CLRS Exercise

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1 7

1.1 7.3

1.1.1 a

This is certain concerning the *Randomized* procedure, the probability of any index i is chosen from $[0, n-1]$is:
$Pr(pivot = i) = \frac{1}{n}$
$E(X_i) = 1 * Pr(pivot = i) + 0 * Pr(pivot \neq i) = \frac{1}{n}$

1.1.2 b

It is certain that if ith element is chosen as pivot, $Random - Parition$ cost $\Theta(n)$ time, and it will call $QuickSort[1, q-1], QuickSort[q+1, n]$ recursively.
Concerning only the first $Parition$, this would be the result:
$E(T(n)) = \Sigma_{i=1}^{n} Pr(pivot = i)(T(i-1) + T(n-i) + \Theta(n))$
$= \Sigma_{i=1}^{n} X_i(T(i-1) + T(n-i) + \Theta(n))$

1.1.3 c

Concerning $X_i = \frac{1}{n}$
$E(T(n)) = \Sigma_{i=1}^{n} \frac{1}{n}(T(i-1) + T(n-i) + \Theta(n))$
$= \Sigma_{i=1}^{n} \frac{1}{n} T(i-1) + \Sigma_{i=1}^{n} \frac{1}{n} T(n-i) + \Sigma_{i=1}^{n} \frac{1}{n} \Theta(n)$
$= \frac{2}{n} \Sigma_{i=1}^{n-1} T(i) + \Theta(n)$

1.1.4 d

$\Sigma_{k=2}^{n-1} klgk$
$\leq lg\frac{n}{2} \Sigma_{k=2}^{\frac{n}{2}} k + lgn \Sigma_{k=\frac{n}{2}}^{n-1} k$
$= lgn \Sigma_{k=2}^{n-1} k - lg2 \Sigma_{k=2}^{\frac{n}{2}} k$
$= lgn \frac{(n+1)(n-2)}{2} - \frac{(\frac{n}{2}+2)(\frac{n}{2}-1)}{2}$
$\leq lgn \frac{n^2}{2} - \frac{n^2}{8}$
by Calculus, we have:
$(\frac{1}{2}x^2 lgx - \frac{1}{4}x^2)|_1^{n-1} \leq E(T(n)) \leq (\frac{1}{2}x^2 lgx - \frac{1}{4}x^2)|_2^n$

1.1.5 e

Proof of $E(T(n)) = O(nlgn)$:
Assume that $\forall k \in [1, n-1], \exists c, E(T(k)) \leq cklgk - \Theta(k)$
For $k = n, E(T(n)) \leq \frac{n}{2}c(lgn\frac{n^2}{2} - \frac{n^2}{4} - \Theta(n^2)) + \Theta(n) \leq cnlgn - \Theta(n)$
Proof of $E(T(n)) = \Omega(nlgn)$:
Assume that $\forall k \in [1, n-1], \exists c, E(T(k)) \geq cklgk + \Theta(k)$
For $k = n, E(T(n)) \geq \frac{n}{2}c(lgn\frac{(n-1)^2}{2} - \frac{(n-1)^2}{4} + \Theta(n^2)) + \Theta(n) \geq cnlgn + \Theta(n)$
$\rightarrow E(T(n)) = \Theta(nlgn)$

1.2 7.5

1.2.1 a

From counting Theorem, it could be noticed that:
$p_i = \frac{(i-1)(n-i)}{C_n^3} = \frac{6(i-1)(n-i)}{n(n-1)(n-2)}$

1.2.2 b

$Pr(i = medium)(normal) = \frac{1}{n}$
$Pr(i = medium)(3part) = \frac{6(\frac{1}{2}n-1)(n-\frac{1}{2}n)}{n(n-1)(n-2)} = \frac{3}{2}\frac{1}{n}$
$Pr(3part) - Pr(normal) = \frac{1}{2}\frac{1}{n}$

1.2.3 c

Consider $f_{diff} = \int_{\frac{n}{3}}^{\frac{2}{3}n}(\frac{6(i-1)(n-i)}{n(n-1)(n-2)} - \frac{1}{n})di$
$= \frac{(-2i^3+3(n+1)i^2-6ni-(n-1)(n-2)i)|_{i=\frac{1}{3}n}^{i=\frac{2}{3}n}}{n(n-1)(n-2)}$
$lim_{n\to\infty}f_{diff} = \frac{4}{27}$

1.2.4 d

Consider we are so lucky that each partition we choose the median:
In the Iteration tree, we have:
$T(n) = \begin{cases} c & n = 1 \\ 2T(\frac{1}{2}n) + n & n > 1 \end{cases}$
The $\Omega(nlgn)$ is kept even in best case.

2 8

2.1 8.1-1

n-1 times, since we need n elements to formulate

2.2 8.1-2

$\Sigma_1^n lgk < \int_1^{n+1} lgkdk = (klgk - k)_1^n = (nlgn - n) - (0 - 1) = nlgn - n + 1$

2.3 8.1-3

$\leftrightarrow$ proof at least half of branch is longer than h
Consider a decision tree with $n!/2$ elements
$\leftrightarrow$ proof at least half of branch is longer than h
Consider a decision tree with $n!/n$ elements
$\leftrightarrow$ proof at least half of branch is longer than h
Consider a decision tree with $n!/2^n$ elements, this is not significant enough and could leave only $\Omega(lg\frac{n!}{2^n}) = \Omega(nlgn - n) = \Omega(nlgn)$ elements

2.4 8.2-4

Consider a trim version of counting sort, build the C map up and query directly:

COUNTING-SORT-TRIM(A, k)

$C[]$
for $i = 0$ **to** k
    $C[i] = 0$
for $j = 1$ **to** $A.length$
    $C[A[j]] + +$
for $m = 1$ **to** k
    $C[m] + = C[m-1]$
return $C[m]$

DIRECT-QUERT(A, k, a, b)

$C =$ COUNTING-SORT-TRIM(A, k)
if $a < 1$
    return $C[b]$
else return $C[b] - C[a-1]$

2.5 8.3-4

First, with $O(n)$ time: convert n numbers k_{10} into k_n which has 3 digits.
Second, with $O(d(n+n))$ time ($Lemma$8.3): Radix sort n 3-digit numbers with each digits take up to n possible values.

DIGITSCONVERT(X)

$result[]$
for $i = 2$ **downto** 0
    $result[i] = X/n^i$
    $X = X \bmod n^i$
return $result$

```
SORT(A, x)
result[]
for each S in A
        S = DIGITSCONVERT(S)
RADIX-SORT(A, x)
```

3 9

3.1 9.1

3.1.1 a

Sorting: MERGE-SORT(A) in worst case $O(nlgn)$
Query: CALL-BY-RANK(A, k) i times in worst case $O(i)$, here we assume manipulating $O(n)$ space cost $O(n)$ time.

3.1.2 b

Building: BUILD-MAP-HEAP(A) in worst case $O(n)$
Query: calling EXTRA-MAX(A, k) i times in worst case $O(ilgn)$

3.1.3 c

Selecting: SELECT(A, i) in worst case $O(n)$
Sorting: MERGE-SORT$(A^{'})$ in worst case $O(ilgi)$

4 11

4.1 11.2

4.1.1 a

Consider for a ball i fall into a specific bucket $Pr(i) = \frac{1}{n}$
Then consider Binomial Distribution, $Pr(k) = C_n^k Pr(i)^k (1 - Pr(i))^{n-k}$

4.1.2 b

Consider random picking a slot, the probability of that slot is maximum is $Pr_{max} = \frac{1}{n}$, and it contains k elements Q_k. for conditional probability, we have:
$P_k = Pr_{i=k|max} = \frac{Pr(i=k \cap max)}{Pr_{max}} \leq \frac{Pr(i=k)}{Pr_{max}} = nQ_k$

4.1.3 c

Proof:
$Q_k = (\frac{1}{n})^k (\frac{n-1}{n})^{n-k} C_n^k$
$= \frac{(n-1)^{n-k}}{n^n} \frac{\Pi_0^{k-1} n-k}{k!}$

$\leq \frac{n^n}{n^n}\frac{1}{k!}$
$= \frac{e^k}{k^k}\frac{1}{k^{\frac{1}{2}}(1+\Theta(\frac{1}{n}))}$
$\leq \frac{e^k}{k^k}$

4.1.4 d

Proof for Q_{k_0}:
$Q_{k_0} = \frac{e^{(\frac{clgn}{lglgn})}}{(\frac{clgn}{lglgn})^{\frac{clgn}{lglgn}}}$
$= \frac{n^{\frac{clg\frac{e}{c}}{lglgn}}}{\frac{n^c}{n^{\frac{clglglgn}{lglgn}}}} = n^{\frac{clg\frac{e}{c}+clglglgn}{lglgn}-c}$
It would not take effort to notice that since $lim_{n\to\infty}\frac{clg\frac{e}{c}+clglglgn}{lglgn} = 0$
$\forall c > 3+\epsilon, Q_{k_0} = O(\frac{1}{n^3})$
And $P_k \leq nQ_k \to P_k = O(\frac{1}{n^2})$

4.1.5 e

$E(M) = \Sigma_{M=1}^{n} MPr(M) < nPr(M > \frac{clgn}{lglgn}) + \frac{clgn}{lglgn}Pr(M \leq \frac{clgn}{lglgn})$
A stronger conclusion to note:
$E(M) = \Sigma_{M=1}^{n} MPr(M) < MPr(M > \frac{clgn}{lglgn}) + \frac{clgn}{lglgn}Pr(M \leq \frac{clgn}{lglgn})$
$\leq \int_{\frac{clgn}{lglgn}}^{\infty}\frac{1}{n}dn + 1 * \frac{clgn}{lglgn}$
$= lg(\frac{clgn}{lglgn}) + \frac{clgn}{lglgn}$
$= O(\frac{clgn}{lglgn})$

5 15

5.1 15.1-1

$2^n - 1 = \Sigma_{j=0}^{n-1}2^j$

5.2 15.1-2

Do not know how!

5.3 15.1-3

See Code

5.4 15.1-4

See Code

Ex 15.2.4

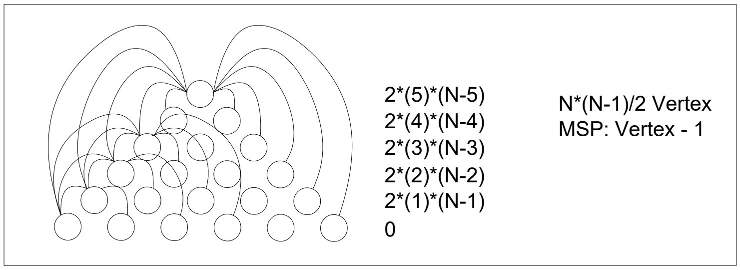

Figure 1: 15.2-4

5.5 15.1-5

See Code

5.6 15.2-1

See Code

5.7 15.2-2

See Code

5.8 15.2-3

Assume that $\forall k \leq n-1, T(k) \geq c2^k$
Then $T(n) = \Sigma_{k=1}^{n-1} T(k)T(n-k) = (n-1)c^2 2^n > c2^n$
So $T(n) = \Omega(n), \omega(n)$

5.9 15.2-4

See Figure 1

5.10 15.2-5

For each level $h(i) = i(n-i)$
For tree $T(n) = 2\Sigma_{i=1}^{n-1} i(n-i)$
$= \frac{3n^3+3n^2}{3} - \frac{2n^3+3n^2+n}{3}$
$= \frac{n^3-n}{3}$

5.11 15.2-6

Assume that $\forall k \leq n-1, N(k) = k-1$
Then $N(n) = N(n-1) + 1$
So $N(n) = n - 1$

5.12 15.3-1

running through: $T(n) = n * P_n^n = n * n! > 4^n$
running recursion: $T(n) = 2\Sigma_{i=1}^{n-1} 4^i + n = \frac{8}{3} 4^{n-1} + n \leq 4^n$
running through takes longer

5.13 15.3-2

no overlapping subproblem call

5.14 15.3-3

Yes

5.15 15.3-4

Do not know how!

5.16 15.4-1

See code

5.17 15.4-2

See code

5.18 15.4-3

See code

5.19 15.1

LSP(s, t, G)

$r = G, size()$
$DPs[r] = 0$
$DPr[r] = path(s, t)$
$max = -\infty$
for $i = 1$ **to** r
  $max(DPs[j] + DPr[r-j] + what)$
return max

5.20 15.1