

Local Planning

$$A_{s_0} = \underset{a}{\operatorname{argmax}} \left(\hat{T}_{s_0}^H U \right) (s_0, a)$$

What does this buy us ??

Localizing computation to s_0

Goal?

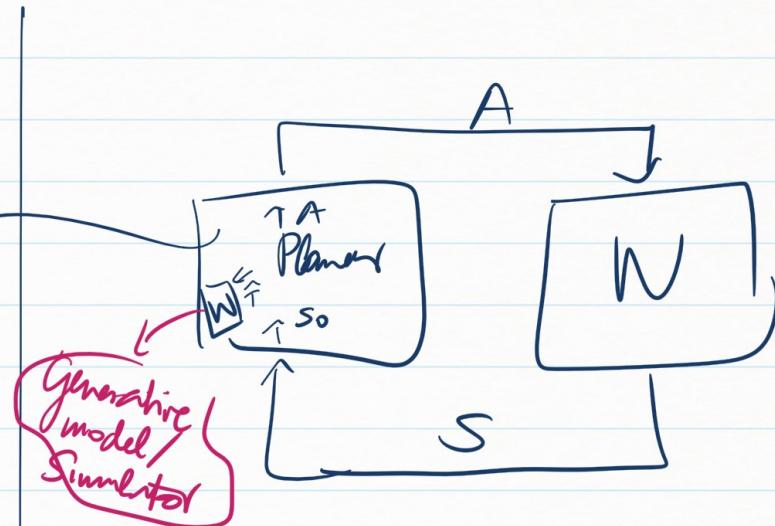
$$| A = a_1$$

① Little compute

② Input : $\delta > 0$

δ -optimality

π

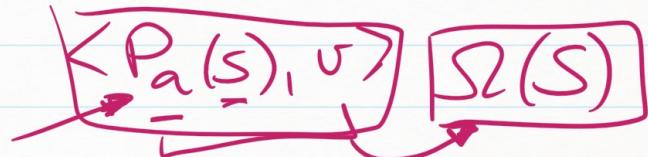


$$\text{Planner} \rightarrow \pi(a | s_0) = \Pr_{s_0}(A_{s_0} = a)$$

$$v^* \geq v^* - \delta I$$

per-state compute cost

$$O((N\delta)^H)$$



$C(s, a) \leftarrow$
 $(s, a) \mapsto \underbrace{S'_1(s, a), \dots, S'_m(s, a)}_{\sim P_a(s)}$ independence
 "call simulator"

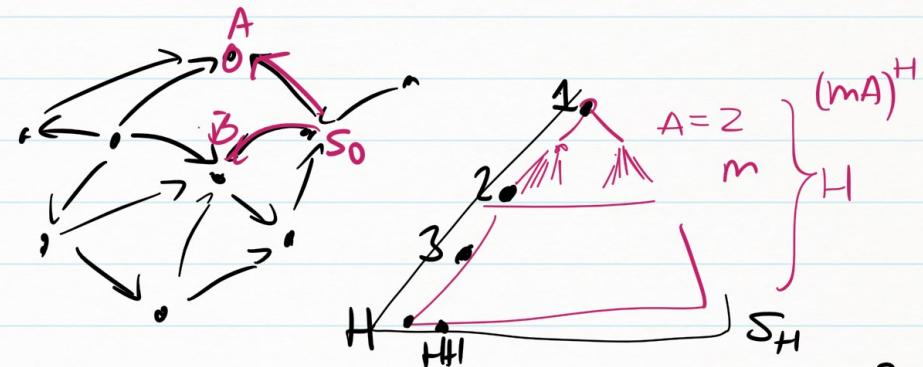
$$C(s) = (C(s, a))_{a \in A} / . = \bigcup_{a \in A} C(s, a)$$

$$\hat{T}: \mathbb{R}^{S \times A} \rightarrow \mathbb{R}^{S \times A}$$

$$\begin{aligned}
 (\hat{T}_q)(s, a) &= v_a(s) + \gamma \frac{1}{m} \sum_{s' \in C(s, a)} \max_{a'} q(s', a') \\
 | \\
 (\hat{T}_q^*)(s, a) &= r_a(s) + \gamma \langle P_a(s), M_q \rangle
 \end{aligned}$$

$$G = (S, \mathcal{E})$$

$$\mathcal{E} = \{(s, s') \mid s' \in C(s)\}$$



$$S_h = \{s \in S \mid \text{dist}(s_0, s) \leq h\}$$

$$h = 0$$

$$S_0 = \{s_0\}$$

$$S_1 = \{s_0, A, B\}$$

$$\vdots -\hat{T}_q^* + \hat{T}_q^*$$

$$\delta_h = \|\hat{T}_h^* - q^*\|_{S_{H-h}}$$

$$\begin{aligned}
 \delta_h &\leq \gamma \delta_h + \underbrace{\|\hat{T}_q^* - T_q^*\|_{S_H}}_{\leq \frac{\varepsilon'}{1-\gamma}} \\
 \delta_0 &\leq \frac{1}{1-\gamma}
 \end{aligned}$$

"whp"

$$\delta_h \leq \gamma \delta_{h-1} + \frac{\varepsilon'}{1-\gamma} \quad , \quad h \geq 1$$

$$\delta_0 \leq \frac{1}{1-\gamma}$$

$$\delta_1 \leq \frac{\gamma + \varepsilon'}{1-\gamma}$$

$$\delta_2 \leq \frac{\gamma(\gamma + \varepsilon') + \varepsilon'}{1-\gamma} = \frac{\gamma^2 + \varepsilon'(1+\gamma)}{1-\gamma}$$

$$\vdots$$

$$\delta_H \leq \frac{\gamma^H + \varepsilon' (1+\gamma + \dots + \gamma^{H-1})}{1-\gamma} \stackrel{\leq \frac{1}{1-\gamma}}{=}$$

$$\leq \left(\gamma^H + \frac{\varepsilon'}{1-\gamma} \right) \frac{1}{1-\gamma} \stackrel{\text{Mgk = r*}}{=} \delta$$

$$\boxed{\varepsilon' = ?}$$

$$\max_{S \in S_H} \max_a \gamma \left[\frac{1}{m} \sum_{j=1}^m v^*(S_j^*(s_a)) - \langle P_a(s), r^* \rangle \right]$$

Hoeffding's inequality

Lemma: $0 \leq X_i \leq 1$ i.i.d.
 $i = 1, \dots, n$

$\forall 0 \leq \zeta < 1 \quad \text{wp at least } 1-\zeta,$

(C) $\left| \frac{1}{n} \sum_{i=1}^n X_i - \mathbb{E} X_i \right| \leq \sqrt{\frac{\log(2/\zeta)}{2n}}$

Lemma: $0 \leq X_i \leq 1, i = 1, \dots, n$
 random element U

(U, X_1, \dots, X_n) jointly distributed.

X_1, \dots, X_n i.i.d. given U

$$[\Pr(X_i \in A_i, i=1..n | U) = \prod_i \Pr(X_i \in A_i | U)]$$

$$\forall 0 \leq \zeta < 1 \quad \text{wp at least } 1-\zeta$$

$$\left| \frac{1}{n} \sum_{i=1}^n X_i - \mathbb{E}[X_i | U] \right| \leq \sqrt{\frac{\log(2/\zeta)}{2n}}$$

(s, a)

$$\|\hat{T}_{q^+} - T_{q^*}\| = \max_{s \in S_H} \max_{a \in A} \left| \frac{1}{m} \sum_{j=1}^m v^*(S_j | s, a) - \langle P_a(s), v^* \rangle \right|$$

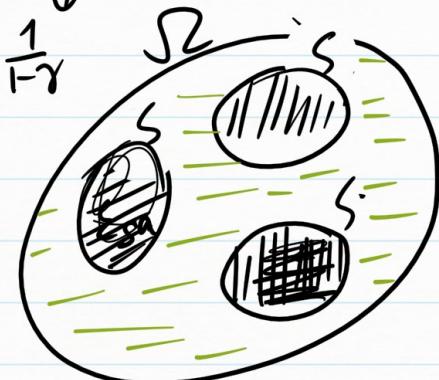
S_H

$\Delta(s, a)$

Union bounds!

$$P(\varepsilon_{sa}) \leq \frac{1}{1-\gamma} \frac{\log(2/\delta)}{m}$$

outside of ε_{sa}



Outside of $\bigcup_{sa} \varepsilon_{sa}$

$$P\left(\left(\bigcup_{sa} \varepsilon_{sa}\right)^c\right) = 1 - P\left(\bigcup_{sa} \varepsilon_{sa}\right) \geq 1 - SA \cdot \gamma$$

wp $1 - \delta_{failure}$

$$\delta_{failure} = SA \gamma$$

$$\gamma = \frac{\delta_{failure}}{SA}$$

$$|S_H| \leq (mA)^{H+1}$$

random
take union bound over these?!

$$\varepsilon' \leq \frac{1}{1-\gamma} \sqrt{\frac{\log(2/\delta_{failure})}{2m}}$$

$H = \frac{1}{(1-\gamma)^2}$

$\delta/3$

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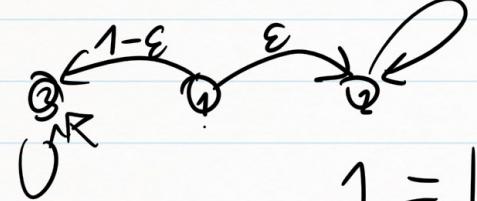
$\log(2/\delta_{failure}) / 2m$

$m \approx H \log(Am)$

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$\frac{2\delta_{failure}}{(1-\gamma)^2} = \frac{\delta}{3}$ target subset.

$m = S_2(\log(S))$



$$S' \sim P_1$$

(S) random

$$1 = \Pr(S' = 2 | S_0) = \underbrace{\Pr(S' = 2)}_{\approx \varepsilon}$$

$$S \doteq S'$$

$\{\hat{S}_1, \dots, \hat{S}_N\}$

$\forall 1 \leq i \leq N$

$$\underline{S_H} = \{S_1, \dots, S_n\}$$

$$n = (Am)^{t+i} - 1 / A - 1$$

S_1, \dots, S_n depth-first

encounters of states in S_H

Lemma:

$\forall i \in [n] \quad \text{wp } 1 - \delta$

$$|\Delta(S_{\tau_i}, a)| \leq \frac{1}{r_i} \sqrt{\frac{\log(\beta/\delta)}{2m}}$$

Proof : Prod i, a

$$\tau_i = \min \{1 \leq k \leq i \mid S_k = S_i\}$$

$$1 \leq \tau_i \leq i$$

$$\Delta(S_i, a) = \frac{1}{m} \sum_{S' \in C(S_i, a)} \underbrace{v^*(S') - \underbrace{P_a(S_i), r^*}_{S_{\tau_i}}}_{S_{\tau_i}}$$

$$\Delta(S_i, a) = \underline{\Delta}(S_{\tau_i}, a)$$

$$C(S_{\tau_i}, a) \perp \underbrace{S_{\tau_i}}_{C(S_1), \dots, C(S_{\tau_i-1})} \text{ via } S_{\tau_i}$$

+ union bond over $i \in [n]$

+ crating together

Thm: Cost = $O((mA)^H)$

$$H \approx H_{\gamma, \eta - \gamma} \delta^2$$

$$m \approx H \log(A)$$

local play goes π :

$$v^\pi \geq v^* - \delta_1$$

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