



# Integer Programming



# *Introduction*

- When one or more variables in an LP problem must assume an integer value we have an Integer Linear Programming (ILP) problem.
  - ILPs occur frequently...
    - Scheduling workers
    - Manufacturing airplanes
  - Integer variables also allow us to build more accurate models for a number of common problems.
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# Relaxation

- Original ILP

$$\text{MAX: } 2X_1 + 3X_2$$

$$\text{S.T.: } X_1 + 3X_2 \leq 8.25$$

$$2.5X_1 + X_2 \leq 8.75$$

$$X_1, X_2 \geq 0$$

$X_1, X_2$  must be integers

- LP Relaxation

$$\text{MAX: } 2X_1 + 3X_2$$

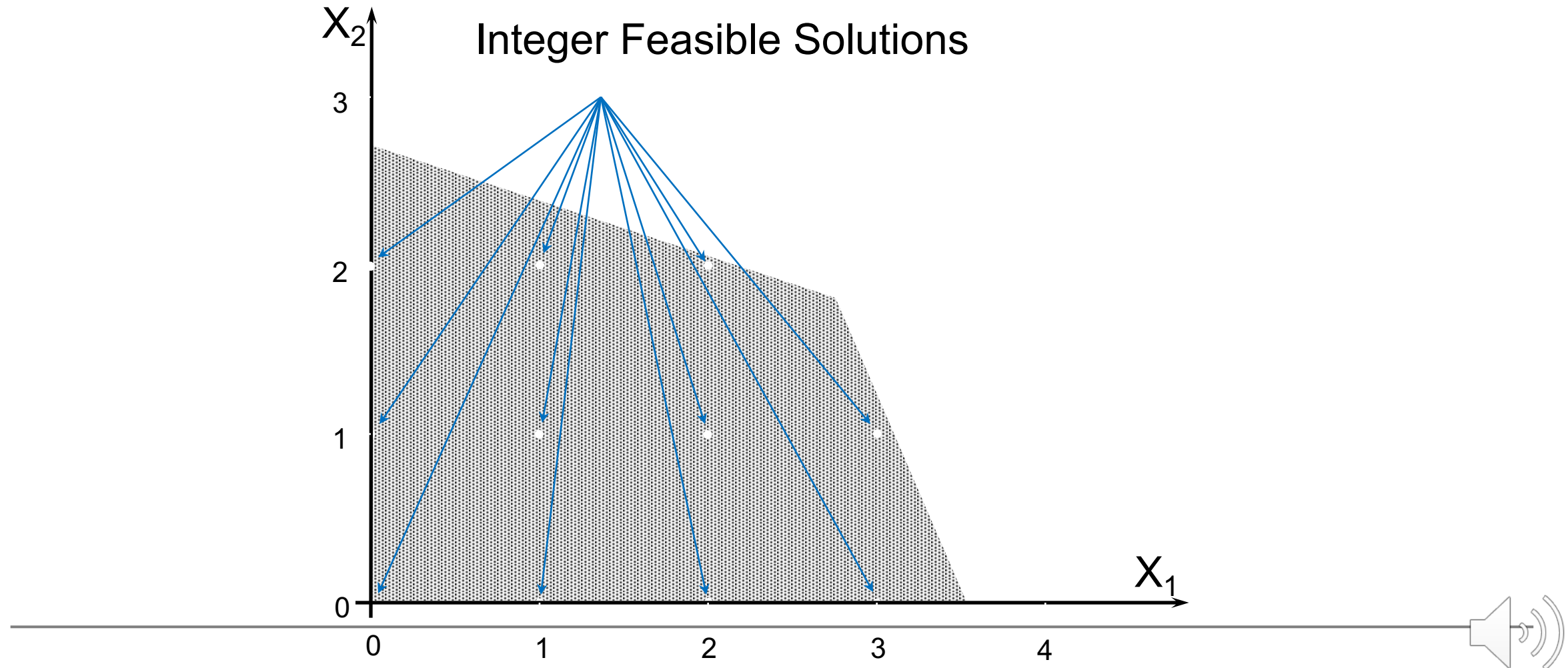
$$\text{S.T.: } X_1 + 3X_2 \leq 8.25$$

$$2.5X_1 + X_2 \leq 8.75$$

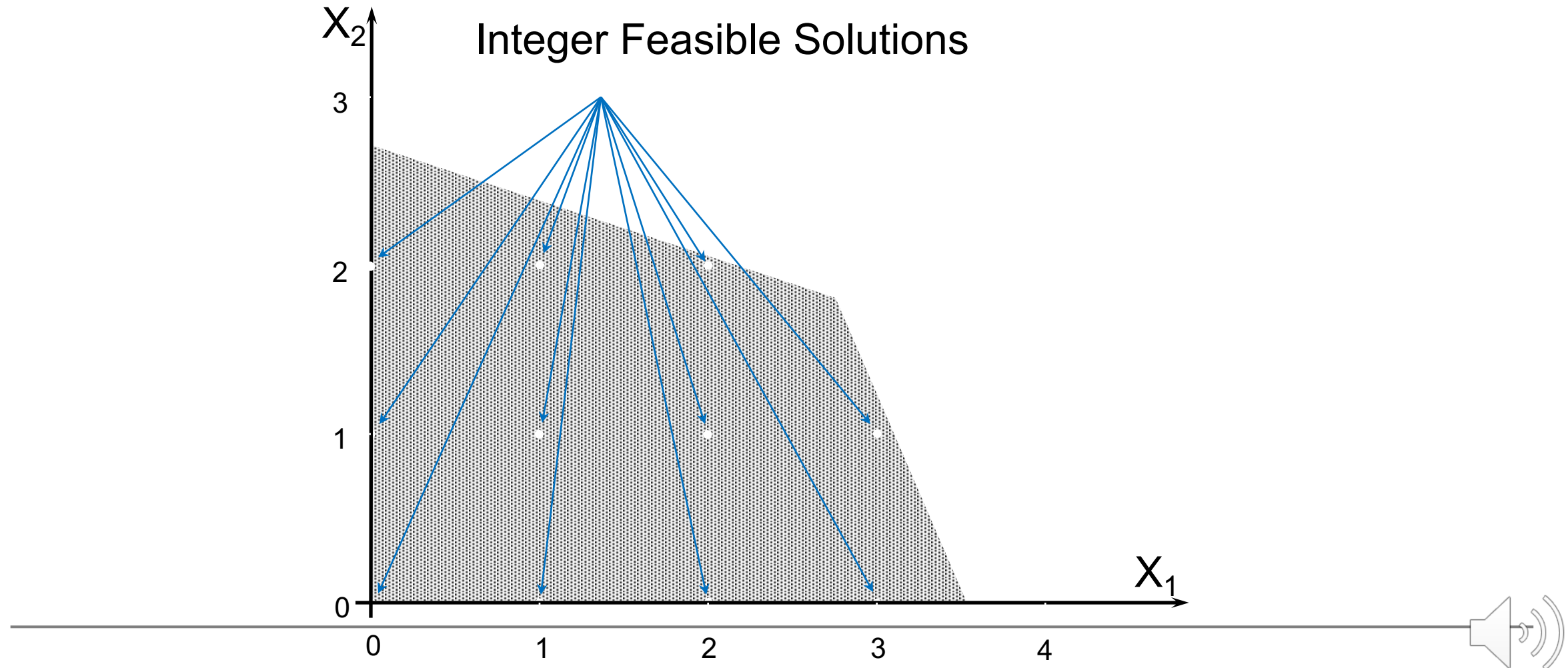
$$X_1, X_2 \geq 0$$



## *Integer Feasible vs. LP Feasible Region*



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## *Integrality Conditions*

$$\begin{array}{ll} \text{MAX: } 350X_1 + 300X_2 & \text{\} profit} \\ \text{S.T.: } 1X_1 + 1X_2 \leq 200 & \text{\} pumps} \\ & 9X_1 + 6X_2 \leq 1566 \quad \text{\} labor} \\ & 12X_1 + 16X_2 \leq 2880 \quad \text{\} tubing} \\ & X_1, X_2 \geq 0 \quad \text{\} nonnegativity} \\ & X_1, X_2 \text{ must be integers} \quad \text{\} integrality} \end{array}$$

Integrality conditions are easy to state but make the problem much more difficult (and sometimes impossible) to solve.



## *Solving ILP Problems*

- When solving an LP relaxation, sometimes you “get lucky” and obtain an integer feasible solution.
- This was the case in the original Blue Ridge Hot Tubs problem in earlier chapters.
- But what if we reduce the amount of labor available to 1520 hours and the amount of tubing to 2650 feet?



## *Integrality Conditions*

$$\begin{array}{ll} \text{MAX: } 350X_1 + 300X_2 & \text{\} \text{ profit} \\ \text{S.T.: } 1X_1 + 1X_2 \leq 200 & \text{\} \text{ pumps} \\ & 9X_1 + 6X_2 \leq 1520 \text{\} \text{ labor} \\ & 12X_1 + 16X_2 \leq 2650 \text{\} \text{ tubing} \\ & X_1, X_2 \geq 0 \text{\} \text{ nonnegativity} \\ & X_1, X_2 \text{ must be integers} \text{\} \text{ integrality} \end{array}$$

Optimal solution of the relaxed problem:

$$X_1 = 116,9444 \quad X_2 = 77,9167$$

Corresponding to a maximum profit of \$64306





# *Rounding*

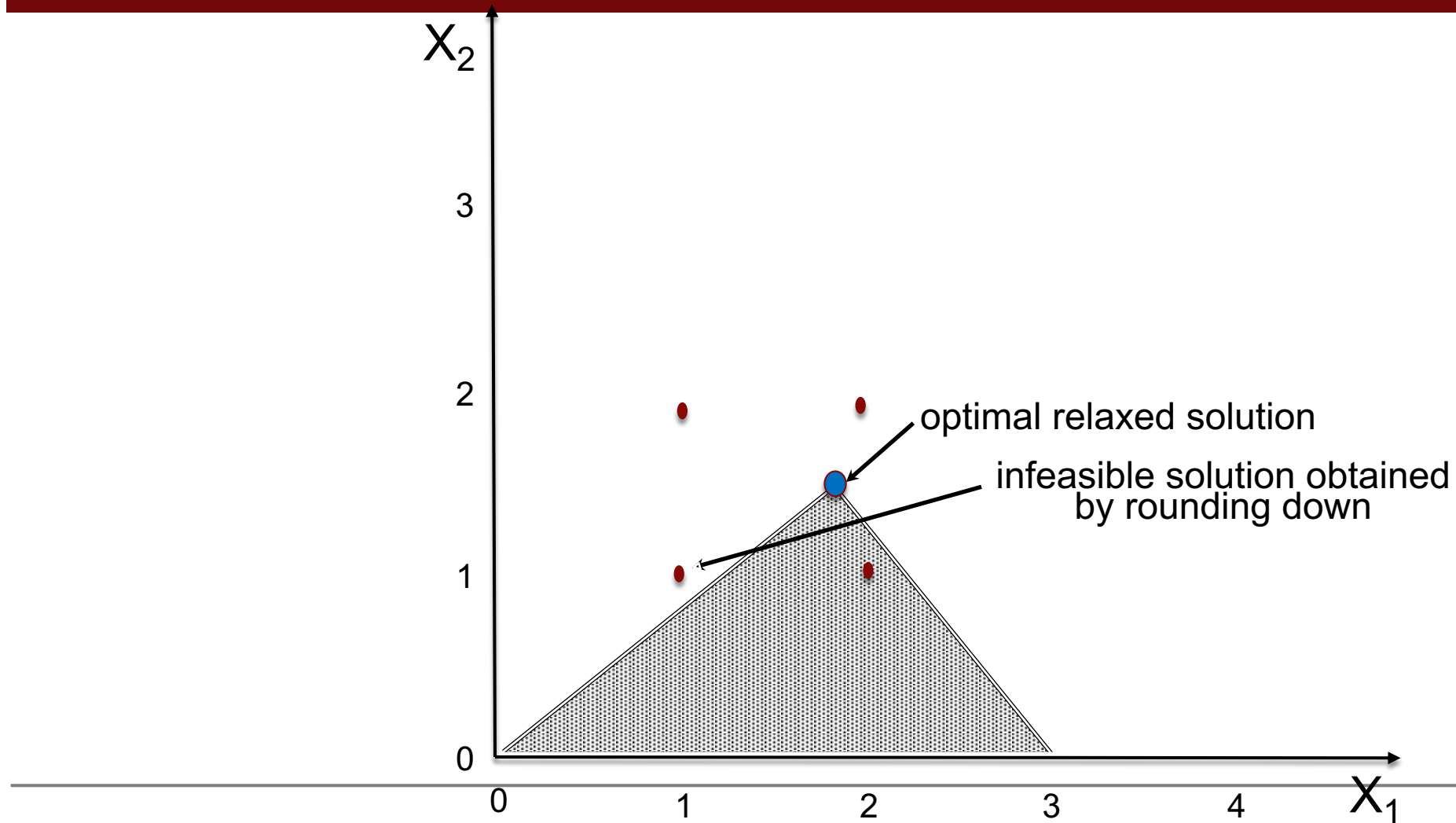
- It is tempting to simply round a fractional solution to the closest integer solution.
- In general, this does not work reliably:
  - The rounded solution may be infeasible.
  - The rounded solution may be suboptimal.



## *How Rounding Down Can Result in an Infeasible Solution*

Blue Ridge Hot Tubs				
	Aqua-Spas	Hydro-Luxes		
Number to Make	117	78	Total Profit	
Unit Profits	\$350	\$300	\$64,350	
Constraints			Used	Available
Pumps Req'd	1	1	195	200
Labor Req'd	9	6	1521	1520
Tubing Req'd	12	16	2652	2650

# *How Rounding Down Can Result in an Infeasible Solution*



# How Rounding Down Can Result in an Infeasible Solution

Blue Ridge Hot Tubs				
	Aqua-Spas	Hydro-Luxes		
Number to Make	116	77	Total Profit	
Unit Profits	\$350	\$300	\$63,700	
Constraints			Used	Available
Pumps Req'd	1	1	193	200
Labor Req'd	9	6	1506	1520
Tubing Req'd	12	16	2624	2650

## *Stopping Rules*

- Because B&B can take so long, most ILP packages allow you to specify a suboptimality tolerance factor.
- This allows you to stop once an integer solution is found that is within some % of the global optimal solution.
- Bounds obtained from LP relaxations are helpful here.
  - Example
    - LP relaxation has an optimal obj. value of \$64,306.
    - 95% of \$64,306 is \$61,090.
    - Thus, an integer solution with obj. value of \$61,090 or better must be within 5% of the optimal solution.



# Bounds

- The optimal solution to an LP relaxation of an ILP problem gives us a *bound* on the optimal objective function value.
- For **maximization** problems, the optimal relaxed objective function value is an *upper bound* on the optimal integer value.
- For **minimization** problems, the optimal relaxed objective function value is a *lower bound* on the optimal integer value.



# Branch-and-Bound

- The Branch-and-Bound (B&B) algorithm can be used to solve ILP problems.
- Requires the solution of a series of LP problems termed “candidate problems”.
- Theoretically, this can solve any ILP.
- Practically, it often takes LOTS of computational effort (and time).

