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$$1) \quad T(n) = \begin{cases} T(1) = O(1) \\ T(n) = 4T\left(\frac{n}{2}\right) + O(n) \end{cases}$$

$$T(n) = 4T\left(\frac{n}{2}\right) + O(n)$$

$$T(n) = 4\left(4T\left(\frac{n}{2^2}\right) + O\left(\frac{n}{2}\right)\right) + O(n) = 4^2 T\left(\frac{n}{2^2}\right) + 4O\left(\frac{n}{2}\right) + O(n) = 4^2 T\left(\frac{n}{2^2}\right) + 2O(n) + O(n)$$

$$T(n) = 4^2\left(4T\left(\frac{n}{2^3}\right) + O\left(\frac{n}{2^2}\right)\right) + 2 \cdot O(n) + O(n) = 4^3 T\left(\frac{n}{2^3}\right) + 2^2 \cdot O(n) + 2^1 \cdot O(n) + O(n)$$

⋮

$$T(n) = 4^K T\left(\frac{n}{2^K}\right) + 2^{K-1} \cdot O(n) + 2^{K-2} \cdot O(n) + \dots + 2^1 \cdot O(n) + 2^0 \cdot O(n)$$

$$T(n) = 4^K T\left(\frac{n}{2^K}\right) + \sum_{i=0}^{K-1} 2^i \cdot O(n) \\ = 4^K T\left(\frac{n}{2^K}\right) + O(n) \sum_{i=0}^{K-1} 2^i$$

A recursão para quando $\frac{n}{2^K} = 1 \Rightarrow n = 2^K \Rightarrow K = \lg n$

$$T(n) = 4^{\lg n} T(1) + O(n) \left[\frac{1(2^{\lg n} - 1)}{2 - 1} \right]$$

$$= n^{\lg 4} \cdot O(1) + O(n) (2^{\lg n} - 1)$$

$$= n^2 \cdot O(1) + O(n) (n^{\lg 2} - 1)$$

$$= n^2 \cdot O(1) + O(n) [n - 1]$$

$$= O(n^2) + O(n^2 - n)$$

$$T(n) = O(n^2)$$

$$2) T(n) = \begin{cases} O(1) & T(1) \\ 3T(\frac{n}{2}) + O(n) & T(n) \end{cases}$$

$$T(n) = 3T(\frac{n}{2}) + O(n)$$

$$= 3[3T(\frac{n}{2^2}) + O(\frac{n}{2})] + O(n) = 3^2 T(\frac{n}{2^2}) + 3O(\frac{n}{2}) + O(n)$$

$$= 3^2[3T(\frac{n}{2^3}) + O(\frac{n}{2^2})] + 3O(\frac{n}{2}) + O(n) =$$

$$= 3^3 T(\frac{n}{2^3}) + 3^2 O(\frac{n}{2^2}) + 3 \cdot O(\frac{n}{2}) + O(n) =$$

$$= 3^3 T\left(\frac{n}{2^3}\right) + \frac{3^2}{2^2} O(n) + \frac{3}{2} \cdot O(n) + O(n)$$

$$= 3^3 [3T(\frac{n}{2^4}) + O(\frac{n}{2^3})] + \left(\frac{3}{2}\right)^2 O(n) + \frac{3}{2} O(n) + O(n) =$$

$$= 3^4 T\left(\frac{n}{2^4}\right) + \left(\frac{3}{2}\right)^3 O(n) + \left(\frac{3}{2}\right)^2 O(n) + \left(\frac{3}{2}\right)^1 O(n) + \left(\frac{3}{2}\right)^0 O(n)$$

⋮

$$= 3^k T\left(\frac{n}{2^k}\right) + \left(\frac{3}{2}\right)^{k-1} O(n) + \dots + \left(\frac{3}{2}\right)^1 O(n) + O(n) =$$

$$= 3^k T\left(\frac{n}{2^k}\right) + \sum_{i=0}^{k-1} \left(\frac{3}{2}\right)^i O(n) =$$

$$= 3^k T\left(\frac{n}{2^k}\right) + O(n) \sum_{i=0}^{k-1} \left(\frac{3}{2}\right)^i$$

Repete até que $\frac{n}{2^k} = 1 \Rightarrow n = 2^k \Rightarrow \lg n = \lg 2^k \Rightarrow k = \lg n$

$$T(n) = 3^{\lg n} T(1) + O(n) \left[\frac{1 \cdot \left[\left(\frac{3}{2}\right)^{\lg n} - 1\right]}{\frac{3}{2} - 1} \right]$$

$$= 3^{\lg n} O(1) + O(n) \left[\frac{\left(\frac{3}{2}\right)^{\lg n} - 1}{\frac{1}{2}} \right]$$

$$= O(n^{\lg 3}) + O(n) [2(n^{\lg \frac{3}{2}} - 1)]$$

$$= O(n^{\lg 3}) + O(n) [2(n^{\lg 3 - \lg 2} - 1)]$$

$$= O(n^{\lg 3}) + O(n) \left[2 \left(\frac{n^{\lg 3}}{n^{\lg 2}} - 1 \right) \right]$$

$$= O(n^{\lg 3}) + O(n) \left[2 \left(\frac{n^{\lg 3}}{n} - 1 \right) \right]$$

$$= O(n^{\lg 3}) + 2 \cdot O \left(n - \frac{n^{\lg 3}}{n} - 1 \cdot n \right)$$

$$= O(n^{\lg 3}) + O(n^{\lg 3} - n)$$

$$T(n) = O(n^{\lg 3})$$