Mattan Raj.
Anumugam Pajkuma.
1002231625 the possible state on when the world

x: ance and that taken on youth A overcount Sketemant x=1, with time operation O(1). Outer loop,

n times grom i=1 to i=n.

inner loop, will not who halor of suffer esponent

ntimes from j=1, to j=n.

Statement $\alpha = x + 1$, givens in times.

Total no. of executions = nan.

Total nuntime, T(n) = nxn (I(n) = n2

$$\frac{2 \operatorname{code} \operatorname{lobation}}{\operatorname{total}} = \sum_{i=1}^{n} \sum_{j=1}^{n} \operatorname{lobation} \operatorname{configuration} \operatorname{configuration}$$

6

6

. Total executions in p2. The nuntime of algorithm = O(n2). 8. The time taken by f(n) is measured, and tested for various n values, ex: n=1,2,3... the measured time is not posted values of n on x-asis and time taken on y cosis. A quadratic pattern plot is received with consistence O(n2). BO, time complexity is quadratic.

in The graph shows a quadratic trend where n increases which validates the time complexity $\Theta(n^2)$.

3. Upper bound (Big O): is specified by light blue that slightly exceeds the polynomial fitted curve.

This indicates that the fr is bounded above O(n2).

This includes time complexity in the worst care
as fast as n2.

lower bound: Orange line thous the actual curve which is bounded by $\Omega(n^2)$. This shows that n^2 to complete

Tight bound (Big theta): Since the upper and lower tounds are quadrate that pollows actual timing data. so, the suntime is $\Theta(n^2)$.

He Identifying No:

The plot quadratic trend no is marked with vertical dashed lines at n = 4800. The data shows raxialrity due to overhead or hoise in system.

No suppresent threshold whose the algorithm's performance aligns with expected no tomplexity.

the malified for will take more time to own because of the added operation y=1+j, within inner loop which is consistent with time O(1), since its executed same no of times like other operations. The asymptotic complexity does not change.

5. No., it won't affect the runtime analysis result sink the operation is O(1) and therefore the runtime remains O(n²) original and modified function perform nested loops for n² times in a constant time.

Therefore,

Poig 0: $O(n^2)$ Big Omega: $O(n^2)$ Poig theta: $O(n^2)$