rightarrow$ Node 11 Successor of Node $11 \rightarrow$ Node 0





Speeding up lookups (2)

- At each step during search if node cannot find the item, route to the node in your routing table with the identifier less then the item identifier
- At each step during routing, the distance between the currrent node and destination is halved (in the worst case), requires O(log₂N) hops at worst

Example: search for item 10 starting at node 15 (item 10 is stored at 13)

Node 15 routes search to node 3 (node 13 is larger than 10!)
Node 3 routes the request to node

