# QCQI: CHAPTER 5 SUMMARY

#### A PREPRINT

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#### **ABSTRACT**

The Fourier Transform and applications.

### 1 The Fourier Transform

## 2 The Phase estimation algorithm

Consider a unitary operator U. Say  $|u\rangle$  is an eigenvector with eigenvalue  $e^{2\pi i \varphi}$ . We would like to find (approximately)  $\varphi$ . We assume that we can query a blackbox to apply  $U^{2^j}$  for any  $j \in \mathbb{Z}_{\geq 0}$ . The key idea is to encode  $\varphi$  into the phase space and then compute the inverse Fourier Transform.

The key chain of events reads as follows (subscripts on the operator describe the qubit(s) it was applied to):

$$|0\rangle|u\rangle\xrightarrow{H_{[n]}^{\otimes n}}\left(\frac{1}{2^{n/2}}\sum_{j=0}^{2^n-1}|j\rangle\right)|u\rangle\xrightarrow{U_{n-i}^{2^i}} \xrightarrow{0\leq i\leq n-1} \xrightarrow{1} \frac{1}{2^{n/2}}\sum_{j=0}^{2^n-1}|j\rangle U^j|u\rangle = \left(\frac{1}{2^{n/2}}\sum_{j=0}^{2^n-1}e^{2\pi i j\varphi}|j\rangle\right)|u\rangle\xrightarrow{\mathrm{IFT}_{[n]}} \approx |\tilde{\varphi}\rangle|u\rangle$$

whereupon measuring the first register gives us the estimate  $0.\tilde{\varphi}$  for  $\varphi$ . The  $\approx$  is for two reasons. One, that  $\varphi$  is possibly more than n bits long, in which case  $\varphi$  is an n-bit approximation to  $\varphi$ . The other reason is that the pther statevectors  $|j\rangle$  could also be received in the measurement with a small probability - their amplitudes are non-zero if  $\varphi$  is longer than n bits (in which case we say that the algorithm failed). Nevertheless, the algorithm is one of most vital importance and use in what follows.

Finally, another note: Preparing the eigenvector  $|u\rangle$  may not be easy. But it may be easy to prepare a superposition of some eigenvectors (for example, in the order-finding algorithm below). In this case, we receive  $\tilde{\varphi}$  with high probability for **one of the eigenvectors in the superposition**.

## 3 Order-finding and Shor's algorithm

Given integers x, N > 0 with gcd(x, N) = 1, we would like to find the order r of x modulo N. Classically, this is hard. Here we show that with high probability we can find it