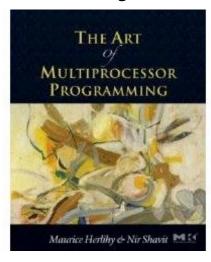
Linked Lists: Locking, Lock-Free, and Beyond ...



Hyungsoo Jung



Recall Three Design Patterns

Fine-grained synchronization

Optimistic Synchronization

Lazy Synchronization



First: Fine-Grained Synchronization

- Instead of using a single lock ...
- Split object into
 - Independently-synchronized components
- Methods conflict when they access
 - The same component …
 - At the same time



Second: Optimistic Synchronization

- Search without locking ...
- If you find it, lock and check
 - OK: we are done
 - Oops: start over
- Evaluation
 - Usually cheaper than locking, but
 - Mistakes are expensive



Third: Lazy Synchronization

- Postpone hard work
- Removing components is tricky
 - Logical removal
 - Mark component to be deleted
 - Physical removal
 - Do what needs to be done



Fourth: Lock-Free Synchronization

- Don't use locks at all
 - Use compareAndSet() & relatives ...
- Advantages
 - No Scheduler Assumptions/Support
- Disadvantages
 - Complex
 - Sometimes high overhead



Reminder: Lock-Free Data Structures

- No matter what ...
 - Guarantees minimal progress in any execution
 - i.e. Some thread will always complete a method call
 - Even if others halt at malicious times
 - Implies that implementation can't use locks



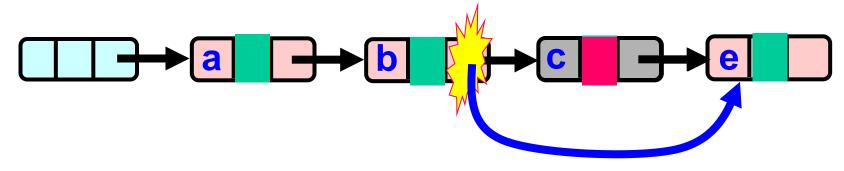
Lock-free Lists

- Next logical step
 - Wait-free contains()
 - lock-free add() and remove()
- Use only compareAndSet()
 - What could go wrong?



Lock-free Lists

Logical Removal



Use CAS to verify pointer is correct

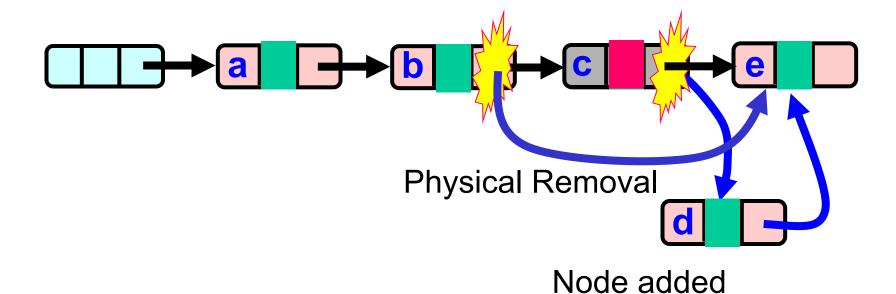
Physical Removal

Not enough!



Problem...

Logical Removal





The Solution: Combine Bit and Pointer

Logical Removal = Set Mark Bit **Physical** Removal Fail CAS: Node not Mark-Bit and Pointer added after logical CAS are CASed together Removal (AtomicMarkableReference)



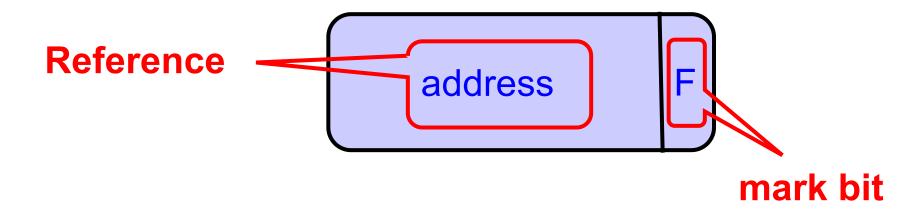
Solution

- Use AtomicMarkableReference
- Atomically
 - Swing reference and
 - Update flag
- Remove in two steps
 - Set mark bit in next field
 - Redirect predecessor's pointer



Marking a Node

- AtomicMarkableReference class
 - Java.util.concurrent.atomic package



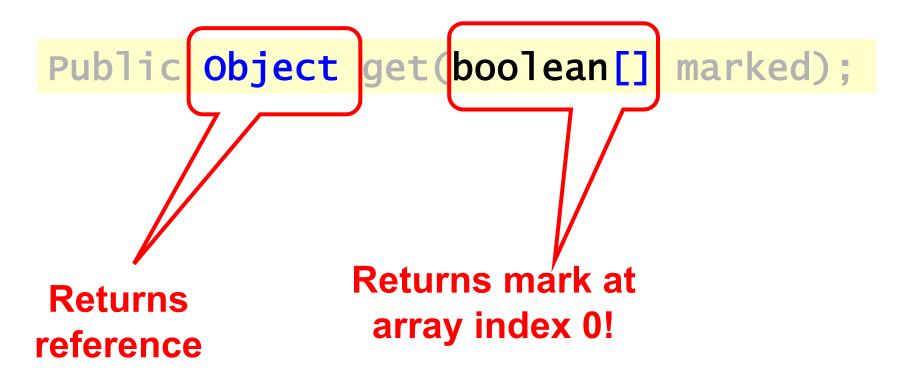


Extracting Reference & Mark

```
Public Object get(boolean[] marked);
```

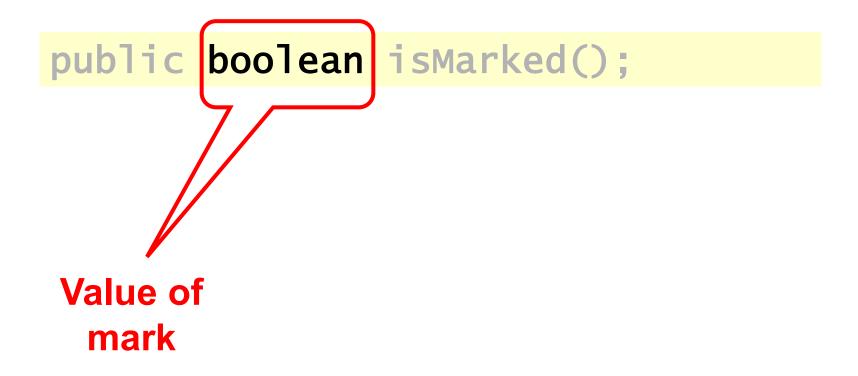


Extracting Reference & Mark





Extracting Mark Only





```
Public boolean compareAndSet(
   Object expectedRef,
   Object updateRef,
   boolean expectedMark,
   boolean updateMark);
```



```
If this is the current
                       reference ...
Public boolean compareAndSet(
  Object expectedRef,
  boolean expectedMark,
  boolean updateMar
                        And this is the
                        current mark ...
```



```
...then change to this
                   new reference ...
Public boolean compareAndSet(
  Object expectedRef,
  Object updateRef,
      ean expectedMark
  boolean updateMark);
                         and this new
                           mark
```



```
public boolean attemptMark(
   Object expectedRef,
   boolean updateMark);
```

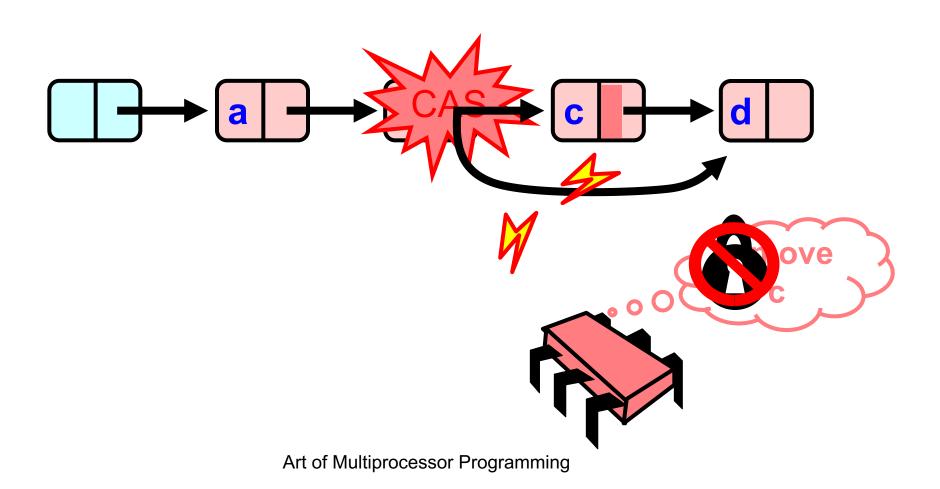


```
public boolean attemptMark(
 Object expectedRef,
  bodleam updateMark);
If this is the current
    reference ...
```

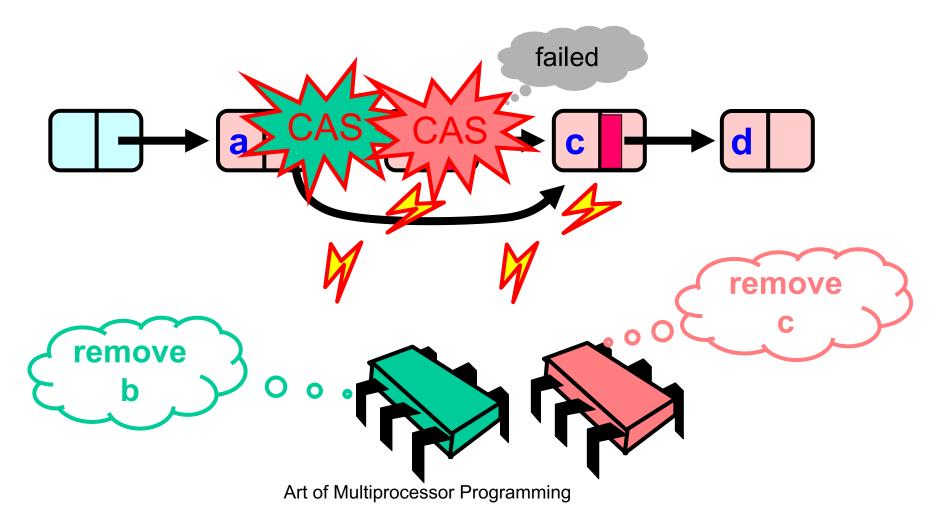


```
public boolean attemptMark(
  Object expectedRef,
 boolean updateMark);
.. then change to
 this new mark.
```

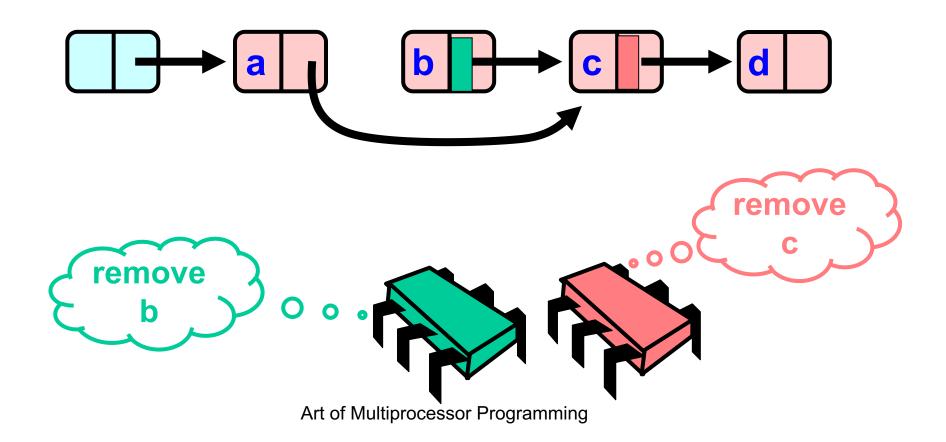




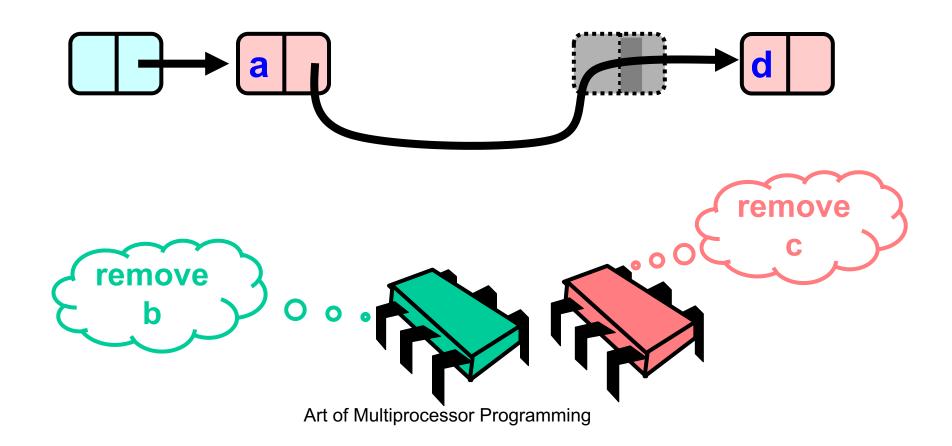












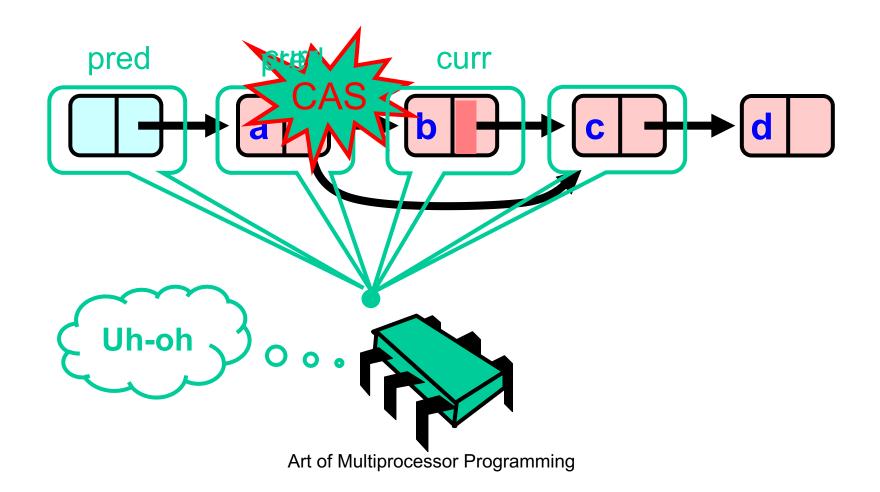


Traversing the List

- Q: what do you do when you find a "logically" deleted node in your path?
- A: finish the job.
 - CAS the predecessor's next field
 - Proceed (repeat as needed)



Lock-Free Traversal (only Add and Remove)





The Window Class

```
class Window {
  public Node pred;
  public Node curr;
  Window(Node pred, Node curr) {
    this.pred = pred; this.curr = curr;
  }
}
```



The Window Class

```
class Window {
  public Node pred;
  public Node curr;
  Window(Node pred, Node curr) {
    this.pred = pred; this.curr = curr;
  }
}
```

A container for pred and current values



Using the Find Method

```
Window window = find(head, key);
Node pred = window.pred;
curr = window.curr;
```



Using the Find Method

```
Window window = find(head, key);
Node pred = window.pred;
curr = window.curr;
Find returns window
```

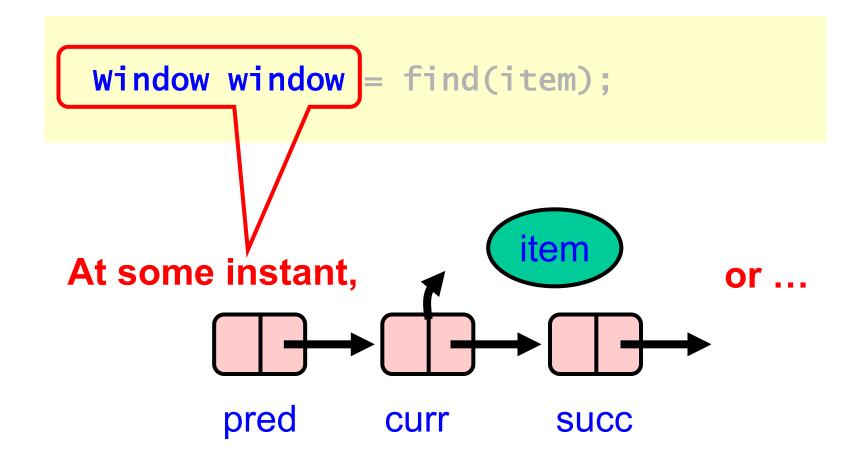


Using the Find Method

```
Window window = find(head, key);
Node pred = window.pred;
curr = window.curr;
Extract pred and curr
```

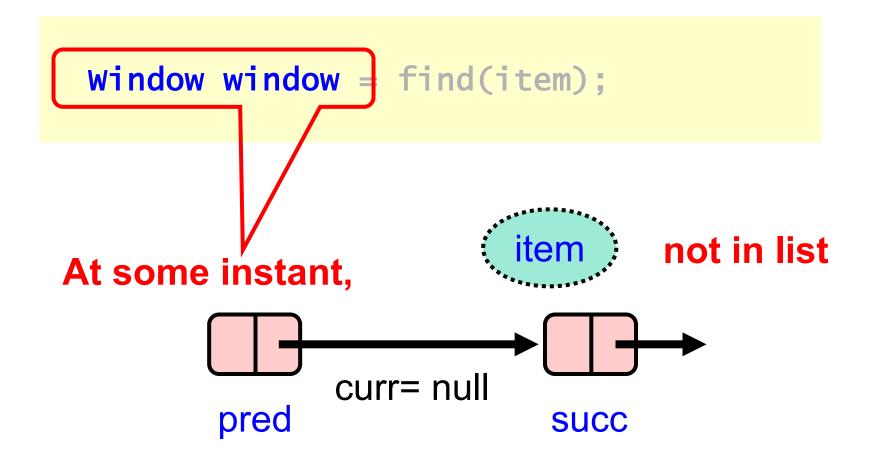


The Find Method





The Find Method





Remove

```
public boolean remove(T item) {
Boolean snip;
while (true) {
Window window = find(head, key);
 Node pred = window.pred, curr = window.curr;
  if (curr.key != key) {
     return false;
  } else {
  Node succ = curr.next.getReference();
  snip = curr.next.compareAndSet(succ, succ, false
true);
  if (!snip) continue;
   pred.next.compareAndSet(curr, succ, false, false);
     return true:
}}}
```



```
public boolean remove(T item) {
Boolean snip:
while (true) {
 window window = find(head, key);
 Node pred = window.pred, curr = window.curr;
  if (curr.key != key) {
     return false;
  } else {
  Node succ = curr.next.getReference();
  snip = curr.next.compareAndSet (succ, succ, false,
true);
  if (!snip) continue;
   pred.next.compareAndSet(curr, succ, false, false);
     return true;
                                Keep trying
}}}
```



```
public boolean remove(T item) {
Boolean snip;
while (true)
 Window window = find(head, key);
 Node pred = window.pred, curr = window.curr;
  if (curr.key != key)
     return false;
  } else {
  Node succ = curr.next.getReference();
  snip = curr.next.compareAndSet\(succ, succ, false,
true);
  if (!snip) continue;
   pred.next.compareAndSet(curr, suct, false, false);
     return true;
                          Find neighbors
}}}
```



```
public boolean remove(T item) {
Boolean snip;
while (true) {
 window window = find(head, key);
 Node pred = window.pred, curr = window.curr;
 if (curr.key != key) {`
     return false;
  } else {
  Node succ = curr.next.getReference();
  snip = curr.next.compareAndSet(succ, succ, false,
true);
  if (!snip) continue;
   pred.next.compareAndSet(curr, succ, false, false);
     return true;
                          She's not there ...
}}}
```



```
public boolean remove(T item) {
Boolean Try to mark node as deleted
while (true) {
 Window window = Find(head, key);
 Node pred = window.pred, curr = window.curr;
  if (curr.key != key) {
     return false;
  Node succ = curr.next.getReference();
  snip = curr.next.compareAndSet(succ, succ, false,
true):
  if (!snip) continue;
   pred.next.compareAndSet(curr, succ, false, false);
     return true;
}}}
```



```
public boolean remove(T item) {
If it doesn't work,
  just retry, if it
                   ind(head, key);
    does, job ndow pred, curr = window.curr;
essentially done y) {
     return taise;
  } else {
  Node succ = curr.next.getReference();
  snip = curr.next.compareAndSet(succ, succ, false,
true).
  if (!snip) continue;
   pred.next.compareAndSet(curr, succ, false, false);
     return true;
}}}
```



```
public boolean remove(T item) {
Boolean snip;
while (true) {
Window window = find(head,
 Node pred = window.pred, curr = window.curr
 if (curr.key != key) {
   Try to advance reference
  (if we don't succeed, someone else did or will).
  snip = curr.next.compareAndSet(succ, succ, false,
true);
  pred.next.compareAndSet(curr, succ, false, false);
     return true:
```



```
public boolean add(T item) {
 boolean splice;
 while (true) {
   Window window = find(head, key);
   Node pred = window.pred, curr = window.curr;
   if (curr.key == key) {
      return false;
   } else {
   Node node = new Node(item);
   node.next = new AtomicMarkableRef(curr, false);
   if (pred.next.compareAndSet(curr, node, false,
false)) {return true;}
}}}
```



```
public boolean add(T item) {
 boolean splice;
 while (true) {
  Window window = find(head, key);
   Node pred - window.pred, curr = window.curr;
  if (curr.key == key) {
      return false;
   Node node = new Node(item);
   node.next = new AtomicMarkableRef(curr, false);
  if (pred.next.compareAndSet(curr, node, false,
false)) freturn true:
                Item already there.
}}}
```



```
public boolean add(T item) {
 boolean splice;
 while (true) {
  Window window = find(head, k
   Node pred = window.pred, curr = window.curr;
   if (curr.key == key) {
      return false;
   } else {
  Node node = new Node(item);
   node.next = new AtomicMarkableRef(curr, false);
   if (pred.next.compareAndSet(curr, node, false,
false)) {return true;}
111
                   create new node
```



```
public boolean add(T item) {
                               Install new node,
 boolean splice;
 while (true) {
  Window window = find(head, key), else retry loop
  Node pred = window.pred, curr = window.curr;
  Node node = new Node(item);
        next - new AtomicMarkableRef(curr, False)
  if (pred.next.compareAndSet(curr, node, false,
false)) {return true;}
```



Wait-free Contains

```
public boolean contains(T item) {
   boolean marked;
   int key = item.hashCode();
   Node curr = this.head;
   while (curr.key < key)
        curr = curr.next;
   Node succ = curr.next.get(marked);
   return (curr.key == key && !marked[0])
}</pre>
```



Wait-free Contains

```
public boolean contains (T item)
    boolean marked; Only diff is that we
    int key = item.hashcoget and check
    Node curr = this.head: marked
    while (curr.key < key)</pre>
      curr = curr.next;
    Node succ = curr.next.get(marked);
    return (curr.key == key && !marked[0])
```



```
public Window find(Node head, int key) {
Node pred = null, curr = null, succ = null;
boolean[] marked = {false}; boolean snip;
 retry: while (true) {
  pred = head;
   curr = pred.next.getReference();
   while (true) {
    succ = curr.next.get(marked);
    while (marked[0]) {
   if (curr.key >= key)
         return new Window(pred, curr);
       pred = curr;
       curr = succ;
}}
```



```
public Window find(Node head, int key) {
Node pred = null, curr = null, succ = null;
 boolean[] marked = {false}; boolean snip;
retry: while (true) {
   pred = head;
  curr = pred.next.getReterence If list changes
   while (true) {
    succ = curr.next.get(marked) while traversed,
    while (marked[0]) {
                             start over
   if (curr.key >= key)
         return new Window(pred, curr);
       pred = curr;
       curr = succ;
```



```
public Window find (Node head int key) {
Node pred = null Start looking from head
boolean[] marked = {talse}; boolean snip;
 retry: while (true) {
  pred = head;
  curr = pred.next.getReference();
   while (true) {
    succ = curr.next.get(marked);
    while (marked[0]) {
   if (curr.key >= key)
         return new Window(pred, curr);
       pred = curr;
       curr = succ;
```



```
public Window find(Node head, int key) {
Node pred = null, curr = null, succ = null;
boolean[] marked = {false}; boolean snip;
 retry: while (true) { Move down the list
   pred = head;
   curr = pred.next.getP
  while (true) {
    succ = curr.next.get(marked);
   while (marked[0]) {
   if (curr.key >= key)
         return new Window(pred, curr);
       pred = curr;
       curr = succ;
```



```
public Window find(Node head, int key) {
Node pred = null, curr = null, succ = null;
boolean[] marked = {false}; boolean snip;
 retry: while (true) {
  pred = head;
   curr = pred.next.getReference();
   while (true) {
   succ = curr.next.get(marked);
          (marked[0
   if (curr.key >= key)
         return new Window (red, curr);
       pred = curr; Get ref to successor and
       curr = succ;
                        current deleted bit
```



```
public Window find(Node head, int key) {
  Node pred = null, curr = null, succ = null;
  boolean[] marked = {false}; boolean snip;
  retry: while (true) {
    pred = head;
    curr = pred.next.getReference();
    while (true) {
     succ = curr.next.get(marked):
     while (marked[0]) {
          return new Window(pr
        nred = curr
Try to remove deleted nodes in
   path...code details soon
```



```
public Window find(Node head, int key) {
Node pred = null, curr = null, succ = null;
boolean[] marked = {false}; boolean snip;
 retry: while (true) {
   pred = head;
   curr - nrad navt natDafaranca():
  If curr key that is greater or
   equal, return pred and curr
   if (curr.key >= key)
         return new Window(pred, curr);
       pred = curr;
       curr = succ:
```



```
public Window find(Node head, int key) {
Node pred = null, curr = null, succ = null;
boolean[] marked = {false}; boolean snip;
 retry: while (true) {
  pred = head;
  curr = pred.next.getReference();
  while (true) {
 Otherwise advance window and
            loop again
   if (curr.key >= key)
         return new Window(pred, curr);
       pred = curr;
```





Try to snip out node

```
retry: while (true) {
   while (marked[0]
     snip = pred.next.compareAndSet(curr,
                          succ, false, false);
        (!snip) continue retry;
     curr = succ;
     succ = curr.next.get(marked);
```



if predecessor's next field changed must retry whole

```
traversal
retry: while (true) {
  while (marked[0]) {
     snip = pred.next.compareAndSet(curr,
     false false)
SUCC
     if (!snip) continue retry;
     succ = curr.next.get(marked);
```



Otherwise move on to check if next node deleted



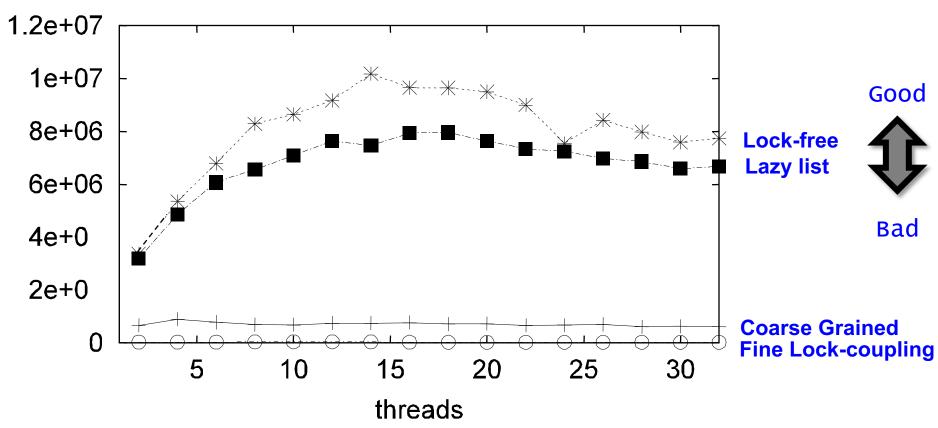
Performance

On 16 node shared memory machine Benchmark throughput of Java List-based Set algs. Vary % of Contains() method Calls.



High Contains Ratio

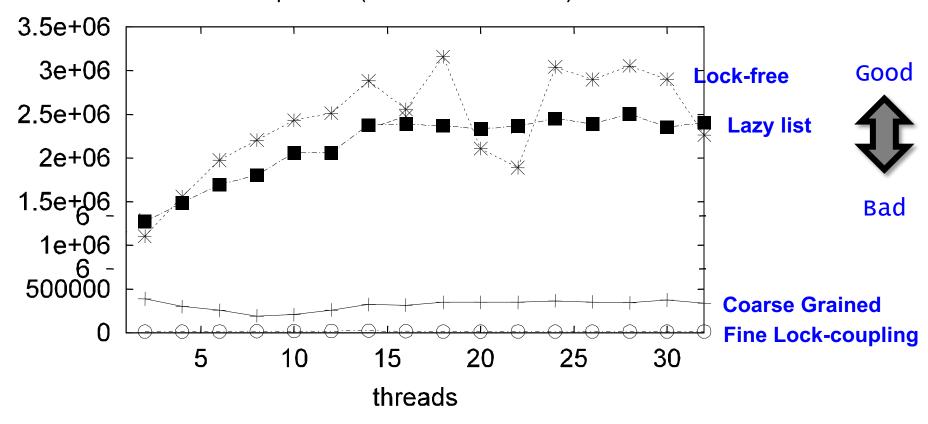
Ops/sec (90% reads/0 load)





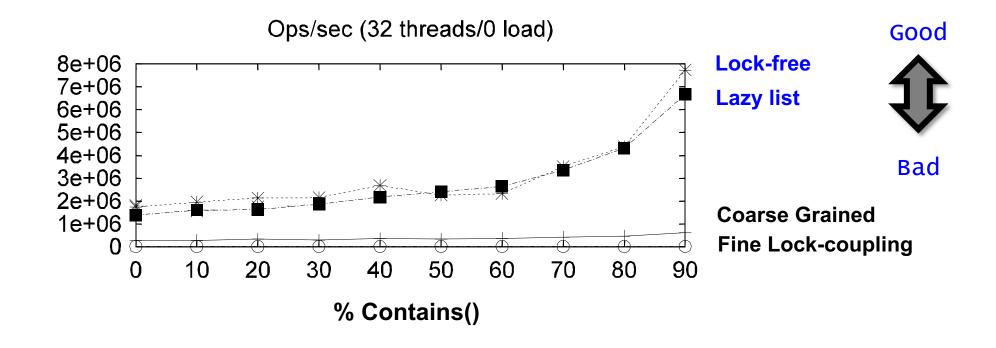
Low Contains Ratio

Ops/sec (50% reads/0 load)





As Contains Ratio Increases





Summary

- Coarse-grained locking
- Fine-grained locking
- Optimistic synchronization
- Lock-free synchronization



"To Lock or Not to Lock"

- Locking vs. Non-blocking: Extremist views on both sides
- The answer: nobler to compromise, combine locking and non-blocking
 - Example: Lazy list combines blocking add() and remove() and a wait-free contains()
 - Remember: Blocking/non-blocking is a property of a method

