# Exponentially many 3-colorings of planar graphs with no short separating cycles

### Mateusz Pach

Based on Carsten Thomassen's article [Tho21]

5 May 2022

# Motivation

# Conjecture [Thomassen 2007]

Every planar triangle-free graph  $\it G$  has exponentially many 3-colorings.

# Results

#### Main Theorem 1.

For infinitely many n, there exists a triangle-free planar graph with separating cycles of length 4 and 5 whose number of proper vertex-3-colorings is  $< 2^{15n/\log_2(n)}$ .

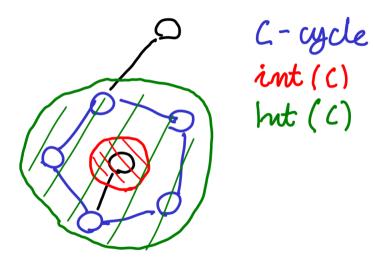
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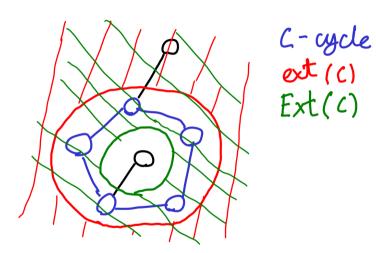
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#### Main Theorem 2.

The number of proper vertex-3-colorings of every triangle-free planar graph with n vertices and with no separating cycle of length 4 or 5 is at least  $2^{n/17700000}$ .





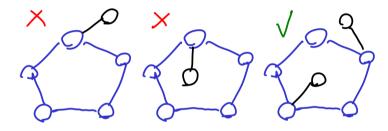


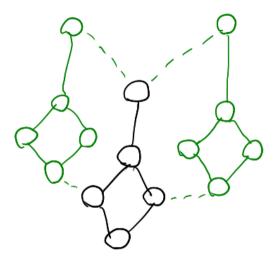
Figure: Seperating cycle

# **Definitions**

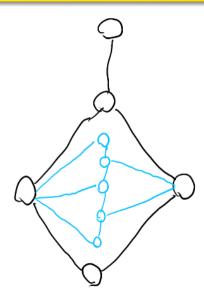
# Definition

If int(C) is empty or ext(C) is empty, then C is facial.

# Planar triangle-free graphs with only subexponentially many 3-colorings



# Planar triangle-free graphs with only subexponentially many 3-colorings



# Exponentially many 3-colorings of planar triangle-free graphs with many facial 5-cycles

#### Theorem 3.

Let G be a connected plane graph of girth at least 4 with an outer cycle C of length 4 or 5. Let q denote the number of facial 5-cycles inside C. Let c be a 3-coloring of C.

If C has length 4, then the number of distinct 3-colorings of G extending c is at least  $2^{3q/20000}$ .

If C has length 5, then the number of distinct 3-colorings of G extending c is at least  $2^{3(q-1)/20000}$ 

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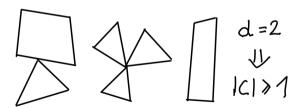
# Intersecting curve systems

#### Theorem 4.

Let d be a natural number  $\geq$  2, and let  $\mathcal C$  be a collection of m simple closed curves in the Euclidean plane such that

- lacktriangledown the interior of a curve in  $\mathcal C$  contains no point of any curve in  $\mathcal C$ ,
- $oldsymbol{@}$  any two curves in  $\mathcal C$  have at most one point common, and
- lacktriangle any curve in  $\mathcal C$  has at most d points of intersection with other curves in  $\mathcal C$ .

Then  $\mathcal{C}$  contains a subset  $\mathcal{C}'$  of at least  $2^{3-2d}3^{1-d}m$  curves such that each connected component  $\mathcal{C}''$  of the union of curves in  $\mathcal{C}'$  contains a point which belongs to all curves in  $\mathcal{C}''$ .



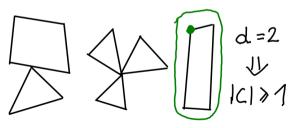
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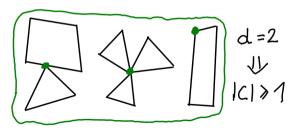
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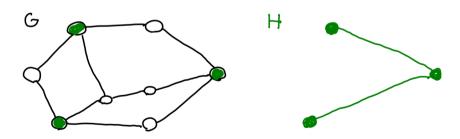
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# Many vertices joined by paths of length 3

#### Theorem 5.

Let G be a graph of girth at least 4, and let A be a set of n pairwise nonadjacent vertices on the outer face boundary. Let H denote the graph with vertex set A such that two vertices in H are neighbors if and only if they are joined by a path of length (precisely) 3 in G. Then H is 16-colorable and has therefore an independent set of at least n/16 vertices.



#### Main Theorem 2.

The number of proper vertex-3-colorings of every triangle-free planar graph with n vertices and with no separating cycle of length 4 or 5 is at least  $2^{n/17700000}$ .

#### Theorem 6.

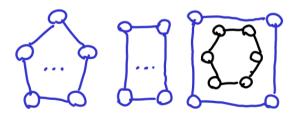
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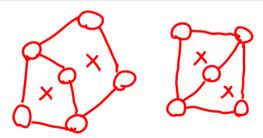


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Assume (reductio ad absurdum) that G is a smallest counterexample for n > 1.

(1): The number of 5-cycles in G is < n/100.

Proof of (1): All 5-cycles in G are facial. If the number of 5-cycles in G is at least n/100, then Theorem 3. implies that G is not a counterexample to Theorem 6, a contradiction.

#### Theorem 3.

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(2): The number of vertices in int(C) which have degree at most 4 and which are not contained in any 5-cycle is at least n/20.

Proof of (2): As G has no triangle, G has n+4 or n+5 vertices and at most 2(n+5)-4 edges, by Euler's formula (F+V-2=E).

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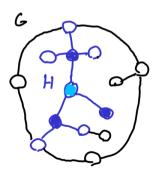
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So, the number described in (2) is at least n - 9n/10 - n/20 = n/20.

- (3): G contains an induced subgraph H such that each component H' of H is a tree containing a vertex v (which we call the root of H') such that
- (i) each H-neighbor x of v is in int(C) and satisfies  $d_G(x) = d_H(x) \le 4$  and is not contained in a 5-cycle, and
- (ii) each vertex in H' has distance (in both G and H) at most 2 to v.

Moreover, at least n/1105920 vertices x in H satisfy  $d_G(x) = d_H(x) \le 4$  and are not contained in a 5-cycle.

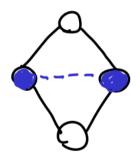


Proof of (3): By (2) and the 4-Color Theorem, G contains an independent set  $I_1$  of at least n/80 vertices, all in int(C), all of degree at most 4, and all not contained in any 5-cycle.

 $\it I_1$  contains a subset  $\it I_2$  with at least  $\it n/320$  vertices such that no two are in a common 4-cycle. (4-Color)

We apply Theorem 4. with d=4 to the family of curves around neighborhood of  $I_2$  vertices and let G' be the subgraph of G consisting of the edges and vertices which are in the resulting subfamily. Then G' contains at least n/276480 vertices x that satisfy  $d_G(x)=d_{G'}(x)\leq 4$  and are not contained in a 5-cycle.

G' may be not induced. Then some edges with ends in G' connect components of G'. Using the 4-Color Theorem we find a subgraph H of G consisting of pairwise nonadjacent components of G' such that H contains at least n/1105920 vertices x in H that satisfy  $d_G(x) = d_H(x) \le 4$  and are not contained in a 5-cycle.

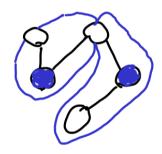


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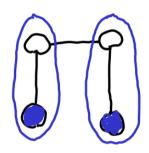


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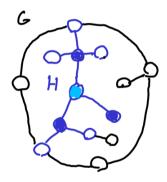


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(4): The subgraph H defined in (3) contains an induced subgraph Q with at least n/17694720 vertices x satisfying  $d_G(x) = d_Q(x) \le 4$  and x not contained in a 5-cycle, such that the contraction of each component Q' of Q into a vertex q results in a triangle-free planar graph.

Proof of (4): H could work, but triangle may occur after contraction.

Such triangle have only one contracted vertex (H is induced).

Such triangle corresponds to path P of length 3 in G.

P must join two vertices  $y_1, y_2$  which are neighbors of distinct vertices  $x_1, x_2$  which are both neighbors of the root v.

We delete v, and for each neighbor x of v we contract x and all (at most 3) neighbors of x in  $G \setminus v$  into a single vertex.

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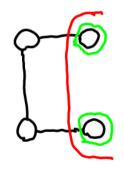
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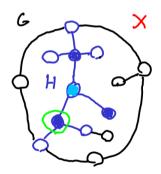
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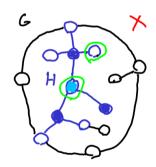
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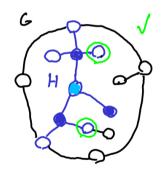
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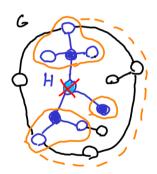
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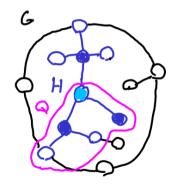
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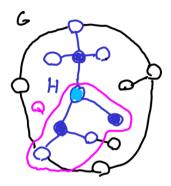


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We color roots and grandchildren in Q' components to the c(q').

Now, there are at least n/17694720 uncolored vertices in  $\it G$  with same color neighbors.

There are at least  $2^{n/17694720}$  color extensions of the coloring of C. Contradiction, G was not a counterexample.



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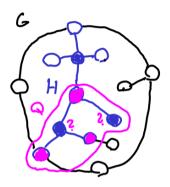
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$$20 = 20$$

Degree  $\leq$  4 and no 5-cycle

$$20 = 20$$
  
 $20 \cdot 4 = 80$ 

$$\mbox{Degree} \leq \mbox{4 and no 5-cycle} \\ \mbox{Independent set}$$

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$$20 \cdot 4 = 80$$

$$20\cdot 4\cdot 4=320$$

Degree  $\leq$  4 and no 5-cycle Independent set No common 4-cycle

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$$20 = 20$$
$$20 \cdot 4 = 80$$
$$20 \cdot 4 \cdot 4 = 320$$
$$20 \cdot 4 \cdot 4 \cdot 864 = 276480$$

 $\label{eq:decomposition} \begin{array}{l} \mbox{Degree} \leq 4 \mbox{ and no 5-cycle} \\ \mbox{Independent set} \\ \mbox{No common 4-cycle} \\ \mbox{Snowflakes} \end{array}$ 

$$20 = 20 \\ 20 \cdot 4 = 80 \\ 20 \cdot 4 \cdot 4 = 320 \\ 20 \cdot 4 \cdot 4 \cdot 864 = 276480$$
 Degree  $\leq 4$  and no 5-cycle Independent set No common 4-cycle Snowflakes Snowflakes Subgraph  $H$ 

### References

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