Team Notebook

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Contents	2	DP		6	5.4	numbertheory	12
1 Algorithms	2	2.1	ConvexHullTrick	6		simplex	
1.1 AP-Bridges	2 3	Geo	0	6	5.0	tricks	14
1.2 Hungarian	2	3.1	ConvexHull	6	6 M	isc	14
1.3 SCC-Tarjans	2	3.2	template	6	6.1	OrderedSet	14
1.4 TernarySearch	$\frac{2}{2}$	~				2 template	
1.5 UnionFind	3 4	Gra	aph	10	0.2		10
1.6 binarylifting-lca	3	4.1	Flow	10	7 D-	-41	11
1.7 centroiddecomposition	3		4.1.1 Dinic	10	·	rthon	16
1.8 dijkstras			4.1.2 MinCostMaxFlow	10	7.1	InputArray	15
1.9 findcicles	4						
1.10 matrix exponentiation	4 5	Ma	${f th}$	11		rings	15
1.11 segtreeIteractive			Miller-Rabin	11	8.1	kmp	15
1.12 segtreeRMQTemplate		5.2	fft	11	8.2	suffix-array	15
1.13 segtreeRSQTemplate	5	5.3	math	12	8.3	8 z	16

1 Algorithms

1.1 AP-Bridges

```
int dfs(int u,int p){
   dfs_num[u] = dfs_low[u] = ++dfs_counter;
   for(auto v : adjList[u]){
       if(dfs_num[v]==0){
           dfs(v.u):
           if(dfs low[v] >= dfs num[u]){
              articulation[u]=true;
           if(dfs_low[v] > dfs_num[u])
              bridge = true;
           dfs_low[u] = min(dfs_low[u],dfs_low[v]);
       } else if(v!=p)
           dfs_low[u] = min(dfs_low[u],dfs_num[v]);
}
int main(){
   memset(dfs_num,0,sizeof(dfs_num));
   memset(dfs low.0.sizeof(dfs low)):
   bridge=false;
   dfs counter=0:
   dfs(0,-1):
   for(int i = 0; i < N; ++i)</pre>
       if(dfs num[i]==0)
           bridge=true;
   puts(bridge ? "Yes" : "No");
   return 0:
```

1.2 Hungarian

```
/*

This will take a matrix a[N][N] and choose one item for each row such that the sum of all items is minimized.

O(n^3)

*/

#define N 107

11 INF = 10000000000000000011;
int n;

11 a[N][N];
```

```
void hungarian(){
   vector<ll> u(n+1), v(n+1), p(n+1), way(n+1);
   for(int i=1; i<=n; i++){</pre>
       i = [0]q
       11 i0 = 0;
       vector<ll> minv(n+1, INF):
       vector<char> used(n+1, false);
       do {
           used[j0] = true;
           11 i0 = p[j0], delta = INF, j1;
           for(int j=1; j<=n; j++){</pre>
              if(!used[i]) {
                  11 cur = a[i0][j]-u[i0]-v[j];
                  if(cur < minv[j])</pre>
                      minv[j] = cur, way[j] = j0;
                  if(minv[j] < delta)</pre>
                      delta = minv[j], j1 = j;
              }
           for(int j=0; j<=n; j++){</pre>
              if(used[i])
                   u[p[j]] += delta, v[j] -= delta;
                   minv[j] -= delta;
           }
           j0 = j1;
       } while(p[j0] != 0);
       do {
           11 i1 = wav[i0];
           p[j0] = p[j1];
           j0 = j1;
       } while(j0);
   vector<int> ans(n+1):
   for(int j=1; j<=n; j++)</pre>
       ans[p[i]] = i;
```

1.3 SCC-Tarjans

```
typedef pair<int, int> ii;
int N,M;
vector<int> adjList[MX_N];
int dfs_num[MX_N],dfs_low[MX_N];
bool vis[MX_N];
stack<int> scc;
```

```
int dfsCounter=1:
int sccIdx=1:
map<int, int> sccMap;
void tarians(int u){
   scc.push(u);
   vis[u]=true;
   dfs_low[u] = dfs_num[u] = dfsCounter++;
   for(int i = 0: i < adiList[u].size(): i++){</pre>
       int v = adjList[u][i];
       if(dfs_num[v]==0){
           tarjans(v);
           dfs_low[u]=min(dfs_low[u],dfs_low[v]);
       } else if(vis[v]){
           dfs_low[u]=min(dfs_low[u],dfs_num[v]);
       }
   }
   if(dfs_low[u] == dfs_num[u]){
       while(1){
           int v = scc.top(); scc.pop();
           sccMap[v]=sccIdx;
           vis[v]=false;
           if(v==u)
              break;
       }
       sccIdx++;
```

1.4 TernarySearch

```
double ternary_search(double 1, double r) {
   double eps = 1e-9;
                                //set the error limit here
   while (r - 1 > eps) {
       double m1 = 1 + (r - 1) / 3;
       double m2 = r - (r - 1) / 3:
       double f1 = f(m1):
                          //evaluates the function at m1
       double f2 = f(m2):
                            //evaluates the function at m2
       if (f1 < f2)
          1 = m1;
       else
   }
   return f(1);
                                //return the maximum of f(x)
         in [1, r]
```

1.5 UnionFind

```
/**
  Union find algorithm
  Complexity O(log n) for Join or Find.
*/
int pai[N];

void init(int n){
  for(int i=1; i<=n; i++){
    pai[i]=i;
  }
}
int find(int i){
  if(pai[i]==i)return i;
  return pai[i]=find(pai[i]);
}

int join(int a, int b){
  a=find(a);
  b=find(b);
  pai[a]=pai[b];
}</pre>
```

1.6 binarylifting-lca

```
#include <bits/stdc++.h>
#define MAXN 100100

typedef long long l1;

using namespace std;

int n, m, s[MAXN], depth[MAXN], anc[MAXN][40];
vector<int> g[MAXN];
bool vis[MAXN];

int dfs(int x, int d, int p){
  vis[x] = true;
  depth[x] = d;
  s[x] = 1;
  anc[x][0] = p;
  for(int i = 1; pow(2, i) <= d; i++){
    anc[x][i] = anc[anc[x][i - 1]][i - 1];</pre>
```

```
for(int i = 0; i < g[x].size(); i++){</pre>
 if(vis[g[x][i]]) continue;
 s[x] += dfs(g[x][i], d + 1, x);
return s[x];
int walk(int x, int d){
int i = 0:
while(d){
 if(d \& 1) x = anc[x][i];
 d /= 2:
 i++;
//cout << "\n":
return x;
int lca(int x, int y){
//cout << x<<y;
if(depth[x] < depth[y]) y = walk(y, depth[y] - depth[x]);</pre>
if(depth[x] > depth[y]) x = walk(x, depth[x] - depth[y]);
//cout << x<<v;
if(x == y) return x;
for(int i = 30: i >= 0: i--){
 if(depth[x] >= pow(2, i) && anc[x][i] != anc[y][i]){
 return lca(anc[x][i], anc[v][i]):
 }
}
return anc[x][0]:
int main(){
ios_base::sync_with_stdio(false);
   cin >> n:
   for(int i = 0; i < n - 1; i++){
    int a. b:
    cin >> a >> b:
    g[a].push_back(b);
    g[b].push_back(a);
   dfs(1, 0, -1);
   cin >> m:
   for(int i = 0; i < m; i++){</pre>
   int a. b:
    cin >> a >> b:
    if(depth[a] > depth[b]) swap(a, b);
```

```
if(a == b) cout << n;
else{
  int l = lca(a, b);
  int d = -2 * depth[l] + depth[a] + depth[b];
  if(d % 2) cout << "0";
else{
   if(depth[a] == depth[b]) cout << s[1] - s[walk(b, d / 2 - 1)];
   else cout << s[walk(b, d / 2)] - s[walk(b, d / 2 - 1)];
}
cout << "\n";
}</pre>
```

1.7 centroiddecomposition

```
#include <bits/stdc++.h>
#define MAXN 100100
typedef long long 11;
using namespace std;
int n. sz[MAXN]:
bool deleted[MAXN], vis[MAXN];
char ch[MAXN];
vector<int> g[MAXN];
void dfs(int x, int p){
if(vis[x]) return;
vis[x] = true;
sz[x] = 1:
for(auto i : g[x]){
 if(i == p || deleted[i]) continue;
 dfs(i.x):
 sz[x] += sz[i];
//cout << x << " " << sz[x] << "\n";
int findCentroid(int x){
memset(vis, 0, sizeof(vis));
dfs(x, -1):
int p = -1, c = sz[x] / 2;
while(true){
bool found = false:
 for(auto i : g[x]){
```

```
if(!deleted[i] && i != p && sz[i] > c){
   found = true:
   p = x;
   x = i:
   break;
 }
 if(!found) return x;
}
void decomp(int x, char c){
int cen = findCentroid(x);
ch[cen] = c:
deleted[cen] = true;
for(auto i : g[cen]){
 if(deleted[i]) continue;
 decomp(i, c + 1);
}
int main(){
   #ifndef ONLINE_JUDGE
 freopen("input.txt", "r", stdin);
#endif
ios_base::sync_with_stdio(false);
cin.tie(NULL):
memset(deleted, 0, sizeof(deleted));
   cin >> n;
   for(int i = 0; i < n - 1; i++){
    int a, b;
    cin >> a >> b;
    g[a].push_back(b);
    g[b].push_back(a);
   //cout << findCentroid(1):</pre>
   decomp(1, 'A');
   for(int i = 1; i <= n; i++){</pre>
    cout << ch[i] << " ";
```

1.8 dijkstras

```
#include <bits/stdc++.h>
#include <utility>
#define MAXN 505
using namespace std;
```

```
typedef long long 11;
typedef pair<int, int> ii;
int n:
vector<pair<int, int> > g[MAXN];
int dist[MAXN];
void dijkstra(int x){
for(int i = 0; i < n; i++){
 dist[i] = 999999999:
priority_queue<pair<int, int>, vector<pair<int, int> >,
     greater<pair<int, int> > pq;
pq.push({0, x});
dist[x] = 0;
while(!pq.empty()){
 pair<int, int> v = pq.top();
 pq.pop();
 for(int i = 0; i < g[v.second].size(); i++){</pre>
  pair<int, int> u = g[v.second][i];
  if(dist[v.second] + u.second < dist[u.first])</pre>
   pg.push({dist[u.first] = dist[v.second] + u.second, u.
        first}):
}
int main(){
   #ifndef ONLINE JUDGE
 freopen("input.txt", "r", stdin);
ios_base::sync_with_stdio(false);
   //cin >> n;
```

1.9 findcicles

```
int n;
vector<vector<int>> adj;
vector<char> color;
vector<int> parent;
int cycle_start, cycle_end; // In O(M)

bool dfs(int v) {
    color[v] = 1;
    for (int u : adj[v]) {
        if (color[u] == 0) {
            parent[u] = v;
        }
}
```

```
if (dfs(u))
              return true:
      } else if (color[u] == 1) {
           cycle_end = v;
           cycle_start = u;
           return true:
      }
   }
   color[v] = 2;
   return false;
void find_cycle() {
   color.assign(n, 0);
   parent.assign(n, -1);
   cycle_start = -1;
   for (int v = 0; v < n; v++) {</pre>
       if (dfs(v))
           break:
   if (cycle_start == -1) {
       cout << "Acvclic" << endl;</pre>
   } else {
       vector<int> cycle;
       cycle.push_back(cycle_start);
       for (int v = cycle_end; v != cycle_start; v = parent[
           cycle.push_back(v);
       cycle.push_back(cycle_start);
       reverse(cycle.begin(), cycle.end());
       cout << "Cycle found: ";</pre>
       for (int v : cycle)
           cout << v << " ":
       cout << endl:
   }
```

1.10 matrixexponentiation

```
#include <bits/stdc++.h>
#define MAXN 100100
#define DIM 2
#define pii pair<int, int>
#define pb push_back
typedef long long ll;
```

```
using namespace std;
11 \mod = 1e9 + 7;
struct matrix{
11 a[DIM][DIM]:
 matrix(){
 memset(a, 0, sizeof(11) * DIM * DIM);
 void init(){
 a[0][0] = 0: a[0][1] = 1:
 a[1][0] = 1: a[1][1] = 1:
 }
 matrix operator*(matrix b){
 matrix c;
 for(int k = 0: k < DIM: k++) {
  for(int i = 0; i < DIM; i++) {</pre>
   for(int j = 0; j < DIM; j++) {</pre>
    c.a[i][j] = (c.a[i][j] + a[i][k] * b.a[k][j]) % mod;
   }
  }
 }
 return c;
 vector<ll> times(vector<ll> v){
 vector<ll> c(DIM, 0):
 for(int i = 0: i < DIM: i++){</pre>
  for(int j = 0; j < DIM; j++){</pre>
   c[i] += v[j] * a[i][j];
   c[i] %= mod;
 return c;
};
matrix pow_matrix(matrix a, ll n) {
 if (n == 1) return a:
 matrix b = pow_matrix(a, n / 2);
 b = b * b:
 if (n \& 1) b = b * a;
 return b;
int n;
int main(){
```

```
#ifndef ONLINE_JUDGE
freopen("input.txt", "r", stdin);
#endif
ios_base::sync_with_stdio(false);
cin.tie(NULL);
    n = 10;
    vector<11> base = {1, 1}, v;
    matrix m;
    m.init();
    m = pow_matrix(m, 4);
    v = m.times(base);
    cout << v[0];
}</pre>
```

1.11 segtreeIteractive

${\bf 1.12}\quad {\bf segtree RMQ Template}$

```
void init(int i, int 1, int r){
lo[i] = 1:
hi[i] = r;
if(1 == r) return;
int m = (1 + r) / 2;
init(2 * i. l. m):
init(2 * i + 1, m + 1, r);
void prop(int i){
delta[2 * i] += delta[i]:
delta[2 * i + 1] += delta[i]:
delta[i] = 0;
void update(int i){
tree[i] = min(tree[2 * i] + delta[2 * i], tree[2 * i + 1] +
      delta[2 * i + 1]);
void inc(int i, int l, int r, int val){
if(r < lo[i] || hi[i] < 1) return:</pre>
if(hi[i] <= r && lo[i] >= 1){
 delta[i] += val:
 return:
//partial cover case
prop(i):
inc(2 * i, 1, r, val);
inc(2 * i + 1, l, r, val):
update(i);
int query(int i, int 1, int r){
if(r < lo[i] || hi[i] < 1) return INT_MAX; //not in range</pre>
if(hi[i] <= r && lo[i] >= 1) return tree[i] + delta[i]: //
     completely in range
prop(i);
int minLeft = query(2 * i, 1, r);
int minRight = query(2 * i + 1, 1, r);
update(i):
return min(minLeft, minRight);
int main(){
init(1, 0, n - 1);
```

1.13 segtreeRSQTemplate

```
#include <bits/stdc++.h>
#define MAXN 505
typedef long long 11;
using namespace std;
int n;
// Range query, range update w/ lazy
11 \log 4 * MAXN + 1, hi [4 * MAXN + 1], tree [4 * MAXN + 1].
    delta[4 * MAXN + 1]:
void init(int i, int 1, int r){
 lo[i] = 1;
 hi[i] = r:
 if(1 == r) return;
 int m = (1 + r) / 2;
 init(2 * i, 1, m):
 init(2 * i + 1, m + 1, r);
void prop(int i){
 delta[2 * i] += delta[i];
 delta[2 * i + 1] += delta[i];
 delta[i] = 0:
void update(int i){
 tree[i] = tree[2 * i] + delta[2 * i] * (hi[2 * i] - lo[2 *
 + tree[2 * i + 1] + delta[2 * i + 1] * (hi[2 * i + 1] - lo
      [2 * i + 1] + 1):
void inc(int i, int l, int r, ll val){
 if(r < lo[i] || hi[i] < 1) return:</pre>
 if(hi[i] <= r && lo[i] >= 1){
 delta[i] += val:
 return;
 //partial cover case
 prop(i);
 inc(2 * i, 1, r, val);
 inc(2 * i + 1, 1, r, val):
 update(i);
7
void simpleinc(int i, ll val){
```

2 DP

2.1 ConvexHullTrick

```
/**
* Source: Simon Lindholm
* Description: Container where you can add lines of the
     form kx+m, and query maximum values at points x.
* Useful for dynamic programming.
* Requires C++ 14
* Use when dp[i] = max(m(j) * i + b(j)) where m and b are
     determined by some i < i
* Negate everything to get min
* Time: O(\log N)
struct Line {
mutable ll k, m, p;
bool operator<(const Line& o) const { return k < o.k; }</pre>
bool operator<(ll x) const { return p < x; }</pre>
};
struct LineContainer : multiset<Line, less<>> {
// (for doubles, use inf = 1/.0, div(a,b) = a/b)
const ll inf = LLONG_MAX;
ll div(ll a, ll b) { // floored division
return a / b - ((a ^ b) < 0 && a % b); }
```

```
bool isect(iterator x. iterator v) {
 if (y == end()) { x->p = inf; return false; }
 if (x->k == y->k) x->p = x->m > y->m ? inf : -inf;
 else x->p = div(y->m - x->m, x->k - y->k);
 return x->p >= y->p;
void add(ll k, ll m) { //slope k, intercept m
 auto z = insert(\{k, m, 0\}), y = z++, x = y;
 while (isect(y, z)) z = erase(z);
 if (x != begin() \&\& isect(--x, y)) isect(x, y = erase(y));
 while ((y = x) != begin() && (--x)->p >= y->p)
 isect(x. erase(v)):
11 query(11 x) { //max value at point x
 assert(!emptv()):
 auto 1 = *lower_bound(x);
 return 1.k * x + 1.m:
}:
```

Geo

3.1 ConvexHull

```
bool cmp(pt a, pt b){return mk(a.y, a.x) < mk(b.y, b.x);}</pre>
vector<pt> convexhull(vector<pt> p){ //counterclockwise, no
    collinear points
sort(p.begin(), p.end(), cmp);
p.erase(unique(p.begin(), p.end()), p.end());
vector<pt> up. dn:
for(pt i : p){
 while(up.size() > 1 and orient(up[up.size() - 2], up.back
      (), i) >= 0) up.pop_back();
 while(dn.size() > 1 and orient(dn[dn.size() - 2], dn.back
      (), i) <= 0) dn.pop_back();
 up.pb(i);
 dn.pb(i);
for(int i = (int) up.size() - 2; i >= 1; i--)
 dn.pb(up[i]):
return dn;
```

3.2 template

```
#include <bits/stdc++.h>
using namespace std;
#define ll long long int
#define pb push_back
#define mk make_pair
#define mt make tuple
#define fi first
#define se second
#define ii pair<int. int>
#define all(x) (x).begin(), (x).end()
#define N 1000007 // 10e6 + 7
struct stableSum {
 Use stableSum to add (positive) elements that are doubles.
 It greatly reduces imprecision.
int cnt = 0:
vector<double> v. pref{0}:
void operator+=(double a) {
 assert(a >= 0):
 int s = ++cnt:
 while (s % 2 == 0) {
  a += v.back():
  v.pop_back(), pref.pop_back();
  s /= 2:
 v.push_back(a);
 pref.push_back(pref.back() + a);
double val() {return pref.back();}
}:
int quadRoots(double a, double b, double c, pair<double,
    double> &out) {
 quadRoots will give the quadratic answer to equation
  x^2*a + x*b + c for a!=0
 Returns how many solutions, place them in out.
assert(a != 0):
double disc = b*b - 4*a*c:
if (disc < 0) return 0:</pre>
double sum = (b >= 0) ? -b-sqrt(disc) : -b+sqrt(disc);
out = \{sum/(2*a), sum == 0 ? 0 : (2*c)/sum\};
```

```
return 1 + (disc > 0):
Tips:
 - Use Integers whenever possible.
 - Minimize division and square root operations.
 - Try to write code that handles many situations at once.
typedef double T;
typedef complex<T> pt;
#define x real()
#define y imag()
// abs(p) = sqrt(x*x + y*y)
T sq(pt p) {return p.x*p.x + p.y*p.y;}
pt translate(pt v, pt p) {
// Translate a point p by a vector v.
return p+v;
pt scale(pt c, double factor, pt p) {
// Scale point p by factor around a center c.
return c + (p-c)*factor;
pt rot(pt p, double a) {
// Rotate point p by an angle a, counterclockwise.
return p * polar(1.0, a);
pt perp(pt p) {
// Rotate point p by 90 degrees, good for integer coords.
return {-p.v, p.x};
T dot(pt v, pt w) {
 v*w = |v|*|w|*cos(angle)
 Check sign of dot product to see if two vectors are going
      in the same dir.
  Positive if angle < pi/2, neg if >, 0 if =
```

```
return v.x*w.x + v.y*w.y;
bool isPerp(pt v, pt w) {return dot(v,w) == 0;}
double angle(pt v, pt w) {
// Angle between two vectors.
return acos(clamp(dot(v,w) / abs(v) / abs(w), -1.0, 1.0));
T cross(pt v. pt w) {
/*
 v*w = |v|*|w|*sin(angle)
 Order of v, w matters! Angle is the ORIENTED angle between
  Positive if 0 < angle < pi, neg if -pi < angle < 0, zero
       if angle = 0 or pi.
return v.x*w.y - v.y*w.x;
T orient(pt a, pt b, pt c) {
// I'll go from a to b to c. If turn left to c, positive.
     Right negative, straight zero.
return cross(b-a,c-a);
double orientedAngle(pt a, pt b, pt c) {
// Return the oriented angle between ab and ac. going from
     b to c.
if (orient(a,b,c) >= 0)
return angle(b-a, c-a):
return 2*M_PI - angle(b-a, c-a);
bool inAngle(pt a, pt b, pt c, pt p) {
// Use this to check if p lies in the angle that ab and ac
assert(orient(a,b,c) != 0):
if (orient(a,b,c) < 0) swap(b,c):
return orient(a,b,p) >= 0 && orient(a,c,p) <= 0;
bool isConvex(vector<pt> p) {
// To check if a polygon is convex, the orientation of all
     three consecutive
// points should be the same.
bool hasPos=false, hasNeg=false;
for (int i=0, n=p.size(); i<n; i++) {</pre>
```

```
int o = orient(p[i], p[(i+1)\%n], p[(i+2)\%n]):
 if (o > 0) hasPos = true:
 if (o < 0) hasNeg = true;</pre>
 return !(hasPos && hasNeg);
bool half(pt p, pt v = \{-1.0, 0.0\}) { // true if in blue
 // Modify v if you want a different starting angle.
 assert(p.x != 0 || p.v != 0); // the argument of (0.0) is
return cross(v,p) < 0 \mid | (cross(v,p) == 0 && dot(v,p) < 0);
void polarSort(vector<pt> &v) {
 This will sort points according to their angle based on
      the origin.
 If I want to do the same thing but with a point not the
      origin, I have
 to subtract that point from all other points.
 If I want to add parameters in the sort such as magnitude,
       just add terms
 to the tuple.
 sort(v.begin(), v.end(), [](pt v, pt w) {
 return make tuple(half(v), 0) < make tuple(half(w), cross(
      v,w));
}):
}
struct line {
pt v; T c:
// From direction vector v and offset c
line(pt v, T c) : v(v), c(c) {}
 // From equation ax+by=c
line(T a, T b, T c) : v(b,-a), c(c) {}
 // From points P and Q
 line(pt p, pt q) : v(q-p), c(cross(v,p)) {}
 // - these work with T = int
 T side(pt p):
 double dist(pt p);
 double sqDist(pt p);
```

```
double slope():
 line perpThrough(pt p);
 bool cmpProj(pt p, pt q);
 line translate(pt t):
 // - these require T = double
 line shiftLeft(double dist):
 pt proj(pt p);
 pt refl(pt p);
T line::side(pt p) {
 // This says what side of the line a point is.
 // Positive side is on the left (remember the line has
      orientation).
 return cross(v.p)-c:
double line::dist(pt p) {
 // Dist point -> line
 return abs(side(p)) / abs(v);
double line::sqDist(pt p) {
 // Dist point -> line squared.
 return side(p)*side(p) / (double)sq(v);
double line::slope(){
 return v.v/v.x;
line line::perpThrough(pt p) {
 // Line that is perpendicular to this line, and goes
      through p.
 return {p, p + perp(v)};
bool line::cmpProj(pt p, pt q) {
 // Use this if you want to sort points through a line.
 return dot(v.p) < dot(v.q):</pre>
line line::translate(pt t) {
 // Translate this line by vector t.
 return {v, c + cross(v,t)};
line line::shiftLeft(double dist) {
 // Shift this line to the left by dist. Note: you gotta
      substitute.
```

```
return {v, c + dist*abs(v)}:
bool inter(line 11, line 12, pt &out) {
// Check if 11 and 12 intersect.
T d = cross(11.v. 12.v):
if (d == 0) return false;
out = (12.v*11.c - 11.v*12.c) / d; // requires floating-
     point coordinates
return true;
pt line::proj(pt p) {
// Projects a point into a line.
return p - perp(v)*side(p)/sq(v);
pt line::refl(pt p) {
// This is the point that is the same distance from line as
      p, but on the other side.
return p - perp(v)*2.0*side(p)/sq(v);
line bisector(line 11, line 12, bool interior) {
// This returns the line that is between 11 and 12.
     dividing the angle in 2.
assert(cross(11.v, 12.v) != 0); // 11 and 12 cannot be
     parallel!
double sign = interior ? 1 : -1:
return {12.v/abs(12.v) + 11.v/abs(11.v) * sign, 12.c/abs(12
     .v) + 11.c/abs(11.v) * sign}:
bool inDisk(pt a, pt b, pt p) {
// Pts a, b are the diameter of a disk, want to know if
     point p is inside.
return dot(a-p, b-p) \le 0:
bool onSegment(pt a, pt b, pt p) {
// Check if point p is in the segment formed by [a, b].
return orient(a,b,p) == 0 && inDisk(a,b,p);
bool properInter(pt a, pt b, pt c, pt d, pt &out) {
// Check if two segments [a, b], [c, d] incercept.
```

```
// The proper interception is an interception that is a
     single point, but not an endpoint.
 double oa = orient(c,d,a),
 ob = orient(c.d.b).
 oc = orient(a,b,c),
 od = orient(a,b,d):
 // Proper intersection exists iff opposite signs
 if (oa*ob < 0 && oc*od < 0) {</pre>
 out = (a*ob - b*oa) / (ob-oa):
 return true:
 return false:
}
// To create sets of points we need a comparison function
struct cmpX {
bool operator()(const pt &a, const pt &b) const{
 return make pair(a.x. a.v) < make pair(b.x. b.v):
}
};
set<pt,cmpX> inters(pt a, pt b, pt c, pt d) {
 // If |set| = 0, no interception.
 // If |set| = 1, point interception.
 // If |set| = 2, segment interception.
 pt out;
 if (properInter(a,b,c,d,out)) return {out};
 set<pt,cmpX> s;
 if (onSegment(c,d,a)) s.insert(a);
 if (onSegment(c,d,b)) s.insert(b);
 if (onSegment(a,b,c)) s.insert(c);
 if (onSegment(a,b,d)) s.insert(d);
 return s;
double segPoint(pt a, pt b, pt p) {
 // Dist of point p to segment [a, b]
 if (a != b) {
 line l(a,b):
 if (1.cmpProj(a,p) && 1.cmpProj(p,b)) // if closest to
      projection
  return l.dist(p): // output distance to line
 return min(abs(p-a), abs(p-b)); // otherwise distance to A
     or B
double segSeg(pt a, pt b, pt c, pt d) {
// Dist of seg [a, b] to seg [c, d]
```

```
pt dummy:
if (properInter(a,b,c,d,dummy))
 return 0:
return min({segPoint(a,b,c), segPoint(a,b,d), segPoint(c,d,
     a), segPoint(c,d,b)});
            ----- Poligons
double areaTriangle(pt a, pt b, pt c) {
return abs(cross(b-a, c-a)) / 2.0:
double areaPolygon(vector<pt> p) {
double area = 0.0:
for (int i = 0, n = p.size(): i < n: i++) {</pre>
 area += cross(p[i], p[(i+1)%n]); // wrap back to 0 if i ==
return abs(area) / 2.0;
bool above(pt a, pt p) {
// True if P at least as high as A (blue part).
return p.v >= a.v:
bool crossesRay(pt a, pt p, pt q) {
// Check if [PQ] crosses ray from A.
return (above(a,q) - above(a,p)) * orient(a,p,q) > 0;
bool inPolygon(vector<pt> p, pt a, bool strict = true) {
// Check if point a is in polygon p.
// If strict, returns false when A is on the boundary.
int numCrossings = 0;
for (int i = 0, n = p.size(); i < n; i++) {</pre>
 if (onSegment(p[i], p[(i+1)%n], a))
 return !strict:
 numCrossings += crossesRay(a, p[i], p[(i+1)%n]);
return numCrossings & 1; // inside if odd number of
     crossings
```

```
pt circumCenter(pt a, pt b, pt c) {
// Gives the center of the circle that goes though a, b, c.
b = b-a, c = c-a; // consider coordinates relative to A
assert(cross(b,c) != 0); // no circumcircle if A,B,C
     aligned
return a + perp(b*sq(c) - c*sq(b))/cross(b,c)/2.0:
template <typename T> int sgn(T k) {
// Return -1, 0, 1 depending on sign of k.
return (T(0) < k) - (k < T(0)):
int circleLine(pt o, double r, line l, pair<pt,pt> &out) {
// Circle-Line intercection (0, 1, 2).
// If only 1 intercection, out.fi == out.se.
double h2 = r*r - 1.sqDist(o);
if (h2 \ge 0)  { // the line touches the circle
 pt p = 1.proj(o); // point P
 pt h = 1.v*sqrt(h2)/abs(1.v); // vector parallel to 1, of
     length h
 out = \{p-h, p+h\}:
return 1 + sgn(h2):
int circleCircle(pt o1, double r1, pt o2, double r2, pair<pt
     ,pt> &out) {
// Circle-Circle intercection (0, 1, 2, inf).
// Similar to circleLine.
pt d=o2-o1; double d2=sq(d);
if (d2 == 0) {assert(r1 != r2); return 0;} // concentric
double pd = (d2 + r1*r1 - r2*r2)/2; // = |0_1P| * d
double h2 = r1*r1 - pd*pd/d2: // = h2
if (h2 >= 0) {
 pt p = o1 + d*pd/d2, h = perp(d)*sqrt(h2/d2);
 out = \{p-h, p+h\};
return 1 + sgn(h2):
int tangents(pt o1, double r1, pt o2, double r2, bool inner,
     vector<pair<pt,pt>> &out) {
```

```
// There can be (0, 1, 2) tangents.
// If 2 tangents, there are two pairs (p1, p2) of points of
      that tangent on the circles.
// If 1 tangent, pairs are equal.
if (inner) r2 = -r2;
pt d = o2-o1:
double dr = r1-r2, d2 = sq(d), h2 = d2-dr*dr;
if (d2 == 0 || h2 < 0) {assert(h2 != 0); return 0;}</pre>
for (double sign : {-1,1}) {
 pt v = (d*dr + perp(d)*sqrt(h2)*sign)/d2;
 out.push back({o1 + v*r1. o2 + v*r2}):
return 1 + (h2 > 0);
pt minEnclosingCircle(vector<pt>v){
// Given a bunch of points, what is the smallest circle
     that contains all of them?
// Return center.
pt p = \{0, 0\}:
for(int i=0; i<v.size(); i++){</pre>
 p+=v[i];
}
if(v.size() == 0)return p;
p/=v.size();
double walk = 0.1;
double d:
for(int i=0: i<30000: i++){</pre>
 int k = 0;
 d = abs(p-v[0]):
 for(int j=1; j<v.size(); j++){</pre>
  if(abs(p-v[i]) > d){
   d = abs(p-v[j]);
   k = j;
 p += (v[k] - p)*walk;
 walk *= 0.999;
// d is the radius
return p:
int main(int argc, char const *argv[]){
pt p = \{3.4, 2.1\}:
cout << p << endl;
return 0:
```

4 Graph

4.1 Flow

4.1.1 Dinic

```
#include<bits/stdc++.h>
using namespace std;
const int MAXN = 2002: //XXX
//Set to the number of nodes in the flow graph.
const int MAXE = 2100012; //XXX
//Number of edges in the flow graph.
int from [MAXE], to [MAXE], cap[MAXE], prv[MAXE], head[MAXN].
    pt[MAXN], ec;
void addEdge(int u, int v, int uv, int vu = 0){
from[ec] = u, to[ec] = v, cap[ec] = uv, prv[ec] = head[u],
     head[u] = ec++:
from[ec] = v, to[ec] = u, cap[ec] = vu, prv[ec] = head[v],
     head[v] = ec++:
int lv[MAXN], q[MAXN];
bool bfs(int source, int sink){
memset(lv, 63, sizeof(lv));
int h = 0, t = 0:
lv[source] = 0:
q[t++] = source;
while (t-h){
 int v = q[h++];
 for (int e = head[v]; "e; e = prv[e])
 if (cap[e] && lv[v] + 1 < lv[to[e]]){</pre>
  lv[to[e]] = lv[v] + 1;
   q[t++] = to[e]:
return lv[sink] < 1e8:
int dfs(int v, int sink, int f = 1e9){
if (v == sink || f == 0)
 return f:
int ret = 0:
for (int &e = pt[v]; ~e; e = prv[e])
 if (lv[v]+1 == lv[to[e]]){
 int x = dfs(to[e], sink, min(f, cap[e]));
  cap[e] -= x;
```

```
cap[e^1] += x;
  ret += x;
  ret += x;
  f -= x;
  if (!f)
    break;
}
return ret;
}
int dinic(int source, int sink){
  int ret = 0;
  while (bfs(source, sink)){
    memcpy(pt, head, sizeof(head));
    ret += dfs(source, sink);
}
return ret;
}
int main(){
  memset(head, -1, sizeof(head));
  return 0;
}
```

4.1.2 MinCostMaxFlow

```
#include<bits/stdc++.h>
using namespace std;
typedef long long 11;
typedef pair<int, int> pii;
#define F first
#define S second
const int MAXN = 1e3 + 10;
const int MAXE = 1e4 + 10:
int from[MAXE], to[MAXE], cap[MAXE], ww[MAXE], prv[MAXE],
    head[MAXN], ec:
//edge from u to v, capacity uv, cost w, vu is backwards
    capacity(not necessary)
void addEdge(int u, int v, int uv, int w, int vu = 0){
from[ec] = u, to[ec] = v, cap[ec] = uv, ww[ec] = w, prv[ec]
      = head[u], head[u] = ec++;
from[ec] = v, to[ec] = u, cap[ec] = vu, ww[ec] = -w, prv[ec
     ] = head[v], head[v] = ec++:
```

```
int lv[MAXN], q[MAXN], rec[MAXN], parE[MAXN];
bool mark[MAXN];
pii spfa(int source, int sink) {
memset(lv, 63, sizeof(lv));
memset(parE, -1, sizeof(parE));
memset(rec, 0, sizeof(rec));
memset(mark, 0, sizeof(mark));
int h = 0, t = 0:
lv[source] = 0;
rec[source] = 1e9;
a[t++] = source:
mark[source] = true;
while (t^h){
 int v = q[h++];
 if (h == MAXN) h = 0;
 mark[v] = false:
 for (int e = head[v]; ~e; e = prv[e])
  if (cap[e] && lv[v] + ww[e] < lv[to[e]]){</pre>
   lv[to[e]] = lv[v] + ww[e]:
   parE[to[e]] = e;
   rec[to[e]] = min(cap[e], rec[v]);
   if (!mark[to[e]]){
    q[t++] = to[e];
    if (t == MAXN) t = 0;
    mark[to[e]] = true;
  }
if (lv[sink] > 1e8) return {0, 0}:
int v = sink;
int totW = 0:
while (~parE[v]){
 int e = parE[v];
 cap[e] -= rec[sink]:
 cap[e^1] += rec[sink];
 totW += ww[e] * rec[sink]:
 v = from[e];
return {rec[sink]. totW}:
//returns (flow, cost)
pii getMCMF(int source, int sink){
pii ret(0, 0), x;
while (true){
 x = spfa(source, sink);
 if (!x.F) break:
 ret.F += x.F:
```

```
ret.S += x.S;
}
return ret;
}
int main(){
  //only need to initialize these
memset(head, -1, sizeof(head));
ec = 0;
return 0;
}
```

5 Math

5.1 Miller-Rabin

```
void factor(ll x, ll& e, ll& k){
   while(x%2LL==OLL){
       x/=2LL;
       ++e:
   }
   k = x:
//increase x for higher certainty, 5 works well
bool is_prime(ll n, int x){
   if(n&2LL==0 || n==1LL)
       return false:
   if(n==2 || n==3 || n==5 || n==7)
      return true:
   ll e. k:
   factor(n-1,e,k);
   while (x-->0) {
      11 a = (rand())\%(n-5LL) + 2LL;
      ll p = mod_exp(a,k,n);
      if(p==1LL || p==n-1LL)
           continue;
       bool all fail = true:
       for(int i = 0; i < e-1; ++i){</pre>
          p = mod_exp(p, 2, n);
           if(p==n-1LL){
              all_fail = false;
              break;
          }
      }
       if(all fail)
           return false:
```

```
return true;
```

5.2 fft.

```
/* emaxx implementation */
/* Multiplication with arbitrary modulos
* use ntt if mod is prime and can be written as 2**k * c
    if not, use Chinese Reminder Theorem
     or transform A(x) = A1(x) + A2(x)*c decompose into A(x)
     /c and A(x)%c
                 B(x) = B1(x) + B2(x)*c
         where c ~= sqrt(mod)
        A * B = A1*B1 + c*(A1*B2 + A2*B1) * c**2(A2*B2)
        with all values < sqrt(mod) subpolynomials have
     coefficientes < mod * N after fft multiply decreasing
     changes of rounding error
const double PI=acos(-1):
typedef complex<double> base;
void fft (vector<base> & a, bool invert) {
int n=(int) a.size():
for (int i=1, j=0; i<n; ++i) {</pre>
 int bit=n>>1:
 for (; j>=bit; bit>>=1)
 j-=bit;
 i+=bit:
 if(i<i)
  swap(a[i],a[j]);
for (int len=2; len<=n; len<<=1) {</pre>
 double ang = 2*PI/len * (invert ? -1 : 1);
 base wlen(cos(ang), sin(ang));
 for (int i=0; i<n; i+=len) {</pre>
  base w(1):
  for (int j=0; j<len/2; ++j) {</pre>
   base u=a[i+j], v=a[i+j+len/2]*w;
   a[i+i]=u+v:
   a[i+j+len/2]=u-v;
   w*=wlen:
if (invert)
 for(int i=0;i<n;++i)</pre>
```

```
a[i]/=n:
// a, b => coefs to multiply, res => resulting coefs
// a[0], b[0], res[0] = coef x^0
// Doesnt work with negative coefs
void multiply (const vector<int> & a, const vector<int> & b,
      vector<int> & res) {
 vector<base> fa (a.begin(), a.end()), fb (b.begin(), b.end
      ()):
 size t n=1:
 while (n<max(a.size(),b.size())) n<<=1:</pre>
 n<<=1:
 fa.resize(n),fb.resize(n);
 fft (fa,false); fft(fb,false);
 for (size_t i=0; i<n; ++i)</pre>
 fa[i]*=fb[i]:
 fft (fa, true);
 res.resize (n):
 // avoid precision errors, mess up with negative values of
      coefs
 for(size_t i=0; i<n; ++i)</pre>
 res[i]=int(fa[i].real() + 0.5);
```

```
5.3 math
Picks theorem
Given a certain lattice polygon with non-zero area.
We denote its area by S. the number of points with integer
    coordinates lying strictly inside the polygon by I and
    the number of points lying on polygon sides by B.
Then, the Pick's formula states:
S=I+B/2 1
Burnsides Lemma
Let G be the finite group of operations we can perform on X
The number of orbits of X is the average of the number of
    fixed points for each g in G
G must be closed, associative, have identity, and inverses
A fixed point wrt g is an element of X such that g.x = x,
    That is, x is unchanged by the group operation.
For an element x, the orbit G.x is the set of all possible
    results of transforming x. Orbits partition X.
```

```
Example: color a square under rotation
G has 4 elements: rotate by x clockwise
X contains 16 colorings if rotations are distinct
Count number of fixed points = 16+2+2+4=24
Number of distinct colorings = 24/4=6
Win if xor of pile sizes is not zero
Sprague-Grundy theorem
Let's consider a state v of a two-player impartial game and
    let vi be the states reachable from it (where i{1,2,,k
    },k0). To this state, we can assign a fully equivalent
    game of Nim with one pile of size x. The number x is
    called the Grundy value or nim-value of state v.
Moreover, this number can be found in the following
    recursive wav:
x=mex \{x1, ,xk\}, where xi is the Grundy value for state vi
    and the function mex (minimum excludant) is the
    smallest non-negative integer not found in the given
Viewing the game as a graph, we can gradually calculate the
    Grundy values starting from vertices without outgoing
    edges. Grundy value being equal to zero means a state
    is losing.
Number Theory
Totient function is the number of positive integers no more
    than n which is coprime with n.
Formula is n * (product over p|n)(1 - 1/p)
Sum (d|n) totient(n) = n
Mobius function
What is mobius function? This function has lots of
    definitions.
However, the main definition is the following.
\mu (n) is 1 if n=1 or n is square-free and has even number
    of prime divisors.
\mu (n) is -1 if n is square-free and has odd number of
    prime divisors.
\mu (n) is 0 if n is not square-free.
```

\mu (n) also has a lot of interesting properties that make \

mu (n) so important.

```
\sum {d|n} \mu(d) = 0 \text{ for all } n>1. \text{ and } mu(n) \text{ is}
    multiplicative.
(A function f is multiplicative if f(mn)=f(m)f(n) for all (m
Enumberating submasks in O(3^n)
for (int m=0; m<(1<<n); ++m)
   for (int s=m: s: s=(s-1)\&m)
... s and m ...
Catalan numbers
The first few numbers Catalan numbers, Cn (starting from
1.1,2,5,14,42,132,429,1430,
C_n = (2n \text{ choose } n)/(n+1)
The Catalan number Cn is the solution for:
Number of correct bracket sequence consisting of n opening
    and n closing brackets.
The number of rooted full binary trees with n+1 leaves (
    vertices are not numbered). A rooted binary tree is
    full if every vertex has either two children or no
    children.
The number of ways to completely parenthesize n+1 factors.
The number of triangulations of a convex polygon with n+2
    sides (i.e. the number of partitions of polygon into
    disjoint triangles by using the diagonals).
The number of ways to connect the 2n points on a circle to
    form n disjoint chords.
The number of non-isomorphic full binary trees with n
    internal nodes (i.e. nodes having at least one son).
The number of monotonic lattice paths from point (0,0) to
    point (n,n) in a square lattice of size nn, which do
    not pass above the main diagonal (i.e. connecting (0,0)
     to (n,n)).
Number of permutations of length n that can be stack sorted
     (i.e. it can be shown that the rearrangement is stack
    sorted if and only if there is no such index i<j<k,
    such that ak<ai<ai).
The number of non-crossing partitions of a set of n elements
The number of ways to cover the ladder 1n using n rectangles
     (The ladder consists of n columns, where ith column
    has a height i).
```

5.4 numbertheory

//Inverse of a mod m by extended euclidean

```
int inv(int a. int m)
    int temp=m, q, t, u=0, v=1;
    if(m==1) return 0:
    while(a>1)
       q=a/m;
       t=m;
       m=a%m;
       a=t;
       t=u:
        u=v-a*u:
        v=t;
    if(v<0) v+=temp;
    return v;
//modular inverse of all numbers in [1,n] in O(n).
int inv[111111], n:
ll mod=1e9+7;
int main(void)
ſ
    cin>>n:
    int i;
    inv[1]=1;
    for(i=2 : i<=n : i++)</pre>
       inv[i]=((mod-mod/i)*inv[mod%i])%mod;
        cout<<inv[i]<<endl:
}
//Totient function sieve
void preprocess(void)
    int i, j;
    eulerphi[1]=1;
    for(i=2 : i<=122000 : i++)
        eulerphi[i]=i;
       primechk[i]=1:
    for(i=2 ; i<=122000 ; i++)</pre>
       if(primechk[i]==1)
           eulerphi[i]-=eulerphi[i]/i;
           for(j=2 ; i*j<=122000 ; j++)</pre>
```

```
primechk[i*j]=0;
               eulerphi[i*j]-=eulerphi[i*j]/i;
       }
   }
//Mobius function sieve
void preprocess(void)
   int i. i:
   for(i=1 ; i<=111100 ; i++)</pre>
       mu[i]=1:
       primechk[i]=1;
   primechk[1]=0;
   for(i=2 : i<=111100 : i++)</pre>
       if(primechk[i]==1)
           mu[i]=-mu[i];
           for(j=2 ; i*j<=111100 ; j++)</pre>
               primechk[i*i]=0;
               if(j%i==0)
               {
                   mu[i*j]=0;
               }
               else
                   mu[i*j]=-mu[i*i];
           }
   }
```

5.5 simplex

```
// Two-phase simplex algorithm for solving linear programs
    of the form
//
// maximize c^T x
// subject to Ax <= b
// x >= 0
//
// INPUT: A -- an m x n matrix
```

```
//
        b -- an m-dimensional vector
11
         c -- an n-dimensional vector
11
        x -- a vector where the optimal solution will be
    stored
//
// OUTPUT: value of the optimal solution (infinity if
11
          above, nan if infeasible)
11
// To use this code, create an LPSolver object with A, b,
    and c as
// arguments. Then, call Solve(x).
typedef long double DOUBLE;
tvpedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9:
struct LPSolver {
 int m. n:
 VI B, N;
 VVD D:
 LPSolver(const VVD &A, const VD &b, const VD &c) :
   m(b.size()), n(c.size()), N(n+1), B(m), D(m+2, VD(n+2)) {
   for (int i = 0; i < m; i++) for (int j = 0; j < n; j++) D
        [i][i] = A[i][i];
   for (int i = 0; i < m; i++) { B[i] = n+i; D[i][n] = -1; D
        [i][n+1] = b[i]; }
   for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c[j];
   N[n] = -1; D[m+1][n] = 1;
 void Pivot(int r, int s) {
   DOUBLE inv = 1.0 / D[r][s];
   for (int i = 0; i < m+2; i++) if (i != r)
     for (int j = 0; j < n+2; j++) if (j != s)
       D[i][i] -= D[r][i] * D[i][s] * inv:
   for (int j = 0; j < n+2; j++) if (j != s) D[r][j] *= inv;</pre>
   for (int i = 0: i < m+2: i++) if (i != r) D[i][s] *= -inv
   D[r][s] = inv;
   swap(B[r], N[s]);
 bool Simplex(int phase) {
   int x = phase == 1 ? m+1 : m;
```

13

```
while (true) {
     int s = -1:
     for (int j = 0; j <= n; j++) {</pre>
       if (phase == 2 && N[j] == -1) continue;
       if (s == -1 || D[x][j] < D[x][s] || D[x][j] == D[x][s]
            1 \&\& N[i] < N[s]) s = i:
     }
     if (s < 0 || D[x][s] > -EPS) return true;
      int r = -1:
     for (int i = 0; i < m; i++) {</pre>
       if (D[i][s] < EPS) continue:
       if (r == -1 || D[i][n+1] / D[i][s] < D[r][n+1] / D[r</pre>
            ll [s][
           D[i][n+1] / D[i][s] == D[r][n+1] / D[r][s] && B[i]
                1 < B[r]) r = i:
     }
     if (r == -1) return false:
     Pivot(r, s);
 }
  DOUBLE Solve(VD &x) {
    int r = 0;
    for (int i = 1; i < m; i++) if (D[i][n+1] < D[r][n+1]) r
    if (D[r][n+1] <= -EPS) {</pre>
     Pivot(r. n):
     if (!Simplex(1) || D[m+1][n+1] < -EPS) return -</pre>
          numeric_limits<DOUBLE>::infinity();
     for (int i = 0; i < m; i++) if (B[i] == -1) {
       int s = -1;
       for (int j = 0; j \le n; j++)
         if (s == -1 || D[i][j] < D[i][s] || D[i][j] == D[i</pre>
              [s] \&\& N[i] < N[s]) s = i;
       Pivot(i, s):
     }
    if (!Simplex(2)) return numeric_limits<DOUBLE>::infinity
         ();
    x = VD(n):
    for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] = D[i][</pre>
    return D[m][n+1]:
 }
};
```

5.6 tricks

```
11 fexp(ll a, int x, ll mod){ // Fast exponenciation
     returns a^x % mod
 if(x==0)return 111;
 if(x\%2==0){
 ll y=fexp(a, x/2, mod);
 return (y*y)%mod;
 return (a*fexp(a, x-1, mod))%mod;
11 divv(11 a, 11 b, 11 mod) { // Division with mod returns a/
 return (a*fexp(b, mod-2, mod))%mod;
11 f[N];
11 fat(11 a, 11 mod){ // Calculates factorial and stores in
     f % mod
 if(a<=1)return 1:</pre>
 return f[a]?f[a]:(f[a]=(a*fat(a-1, mod))%mod);
ll choose(ll n, ll k, ll mod){ // Returns n choose k % mod
 return divv(fat(n, mod), (fat(k, mod)*fat(n-k, mod))%mod,
     mod)%mod:
11 gcd(ll a, ll b){ // Greatest common divisor
 return b?gcd(b, a%b):a;
11 lcm(ll a, ll b){ // Least common multiple
 return (a*b)/gcd(a, b);
/* Fast factorization */
int p[N];
void start_fast(int MAX){ // Runs O(nlog(n)) Needs to be
     called to use fast fact or ammount of divisors.
 for(int i=2; i<=MAX; i++){</pre>
 if(p[i]==0){
  for(int j=i; j<=MAX; j+=i){</pre>
   p[j]=i;
  }
 }
```

```
vector<int>fast_fact(int x){ // Fast factorization in O(log2
vector<int>ret;
while(x>1){
 ret.pb(p[x]);
 x/=p[x];
return ret;
int amount_of_divisors(int x){ // Calculate the ammount of
    divisors of a number in O(log2(x)) assume already ran
    start fast.
if(x==1)return 1;
vector<int>v=fast fact(x):
int ret=1, curr=2;
for(int i=1: i<v.size(): i++){</pre>
 if(v[i]==v[i-1])curr++;
 else{
  ret*=curr:
  curr=2;
return ret*curr;
```

6 Misc

6.1 OrderedSet

```
OrderedSet st;
st.insert({1, 22});
st.insert({1, 33});
st.insert({1, 44});
st.insert({1, 55});
cout << st.order_of_key({1,33}) << endl;
cout << st.order_of_key({1,35}) << endl; // Where would it
be?
cout << (*st.find_by_order(2)).S << endl;
return 0;</pre>
```

6.2 template

```
#include <bits/stdc++.h>
using namespace std;

typedef long long ll;
typedef pair<int, int> pii;

#define F first
#define S second
#define se second
#define fi first
#define pb push_back
#define eb emplace_back
#define mk make_pair

#define N 1000007 //10e6 +7

int main(){
  ios::sync_with_stdio(false);
}
```

7 Python

7.1 InputArray

```
n = int(input())
a = list(map(int, input().split()));
ans = 1
v = []
for i in range(1, n):
    if(a[i] == 1):
```

```
ans += 1
v.append(a[i - 1])
v.append(a[n - 1])
print(len(v))
print(' '.join(map(str, v)));
```

8 Strings

8.1 kmp

```
/*
border = proper prefix that is suffix
p[i] = length of longest border of prefix of length i, s
     [0...i-1]
typedef long long 11;
typedef pair<int, int> ii;
const int INF = 0x3f3f3f3f;
const double PI = acos(-1.0):
const int N = 1e6 + 6;
int pi[N];
string p, t;
void pre () {
p += '#';
pi[0] = pi[1] = 0;
for (int i = 2; i <= (int)p.size(); i++) {</pre>
 pi[i] = pi[i-1];
 while (pi[i] > 0 \text{ and } p[pi[i]] != p[i-1])
  pi[i] = pi[pi[i]];
 if (p[pi[i]] == p[i-1])
  pi[i]++;
void report (int at) {
void KMP () {
pre ();
int k = 0;
```

```
int m = p.size() - 1;

for (int i = 0; i < (int)t.size(); i++) {
  while (k > 0 and p[k] != t[i])
  k = pi[k];

if (p[k] == t[i])
  k++;
  if (k == m)
   report (i - m + 1);
}

int main (void) {
  ios_base::sync_with_stdio(false);

return 0;
}
```

8.2 suffix-array

```
#include<bits/stdc++.h>
using namespace std;
typedef long long 11;
typedef pair<int, int> pii;
#define rank asdkflj
#define F first
#define S second
const int MAXN = 1e6 + 10;
const int LOG = 21;
rank[i][j] = position of lex smallest (not continuous)
     string of len 2<sup>i</sup> starting in position j
2<sup>i</sup> = len of string, j = starting position
sa[i] = index where it starts
i = ith smallest suffix
lc[i] = longest common prefix between i, i+1 smallest
Note: This is not the ith guy, but the ith smallest guy.
```

```
call lcp(i, i) to get the longest common prefix between the
       suffix starting
in i and the suffix starting in j
O(nlogn)
*/
// IMPORTANT: See if need to change to long long!
string s;
int rank[LOG][MAXN], n, lg;
pair<pair<int, int>, int> sec[MAXN];
int sa[MAXN]. lc[MAXN]:
int lcp(int a, int b){
int a = a:
for (int w = lg-1; ~w && max(a, b) < n; w--)</pre>
 if (max(a, b) + (1<<w) <= n && rank[w][a] == rank[w][b])</pre>
  a += 1 << w, b += 1 << w:
return a - _a;
}
int cnt[MAXN]:
pair<pii, int> gec[MAXN];
void srt() {
memset(cnt, 0, sizeof(cnt));
for (int i = 0; i < n; i++) cnt[sec[i].F.S+1]++;</pre>
 for (int i = 1; i < MAXN; i++) cnt[i] += cnt[i-1];</pre>
 for (int i = 0: i < n: i++) gec[--cnt[sec[i].F.S+1]] = sec[
     i];
 memset(cnt, 0, sizeof(cnt));
 for (int i = 0; i < n; i++) cnt[gec[i].F.F+1]++;</pre>
 for (int i = 1; i < MAXN; i++) cnt[i] += cnt[i-1];</pre>
 for (int i = n-1; ~i; i--) sec[--cnt[gec[i].F.F+1]] = gec[i
     1:
```

```
void build() {
n = s.size():
 int cur = 1: lg = 0:
 while (cur < n){</pre>
  lg++;
  cur <<= 1;
 lg++;
for (int i = 0; i < n; i++) rank[0][i] = s[i];</pre>
for (int w = 1; w < lg; w++){
 for (int i = 0; i < n; i++)</pre>
  if (i + (1 << w-1) >= n)
   sec[i] = \{\{rank[w-1][i], -1\}, i\};
   sec[i] = \{\{rank[w-1][i], rank[w-1][i+(1<< w-1)]\}, i\};
 //sort(sec, sec + n); This makes the code n*log^2, faster
      to code.
 srt(); // Comment this if using above code
 rank[w][sec[0].S] = 0;
 for (int i = 1; i < n; i++)</pre>
  if (sec[i].F == sec[i-1].F)
   rank[w][sec[i].S] = rank[w][sec[i-1].S];
   rank[w][sec[i].S] = i:
for (int i = 0; i < n; i++)</pre>
 sa[rank[lg-1][i]] = i:
for (int i = 0; i + 1 < n; i++)
 lc[i] = lcp(sa[i], sa[i+1]);
```

```
}
```

8.3 z

```
\{0, if i = 0\}
   z[i] = {length longest commom prefix of s and s[i...n-1]
typedef long long 11;
typedef pair<int, int> ii;
const int INF = 0x3f3f3f3f;
const double PI = acos(-1.0);
const int N = 2e5 + 5;
string s;
int z[N];
void go () {
int 1 = 0, r = 0;
int n = s.size();
memset (z. 0. sizeof z):
for (int i = 1: i < n: i++) {
 if (i <= r)
 z[i] = min (z[i-1], r - i + 1);
 while (z[i] + i < n \text{ and } s[z[i] + i] == s[z[i]])
 if (r < i + z[i] - 1) {
 1 = i:
 r = i + z[i] - 1;
}
```