Investigating Higher-order Simplification for SMT Solving

work in progress

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What this is about

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The Isabelle simplifier provides a basic form of automatization.

Examples

```
fun add :: "nat ⇒ nat ⇒ nat" where
  "add 0 n = n" |
  "add (Suc m) n = Suc(add m n)"

lemma add_02: "add m 0 = m"
  apply(induction m)
  apply(simp)
  apply(simp)
  done
```

What we are looking at

- ► Isabelle interfaces with automatic systems via the Sledgehammer tool
- ▶ How does the simplifier compare to Sledgehammer?

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- Isabelle interfaces with automatic systems via the Sledgehammer tool
- How does the simplifier compare to Sledgehammer?
- ► The *Mirabelle* tool is a tool included with the Isabelle distribution for experiments with Sledgehammer
- ► Goals are considered trivial if they are also solved by try0
- We added an additional parameter to run the simplifier only

Numbers

- ▶ We ran this on 64 randomly chosen AFP entries
- ► Configuration: Timeout of 20s, default provers
- Overall 1973 Sledgehammer calls
- ► Sledgehammer succeeded 1223 (62%) times
- ► The simplifier solved 344 (17%) problems and 45 (2.3%) exclusively

Numbers with comments

- ▶ In an earlier test we identified 24 simp only problems and investigated them further
- ▶ 13 were solved when reproducing locally with default Sledgehammer settings
- Many were some form of computation (i.e. in theories about cryptography or functional programming)
- Example problem Tycon/Lazy_List_Monad.thy:coerce_LNi¹:
 - ► Fails when naively applying Sledgehammer
 - The simplifier uses four rules
 - Sledgehammer succeeds when adding these rules, but took a while
 - Only exporting those rules results in a easy problem

¹AFP 2018-05-09

The Isabelle simplifier (Nipkow 1989)

- User defined set of rewriting rules
 - automatically from e.g. function definitions
 - ► Lemmas annotated with [simp]
 - When calling the simplifier (also allows restricting the set)
- ▶ This should be confluent and terminating
 - In practice both often does not hold
 - but it doesn't matter much for the user
- Workings in HOL: no encoding needed.
- Some extensions: Conditional, local assumptions, simp progs

The Isabelle simplifier

Main task of the Isabelle simplifier: Express rewriting in the LCF-Style framework of Isabelle.

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Utilizes lazy lists to express the possible infinite number of unifiers.

(First-order) rewriting and SMT

- Exported simp annotations have been used in the superposition prover SPASS (Blanchette et al. 2012)
- ► SMT however is not fundamentally based on rewriting, but ground reasoning + instantiation
- Some current usage:
 - 1. as pre-processing
 - 2. as part of the theory reasoner (see next slide)
 - 3. Z3 supports explicit simplification commands (Moura and Passmore 2013)

•
$$\{z = y \times y, y = w + 2, w = 1\}$$

- $M = \{z = x, y = w + 2, w = 1\}, X(M) = \{x = y \times y\}$

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- ▶ $M \vDash_A y = 3$ and we get $\sigma = \{y \mapsto 3\}$

Solve some extension of a theory by using qualities entailed by the base theories together with simplification rules.

$$M = \{z = x, y = w + 2, w = 1\}, X(M) = \{x = y \times y\}$$

▶
$$M \vDash_A y = 3$$
 and we get $\sigma = \{y \mapsto 3\}$

$$(y \times y) \sigma \downarrow = (3 \times 3) \downarrow = 9$$

Then apply the solver for the base theory.

Onward

Some ideas:

- Naive: Use some rules as pre-processing and get some potentially helpful new rules
- ▶ Handling quantifiers in SMT is still an area of active research
- Even more important as we move to HOL

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Can we separate some quantified input formulas into a *simpset* and use this for some form of rewriting instead of instantiation?

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