

course: [CSC 135 - Computing Theory and Programming Languages](#)

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related notes: [2022-04-12](#) [2022-04-07-CSC135-01-LEC-parsing](#)

# Pushdown Automata Parsing

W15.2 | Tuesday, April 12, 2022 | 10:27 AM

## PDA Parsing Example01

$$S \rightarrow aS$$

$$S \rightarrow T$$

$$T \rightarrow bT$$

$$T \rightarrow R$$

$$R \rightarrow cR$$

$$R \rightarrow \lambda$$

### Step01: Find which are Nullable?

All are Nullable:  $S, T, R$

### Step02: What is our set constraints?

Grammar		
$S \rightarrow aS$	$a \in \text{first}(S)$	
$S \rightarrow T$	$\text{First}(T) \leq \text{First}(S)$	$\text{First}(S) \leq \text{First}(T)$
$T \rightarrow bT$	$b \in \text{first}(T)$	
$T \rightarrow R$	$\text{First}(R) \leq \text{First}(T)$	$\text{First}(T) \leq \text{First}(R)$
$R \rightarrow cR$	$c \in \text{first}(R)$	
$R \rightarrow \lambda$		
$R \rightarrow \$\$$	$\$ \in \text{Follows}(S)$	

$S \rightarrow aS \rightarrow aT \rightarrow abT \rightarrow abbT \rightarrow abbR \rightarrow abbc$

- $abbR \rightarrow R \rightarrow \lambda \rightarrow abbc \rightarrow$

## Step03: Build out our sets

	FIRST	Follow
$S$	$a, b, c$	$\$$
$T$	$b, c$	$\$$
$R$	$c$	$\$$

## STEP04: Build out our prediction table

Grammar	FIRST RHS	IF RHS Nullable FOLLOWS LHS	PREDICTORS
$S \rightarrow aS$	$a$	—	$a$
$S \rightarrow T$	$b, c$	$\$$	$b, c, \$$
$T \rightarrow bT$	$b$	—	$b$
$T \rightarrow R$	$c$	$\$$	$c, \$$
$R \rightarrow cR$	$c$	—	$c$
$R \rightarrow \lambda$	—	$\$$	$\$$

## STEP05: Is this grammar suitable for LL(1) parsing?

Check for three properties - if all are "NO" then yes the grammar is suitable for LL(1) parsing - Need prediction table

1. Is ambiguous grammar? - **has to be "NO"**
  1. Two left most \_\_\_\_
  2. Does it have any left recursion? - Is it left recursive? - **has to be "NO"**
    1.  $A \rightarrow fw$
  3. Are there conflicting predictors for any non-terminal? - **has to be "NO"**

## STEP 06: now write the parser...



### Code writing the parser - Pushdown Automata Parsing

```
def parse(input):  
    toks = scanner(input)  
    stack = ['S']  
    while len(stack) > 0:  
        top = stack.pop()
```

# Always pop top of stack

```

tok = toks.next()                # None indicates token stream
empty
if top in ('a', 'b', 'c'):        # Try input/stack match
    toks.match(top)
elif top == 'S' and tok == 'a':
    stack.append('S')
    stack.append('a')
elif top == 'S' and (tok == None or tok in ('b', 'c')):
    stack.append('T')
elif top == 'T' and (tok == None or tok in ('b')):
    stack.append('R')            # production to follow here
elif top == 'R' and tok == 'c':
    stack.append('R')
    stack.append('c')            # production to follow here
elif top == 'R' and tok == None:
    pass # Push nothing
else:
    raise Exception              # Unrecognized top/tok
combination
if toks.next() != None:
    raise Exception

```

## Recursive Decent Parsing

If you have a production

Production	
$A \rightarrow bCD$	in: b

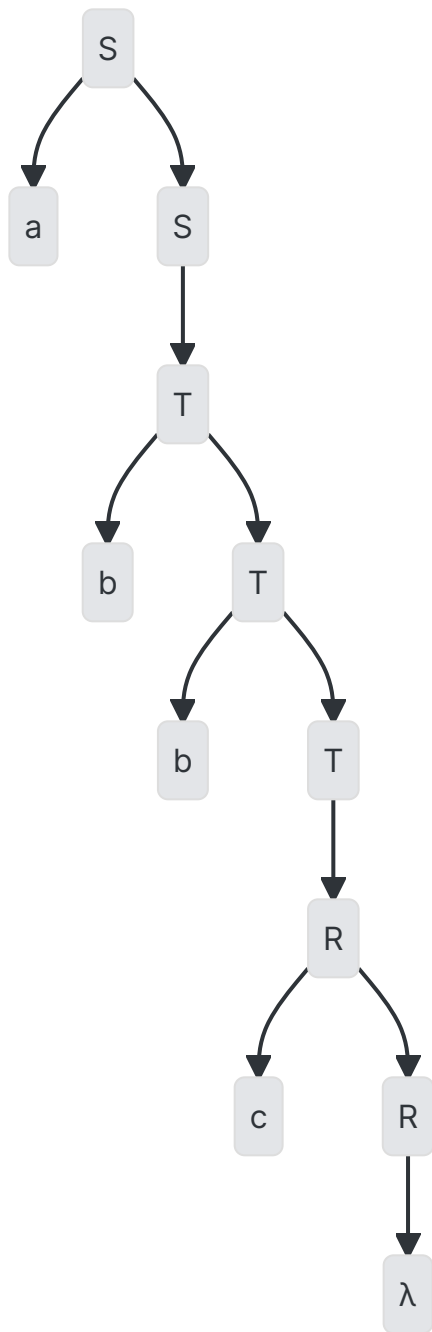
b

C

D

STACK

where b, C, D is our to-do list



### CODE: Recursive Decent Parsing

```
def parseT(toks):
    toks = toks.next()
    if tok == 'b':
        toks.match(toks)
        parseS(toks)
    elif tok == None or tok == 'c':
        parseR(toks)
    else:
        raise Exception
```

```
def parseS(toks):
    tok = toks.next()
    if tok == 'a':
        toks.match(toks)
        parseS(toks)
    elif tok == None or tok in ('b', 'c'):
        parseT(toks)
    else:
        raise Exception
```

```
def recursive_parse(input):
    toks = scanner(input)
    parseS(toks)
    if toks.next() != None:
        raise Exception
```

```
try:
    recursive_parse("aabbcc")
except:
    print("Reject")
else:
    print("Accept")
```