```
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T) root
   0110
multiply (1001, 0110)
      X= 1001 A=0110
      m = 4
     XL= LO XR=OL
      YL=OL YR= LO
      PL= multiply (10,01)
             X=10 Y=01
             XT=T XB=0
             YL=0 YR=1
             P1= multiply (1,0)
                   X=T Y=0
                    m = 1
                  return 1.000
             PL=O
             P2 = multiply (0,1)
                    X=0 Y=1
                  return 0.1:0
             P2=0
             P3 = multiply (1+0, 0+1)
                    Xzl Yzl
                 return 1.121
             P_3 = 1
return 0 \times 2^2 + (1 - 0 - 0) \times 2^{2} + 0
```

```
return 000 + 10 +0
      return 010
PL=010
P2= multiply (O1, LO)
      X=01 Y=10
       m = 2
      XL=O XR=1
      YL=L YR=0
       PL= multiply (0, 1)
            X=0 X=1
            return 0-1 =0
       P L = 0
       P2 = multiply (1,0)
            x= L y=0
            return 1.0 20
      P2=0
       P3 = multiply (0+1, 1+0)
             X=1 Y=1
             return L.1=1
      return 0 × 2 + (1-0-0) × 2 2 + 0
      return 000 + 10 + 6
      return 010
P2=010
P3=multiply (10+01,01+10)
      X= TT A= TT
```

XC= L XB= L

```
YL=1 YR=1
PL= multiply (L, L)
     X=T X=T
     return L. L=L
pr= T
P2= multiply (1,1)
     ×=1 7=1
     return 1.1=1
Pz=L
P3 = multiply (1+1, 121)
     X = LO Y= LO
     XL=L XR=0
     YL=L YR=0
     PL= multiply (LA)
         X=T Y=T
         return L. L=1
     PL= L
     P2= multiply (0,0)
          X20 Y20
            return 0-0=0
     P2=0
      P3 = multiply (1+0, 1+0)
           X=T X=T
            return 1-1=1
```

$$\begin{array}{c} P_{3}=1 \\ \text{ within } 100 + 00 + 0 \\ \text{ within } 100 + 00 + 0 \\ \text{ within } 100 + 100 + 0 \\ \text{ within } 100 + 100 + 1 \\ \text{ within } 100$$

$$m \cdot k = 0 - m = k$$

$$= 2^{m} 7(0) + \frac{1(2^{m} - 1) \cdot O(1)}{2 - 1}$$

$$= O(2^{n}) + O(2^{n} - 1) O(1)$$

$$= O(2^{n}) + O(2^{n} - 1)$$

$$= O(2^{n}) + O(2^{n} - 1)$$

c) 
$$T(n) = g(\frac{\pi}{3}) + O(n^2)$$
 $a = 9$ 
 $b = 3$ 
 $2 = 2$ 
 $d = 2$ 
 $0 = 2$ 

$$a:T(n)=O(n \log_2 5) \approx O(n^{2,32})$$
  
 $b:T(n)=O(2^n)$   
 $c:T(n)=O(n^2 \log_2 n)$ 

blersa forma, n log 2 = co (n² log n), logo, a purção c é o limite injerios pora o função a.

Assim, o algoritmo c dene rer oscolhido já que ele possur a menos complexidade dentre os três