

Construção de Compiladores

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3ª Etapa do Projeto

Uberlândia - MG

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Parte 1

Código em c:

```
#include <stdio.h>
#include <stdlib.h>
```

```
int main() {
    return 0;
}
```

Código do IR:

```
; ModuleID = 'test.c'
source_filename = "test.c"
target datalayout = "e-m:e-i64:64-f80:128-n8:16:32:64-S128"
target triple = "x86_64-pc-linux-gnu"

; Function Attrs: noinline nounwind optnone sspstrong uwtable
define dso_local i32 @main() #0 {
    %1 = alloca i32, align 4
    store i32 0, i32* %1, align 4
    ret i32 0
}

attributes #0 = { noinline nounwind optnone sspstrong uwtable
"correctly-rounded-divide-sqrt-fp-math"="false" "disable-tail-calls"="false"
"less-precise-fpmad"="false" "min-legal-vector-width"="0" "no-frame-pointer-elim"="true"
"no-frame-pointer-elim-non-leaf" "no-infs-fp-math"="false" "no-jump-tables"="false"
"no-nans-fp-math"="false" "no-signed-zeros-fp-math"="false" "no-trapping-math"="false"
"stack-protector-buffer-size"="8" "target-cpu"="x86-64"
"target-features"="+fxsr,+mmx,+sse,+sse2,+x87" "unsafe-fp-math"="false"
"use-soft-float"="false" }

!llvm.module.flags = !{!0, !1, !2}
!llvm.ident = !{!3}

!0 = !{i32 1, !"wchar_size", i32 4}
!1 = !{i32 7, !"PIC Level", i32 2}
!2 = !{i32 7, !"PIE Level", i32 2}
!3 = !{!"clang version 8.0.1 (tags/RELEASE_801/final)"}
```

Nesse código dá pra perceber que ele mapeia a função main() alocando um inteiro de 32 bits, armazenando o valor com zero e depois retornando o valor alocado

Código em c:

```
#include <stdio.h>
#include <stdlib.h>
```

```
int main() {
    int x = 4;
    float y = 5;
    return 0;
}
```

Código do IR:

```
; ModuleID = 'test.c'
source_filename = "test.c"
target datalayout = "e-m:e-i64:64-f80:128-n8:16:32:64-S128"
target triple = "x86_64-pc-linux-gnu"
```

```
; Function Attrs: noline nounwind optnone sspstrong uwtable
define dso_local i32 @main() #0 {
    %1 = alloca i32, align 4
    %2 = alloca i32, align 4
    %3 = alloca float, align 4
    store i32 0, i32* %1, align 4
    store i32 4, i32* %2, align 4
    store float 5.000000e+00, float* %3, align 4
    ret i32 0
}
```

```
attributes #0 = { noline nounwind optnone sspstrong uwtable
"correctly-rounded-divide-sqrt-fp-math"="false" "disable-tail-calls"="false"
"less-precise-fpmad"="false" "min-legal-vector-width"="0" "no-frame-pointer-elim"="true"
"no-frame-pointer-elim-non-leaf" "no-infs-fp-math"="false" "no-jump-tables"="false"
"no-nans-fp-math"="false" "no-signed-zeros-fp-math"="false" "no-trapping-math"="false"
"stack-protector-buffer-size"="8" "target-cpu"="x86-64"
"target-features"="+fxsr,+mmx,+sse,+sse2,+x87" "unsafe-fp-math"="false"
"use-soft-float"="false" }
```

```
!llvm.module.flags = !{!0, !1, !2}
!llvm.ident = !{!3}
```

```
!0 = !{i32 1, !"wchar_size", i32 4}
!1 = !{i32 7, !"PIC Level", i32 2}
!2 = !{i32 7, !"PIE Level", i32 2}
!3 = !{!"clang version 8.0.1 (tags/RELEASE_801/final)"}
```

Nesse já dá pra perceber que as variáveis criadas são alocadas dentro do define da função main e são dados números sequenciais como IDs das variáveis

```
#include <stdio.h>
#include <stdlib.h>
```

Código do IR:

```
; Function Attrs: noline nounwind optnone sspstrong uwtable
```

```
define dso_local i32 @main() #0 {  
    %1 = alloca i32, align 4  
    %2 = alloca i32, align 4  
    %3 = alloca float, align 4  
    store i32 0, i32* %1, align 4  
    store i32 4, i32* %2, align 4  
    store float 5.000000e+00, float* %3, align 4  
    %4 = load float, float* %3, align 4  
    %5 = fcmp oeq float %4, 5.000000e+00  
    br i1 %5, label %6, label %8  
  
; <label>:6:                                ; preds = %0  
    %7 = call i32 (@i8*, ...) @printf(i8* getelementptr inbounds ([6 x i8], [6 x i8]* @.str, i32 0,  
i32 0))  
    br label %8  
  
; <label>:8:                                ; preds = %6, %0  
    ret i32 0  
}
```

```
declare i32 @printf(i8*, ...) #1
```

```
attributes #0 = { noline nounwind optnone sspstrong uwtable
"correctly-rounded-divide-sqrt-fp-math"="false" "disable-tail-calls"="false"
"less-precise-fpmad"="false" "min-legal-vector-width"="0" "no-frame-pointer-elim"="true"
"no-frame-pointer-elim-non-leaf" "no-infs-fp-math"="false" "no-jump-tables"="false"
"no-nans-fp-math"="false" "no-signed-zeros-fp-math"="false" "no-trapping-math"="false"
"stack-protector-buffer-size"="8" "target-cpu"="x86-64"
"target-features"="+fxsr,+mmx,+sse,+sse2,+x87" "unsafe-fp-math"="false"
"use-soft-float"="false" }
attributes #1 = { "correctly-rounded-divide-sqrt-fp-math"="false" "disable-tail-calls"="false"
"less-precise-fpmad"="false" "no-frame-pointer-elim"="true" "no-frame-pointer-elim-non-leaf"
"no-infs-fp-math"="false" "no-nans-fp-math"="false" "no-signed-zeros-fp-math"="false"
"no-trapping-math"="false" "stack-protector-buffer-size"="8" "target-cpu"="x86-64"
"target-features"="+fxsr,+mmx,+sse,+sse2,+x87" "unsafe-fp-math"="false"
"use-soft-float"="false" }
```

```
!llvm.module.flags = !{!0, !1, !2}
```

```
!llvm.ident = !{!3}
```

```
!0 = !{i32 1, !"wchar_size", i32 4}
```

```
!1 = !{i32 7, !"PIC Level", i32 2}
```

```
!2 = !{i32 7, !"PIE Level", i32 2}
```

```
!3 = !{!"clang version 8.0.1 (tags/RELEASE_801/final)"}
```

Aqui dá pra perceber que ao comparar dois numeros, ele aloca um numero para guardar o y, e outro para alocar o 5. Depois usa um comando de comparacao com algo parecido com um goto e redireciona dependendo do resultado

Código em c:

```
#include <stdio.h>
```

```
#include <stdlib.h>
```

```
int main() {
```

```
    int x = 4;
```

```
    float y = 5;
```

```
    if (y == 5) {
        printf("Hello");
    }
```

```
    while (x < 10) {
        x++;
    }
```

```
return 0;
}
```

Código do IR:

```
; ModuleID = 'test.c'
source_filename = "test.c"
target datalayout = "e-m:e-i64:64-f80:128-n8:16:32:64-S128"
target triple = "x86_64-pc-linux-gnu"

@.str = private unnamed_addr constant [6 x i8] c"Hello\00", align 1

; Function Attrs: noinline nounwind optnone sspstrong uwtable
define dso_local i32 @main() #0 {
    %1 = alloca i32, align 4
    %2 = alloca i32, align 4
    %3 = alloca float, align 4
    store i32 0, i32* %1, align 4
    store i32 4, i32* %2, align 4
    store float 5.000000e+00, float* %3, align 4
    %4 = load float, float* %3, align 4
    %5 = fcmp oeq float %4, 5.000000e+00
    br i1 %5, label %6, label %8

; <label>:6:                                     ; preds = %0
    %7 = call i32 @printf(i8* getelementptr inbounds ([6 x i8], [6 x i8]* @.str, i32 0,
i32 0))
    br label %8

; <label>:8:                                     ; preds = %6, %0
    br label %9

; <label>:9:                                     ; preds = %12, %8
    %10 = load i32, i32* %2, align 4
    %11 = icmp slt i32 %10, 10
    br i1 %11, label %12, label %15

; <label>:12:                                    ; preds = %9
    %13 = load i32, i32* %2, align 4
    %14 = add nsw i32 %13, 1
    store i32 %14, i32* %2, align 4
    br label %9

; <label>:15:                                    ; preds = %9
```

```

ret i32 0
}

declare i32 @printf(i8*, ...) #1

attributes #0 = { noline nounwind optnone sspstrong uwtable
"correctly-rounded-divide-sqrt-fp-math"="false" "disable-tail-calls"="false"
"less-precise-fpmad"="false" "min-legal-vector-width"="0" "no-frame-pointer-elim"="true"
"no-frame-pointer-elim-non-leaf" "no-infs-fp-math"="false" "no-jump-tables"="false"
"no-nans-fp-math"="false" "no-signed-zeros-fp-math"="false" "no-trapping-math"="false"
"stack-protector-buffer-size"="8" "target-cpu"="x86-64"
"target-features"="+fxsr,+mmx,+sse,+sse2,+x87" "unsafe-fp-math"="false"
"use-soft-float"="false" }
attributes #1 = { "correctly-rounded-divide-sqrt-fp-math"="false" "disable-tail-calls"="false"
"less-precise-fpmad"="false" "no-frame-pointer-elim"="true" "no-frame-pointer-elim-non-leaf"
"no-infs-fp-math"="false" "no-nans-fp-math"="false" "no-signed-zeros-fp-math"="false"
"no-trapping-math"="false" "stack-protector-buffer-size"="8" "target-cpu"="x86-64"
"target-features"="+fxsr,+mmx,+sse,+sse2,+x87" "unsafe-fp-math"="false"
"use-soft-float"="false" }

!llvm.module.flags = !{!0, !1, !2}
!llvm.ident = !{!3}

!0 = !{i32 1, !"wchar_size", i32 4}
!1 = !{i32 7, !"PIC Level", i32 2}
!2 = !{i32 7, !"PIE Level", i32 2}
!3 = !{!"clang version 8.0.1 (tags/RELEASE_801/final)"}

```

Com o while, ele também faz igual a um if, faz a comparação que direciona a blocos (<label>) diferentes dependendo do resultado. A diferença é que sempre tem uma opção que volta na primeira <label> do while. E também quando o x é adicionado de 1 unidade, é usado uma variável a mais para calcular a soma e depois salvar na variável original

Código em c:

```

#include <stdio.h>
#include <stdlib.h>

int main() {
    int i = 0;
    for (i = 0; i < 5; i++) {
        printf("%d\n", i);
    }

    return 0;
}

```

```
}
```

Código do IR:

```
; ModuleID = 'test.c'
source_filename = "test.c"
target datalayout = "e-m:e-i64:64-f80:128-n8:16:32:64-S128"
target triple = "x86_64-pc-linux-gnu"

@.str = private unnamed_addr constant [4 x i8] c"%d\0A\00", align 1

; Function Attrs: noline nounwind optnone sspstrong uwtable
define dso_local i32 @main() #0 {
    %1 = alloca i32, align 4
    %2 = alloca i32, align 4
    store i32 0, i32* %1, align 4
    store i32 0, i32* %2, align 4
    store i32 0, i32* %2, align 4
    br label %3

; <label>:3:                                ; preds = %9, %0
    %4 = load i32, i32* %2, align 4
    %5 = icmp slt i32 %4, 5
    br i1 %5, label %6, label %12

; <label>:6:                                ; preds = %3
    %7 = load i32, i32* %2, align 4
    %8 = call i32 @printf(i8* getelementptr inbounds ([4 x i8], [4 x i8]* @.str, i32 0,
i32 0), i32 %7)
    br label %9

; <label>:9:                                ; preds = %6
    %10 = load i32, i32* %2, align 4
    %11 = add nsw i32 %10, 1
    store i32 %11, i32* %2, align 4
    br label %3

; <label>:12:                               ; preds = %3
    ret i32 0
}

declare i32 @printf(i8*, ...) #1

attributes #0 = { noline nounwind optnone sspstrong uwtable
"correctly-rounded-divide-sqrt-fp-math"="false" "disable-tail-calls"="false"
"less-precise-fpmad"="false" "min-legal-vector-width"="0" "no-frame-pointer-elim"="true"
```



```

"no-frame-pointer-elim-non-leaf" "no-infs-fp-math"="false" "no-jump-tables"="false"
"no-nans-fp-math"="false" "no-signed-zeros-fp-math"="false" "no-trapping-math"="false"
"stack-protector-buffer-size"="8" "target-cpu"="x86-64"
"target-features"="+fxsr,+mmx,+sse,+sse2,+x87" "unsafe-fp-math"="false"
"use-soft-float"="false" }
attributes #1 = { "correctly-rounded-divide-sqrt-fp-math"="false" "disable-tail-calls"="false"
"less-precise-fpmad"="false" "no-frame-pointer-elim"="true" "no-frame-pointer-elim-non-leaf"
"no-infs-fp-math"="false" "no-nans-fp-math"="false" "no-signed-zeros-fp-math"="false"
"no-trapping-math"="false" "stack-protector-buffer-size"="8" "target-cpu"="x86-64"
"target-features"="+fxsr,+mmx,+sse,+sse2,+x87" "unsafe-fp-math"="false"
"use-soft-float"="false" }

```

```

!llvm.module.flags = !{!0, !1, !2}
!llvm.ident = !{!3}

```

```

!0 = !{i32 1, !"wchar_size", i32 4}
!1 = !{i32 7, !"PIC Level", i32 2}
!2 = !{i32 7, !"PIE Level", i32 2}
!3 = !{"clang version 8.0.1 (tags/RELEASE_801/final)"}

```

Pelo que foi percebido os comandos de iteracao, trabalham da mesma forma, usando goto para ir a blocos <label> diferentes de acordo com a comparacao. O for nao foge da regra e trabalha da mesma forma, porem com 1 operacao a mais de adicao. Já o printf é um comando diferente que recebe as variaveis alocadas como argumento.

Parte 2

Gramática que especifica a Linguagem:

$G = (V, T, P, S)$

$V = \{ S, \text{bloco}, \text{declaracao}, \text{comandos}, \text{tipo}, \text{condicao}, \text{expressao}, \text{termo}, \text{relop}, \text{artop}, \text{atrop} \}$

$T = \{ \text{int}, \text{char}, \text{real}, \text{se}, \text{enquanto}, ==, <, >, <=, >=, +, -, *, /, =, \text{id}, \epsilon, \text{numero}, ;, . \}$

$P = \{ S \rightarrow \text{programa bloco}, \text{bloco} \rightarrow \text{inicio declaracao comandos fim}, \text{declaracao} \rightarrow \text{tipo id};$
 $\text{declaracao} \mid \epsilon, \text{tipo} \rightarrow \text{int} \mid \text{char} \mid \text{real}, \text{comandos} \rightarrow \text{se} (\text{condicao}) \text{bloco comandos} \mid$
 $\text{enquanto} (\text{condicao}) \text{bloco comandos} \mid \text{id atop expressao}; \text{comandos} \mid \epsilon, \text{condicao} \rightarrow$
 $\text{termo relop termo}, \text{expressao} \rightarrow \text{expressao artop termo} \mid \text{expressao atop termo} \mid \text{termo},$
 $\text{termo} \rightarrow \text{id} \mid \text{numero} \mid (\text{expressao}),$
 $\text{relop} \rightarrow == \mid < \mid > \mid <= \mid >=,$
 $\text{artop} \rightarrow + \mid - \mid * \mid /,$
 $\text{atrop} \rightarrow = \}$

Tokens:

ID (identificadores -> contem letras, digitos e _)

numero (constantes numericas)

programa

inicio

fim

enquanto

relop (operadores relacionais -> ==, <, >, <=, >=)

artop (operadores aritmeticos -> +, -, *, /)

atrop (operadores de atribuicao -> =)

se

tipo (tipos de variaveis -> int, char, real)

simbolos (simbolos usados no codigo -> (,))

Expressões Regulares dos Lexemas:

letra -> [a-zA-Z]

digito -> [0-9]

digitos -> digito+

numero -> digitos ($\epsilon \mid (.digitos)$) ($\epsilon \mid (E (+ \mid - \mid \epsilon) digitos)$)

id -> letra (letra \mid _ \mid digito) *

programa -> programa

inicio -> inicio

fim -> fim

enquanto -> enquanto

relop -> == \mid < \mid > \mid <= \mid >=

atrop -> =

artop -> + \mid - \mid * \mid /

tipo -> int | char | real

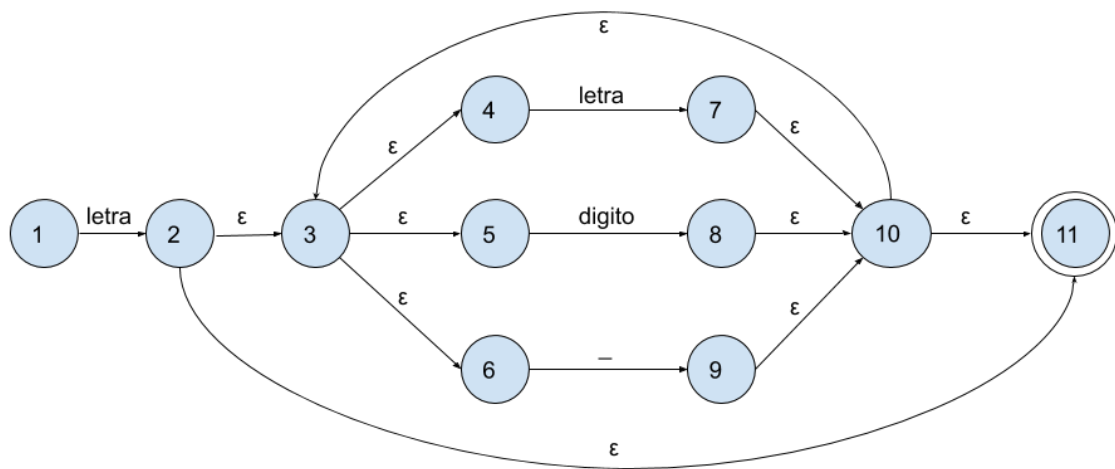
se -> se

simbolos -> (|)

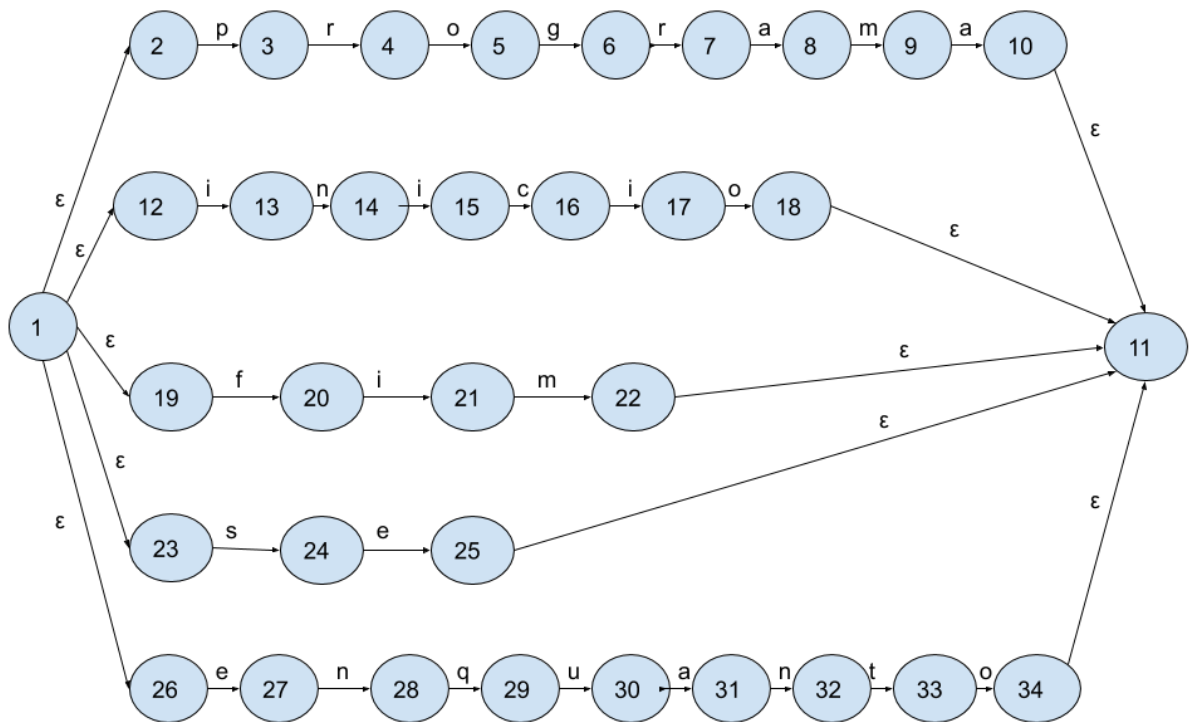
Parte 3

Diagramas

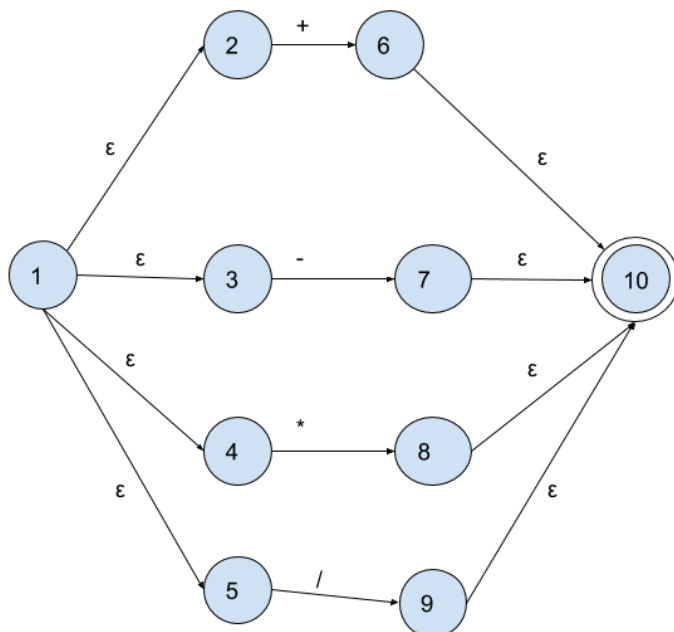
ID



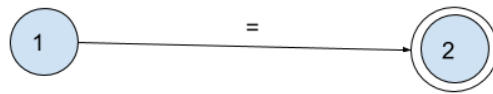
Palavras Reservadas



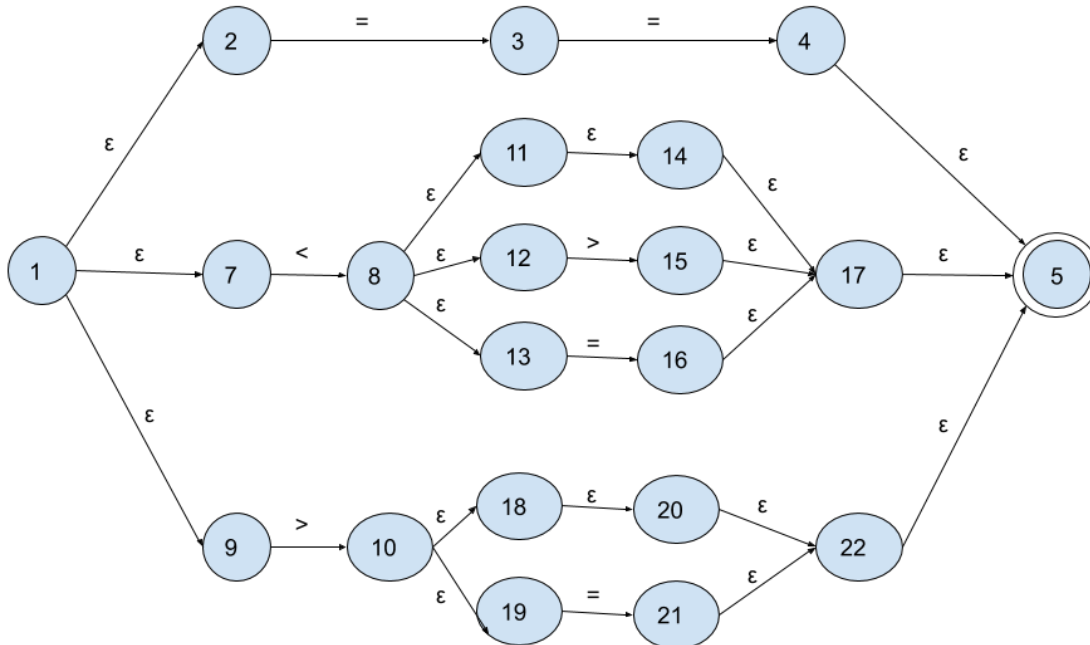
Artop (Operadores Aritmeticos)



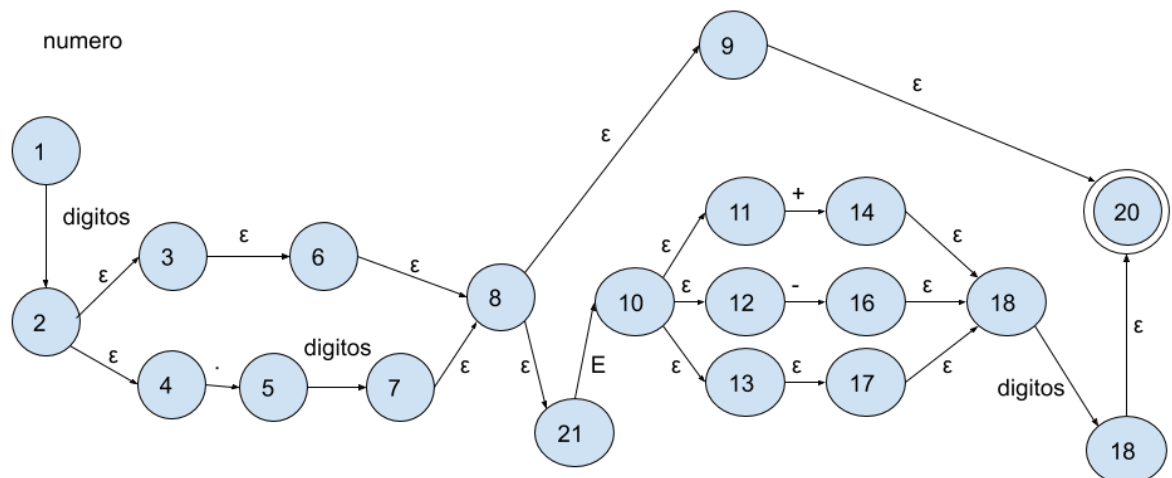
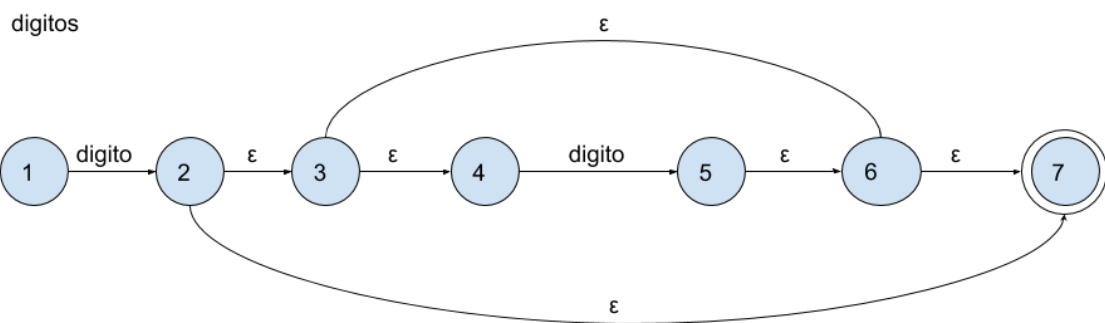
Atrop (Operador de Atribuicao)



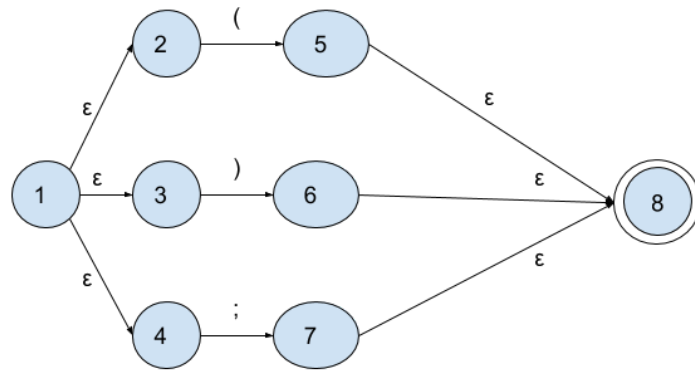
Relop (Operadores relacionais)



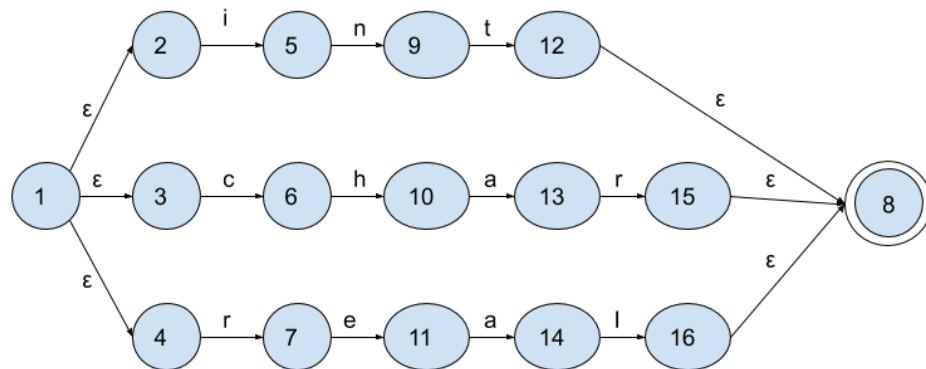
Numero



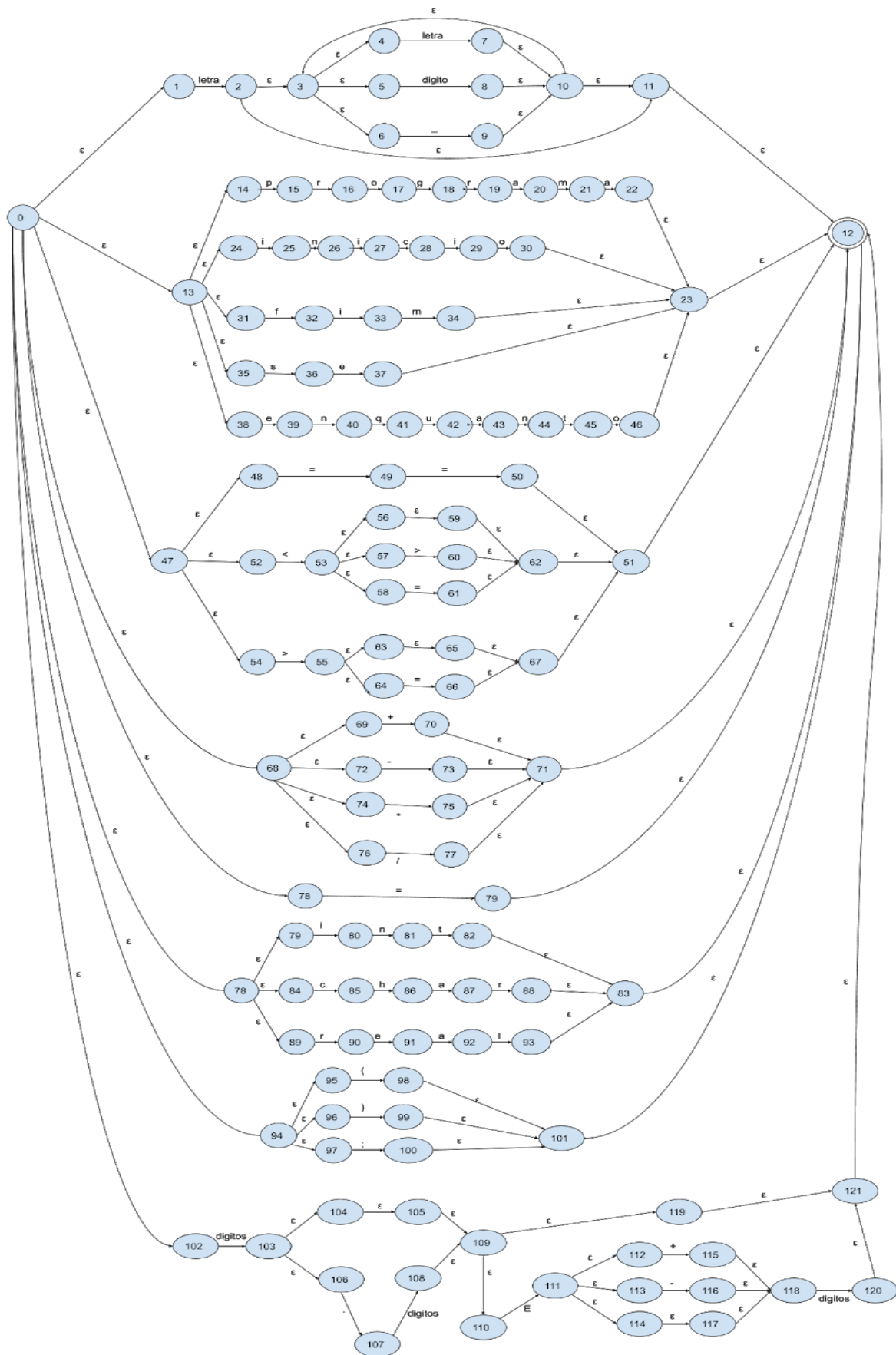
Simbolos



Tipo

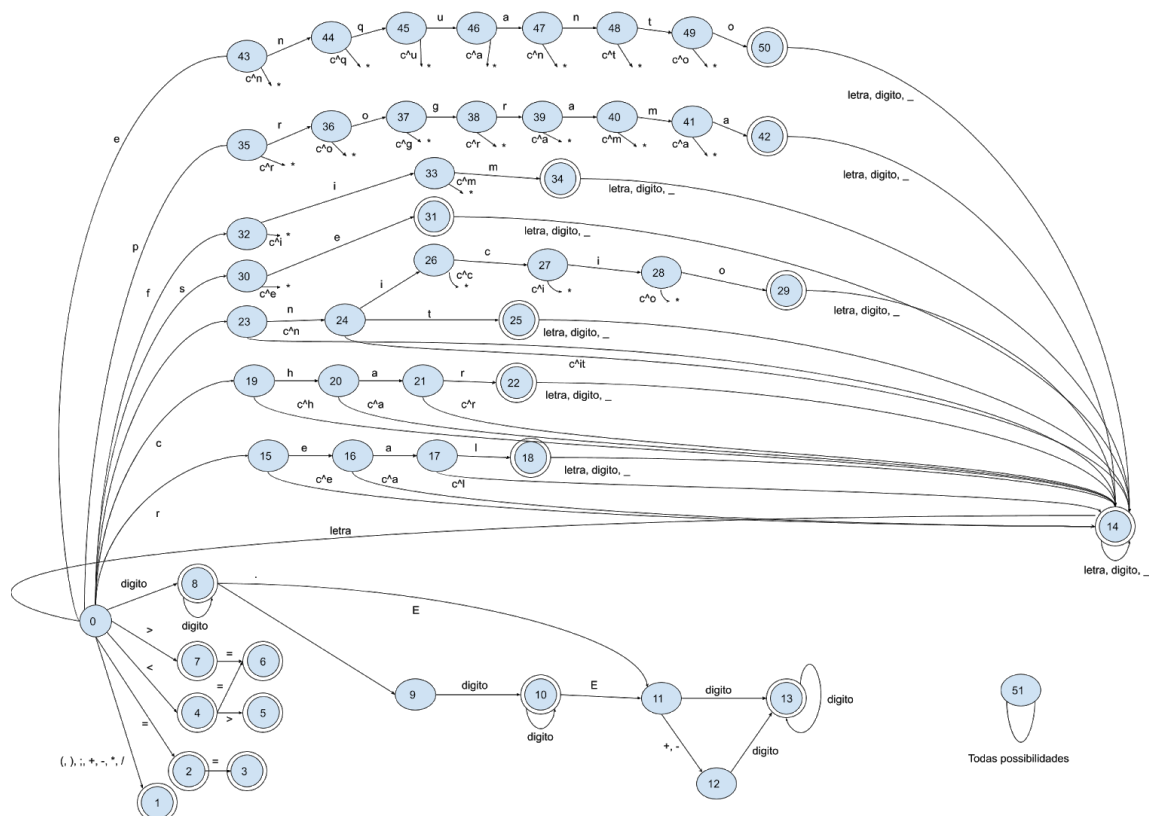


Afnd Uniao



Afd da Uniao

As possíveis saídas dos nos que não estão aparecendo estão ligadas ao nó 51 (Foi para poupar espaço, senão ia ficar muito bagunçado)
 As saídas que tem * significa conectar ao nó 14
 -> [letra digito _ ^z] = c^z



Parte 4

Gramática sem ambiguidade e recursão à esquerda:

$G = (V, T, P, S)$

$V = \{S, \text{bloco}, \text{declaracao}, \text{comandos}, \text{tipo}, \text{condicao}, \text{expressao}, \text{termo}, \text{relop}, \text{artop}, \text{atrop}\}$

$T = \{\text{int}, \text{char}, \text{real}, \text{se}, \text{enquanto}, ==, <, >, <=, >=, +, -, *, /, =, \text{id}, \epsilon, \text{numero}\}$

$P = \{S \rightarrow \text{programa bloco}, \text{bloco} \rightarrow \text{inicio declaracao comandos fim}, \text{declaracao} \rightarrow \text{tipo id};$
 $\text{declaracao} \mid \epsilon, \text{tipo} \rightarrow \text{int} \mid \text{char} \mid \text{real}, \text{comandos} \rightarrow \text{se} (\text{condicao}) \text{bloco comandos} \mid \text{enquanto}$
 $(\text{condicao}) \text{bloco comandos} \mid \text{id atrop expressao comandos} \mid \epsilon, \text{condicao} \rightarrow \text{termo relop termo},$
 $\text{expressao} \rightarrow \text{termo expressao}', \text{expressao}' \rightarrow \text{artop termo expressao}' \mid \text{atrop termo expressao}'$
 $\mid \epsilon, \text{termo} \rightarrow \text{id} \mid \text{numero}, \text{relop} \rightarrow == \mid < \mid > \mid <= \mid >=, \text{artop} \rightarrow + \mid - \mid * \mid / , \text{atrop} \rightarrow =\}$

FIRST e FOLLOW

$\text{FIRST}(S) = \{\text{programa}\}$

$\text{FIRST}(\text{bloco}) = \{\text{inicio}\}$

$\text{FIRST}(\text{declaracao}) = \{\text{tipo}, \epsilon\}$

$\text{FIRST}(\text{tipo}) = \{\text{int}, \text{char}, \text{real}\}$

$\text{FIRST}(\text{comandos}) = \{\text{se}, \text{enquanto}, \text{id}, \epsilon\}$

$\text{FIRST}(\text{condicao}) = \{\text{id}, \text{numero}\}$

$\text{FIRST}(\text{expressao}) = \{\text{id}, \text{numero}\}$

$\text{FIRST}(\text{expressao}') = \{+, -, *, /, =, \epsilon\}$
 $\text{FIRST}(\text{termo}) = \{\text{id}, \text{numero}, ()\}$
 $\text{FIRST}(\text{relop}) = \{==, <, <=, >, >=\}$
 $\text{FIRST}(\text{artop}) = \{+, -, *, /\}$
 $\text{FIRST}(\text{atrop}) = \{=\}$

$\text{FOLLOW}(S) = \{\$ \}$
 $\text{FOLLOW}(\text{bloco}) = \text{FOLLOW}(S) + \text{FIRST}(\text{comandos}) + \text{FOLLOW}(\text{comandos}) = \{\$, \text{se}, \text{enquanto}, \text{id}, \epsilon, \text{fim}\}$
 $\text{FOLLOW}(\text{declaracao}) = \{\text{se}, \text{enquanto}, \text{id}, \text{fim}\}$
 $\text{FOLLOW}(\text{tipo}) = \{\text{id}\}$
 $\text{FOLLOW}(\text{comandos}) = \{\text{fim}\}$
 $\text{FOLLOW}(\text{condicao}) = \{\}$
 $\text{FOLLOW}(\text{expressao}) = \{;,)\}$
 $\text{FOLLOW}(\text{expressao}') = \text{FOLLOW}(\text{expressao}) = \{;,)\}$
 $\text{FOLLOW}(\text{termo}) = \text{FIRST}(\text{relop}) + \text{FOLLOW}(\text{condicao}) + \text{FIRST}(\text{expressao}') +$
 $\text{FOLLOW}(\text{expressao}) + \text{FOLLOW}(\text{expressao}') = \{==, <, <=, >, >=,), +, -, *, /, =, \epsilon, ;\}$
 $\text{FOLLOW}(\text{relop}) = \text{FIRST}(\text{termo}) = \{\text{id}, \text{numero}\}$
 $\text{FOLLOW}(\text{artop}) = \text{FIRST}(\text{termo}) = \{\text{id}, \text{numero}\}$
 $\text{FOLLOW}(\text{atrop}) = \text{FIRST}(\text{termo}) = \{\text{id}, \text{numero}\}$

Verificação de Gramática LL

Regra 1:

$\text{FIRST}(\text{tipo id ; declaracao}) \cap \text{FIRST}(\epsilon) = \emptyset$
 $\text{FIRST}(\text{se (condicao) bloco comandos}) \cap \text{FIRST}(\text{enquanto (condicao) bloco comandos}) \cap$
 $\text{FIRST}(\text{id atrop expressao ; comandos}) \cap \text{FIRST}(\epsilon) = \emptyset$
 $\text{FIRST}(\text{artop termo expressao'}) \cap \text{FIRST}(\text{atrop termo expressao'}) \cap \text{FIRST}(\epsilon) = \emptyset$
 $\text{FIRST}(\text{id}) \cap \text{FIRST}(\text{numero}) \cap \text{FIRST}((\text{expressao})) = \emptyset$
 $\text{FIRST}(==) \cap \text{FIRST}(<) \cap \text{FIRST}(<=) \cap \text{FIRST}(>) \cap \text{FIRST}(>=) = \emptyset$
 $\text{FIRST}(+) \cap \text{FIRST}(-) \cap \text{FIRST}(*) \cap \text{FIRST}(/) = \emptyset$

Regra 2:

$\text{FIRST}(\text{tipo id ; declaracao}) \cap \text{FOLLOW}(\text{declaracao}) = \emptyset$
 $\text{FIRST}(\text{se (condicao) bloco comandos}) \cap \text{FOLLOW}(\text{comandos}) = \emptyset$
 $\text{FIRST}(\text{enquanto (condicao) bloco comandos}) \cap \text{FOLLOW}(\text{comandos}) = \emptyset$
 $\text{FIRST}(\text{id atrop expressao ; comandos}) \cap \text{FOLLOW}(\text{comandos}) = \emptyset$

$\text{FIRST}(\text{artop termo expressao}') \cap \text{FOLLOW}(\text{expressao}') = \emptyset$

$\text{FIRST}(\text{atrop termo expressao}') \cap \text{FOLLOW}(\text{expressao}') = \emptyset$

Como a Gramática obedece as 2 regras então ela é LL

Ordem das Produções:

- 1 - S -> programa bloco
- 2 - bloco -> inicio declaracao comandos fim
- 3 - declaracao -> tipo id ; declaracao
- 4 - declaracao -> ϵ
- 5 - tipo -> int
- 6 - tipo -> char
- 7 - tipo -> real
- 8 - comandos -> se (condicao) bloco comandos
- 9 - comandos -> enquanto (condicao) bloco comandos
- 10 - comandos -> id atrop expressao ; comandos
- 11 - comandos -> ϵ
- 12 - condicao -> termo relop termo
- 13 - expressao -> termo expressao'
- 14 - expressao' -> artop termo expressao'
- 15 - expressao' -> atrop termo expressao'
- 16 - expressao' -> ϵ
- 17 - termo -> id
- 18 - termo -> numero
- 19 - termo -> (expressao)
- 20 - relop -> ==
- 21 - relop -> <>
- 22 - relop -> <
- 23 - relop -> >
- 24 - relop -> <=
- 25 - relop -> >=
- 26 - artop -> +
- 27 - artop -> -
- 28 - artop -> *
- 29 - artop -> /
- 30 - atrop -> =

Tabela de Análise Preditiva

	progr ama	inic io	fi m	id	;	in t	cha r	re al	s e	()	enqu anto	=	<	<	>	>=	+	-	*	/	=	\$	numero
S	1																						
bloco		2																					
declar acao			4	4		3	3	3	4		4												
tipo						5	6	7															
coman dos			11	10					8		9												
condic ao				12						12													12
expres sao				13						13													13
expres sao'				1 6						1 6							1 4	1 4	1 4	1 4	1 5		
termo				17						19													18
relop												2 0	2 1	2 2	2 4	2 3	25						
artop																	2 6	2 7	2 8	2 9			
atrop																					3 0		