Inferencia de facetas de declasificación Propuesta de solución

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- 2. Consideraciones importantes
- 3. Propuesta: Proceso iterativo de inferencia asistida
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Descripción y motivación del

problema

#### Descrpción y motivación del problema

• Se quiere realizar inferencia de facetas públicas para facilitarle el trabajo al programador.

#### Descrpción y motivación del problema

- Se quiere realizar inferencia de facetas públicas para facilitarle el trabajo al programador.
- Dado que Dart provee la inferencia de facetas privadas, el problema se denomina "Inferencia de facetas públicas (o de declasificación) dadas las facetas privadas".

La inferencia global es un problema no decidible, y la presencia de subtyping no mejora la situación.

```
class Person {
   Person foo(Person a, Person b) {
    return a.bar(b);
  }
  Person bar(Person a) {
   return a.foo(a, this);
  }
}
```

La inferencia global es un problema no decidible, y la presencia de subtyping no mejora la situación.

```
class Person {
    @???? Person foo(@PersonBar Person a, @PersonFoo Person b) {
    return a.bar(b);
  }
    @???? Person bar(@PersonFoo Person a) {
    return a.foo(a, this);
  }
}
```

Entonces, mejor le pedimos al usuario que especifique las facetas de declasificación de la firma de los métodos...

# Propuesta: Proceso iterativo de

inferencia asistida

• Paso 1: Inferir considerando solo las facetas de la firma de un metodo que están definidas.

- Paso 1: Inferir considerando solo las facetas de la firma de un metodo que están definidas.
- Pasos sucesivos: Utilizando las facetas inferidas en el paso 1 y las modificaciones del usuario, refinar las facetas cuando sea pertinente.

```
class Person {
 void foo(Person a, Person b) {
   a.bar(b);
 void bar(Person p) {
   p.baz(p);
 void baz(Person p) {
   print(p);
```

```
class Person {
 void foo(Person a, Person b) {
   a.bar(b);
 void bar(Person p) {
   p.baz(p);
 void baz(Person p) {
   print(p);
```



Inferir facetas no definidas.

```
class Person {
 @void void foo(@PersonBar Person a,
     @Person Person b) {
   a.bar(b);
 @void void bar(@PersonBaz Person p) {
   p.baz(p);
 @void void baz(@Person Person p) {
   print(p);
```

```
class Person {
 @void void foo(@PersonBar Person a,
     @Person Person b) {
   a.bar(b);
 @void void bar(@PersonBaz Person p) {
   p.baz(p);
 @void void baz(@Person Person p) {
   print(p);
```



Flujo no permitido. Refinar facetas.

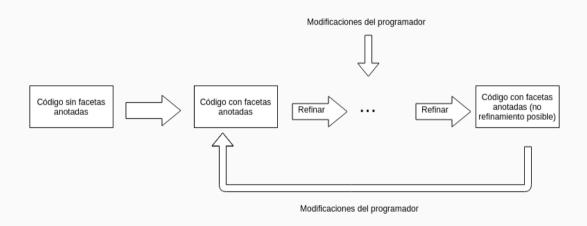
```
class Person {
 @void void foo(@PersonBar Person a,
     @PersonBaz Person b) {
   a.bar(b);
 @void void bar(@Person Person p) {
   p.baz(p);
 @void void baz(@Person Person p) {
   print(p);
```

```
class Person {
 @void void foo(@PersonBar Person a,
     @PersonBaz Person b) {
   a.bar(b);
 @void void bar(@Person Person p) {
   p.baz(p);
 @void void baz(@Person Person p) {
   print(p);
```



Flujo no permitido. Refinar facetas.

```
class Person {
 @void void foo(@PersonBar Person a,
     @Person Person b) {
   a.bar(b);
 @void void bar(@Person Person p) {
                                                  No hay errores:)
   p.baz(p);
 @void void baz(@Person Person p) {
   print(p);
```



Primero, se extiende el lenguaje base de Type-based relaxed noninterference para incluir referencias, instrucciones condicionales y secuencias de instrucciones.

Luego, se escriben las reglas de inferencia, que generan un set de constraints para cada tipo de expresión.

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$$\frac{\Gamma, M, pc, pt \vdash e1: t1 \mid C1 \qquad \Gamma, M, pc, pt \vdash e1: t1 \mid C2}{\Gamma, M, pc, pt \vdash e1: e2: t2 \mid C2 \cup C1}$$

Luego, se escriben las reglas de inferencia, que generan un set de constraints para cada tipo de expresión.

(if)

 $\Gamma, M, pc, pt \vdash e1 : t1 \mid C1$   $\Gamma, M, pc1, pt \vdash e2 : t2 \mid C2$  $\Gamma, M, pc1, pt \vdash e3 : t3 \mid C3$ 

 $\overline{\Gamma, M, pc, pt \vdash if \ e1 \ then \ e2 \ else \ e3: t \ | \ \{t2 <: t\} \cup \{t3 <: t\} \cup C3 \cup C2 \cup \{t1 <: pc1\} \cup C1}$ 

#### Sections

Sections group slides of the same topic

\section{Elements}

for which metropolis provides a nice progress indicator  $\dots$ 

# Titleformats

#### Metropolis titleformats

metropolis supports 4 different titleformats:

- Regular
- $\bullet$  Smallcaps
- allsmallcaps
- ALLCAPS

They can either be set at once for every title type or individually.

#### Small caps

This frame uses the smallcaps titleformat.

#### Potential Problems

Be aware, that not every font supports small caps. If for example you typeset your presentation with pdfTeX and the Computer Modern Sans Serif font, every text in smallcaps will be typeset with the Computer Modern Serif font instead.

#### all small caps

This frame uses the allsmallcaps titleformat.

#### Potential problems

As this titleformat also uses smallcaps you face the same problems as with the smallcaps titleformat. Additionally this format can cause some other problems. Please refer to the documentation if you consider using it.

As a rule of thumb: Just use it for plaintext-only titles.

#### ALL CAPS

This frame uses the allcaps titleformat.

#### Potential Problems

This titleformat is not as problematic as the allsmallcaps format, but basically suffers from the same deficiencies. So please have a look at the documentation if you want to use it.

Elements

#### Typography

The theme provides sensible defaults to \emph{emphasize} text, \alert{accent} parts or show \textbf{bold} results.

becomes

The theme provides sensible defaults to emphasize text, accent parts or show bold results.

#### Font feature test

- Regular
- Italic
- $\bullet$  SmallCaps
- Bold
- Bold Italic
- Bold SmallCaps
- Monospace
- Monospace Italic
- Monospace Bold
- Monospace Bold Italic

#### Lists

#### ${\rm Items}$

- Milk
- Eggs
- Potatos

#### Enumerations

- 1. First,
- 2. Second and
- 3. Last.

#### Descriptions

PowerPoint Meeh.

Beamer Yeeeha.

#### Animation

 $\bullet\,$  This is important

#### Animation

- ullet This is important
- Now this

# Animation

- $\bullet$  This is important
- Now this
- And now this

# Animation

- ullet This is really important
- Now this
- And now this

# Figures

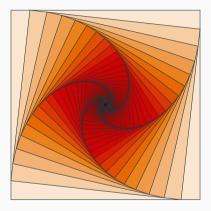


Figure 1: Rotated square from texample.net.  $\,$ 

# Tables

Table 1: Largest cities in the world (source: Wikipedia)

City	Population
Mexico City	20,116,842
Shanghai	19,210,000
Peking	15,796,450
Istanbul	$14,\!160,\!467$

### Blocks

Three different block environments are pre-defined and may be styled with an optional background color.

Default Block co

Block content.

Alert

Block content.

Example

Block content.

Default

Block content.

Alert

Block content.

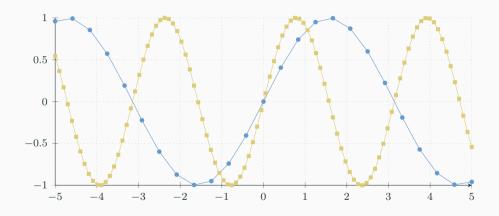
Example

Block content.

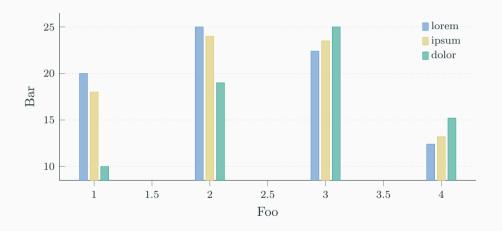
# Math

$$e = \lim_{n \to \infty} \left( 1 + \frac{1}{n} \right)^n$$

# Line plots



# Bar charts



# Quotes

Veni, Vidi, Vici

### Frame footer

metropolis defines a custom beamer template to add a text to the footer. It can be set via

\setbeamertemplate{frame footer}{My custom footer}

My custom footer 32

### References

Some references to showcase [allowframebreaks] [4, 2, 5, 1, 3]

Conclusion

## Summary

Get the source of this theme and the demo presentation from

github.com/matze/mtheme

The theme itself is licensed under a Creative Commons Attribution-ShareAlike 4.0 International License.



Questions?

# Backup slides

Sometimes, it is useful to add slides at the end of your presentation to refer to during audience questions.

The best way to do this is to include the appendixnumberbeamer package in your preamble and call \appendix before your backup slides.

metropolis will automatically turn off slide numbering and progress bars for slides in the appendix.

### References i



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### References ii



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