

3. Let a_0, a_1, a_2, \dots be the sequence defined by the 2nd-order linear recursion relation

$$a_n = 6a_{n-1} - 5a_{n-2}, \quad \text{for } n \geq 2, \quad \text{with } a_0 = 0, a_1 = 4.$$

Take $P(n): a_n = 5^n - 1$. Then $\forall n \in \mathbb{N}$, $P(n)$ (use strong mathematical induction).

P_n is the claim that the formula $a_n = 5^n - 1$ is closed formula for the n^{th} term of the sequence solving our recurrence relation.

Base case: $n=2$ $P_2: a_2 = 5^2 - 1$ satisfies

$$a_2 = 6(a_1) - 5(a_0) \quad \text{or} \quad 24 = a_2 \stackrel{?}{=} 6(4) - 5(0) \quad \checkmark$$

strong induction: Assume P_2, P_3, \dots, P_k hold for some $k \geq 2$.

Thus

$$\begin{aligned} a_{k+1} &= 6(\underline{a_k}) - 5(\underline{a_{k-1}}) = 6(5^k - 1) - 5(5^{k-1} - 1) \\ &= 6(5^k) - 6 - 5(5^{k-1}) + 5 = 6(5^k) - 5^k - 1 = 5^k(6-1) - 1 \\ &= \underline{5^{k+1} - 1} \quad \text{destination (hoped-for)} \end{aligned}$$

4. Use strong mathematical induction to show the product of n numbers requires $n - 1$ multiplications, regardless of grouping.

To multiply $r_1 r_2 r_3 \dots r_n$ requires $n-1$ multiplications

What domain for n ? $n \geq 1$ \checkmark

Base step: list of 1 number requires 0 multiplications

Suppose for inductive step that $\underline{P_1, P_2, \dots, P_k}$ holds for some $k \geq 1$.
I.H.

Now take a list of $k+1$ numbers

$(r_1 r_2 r_3 \dots r_k r_{k+1})$ last (final outer set of parens) group

This final grouping produces 2 sublists, one of size m , the other of size j

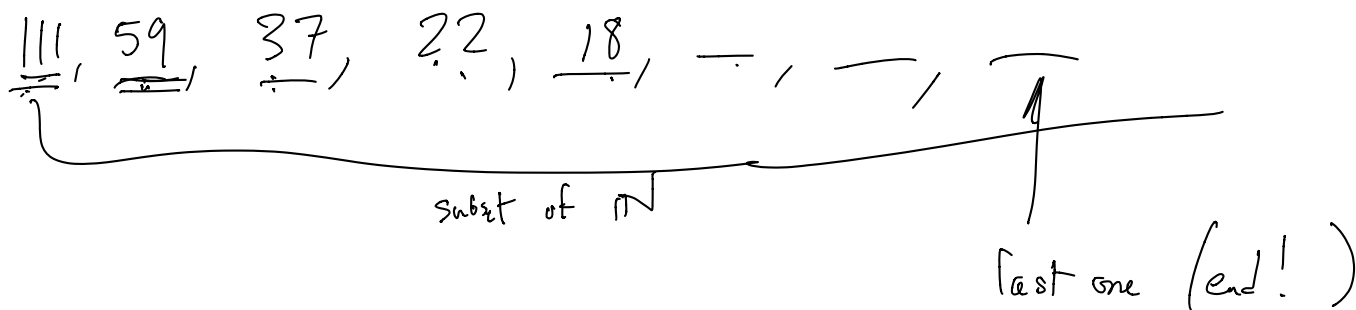
where $\left. \begin{array}{l} 1 \leq m < k+1 \\ 1 \leq j < k+1 \end{array} \right\} m+j = k+1$

By I.H., 1st sublist requires $(m-1)$ multiplies, 2nd requires $(j-1)$ multiplies

total # of multiplies:
 $(m-1) + (j-1) + 1 = k$ multiplies

5. A simple polygon with $n \geq 3$ sides can be triangulated into $n - 2$ triangles (use strong mathematical induction, and the fact that every simple polygon with at least four sides has an interior diagonal).

6. Given a strictly decreasing sequence of positive integers r_1, r_2, r_3, \dots (so $r_{i+1} < r_i$ for each i),
 For instance the sequence terminates (use the well-ordering principle).



Ex.] $n = 57, d = 5$
 $\Rightarrow q = 11, r = 5$?

Ex.] $n = -33, d = 4$
 $\Rightarrow q = -9, r = 3$?

$\{ \dots, 47, 52, 57, 62, 67, \dots \}$

Division Algorithm

7. Given any integer n and any positive integer d , there exist integers q and r such that $n = dq + r$ and $0 \leq r < d$ (use the well-ordering principle).

Statement to prove: Given integers n and d , $d > 0$, there exist unique integers q and r satisfying:

- $n = qd + r$
- $0 \leq r < d$

$$\begin{array}{r} \text{quotient} \\ \downarrow \\ q \\ d \overline{) 73} \\ \underline{-(q \cdot d)} \\ r \\ \text{remainder} \end{array}$$

prove using well-ordering principle:

Consider the set $A = \{ n - qd \mid q \in \mathbb{Z} \} \cap \mathbb{N}$

This $A \subseteq \mathbb{N}$, so by well-ordering principle, A has a smallest element, call it

$$r = n - qd \text{ for some } q \in \mathbb{Z}.$$

Now we have r and q . Note that, by construction,

$$n = r + qd$$

And $r \geq 0$ since r comes from A (containing only nat. nos.)

and either $r < d$ or it isn't.

But, if $r \geq d$ (or $d \leq r$), we'd have

$$d \leq r = n - qd$$

so subtracting d from all these expressions

$$0 \leq r - d = n - (q+1)d$$

which makes $r-d$ another element inside A , even smaller than r , and that's impossible, since r is already smallest.