# **CSC 138 Final Study Guide**

Sunday, December 12, 2021 | 02:38 AM

# **Chapter 01**

#### What's Internet protocol

Are specific messages sent and specific actions we take in response to the received reply messages or events; on the whole, protocols define how the format, the order, of the messages sent, and actions taken on transmission or received, receipt, message.

#### **Network Core**

**The Network Core**: The mesh of packet switches and links that interconnects the Internet's end systems.

**Circuit switching networks VS Packet-switching networks** have both traditionally occupied different spaces within corporations.

- Circuit-switched networks are ideal for communications which require data to be transmitted in real-time
  - · Are normally used for traditional telephones calls
- Packet-switched networks are more flexible and more efficient if some amount of delay is acceptable
  - Are normally used to handle digital data

#### **Packet Switching**

- 1. In a network application end systems exchange **messages** with each other
  - 1. **Messages** can contain anything the application designer wants
  - 2. Messages can perform control functions ("handshake" between systems)
  - 3. Messages can contain data (e-mail message, JPEG image, or MP3 audio file)
- 2. To send a **message** between two systems (source and destination end system)
  - 1. Packets: The source system breaks messages into smaller chunks of data
  - 2. In travel each packet goes though communication links and packet switches
    - 1. Packet Switches: Two main types routers and link-layer switches
  - 3. Packets pass over each comm link equal to full transmission rate of the link
    - 1. If a source system or packet switch is sending L-bits over a link with transmission rate R bits/sec; then, the time to transmit the packet is

$$\frac{L}{R}seconds$$

#### **Store and Forward**

The entire packet must arrive at the router before it can be transmitted to the next link.

#### **End-End Delay -**

It takes

$$\frac{L}{R}seconds$$

to transmit, send-out, L-bit packet over communication link at a rate of

$$R\frac{bits}{sec}$$

#### **Queuing and Loss -**

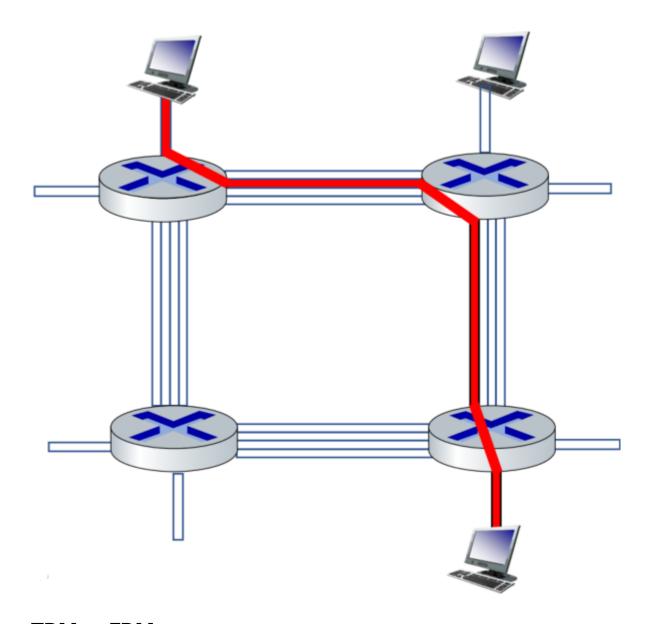
If the arrival rate, in bits/second, to the communication link exceeds transmission rate (bits/second) of the link for some period of time:

- 1. Packets will queue waiting to be transmitted to the output link
- 2. Packets can be dropped (loss) when the memory (buffer) in the router fills up

## **Circuit Switching**

Circuit switching is a type of switching in which a dedicated channel or circuit is established for the duration of communications. A method used by the old traditional telephone call where the call/signal is carried over the **Public Switched Telephone Network** (PSTN), or also known as the **Plain Old Telephone Service** (POTS).

**In short**: An end-to-end resources allocated, reserved (not shared), for the "call" between source and destination. *It works like a telephone switchboard* 



#### **TDM vs FDM**

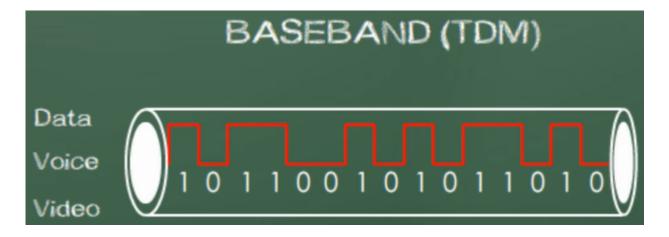
When do we use Time Division Multiplexing (TDM)? When do we use Frequency Division Multiplexing (FDM)?

To answer the questions above, we need to know about baseband and broadband

#### **BASEBAND (TDM)**

Baseband deals with digital signals. Baseband systems can transmit only one signal at a time. Digital signals occupy the entire bandwidth, which means baseband requires an exclusive use of the channel's capacity.

- When a user is sending a signal, no one else can use it at the same time.
  - For example: Ethernet



#### **Baseband: Time Division Multiplexing (TDM)**

TDM transmit different users' signals, the "call", over one link (the maximum rate of the frequency band) by dividing time into slots or intervals and assigning them equally among the users.

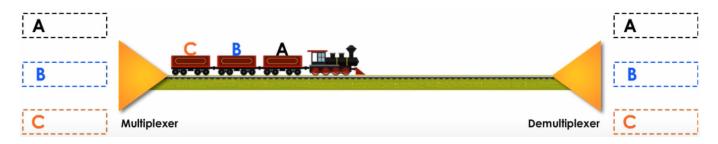
**Main Feature quality of TDM**: Every user is allocated with the same amount of time, but equality is not efficient in many situations.

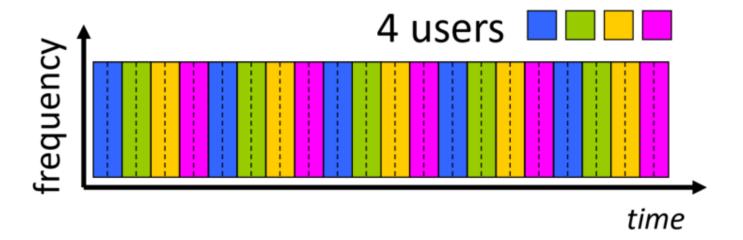
**In summary**: TDM, Time Division Multiplexing, is about time allocation. Every user is given the same amount of time in a "call".

#### **Example:**

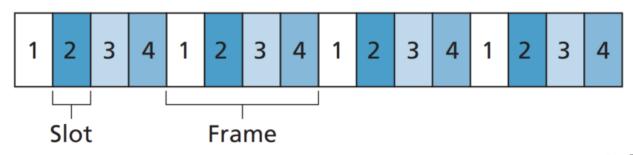
1. **Ethernet** is an example of baseband system.

Some users may have more data or a higher priority than others. Giving the same amount of time regardless of need or priority is not cost effective. As a result Statistical Time Division Multiplexing, or Statistical Multiplexing, is recommended.





#### **TDM**



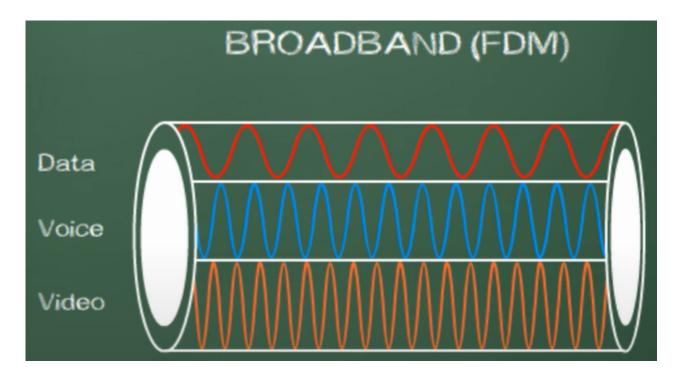
Key:

2

All slots labeled "2" are dedicated to a specific sender-receiver pair.

#### **BROADBAND (FDM)**

Broadband transmissions deal with analog signals which are in the form of optical or Radio Frequency, RF, waves. Broadband allows multiple transmissions at the same time. They can share one channel by using different frequency range. Thus, FDM, Frequency Division Multiplexing, is used in the broadband system.



#### **Broadband: Frequency Division Multiplexing (FDM)**

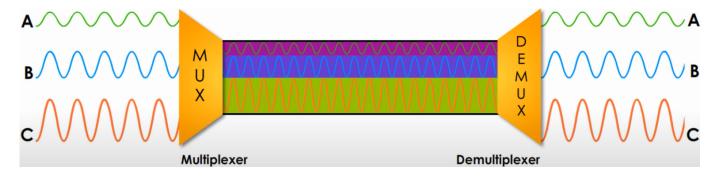
The optical electromagnetic frequencies, analog signals (not digital signals), are divided (narrowed) into frequency bands; where, each call gets its own band, and can only transmit at the max rate of the narrow band.

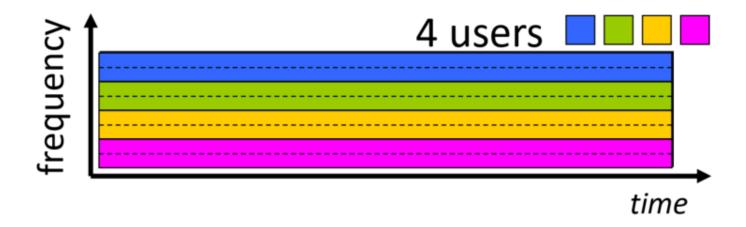
- 1. Optical electromagnetic frequencies are divided (narrowed) into frequency bands
- Each call gets its own band and can only transmit at the max rate of the narrow band

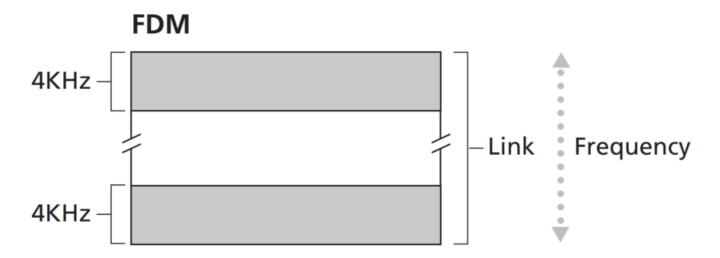
**In Summary**: FDM, Frequency Division Multiplexing, is about frequency channel and each user used their own non-overlapping frequency channel.

#### **Example:**

- 1. FM radio stations also use Frequency Division Multiplexing (FDM) to share the frequency spectrum between 88 MHz and 108 MHz, with each station being allocated a specific frequency band.
- 2. The broadband system is an example of Frequency Division Multiplexing (FDM)







#### **Packet Delay -**

#### Four Sources of Packet Delay -

- 1. Nodal Processing
- 2. Queuing Delay
- 3. Transmission Delay
- 4. Propagation Delay

# **Packet Loss -**

If a router is queuing packets because of high congestion, and the queue is full, the packets will be dropped.

## **Throughput**

The rate (bits per second / time unit) at which bits are being sent from sender to receiver.

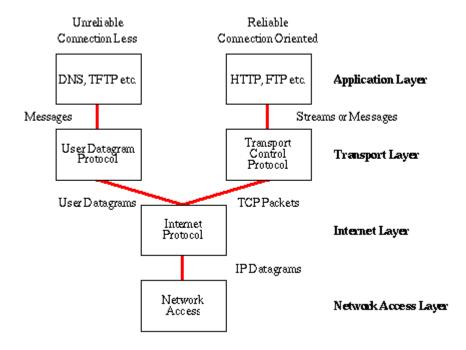
- 1. Instantaneous: rate at given point in time
- 2. Average: rate over longer period of time

#### **Protocol Layer**

- 1. Network designers organize protocols including the hardware and software that implement the protocols in **layers**, so to provide structure.
- 2. Each protocol belongs to one of the layers

#### **Internet Protocol Stack**

- 1. Application Layer (**message**): Applications implement services of the Transport Layer
- 2. Packets of information at the Application Layer are referred to as messages
- 3. Transport Layer (**segment**): Encapsulates the application layer message and integrates the protocol type (UDP, TCP, etc...)
- 4. The book refers to a transport-layer packet as a segment
- 5. Network Layer (**datagrams**): Transfers messages from one host to another using link layer.
- 6. The internet's network layer is responsible for moving network-layer packets known as **datagrams** from one host to another.
- Link Layer (frames): Transfers datagrams from host to neighboring host using the network layer
- 8. The book refers to the link-layer packets as **frames**
- Physical Layer: The physical layer includes the physical medium of transfer (twistedpair copper wire, single-mode fiber optics, coaxial cable) and protocols to handle that transfer



Messa			ge	M	Application
Segment			H <sub>t</sub>	M	Transport
Datagram		H <sub>n</sub>	$H_{t}$	M	Network
Frame	$H_{\mathrm{I}}$	$H_{n}$	H <sub>t</sub>	M	Link
					Physical

Five-layer Internet protocol stack

Figure 1.23 ◆ The Internet protocol stack

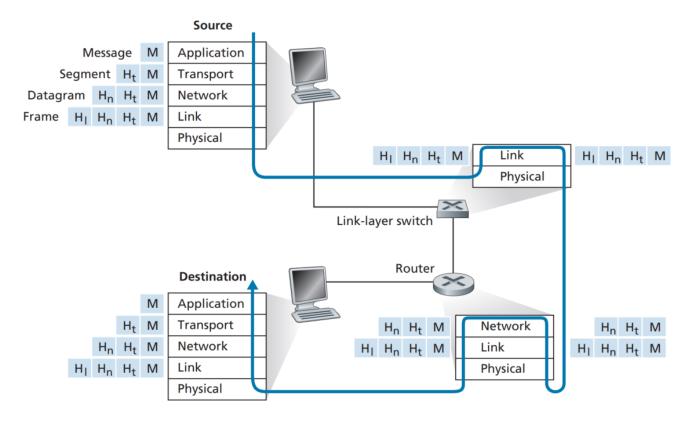


Figure 1.24 • Hosts, routers, and link-layer switches; each contains a different set of layers, reflecting their differences in functionality

## **Network Security**

#### Denial-of-Service (DoS) Attacks / Distributed DoS (DDoS) Attacks

Purposeful overloading of a server with requests (usually through multiple machines controlled by one attacker) to prevent legitimate access

#### **Categories of DoS attacks**

- 1. **Vulnerability attack**: Sending a few well-crafted messages to a vulnerable application or Operating System running on the targeted system, so if the right sequence of packets sent will stop service or crash the host system
- 2. **Bandwidth flooding**: Send a deluge of packets to the targeted host, so that the target's access link becomes clogged
- 3. **Connecting flooding**: Establishes a large number of half-open or fully open TCP connections at the target host, so the host system will become bogged down with bogus connections it will stop accepting ones from legitimate connections

#### **Packet Sniffing**

Intercept and copy packets via an application where unencrypted info can be displayed and inspected

# **IP Spoofing**

Creation of packets with a modified source address to hide the identity of the sender or impersonate another sender

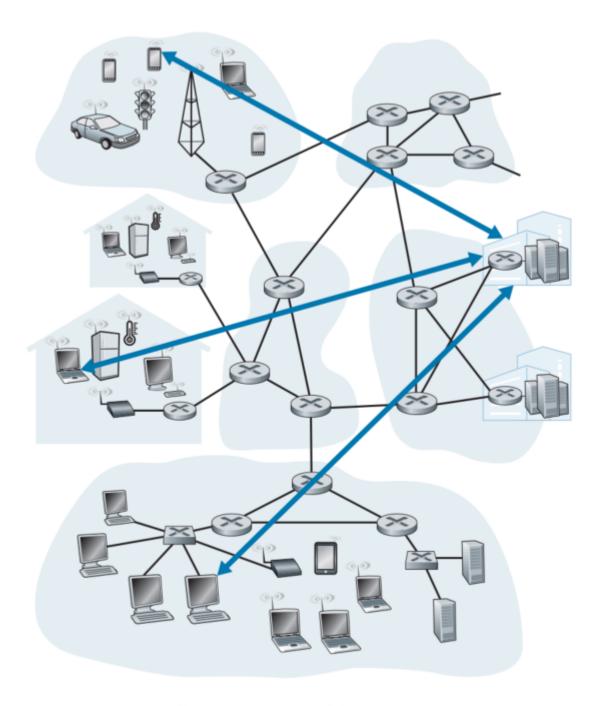
# **Chapter 02**

# **Application Models**

#### **Client-server**

Clients do not directly talk to each other and is mediated by an always-on host (a server)

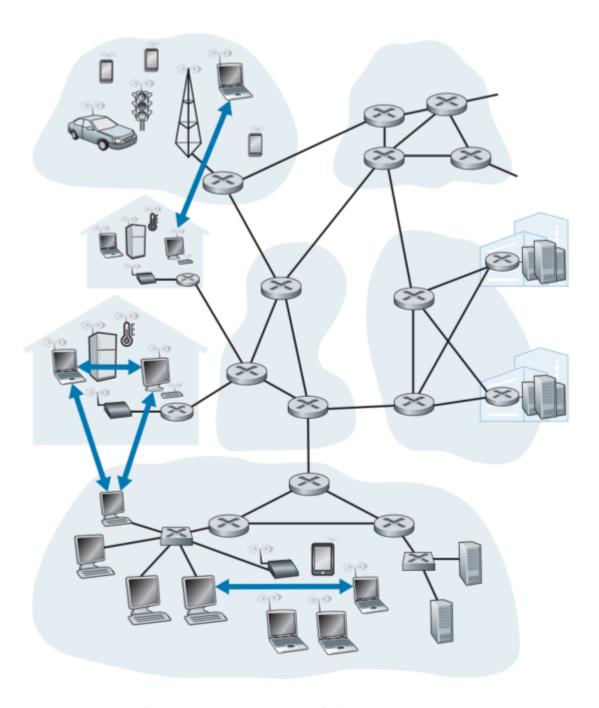
Managing a client-server connection is far simpler than Peer-to-Peer (P2P)



#### a. Client-server architecture

# Peer-to-Peer (P2P)

In a P2P architecture, there is minimal (or no) reliance on dedicated servers in data centers. Instead the application exploits direct communication between pairs of intermittently connected hosts, called peers. The peers are not owned by the service provider, but are instead desktops and laptops controlled by users, with most of the peers residing in homes, universities, and offices.



b. Peer-to-peer architecture

# **Transport protocols -**

There are two transport protocols **TCP** and **UDP**, and either of which can transport application-layer messages.

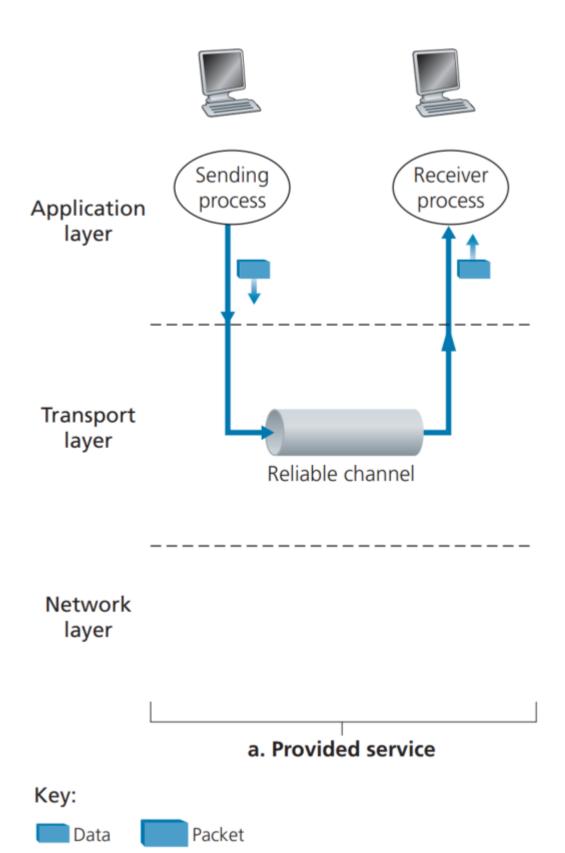
#### **TCP**

- 1. **TCP** provides a connection-oriented service to its applications.
- 2. **TCP** guaranties delivery of application-layer messages to the destination and flow control (sender/receiver speed matching)

- 3. **TCP** Breaks long messages into shorter segments and provides a congestion control mechanism
  - 1. so the source throttles its transmission rate when the network is congested

#### TL;DR

- Reliable transport
- Flow control; sender will not overwhelm receiver
- Congestion control; sender can be throttled
- No timing, security, or minimum throughput
- Setup is required to establish connection



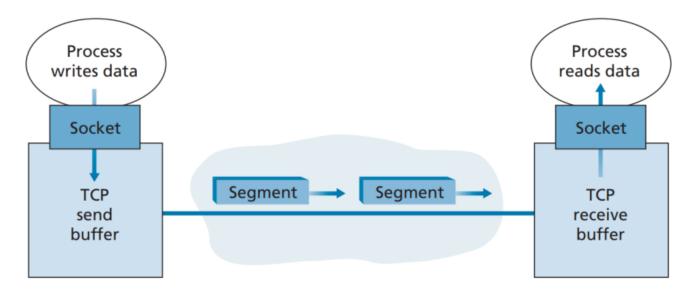


Figure 3.28 → TCP send and receive buffers

#### **UDP**

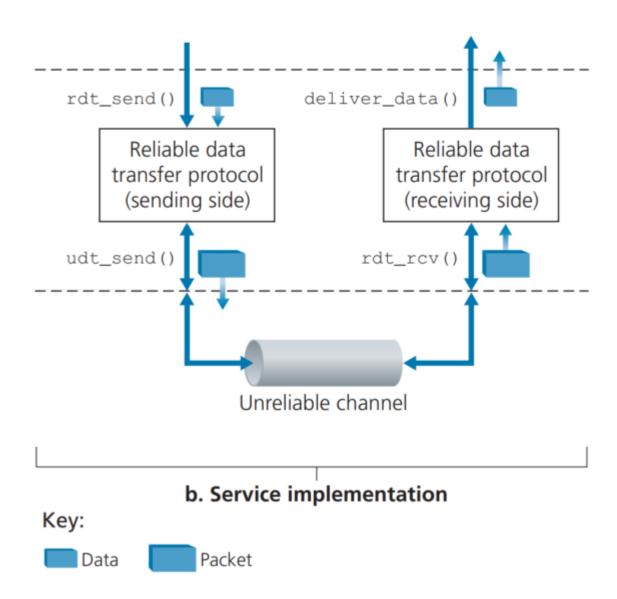
- 1. **UDP** provides a connectionless service to its applications
- 2. UDP DOES NOT provide reliability, flow control, and congestion control

#### TL;DR

- Unreliable transport
- No flow control, congestion control, timing, minimum throughput, security, or setup







**Web and HTTP-**

**HTTP protocol-**

**Persistence and non-persistence HTTP-**

Cookies-
Web caches-
Conditional GET-
Electronic mail
Three major components
SMTP
Mail access protocols
SMTP vs HTTP
DNS
Definition
Services
Different Tiers of servers
Video Steaming
DASH: server and client
<b>Content distribution networks</b>
Chapter 03
Transport-layer services
Multiplexing and demultiplexing
UDP
Internet checksum
Rdt1.0 to 3.0

Pipelined protocols
Go-Back N
Selective repeat
TCP overview
Seq and ack
RTT
Retransmission
Flow control
3-way handshake
TCP congestion control
AIMD
RENO and Tahoe
ECN
Chapter 04 and Chapter 05
Forwarding and routing
Data plane and control plane
What's inside a router
Longest prefix matching
Switching fabrics
IP datagram
IP addressing

Subnets
DHCP
NAT
IPv6 Format
Tunneling
Routing protocols
Classification
Link-state routing
Distance vector
Intra-AS routing (RIP, OSPF, IS-IS, IGRP)
Inter-AS routing (BGP)
Chapter 06
Link Layer
Error Detection (CRC)
Multiple Access protocols
TDMA and FDMA
Slotted ALOHA, Pure ALOHA, CSMA/CD
MAC addresses and ARP
Ethernet (switch)
Putting all together
Chapter 07

#### Wireless

Wireless links and network characteristics

WiFi

Cellular network: 4G LTE

# **Chapter 08**

**Network Security** 

**Principles of cryptography** 

Symmetric key cryptography

**Public Key Cryptography** 

**Authentication, message integrity** 

**Securing e-mail** 

# Part 1: 15 Short-answer questions (carefully, briefly and clearly)

**EXAMPLE** 

1. What is the main reason that packet loss may occur in computer network?

# Part 2: 10 Multiple-choice questions (Single answer) please circle it

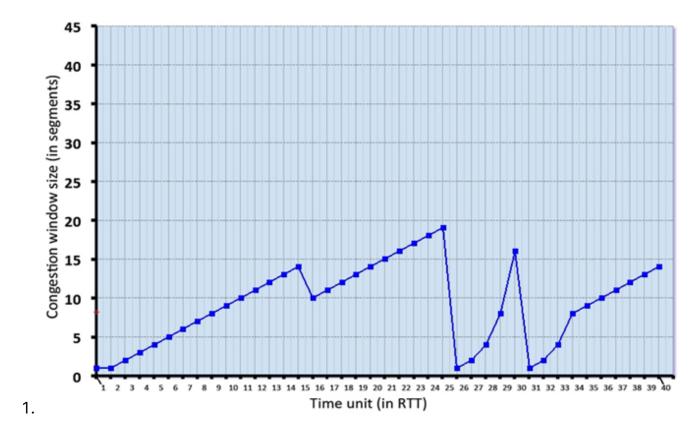
**EXAMPLE** 

- 5. Which of the following about SMTP is false?
- (a) It enables a mail server to send email messages to another mail server
- (b) User agent cannot use it to retrieve email
- (c) It typically uses a connectionless protocol such as UDP
- (d) Unlike FTP it uses a single port number

1

# Part 3: 5 Long-answer questions

**EXAMPLE** 



# 2) Consider the Cyclic Redundancy Check (CRC) algorithm discussed in Section 5.2.3 of the text. Suppose that the 4-bit generator (G) is 1001, that the data payload (D) is 10011101 and that r=3. What are the CRC bits R?

Recall from the text that in calculating the CRC bits, all CRC calculations are done in modulo-2 arithmetic without carries in addition or borrows in subtraction. This means that addition and subtraction are identical and both equal to the bitwise exclusive-OR (XOR) of the operands.

To compute the CRC, we begin by taking the value of D, 10011101, and multiplying it by 2^3, giving 10011101000. We then divide this number by the generator bits [G] = 1001, using modulo-2 arithmetic. The final remainder, R, after this division are then the CRC bits. Here is that calculation:

```
10001100

1001 ) 10011101000

1001

----

1101

1001

----

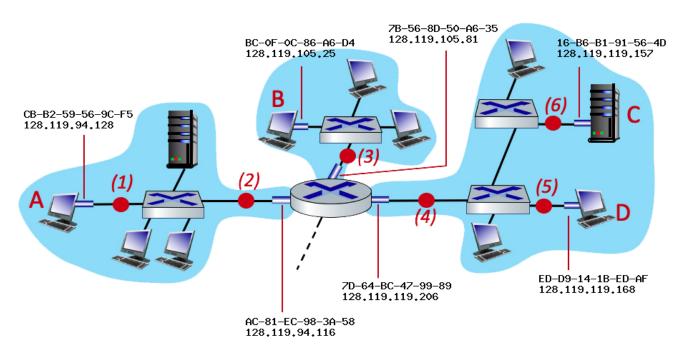
1000

1001

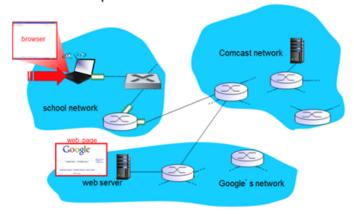
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```

We've computed the remainder as R = 100 and the quotient n = 10001100. You should verify that n\*G XOR R is indeed equal to 10011101000. You can use <u>this calculator</u> to do the modulo-2 arithmetic if you don't want to do it by hand.

3) The IP and MAC addresses are shown for nodes A, B, C and D, as well as for the router interfaces. Consider an IP datagram being sent from NODE C to NODE B. Give the source and destination Ethernet address, as well as the source and destination address of the IP datagram encapsulated within the Ethernet frame at points (6), (4), and (3) in the figure



Let's use the following scenario to journey down the protocol stack completely: a student attaches laptop to campus network, requests the homepage from www.google.com web server. Suppose Comcast is the Internet Service Provider of the campus network.



5.

Bob chooses 
$$p=5$$
,  $q=7$ . Then  $n=35$ ,  $z=24$ .  
 $e=5$  (so e, z relatively prime).  
 $d=29$  (so ed-1 exactly divisible by z).  
encrypting 8-bit messages.

encrypt: bit pattern m me c = me mod n of m = cd mod n decrypt: c = cd = cd

481968572106750915091411825223071697

6.

#### course:

• CSC 138-06 - Computer Network Fundamentals

#### **HELPFUL RESOURCE:**

https://gaia.cs.umass.edu/kurose\_ross/interactive/index.php