

Optimality Conditions for Linearly Constrained Optimization

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General Optimization Problems

Let the problem have the general mathematical programming (MP) form:

$$\begin{array}{ll} \text{(P)} & \text{minimize} \quad f(\mathbf{x}) \\ & \text{subject to} \quad \mathbf{x} \in \mathcal{F}. \end{array}$$

In all forms of MP, a **feasible solution** of a given problem is a vector that satisfies the constraints of the problem, that is, in \mathcal{F} .

Question: how does one recognize or certify an optimal solution to a **generally constrained and objectived** optimization problem?

Answer: **Optimality Condition Theory** again.

Descent Direction

Let f be a differentiable function on R^n . Given a point $\bar{\mathbf{x}} \in R^n$, if there is a vector \mathbf{d} such that

$$\nabla f(\bar{\mathbf{x}})\mathbf{d} < 0,$$

then there exists a scalar $\bar{\tau} > 0$ such that

$$f(\bar{\mathbf{x}} + \tau\mathbf{d}) < f(\bar{\mathbf{x}}) \text{ for all } \tau \in (0, \bar{\tau}).$$

The vector \mathbf{d} is called a **descent direction** at $\bar{\mathbf{x}}$. If $\nabla f(\bar{\mathbf{x}}) \neq 0$, then $\nabla f(\bar{\mathbf{x}})$ is the direction of **steepest ascent** and $-\nabla f(\bar{\mathbf{x}})$ is the direction of **steepest descent** at $\bar{\mathbf{x}}$.

Let's denote the set of descent directions at $\bar{\mathbf{x}}$ by $\mathcal{D}_{\bar{\mathbf{x}}}^d$, i.e.,

$$\mathcal{D}_{\bar{\mathbf{x}}}^d := \{\mathbf{d} \in R^n : \nabla f(\bar{\mathbf{x}})\mathbf{d} < 0\}.$$

Feasible Direction

At feasible point $\bar{\mathbf{x}}$, the set of feasible directions is given by

$$\mathcal{D}_{\bar{\mathbf{x}}}^f := \{\mathbf{d} \in R^n : \mathbf{d} \neq \mathbf{0}, \bar{\mathbf{x}} + \lambda \mathbf{d} \in \mathcal{F} \text{ for all small } \lambda > 0\}.$$

Examples:

$$\mathcal{F} = R^n \Rightarrow \mathcal{D}^f = R^n.$$

$$\mathcal{F} = \{\mathbf{x} : A\mathbf{x} = \mathbf{b}\} \Rightarrow \mathcal{D}^f = \{\mathbf{d} : A\mathbf{d} = \mathbf{0}\}.$$

$$\mathcal{F} = \{\mathbf{x} : A\mathbf{x} \geq \mathbf{b}\} \Rightarrow \mathcal{D}^f = \{\mathbf{d} : A_i \mathbf{d} \geq 0, \forall i \in \mathcal{A}(\bar{\mathbf{x}})\},$$

where the **active** or **binding** constraint set $\mathcal{A}(\bar{\mathbf{x}}) := \{i : A_i \bar{\mathbf{x}} = b_i\}$.

Optimality Conditions

Optimality Conditions: given a feasible solution or point $\bar{\mathbf{x}}$, what are the **necessary conditions** for $\bar{\mathbf{x}}$ to be a local optimizer?

A general answer would be: there exists no direction at $\bar{\mathbf{x}}$ that is both **descent and feasible**. Or the **intersection** of $\mathcal{D}_{\bar{\mathbf{x}}}^d$ and $\mathcal{D}_{\bar{\mathbf{x}}}^f$ must be **empty**.

Unconstrained Problems

Consider the following **unconstrained** problem where f is differentiable on R^n :

$$\begin{array}{ll} \text{(UP)} & \text{minimize} \quad f(\mathbf{x}) \\ & \text{subject to} \quad \mathbf{x} \in R^n. \end{array}$$

$\mathcal{D}_{\bar{\mathbf{x}}}^f = R^n$, so that $\mathcal{D}_{\bar{\mathbf{x}}}^d = \{\mathbf{d} \in R^n : \nabla f(\bar{\mathbf{x}})\mathbf{d} < 0\} = \emptyset$.

Theorem 1 Let \bar{x} be a (local) minimizer of (UP). If the functions f is continuously differentiable at \bar{x} , then

$$\nabla f(\bar{\mathbf{x}}) = \mathbf{0}.$$

Linear Equality-Constrained Problems

Consider the following **linear equality-constrained** problem where f is differentiable on R^n :

$$\begin{array}{ll} \text{(LEP)} & \text{minimize} \quad f(\mathbf{x}) \\ & \text{subject to} \quad A\mathbf{x} = \mathbf{b}. \end{array}$$

Theorem 2 (the Lagrange Theorem) Let $\bar{\mathbf{x}}$ be a (local) minimizer of LEP. If the functions f is continuously differentiable at $\bar{\mathbf{x}}$, then

$$\nabla f(\bar{\mathbf{x}}) = \mathbf{y}^T A$$

for some $\mathbf{y} = (\bar{y}_1; \dots; \bar{y}_m) \in R^m$, which are called **Lagrange or dual multipliers**.

Geometric interpretation: the objective gradient vector is **perpendicular** to or the objective level set **tangents** the constraint hyperplanes.

Proof

Consider feasible direction space

$$\mathcal{F} = \{\mathbf{x} : A\mathbf{x} = \mathbf{b}\} \Rightarrow \mathcal{D}_{\bar{\mathbf{x}}}^f = \{\mathbf{d} : A\mathbf{d} = 0\}.$$

If $\bar{\mathbf{x}}$ is a local optimizer, then the intersection of the descent and feasible direction sets at $\bar{\mathbf{x}}$ must be empty or

$$A\mathbf{d} = \mathbf{0}, \nabla f(\bar{\mathbf{x}})\mathbf{d} \neq 0$$

has no feasible solution for \mathbf{d} . By the Alternative System Theorem it must be true that its alternative system has a solution, that is, there is $\mathbf{y} \in R^n$ such that

$$\nabla f(\bar{\mathbf{x}}) = \mathbf{y}^T A = \sum_{i=1}^m \bar{y}_i A_i.$$

The Barrier Function Problem

Consider the problem

$$\begin{array}{ll}\text{minimize} & -\sum_{j=1}^n \log x_j \\ \text{subject to} & A\mathbf{x} = \mathbf{b}, \\ & \mathbf{x} \geq \mathbf{0}\end{array}$$

The non-negativity constraint can be removed if the feasible region has an "interior". If a minimizer exists, then

$$-\mathbf{e}^T X^{-1} = \mathbf{y}^T A = \sum_{i=1}^m \bar{y}_i A_i.$$

Linear Inequality-Constrained Problems

Let us now consider the inequality-constrained problem

$$\begin{array}{ll} \text{(LIP)} & \text{minimize} \quad f(\mathbf{x}) \\ & \text{subject to} \quad A\mathbf{x} \geq \mathbf{b}. \end{array}$$

Theorem 3 (*the KKT Conditions*) Let $\bar{\mathbf{x}}$ be a (local) minimizer of LIP). If the functions f is continuously differentiable at $\bar{\mathbf{x}}$, then

$$\nabla f(\bar{\mathbf{x}}) = \mathbf{y}^T A, \quad \mathbf{y} \geq \mathbf{0}$$

for some $\mathbf{y} = (\bar{y}_1; \dots; \bar{y}_m) \in R^m$, which are called *Lagrange or dual multipliers*, and $\bar{y}_i = 0$, if $i \notin \mathcal{A}(\bar{\mathbf{x}})$.

Geometric interpretation: the objective gradient vector is in the **cone** generated by the **normal directions** of the active-constraint hyperplanes.

Proof

$$\mathcal{F} = \{\mathbf{x} : A\mathbf{x} \geq \mathbf{b}\} \Rightarrow \mathcal{D}_{\bar{\mathbf{x}}}^f = \{\mathbf{d} : A_i \mathbf{d} \geq 0, \forall i \in \mathcal{A}(\bar{\mathbf{x}})\},$$

or

$$\mathcal{D}_{\bar{\mathbf{x}}}^f = \{\mathbf{d} : \bar{A}\mathbf{d} \geq \mathbf{0}\},$$

where \bar{A} corresponds to those active constraints. If $\bar{\mathbf{x}}$ is a local optimizer, then the **intersection** of the **descent and feasible** direction sets at $\bar{\mathbf{x}}$ must be empty or

$$\bar{A}\mathbf{d} \geq \mathbf{0}, \nabla f(\bar{\mathbf{x}})\mathbf{d} < 0$$

has no feasible solution. By **the Alternative System Theorem** it must be true that its alternative system has a solution, that is, there is $\bar{\mathbf{y}} \geq \mathbf{0}$ such that

$$\nabla f(\bar{\mathbf{x}}) = \bar{\mathbf{y}}^T \bar{A} = \sum_{i \in \mathcal{A}(\bar{\mathbf{x}})} \bar{y}_i A_i.$$

Let $\bar{y}_i = 0$ for all remaining inactive constraints. Then we prove the theorem.

Optimization with Mixed Constraints

We now consider optimality conditions for problems having both **inequality** and **equality** constraints. That is,

$$\begin{array}{ll} \text{minimize} & f(\mathbf{x}) \\ \text{subject to} & A\mathbf{x} = \mathbf{b} \\ & \mathbf{x} \geq \mathbf{0} \end{array} \quad (\text{P})$$

For any feasible point $\bar{\mathbf{x}}$ of (P) we have the sets

$$\mathcal{A}(\bar{\mathbf{x}}) = \{j : \bar{x}_j = 0\}$$

$$\mathcal{D}_{\bar{\mathbf{x}}}^d = \{\mathbf{d} : \nabla f(\bar{\mathbf{x}})\mathbf{d} < 0\}.$$

The KKT Conditions again

Theorem 4 Let $\bar{\mathbf{x}}$ be a local minimizer for (P). Then there exist multipliers \mathbf{y} and \mathbf{s} such that

$$\left\{ \begin{array}{ll} \nabla f(\bar{\mathbf{x}}) &= \mathbf{y}^T A + \mathbf{s}^T \\ \mathbf{s} &\geq \mathbf{0} \\ s_j &= 0 \text{ if } j \notin \mathcal{A}(\bar{\mathbf{x}}). \end{array} \right.$$

Optimality and Complementarity Conditions

$$\left\{ \begin{array}{rcl} x_j (\nabla f(\mathbf{x}) - \mathbf{y}^T A)_j & = & 0, \quad \forall j = 1, \dots, n \\ A\mathbf{x} & = & \mathbf{b} \\ \nabla f(\mathbf{x}) - \mathbf{y}^T A & \geq & \mathbf{0} \\ \mathbf{x} & \geq & \mathbf{0}. \end{array} \right.$$

$$\left\{ \begin{array}{rcl} x_j s_j & = & 0, \quad \forall j = 1, \dots, n \\ A\mathbf{x} & = & \mathbf{b} \\ \nabla f(\mathbf{x}) - \mathbf{y}^T A - \mathbf{s}^T & = & \mathbf{0} \\ \mathbf{x}, \mathbf{s} & \geq & \mathbf{0} \end{array} \right.$$

Sufficient Optimality Conditions

Theorem 5 If f is a differentiable *convex* function in the feasible region and the feasible region is a convex set, then the (first-order) KKT optimality conditions are *sufficient* for the *global optimality* of a feasible solution.

Corollary 1 If f is differentiable *convex* function in the feasible region, then the (first-order) KKT optimality conditions are *sufficient* for the *global optimality* of a feasible solution for linearly constrained optimization.

How to check convexity, say $f(x) = x^3$?

- Hessian matrix is PSD in the feasible region.
- Epigraph is a convex set.

LCCP Examples: Linear Optimization

$$\begin{aligned} (LP) \quad & \text{minimize} \quad \mathbf{c}^T \mathbf{x} \\ & \text{subject to} \quad A\mathbf{x} = \mathbf{b}, \mathbf{x} \geq \mathbf{0}. \end{aligned}$$

For any feasible \mathbf{x} of LP, it's optimal if for some \mathbf{y}, \mathbf{s}

$$\begin{aligned} x_j s_j &= 0, \forall j = 1, \dots, n \\ A\mathbf{x} &= \mathbf{b} \\ \nabla(\mathbf{c}^T \mathbf{x}) = \mathbf{c}^T &= \mathbf{y}^T A + \mathbf{s}^T \\ \mathbf{x}, \mathbf{s} &\geq \mathbf{0}. \end{aligned}$$

Here, \mathbf{y} are Lagrange multipliers of equality constraints, and \mathbf{s} (reduced cost or dual slack vector in LP) are Lagrange multipliers for $\mathbf{x} \geq \mathbf{0}$.

LCCP Examples: Barrier Optimization

$$f(\mathbf{x}) = \mathbf{c}^T \mathbf{x} - \mu \sum_{j=1}^n \log(x_j),$$

for some fixed $\mu > 0$. Assume that interior of the feasible region is not empty:

$$\left\{ \begin{array}{rcl} A\mathbf{x} & = & \mathbf{b} \\ c_j - \frac{\mu}{x_j} - (\mathbf{y}^T A)_j & = & 0, \quad \forall j = 1, \dots, n \\ \mathbf{x} & > & \mathbf{0}. \end{array} \right.$$

$$\left\{ \begin{array}{rcl} x_j s_j & = & \mu, \quad \forall j = 1, \dots, n, \\ A\mathbf{x} & = & \mathbf{b}, \\ A^T \mathbf{y} + \mathbf{s} & = & \mathbf{c}, \\ (\mathbf{x}, \mathbf{s}) & > & \mathbf{0}. \end{array} \right.$$

Proof of Uniqueness

Solution pair of (\mathbf{x}, \mathbf{s}) of the barrier optimization problem is unique. Suppose there two different pair $(\mathbf{x}^1, \mathbf{s}^1)$ and $(\mathbf{x}^2, \mathbf{s}^2)$. Then we have

$$(\mathbf{s}^1 - \mathbf{s}^2)^T (\mathbf{x}^1 - \mathbf{x}^2) = 0.^a$$

Thus, there is j such that

$$(s_j^1 - s_j^2)(x_j^1 - x_j^2) > 0.$$

If $x_j^1 > x_j^2$, then $s_j^1 < s_j^2$ since $x_j^1 s_j^1 = x_j^2 s_j^2 = \mu > 0$, which leads to $(s_j^1 - s_j^2)(x_j^1 - x_j^2) < 0$ which is a contradiction. Similarly, one cannot have $x_j^1 < x_j^2$.

^aTo see this, the second condition gives $A\mathbf{x}_1 = \mathbf{b}_1$ and $A\mathbf{x}_2 = \mathbf{b}_3$ and the third condition gives $A^T \mathbf{y}_1 + \mathbf{s}_1 = \mathbf{c}$ and $A^T \mathbf{y}_2 + \mathbf{s}_2 = \mathbf{c}$. Which imply $A(\mathbf{x}_1 - \mathbf{x}_2) = 0$ and $A^T(\mathbf{y}_1 - \mathbf{y}_2) = \mathbf{s}_2 - \mathbf{s}_1$ and thus give $0 = (\mathbf{x}_1 - \mathbf{x}_2)^T A^T(\mathbf{y}_1 - \mathbf{y}_2) = (\mathbf{x}_1 - \mathbf{x}_2)^T (\mathbf{s}_2 - \mathbf{s}_1)$.

KKT Applications: Fisher's Equilibrium Price

Player $i \in B$'s optimization problem for given prices $p_j, j \in G$.

$$\begin{aligned} &\text{maximize} && \mathbf{u}_i^T \mathbf{x}_i := \sum_{j \in G} u_{ij} x_{ij} \\ &\text{subject to} && \mathbf{p}^T \mathbf{x}_i := \sum_{j \in G} p_j x_{ij} \leq w_i, \\ &&& x_{ij} \geq 0, \quad \forall j, \end{aligned}$$

Assume that the amount of each good is s_j . The equilibrium price vector is the one that for all $j \in G$

$$\sum_{i \in B} x(\mathbf{p})_{ij} = s_j$$

Example of Fisher's equilibrium price

There two goods, x and y , each with 1 unit on the market. Buyer 1, 2's optimization problems for given prices p_x, p_y .

$$\begin{aligned} &\text{maximize} && 2x_1 + y_1 \\ &\text{subject to} && p_x \cdot x_1 + p_y \cdot y_1 \leq 5, \end{aligned}$$

$$x_1, y_1 \geq 0;$$

$$\begin{aligned} &\text{maximize} && 3x_2 + y_2 \\ &\text{subject to} && p_x \cdot x_2 + p_y \cdot y_2 \leq 8, \end{aligned}$$

$$x_2, y_2 \geq 0.$$

Given $(p_x, p_y) = \left(\frac{26}{3}, \frac{13}{3}\right)$, the optimal quantities are found to be

$$(x_1, y_1) = \left(\frac{1}{13}, 1\right), \quad (x_2, y_2) = \left(\frac{12}{13}, 0\right).$$

Equilibrium Price Conditions

Player $i (\in B)$'s dual problem for given prices $p_j, j \in G$.

$$\begin{array}{ll} \text{minimize} & w_i y_i \\ \text{subject to} & \mathbf{p} y_i \geq \mathbf{u}_i, \quad y_i \geq 0 \end{array}$$

The necessary and sufficient conditions for an equilibrium point \mathbf{x}_i, \mathbf{p} are:

$$\left\{ \begin{array}{ll} \mathbf{p}^T \mathbf{x}_i \leq w_i, \quad \mathbf{x}_i \geq \mathbf{0}, & \forall i, \\ p_j y_i \geq u_{ij}, \quad y_i \geq 0, & \forall i, j, \\ \mathbf{u}_i^T \mathbf{x}_i = w_i y_i, & \forall i, \\ \sum_i x_{ij} = s_j, & \forall j. \end{array} \right.$$

Equilibrium Price Conditions (cont'd)

These conditions can be represented by

$$\left\{ \begin{array}{ll} \sum_j s_j p_j \leq \sum_i w_i, \mathbf{x}_i \geq \mathbf{0}, & \forall i, \\ \frac{\mathbf{u}_i^T \mathbf{x}_i}{w_i} \cdot p_j \geq u_{ij}, & \forall i, j, \\ \sum_i x_{ij} = s_j, & \forall j. \end{array} \right.$$

since from the second inequality (after multiplying x_{ij} to both sides and take sum over j) we have

$$\mathbf{p}^T \mathbf{x}_i \geq w_i, \forall i.$$

Then, from the rest conditions

$$\sum_i w_i \geq \sum_j s_j p_j = \sum_i \mathbf{p}^T \mathbf{x}_i \geq \sum_i w_i.$$

Thus, these conditions imply $\mathbf{p}^T \mathbf{x}_i = w_i, \forall i$.

Equilibrium Price Property

If u_{ij} has at least one positive coefficient for every j , then we must have $p_j > 0$ for every j at every equilibrium. Moreover, The second inequality can be rewritten as

$$\log(\mathbf{u}_i^T \mathbf{x}_i) + \log(p_j) \geq \log(w_i) + \log(u_{ij}), \quad \forall i, j, u_{ij} > 0.$$

The function on the left is (strictly) concave in \mathbf{x}_i and p_j . Thus,

Theorem 6 *The equilibrium set of the Fisher Market is convex, and the equilibrium price vector is unique.*

Aggregate Social Optimization

$$\begin{aligned} &\text{maximize} && \sum_{i \in B} w_i \log(\mathbf{u}_i^T \mathbf{x}_i) \\ &\text{subject to} && \sum_{i \in B} x_{ij} \leq s_j, \quad \forall j \in G \\ &&& x_{ij} \geq 0, \quad \forall i, j, \end{aligned}$$

Theorem 7 (Eisenberg and Gale 1959) Optimal dual (Lagrange) multiplier vector of equality constraints is an *equilibrium price vector*.

Optimality Conditions of the Aggregated Problem

$$\left\{ \begin{array}{ll} w_i \frac{u_{ij}}{\mathbf{u}_i^T \mathbf{x}_i} & \leq p_j, \quad \forall i, j \\ w_i \frac{u_{ij} x_{ij}}{\mathbf{u}_i^T \mathbf{x}_i} & = p_j x_{ij}, \quad \forall i, j \\ \sum_i x_{ij} & \leq s_j, \quad \forall j \\ p_j \sum_i x_{ij} & \leq p_j s_j, \quad \forall j \\ \mathbf{x}_i, \mathbf{p} & \geq \mathbf{0}. \end{array} \right.$$

Let $y_i = \mathbf{u}_i^T \mathbf{x}_i / w_i$. Then, these conditions are **identical** to the equilibrium price conditions, since

$$y_i = \frac{\mathbf{u}_i^T \mathbf{x}_i}{w_i} \geq \frac{u_{ij}}{p_j}, \quad \forall i, j.$$

Rewriting the Aggregate Social Optimization

$$\begin{aligned} &\text{maximize} && \sum_{i \in B} w_i \log u_i \\ &\text{subject to} && \sum_{j \in G} u_{ij}^T x_{ij} - u_i = 0, \quad \forall i \in B \\ &&& \sum_{i \in B} x_{ij} \leq s_j, \quad \forall j \in G \\ &&& x_{ij} \geq 0, \quad s_i \geq 0, \quad \forall i, j, \end{aligned}$$

This is called the **weighted analytic center** problem.

Question: Is the price vector **p** **unique** when at least one $u_{ij} > 0$ among $i \in B$ and $u_{ij} > 0$ among $j \in G$.

Aggregate Example

$$\begin{aligned} \text{maximize} \quad & 5 \log(2x_1 + y_1) + 8 \log(3x_2 + y_2) \\ \text{subject to} \quad & x_1 + x_2 = 1, \\ & y_1 + y_2 = 1, \\ & x_1, x_2, y_1, y_2 \geq 0. \end{aligned}$$

Or

$$\begin{aligned} \text{maximize} \quad & 5 \log(u_1) + 8 \log(u_2) \\ \text{subject to} \quad & 2x_1 + y_1 - u_1 = 0, \\ & 3x_2 + y_2 - u_2 = 0, \\ & x_1 + x_2 = 1, \\ & y_1 + y_2 = 1, \\ & x_1, x_2, y_1, y_2 \geq 0. \end{aligned}$$