

## Examples of recurrences

$$1. \quad a_n = 5a_{n-1} - 6a_{n-2}, \quad a_0 = 2, \quad a_1 = 3$$

linear, 2nd degree, const-coeff., homogeneous

Our approach  $a_n = r^n$

$$r^n = 5r^{n-1} - 6r^{n-2}$$

$$r^{n-2}(r^2 - 5r + 6) = 0$$

$$r^2 - 5r + 6 = 0$$

$$(r-3)(r-2) = 0 \rightarrow r = 2, 3$$

$$a_n = \alpha_1 2^n + \alpha_2 3^n \quad (*)$$

Choose  $\alpha_1, \alpha_2$  in order to satisfy IC

$2 = a_0$	$= \alpha_1 \cdot 2^0 + \alpha_2 \cdot 3^0 = \alpha_1 + \alpha_2$
$3 = a_1$	$= \alpha_1 2^1 + \alpha_2 3^1 = 2\alpha_1 + 3\alpha_2$

2 eqns in  $\alpha_1, \alpha_2$

$$\alpha_1 + \alpha_2 = 2 \rightarrow$$

$$2\alpha_1 + 3\alpha_2 = 3$$

$$\alpha_1 = 2 - \alpha_2$$

Subst. into 2nd eqn.

$$2(2 - \alpha_2) + 3\alpha_2 = 3$$

$$4 + \alpha_2 = 3 \rightarrow \alpha_2 = -1$$

$$\alpha_1 = 3$$

Our recurrence soln. (also satisfying the ICs) is

$$a_n = 3(2^n) - 3^n$$

2.  $a_n = 4a_{n-1} - 4a_{n-2}, \quad a_0 = 1, \quad a_1 = 3$

linear, 2<sup>nd</sup> degree, const coeff., homog.

Take  $a_n = r^n$

$$r^n - 4r^{n-1} + 4r^{n-2} = 0$$

$$r^{n-2}(r^2 - 4r + 4) = 0$$

char. eqn.  $r^2 - 4r + 4 = 0$

$$(r-2)(r-2) = 0 \rightarrow \begin{matrix} \text{roots} \\ (\text{repeated}) \end{matrix} r = 2,$$

Have

seq.

$$2^n: 1, 2, 2^2, 2^3, \dots$$

$2^n \leftarrow$  usually  $2^n$  seq. is different (different root)

Repeated root makes it impossible to combine

Bad news  $\alpha_1 \cdot 2^n + \alpha_2 \cdot 2^n$  to satisfy the ICs.

Good News

It's as simple as adding  $n$  as a factor to the  $2^n$  sequence

$$\begin{array}{c} \text{1st seq.} \\ \text{---} \\ n=0 & n=1 & n=2 & n=3 \\ \text{2}^n: & 1, & 2, & 2^2, & 2^3, \dots \end{array}$$

$2^{nd}$  seq

$$n2^n: 0, 2, 8, 24, \dots$$

↑  
added factor of  $n$

- a trick to produce a  $2^{nd}$  sequence which also satisfies the recurrence (not necessarily the ICs)
- a different seq. than the first
- only applies if there is a repeated root underneath

Now use just as in other examples  
mixtures

$$a_n = \alpha_1 2^n + \alpha_2 n \cdot 2^n \quad (\star)$$

Still must solve for  $\alpha_1, \alpha_2$

$$\begin{aligned} 1 &= a_0 = \alpha_1 \cdot 2^0 + \alpha_2 \cdot 0 && \xrightarrow{\text{2 eqs. in } \alpha_1, \alpha_2} \\ 3 &= a_1 = \alpha_1 \cdot 2 + \alpha_2 \cdot 1 \cdot 2^1 && \end{aligned}$$

$$\begin{cases} 1 = \alpha_1 \\ 3 = 2\alpha_1 + 2\alpha_2 \end{cases}$$

$$\text{get } \alpha_1 = 1$$

$$\alpha_2 = 1/2$$

Soln. is (A) rewritten w/ these alphas:

$$a_n = 2^n + \frac{1}{2} \cdot n \cdot 2^n = 2^n \left( 1 + \frac{n}{2} \right).$$

### 8.3: Divide-and-conquer algorithms

#### A. Binary search algorithm

Have sorted list of size  $n$

Divide list in 2 halves

make 2 comparisons

- key w/ the element at end of 1<sup>st</sup> sublist  
(tells which sublist might contain key)
- see if sublist has 2 or more elements in it

Have list (still sorted, half size of original) to search

Say, for a list of size  $n$ ,

$f(n) = \# \text{ of comparisons to search a sorted list of size } n$   
using this alg. for your key

If  $n$  is even

$$f(n) = 2 + f(\frac{n}{2})$$

divide-and-conquer  
recurrence  
relation

#### B. Fast-integer multiplication

Idea: Say I'm stuck using 4-digit decimal nos.

$$d = (d_3 d_2 d_1 d_0)_{10} = d_3 \times 10^3 + d_2 \times 10^2 + d_1 \times 10^1 + d_0 \times 10^0$$

Task: multiply two such nos.  $d$  and  $s = (\delta_3 \delta_2 \delta_1 \delta_0)$

Define  $f(n) = \# \text{ of operations required to multiply } n\text{-dig.t decimal nos.}$

Notice:

$$\begin{aligned} d &= (d_3 \times 10^3 + d_2 \times 10^2) \times 10^1 + (d_1 \times 10^1 + d_0 \times 10^0) \\ &= (d_3 d_2)_{10} \times 10^2 + (d_1 d_0)_{10} \quad \left( \begin{array}{l} 2791 = 27 \times 10^2 + 91 \\ = (27) \times 10^2 + (91) \end{array} \right) \\ &= D_1 \times 10^2 + D_0 \end{aligned}$$

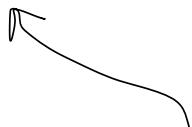
Can do the same w/  $s$ : Define  $\Delta_1 = (\delta_3 \delta_2)_{10}, \Delta_0 = (\delta_1 \delta_0)_{10}$

$$\text{so } s = \Delta_1 \times 10^1 + \Delta_0$$

Intent: Multiply

$$\begin{aligned} ds &= (D_1 \cdot 10^2 + D_0)(\Delta_1 \cdot 10^2 + \Delta_0) \\ &= D_1 \Delta_1 \cdot 10^4 + 10^2(D_1 \Delta_0 + D_0 \Delta_1) + D_0 \Delta_0 \\ &= \text{Same expression as last} + 10^2 D_1 \Delta_1 - 10^2 D_1 \Delta_1 + 10^2 D_0 \Delta_0 - 10^2 D_0 \Delta_0 \end{aligned}$$

$$\text{you verify } = (10^4 + 10^2) D_1 \Delta_1 + 10^2 (D_1 - D_0)(\Delta_0 - \Delta_1) + (10^2 + 1) D_0 \Delta_0$$



Note: Each of  $D_1, D_0, \Delta_1, \Delta_0$  are 2-digit nos. }  
If I multiply ds in the manner prescribed here

I replace a single 4-digit multiplication problem

(w/  $f(4)$  operations) by 3 multiplications of

2 digit nos w/  $3 \cdot f(2)$  operations

also picking up 4 operations (add/subtracts)

So

$$f(4) = 3f(2) + 4$$

More generally

$$f(n) = 3f(\frac{n}{2}) + 4$$

Examples: Solving linear recurrences

$$1. \quad a_n = 5a_{n-1} - 6a_{n-2}, \quad a_0 = 2, \quad a_1 = 3$$

Substitute  $a_n = r^n$

$n=0$	$n=1$	$n=2$
2: 1, 2, 4, ...		
3: 1, 3, 9, ...		

$$r^n = 5r^{n-1} - 6r^{n-2}$$

$$r^{n-2}(r^2 - 5r + 6) = 0$$

characteristic eqn:  $r^2 - 5r + 6 = 0$   
 $(r-3)(r-2) = 0 \rightarrow r = 2, 3$

Need a mixture

$$a_n = \alpha_1 2^n + \alpha_2 3^n \quad (\dagger)$$

Get  $\alpha_1, \alpha_2$  applying ICS

$$\underline{\text{ICS}} \quad 2 = a_0 = \alpha_1 \cdot 2^0 + \alpha_2 \cdot 3^0 = \alpha_1 + \alpha_2$$

$$3 = a_1 = \alpha_1 \cdot 2^1 + \alpha_2 \cdot 3^1 = 2\alpha_1 + 3\alpha_2$$

To finish, must solve 2 eqns. in unknowns  $\alpha_1, \alpha_2$

$$\left. \begin{array}{l} \alpha_1 + \alpha_2 = 2 \\ 2\alpha_1 + 3\alpha_2 = 3 \end{array} \right\} \rightarrow \begin{array}{l} \alpha_1 = 3 \\ \alpha_2 = -1 \end{array}$$

Rewrite  $(\dagger)$  as a final soln. using this ↑

$$\boxed{a_n = 3 \cdot 2^n - 3^n}$$

2. (contains a new element)

$$a_n = 4a_{n-1} - 4a_{n-2}, \quad a_0 = 1, \quad a_1 = 3$$

$$\text{Insert } a_n = r^n$$

$$r^n - 4r^{n-1} + 4r^{n-2} = 0$$

$$r^{n-2}(r^2 - 4r + 4) = 0$$

$$\text{char. eqn.} \quad r^2 - 4r + 4 = 0$$

$$(r-2)(r-2) = 0 \rightarrow \begin{matrix} \text{repeated} \\ \text{root} \end{matrix} \quad r=2$$

(New!)

In other examples  $n=0, n=1, n=2, n=3, \dots$

$$\text{root 1} \quad r_1 : \quad | \quad r_1 \quad r_1^2 \quad r_1^3 \quad \dots$$

$$\text{root 2} \quad r_2 : \quad | \quad r_2 \quad r_2^2 \quad r_2^3 \quad \dots$$

This time

$$z : \quad | \quad 1 \quad 2 \quad z^2 \quad z^3 \quad \dots$$

$$\text{identical} \rightarrow z : \quad | \quad 1 \quad 2 \quad z^2 \quad z^3$$

*Bad news*  $\left\{ \begin{array}{l} \text{Repeated root} \rightarrow \text{impossible to combine} \\ \alpha_1 (\text{first}) + \alpha_2 (r_2^{n^*}) \\ \text{and satisfy the ICS} \end{array} \right.$

Good news: ready replacement for that  $2^n$  sequence

i.e. repeated

$$2^n : 2^0, 2^1, 2^2, \dots$$

Replace w/

$$n2^n : 0 \cdot 2^0, 1 \cdot 2^1, 2 \cdot 2^2, 3 \cdot 2^3, \dots$$

Note: This sort of replacement for a soln. seq. does not work in general, but it does when a repeated root underlies things.

Plan of action:

$$a_n = \alpha_1 2^n + \alpha_2 n 2^n (\star)$$

Choosing  $\alpha_1, \alpha_2$  via ICs as before.

ICs

$$1 = a_0 = \alpha_1 + \alpha_2 \cdot 0 \rightarrow \alpha_1 = 1$$

$$3 = a_1 = 2\alpha_1 + 2\alpha_2 \rightarrow 2\alpha_1 + 2\alpha_2 = 3$$

$$\text{get } \alpha_1 = 1, \alpha_2 = \frac{1}{2}$$

So far.

$$a_n = 2^n + \frac{1}{2} n \cdot 2^n = 2^n \left( 1 + \frac{n}{2} \right)$$

### 8.3 Divide-and-conquer algorithms.

#### A. Binary search

Sorted list of size  $n$

Write:  $f(n) = \# \text{ of comparisons required in searching}$   
 this list of size  $n$  for a key  
 using binary search.

$$f(n) = 2 + f\left(\frac{n}{2}\right) \quad \text{if } n \text{ is even}$$

division eqn. - recurrence relation

$$f_n = 2 + f_{n/2}$$

#### B. Fast-integer multiplication

Restrained to 4-digit decimal nos

$$\begin{aligned} d &= (d_3 d_2 d_1 d_0)_{10} = d_3 \times 10^3 + d_2 \times 10^2 + d_1 \times 10^1 + d_0 \times 10^0 \\ &= (d_3 d_2)_{10} \times 10^2 + (d_1 d_0)_{10} \\ &\quad \uparrow \qquad \text{like saying } 4279 = (42) \times 10^2 + 79 \\ &= D_1 \times 10^2 + D_0. \end{aligned}$$

Task: Multiply 2 4-digit nos.  $d = (d_3 d_2 d_1 d_0)_{10}$  and  $\delta = (\delta_3 \delta_2 \delta_1 \delta_0)_{10}$

Similarly,  $\delta = \Delta_1 \times 10^2 + \Delta_0$

$$d\delta = (D_1 \cdot 10^2 + D_0)(\Delta_1 \cdot 10^2 + \Delta_0)$$

Note:  
 $D_1, D_0, \Delta_1, \Delta_0$   
 are all  
 2-digit nos.

$$\begin{aligned}
 &= D_1 \Delta_1 \cdot 10^4 + 10^2 (D_1 \Delta_0 + D_0 \Delta_1) + D_0 \Delta_0 \\
 &= \underset{\text{same as above}}{+} 10^2 D_1 \Delta_1 - 10^2 D_1 \Delta_1 + 10^2 D_0 \Delta_0 - 10^2 D_0 \Delta_0 \\
 \text{you verify} \quad &= (10^4 + 10^2) D_1 \Delta_1 + 10^2 (D_1 - D_0)(\Delta_0 - \Delta_1) + (10^2 + 1) D_0 \Delta_0
 \end{aligned}$$

Why? !!!!

Let  $f(n) =$  # of operations required to multiply  
2 n-digit decimal nos.

Our work shows adds/subtracts

$$f(4) = 3f(2) + 4$$

More generally, if  $n$  is even

$$f(n) = 3f(n/2) + 4 \quad \begin{matrix} \text{divide-and-conq.} \\ \text{recurrence eqn.} \end{matrix}$$

$$\begin{aligned}
 (10^2 + 1) \cdot 5^9 &= (5900) + 5^9 \\
 &= 5959
 \end{aligned}$$

$$\begin{aligned}
 (2^4 + 2^2 + 1) \cdot (10)_2 &= (100000)_2 + (1000)_2 + (10)_2
 \end{aligned}$$

$$= \begin{pmatrix} 0 & 1 & 0 & 1 & 0 \end{pmatrix}_2$$