

Fast On-Line Index Construction by Geometric Partitioning

Lester, Nicholas, Alistair Moffat, and Justin Zobel

Proceedings of the 14th ACM international conference on Information and knowledge management. ACM, 2005.

Hybrid Index Maintenance for Growing Text Collections

Büttcher, Stefan, Charles LA Clarke, and Brad Lushman

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Presented by Matt Chaney

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Inverted Indexes

- Similar to index in book
- Composed of two main parts
 - Vocabulary
 - Set of inverted lists, some with extra info
 - Frequencies
 - Locations
- Normally stored on disk as a single, contiguous partition

and	1:1		only	2:1	
aquarium	3:1		pigmented	4:1	
are	3:1	4:1	popular	3:1	
around	1:1		refer	2:1	
as	2:1		referred	2:1	
both	1:1		requiring	2:1	
bright	3:1		salt	1:1	4:1

Fig. 5.5. Inverted Index with Counts.
Yang, Christopher C. "Search engines information retrieval in practice."
(2010).

Problem

- Adding new documents challenging
- Two basic strategies
 - Off-line
 - Index creation simple process
 - Lag between document discovery and indexing
 - **On-line**
 - When delay cannot be tolerated
 - Index always queryable and up to date
 - Adds complexity to system
- Merge process
 - In-memory for new documents
 - On disk for existing index

Merge-Based Index Maintenance Algorithms

- Rebuild
 - Simple, rebuild on disk and switch
 - 10^{10} total items with 10^8 in-memory requires 17.1 trillion operations
- Remerge
 - Uses *bufferloads* to avoid redundant sorting
 - Rewrites on-disk, inserting in-memory postings during merge
 - 10^{10} total items with 10^8 in-memory require 0.8 trillion operations

Merge-Based Index Maintenance Algorithms

- In-place Update
 - Write postings when discovered
 - In-memory appended to on-disk during merge
 - Over-allocates space to avoid excessive file moves
 - 10^{10} total items with 10^8 in-memory require 0.3 trillion operations, *but...*
 - Metric is misleading
 - Memory overhead - 60% allocated and not used
 - Only counts *data movement* - Each term generates disc op → +1 trillion more ops
- Multiple partitions - “No Merge”
 - Don't merge, adds fragments somewhere on disk
 - Partition reference overhead

Geometric Partitioning

- Blends the Remerge and Multiple Partitions approach
 - Better construction time than Remerge
 - Better query processing time than Multiple Partitions
- Two approaches
 - Fixed partition size, r
 - Fixed number of partitions, p
- New documents in smaller partitions
- Old documents in larger partitions

Geometric Partitioning

- Controlled size, static r
 - Ex. partition 3 is 2x size of partition 2, which is 2x size of partition 1
 - Most recent documents go into smallest partition
- Leads to a hierarchical merging scheme

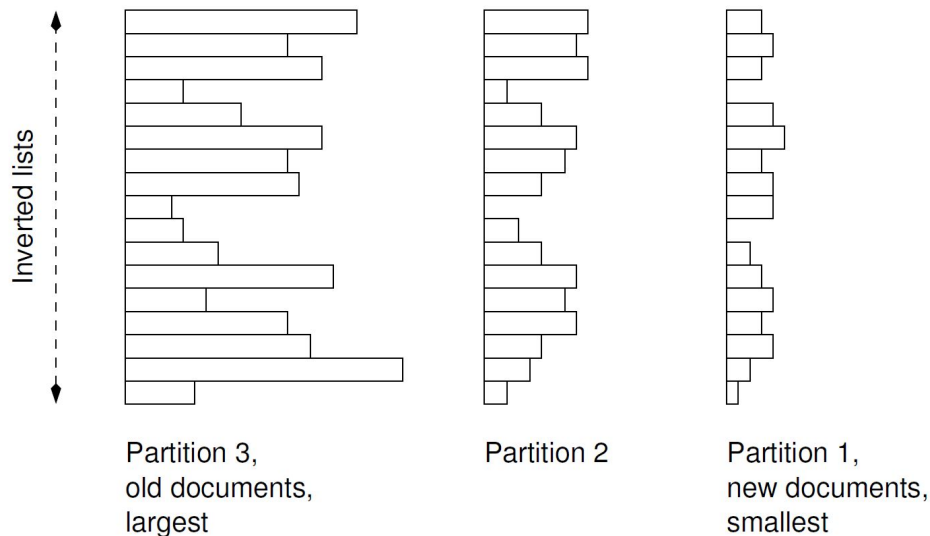


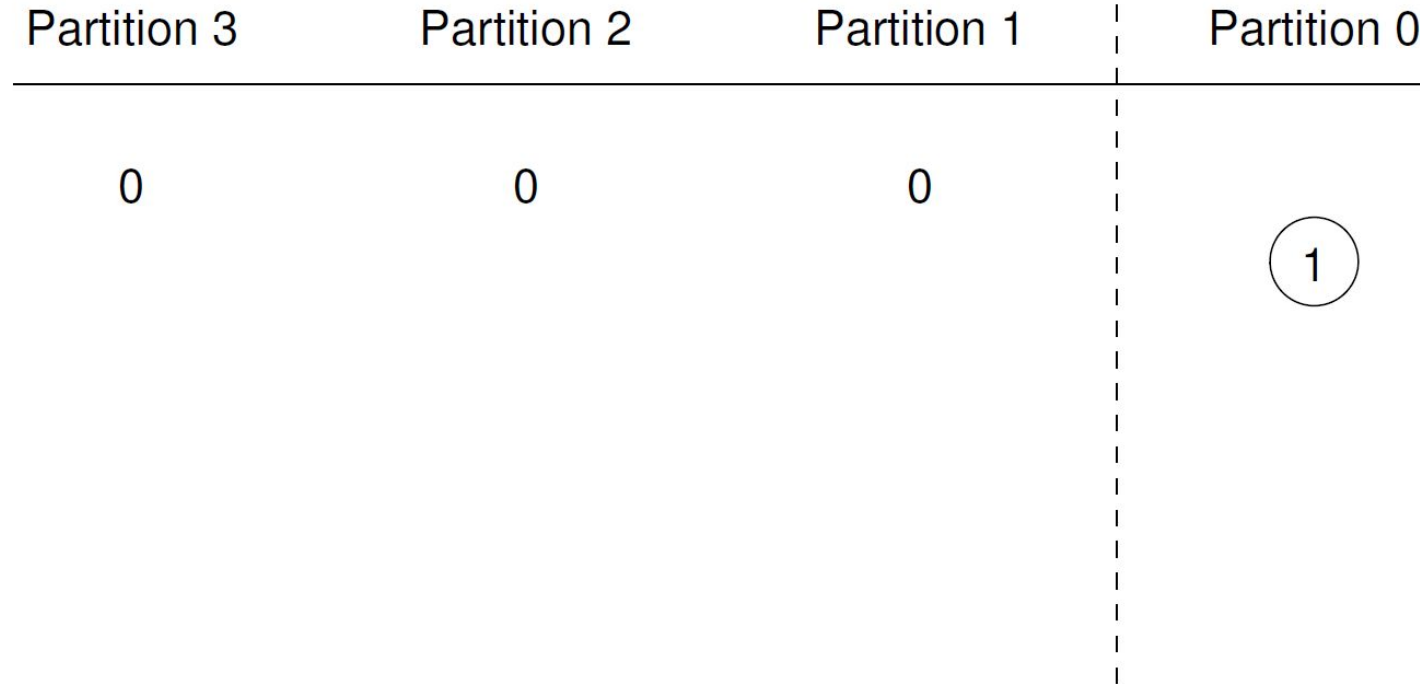
Fig. 1. Geometric Partitioning example.

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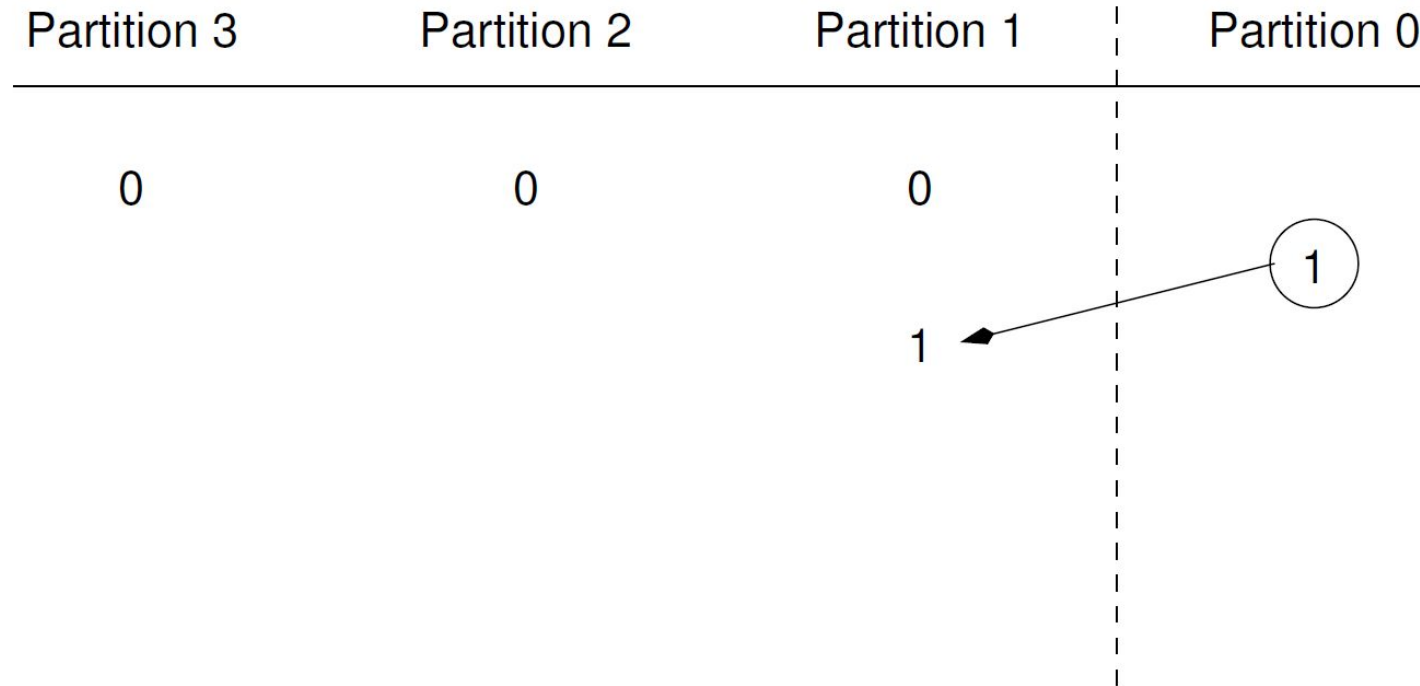
Hierarchical Merging, $r = 3$

Partition 3	Partition 2	Partition 1	Partition 0
0	0	0	

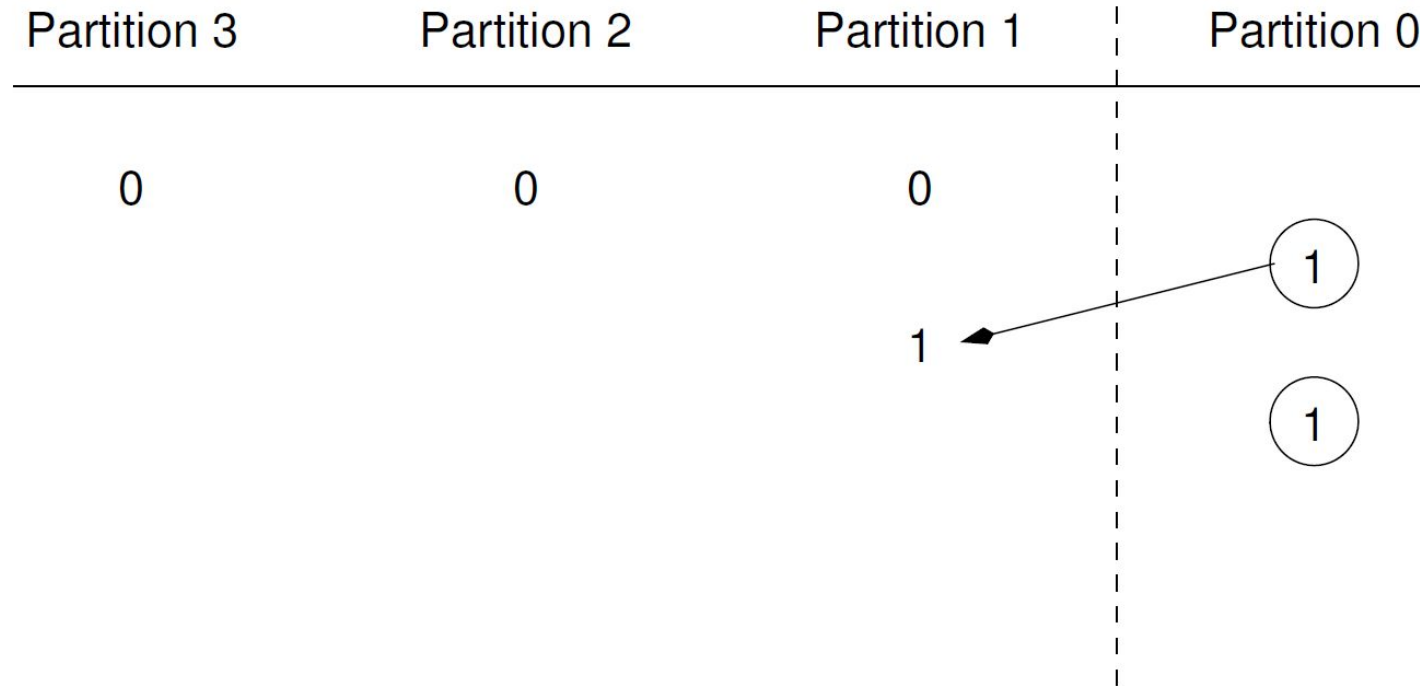
Hierarchical Merging, $r = 3$



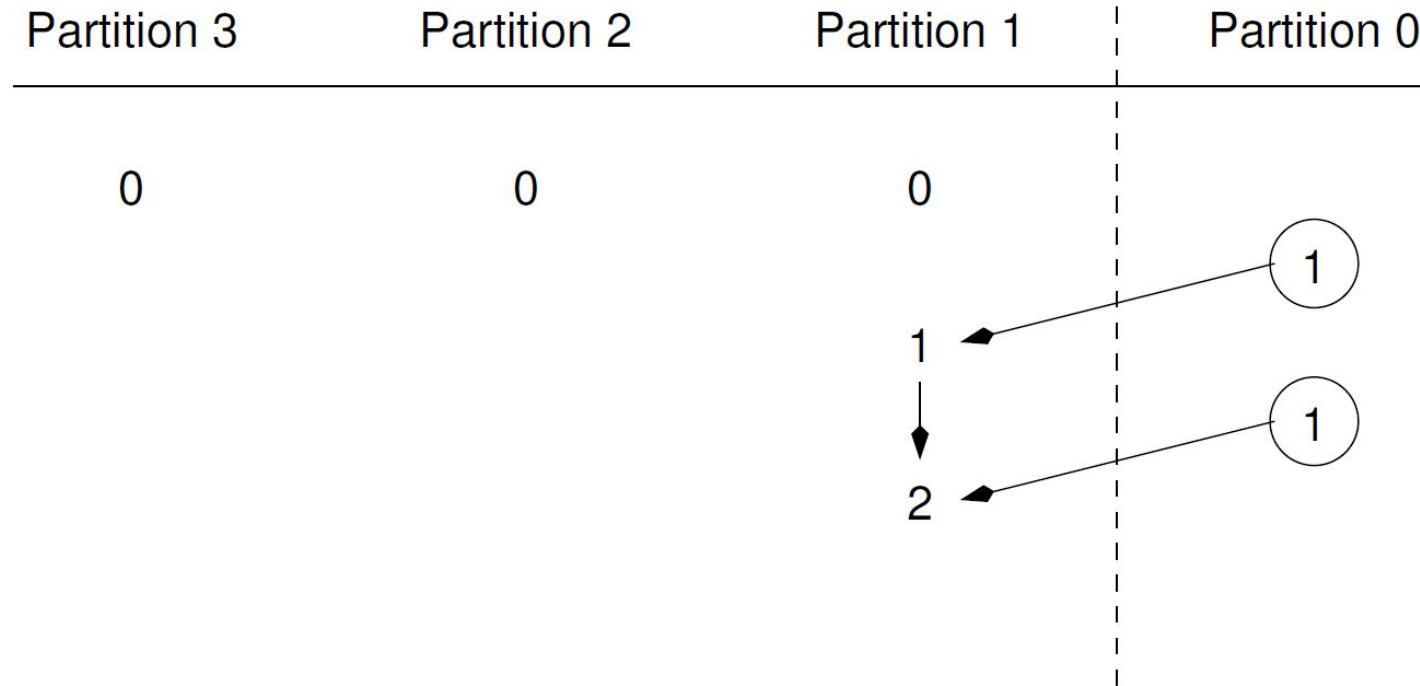
Hierarchical Merging, $r = 3$



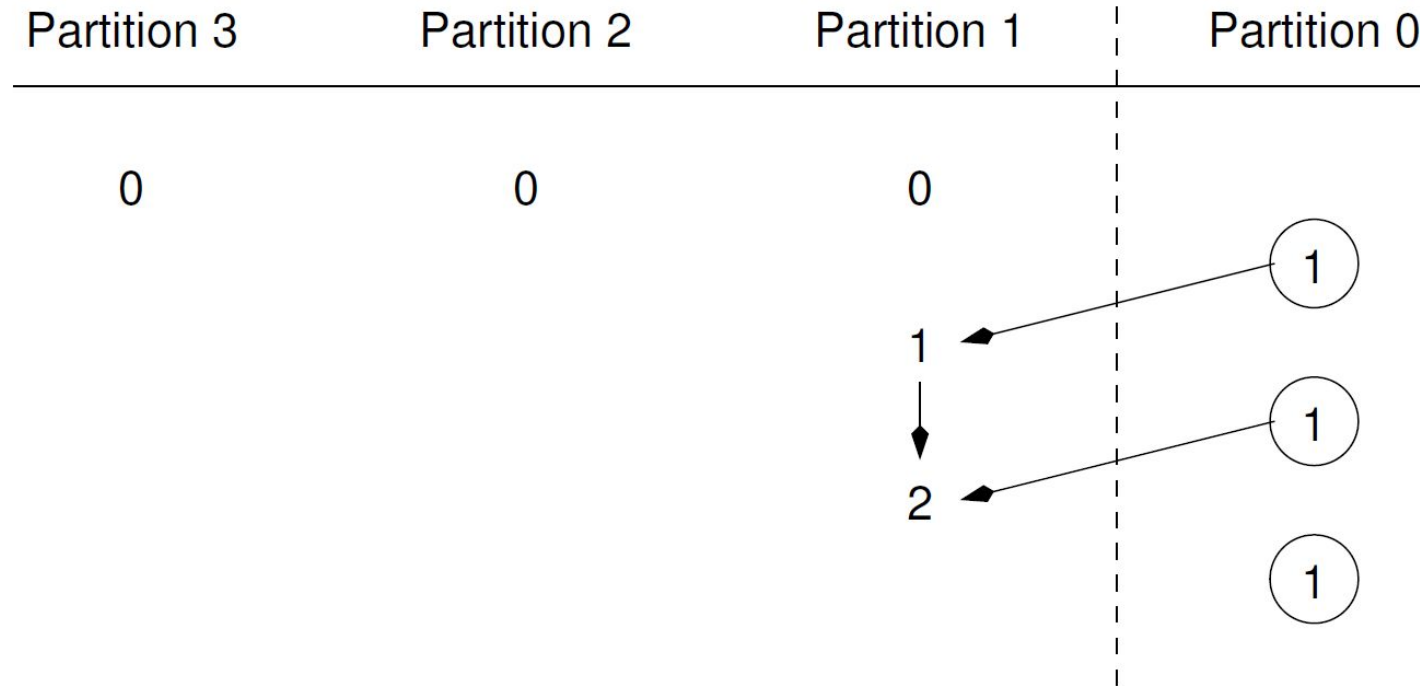
Hierarchical Merging, $r = 3$



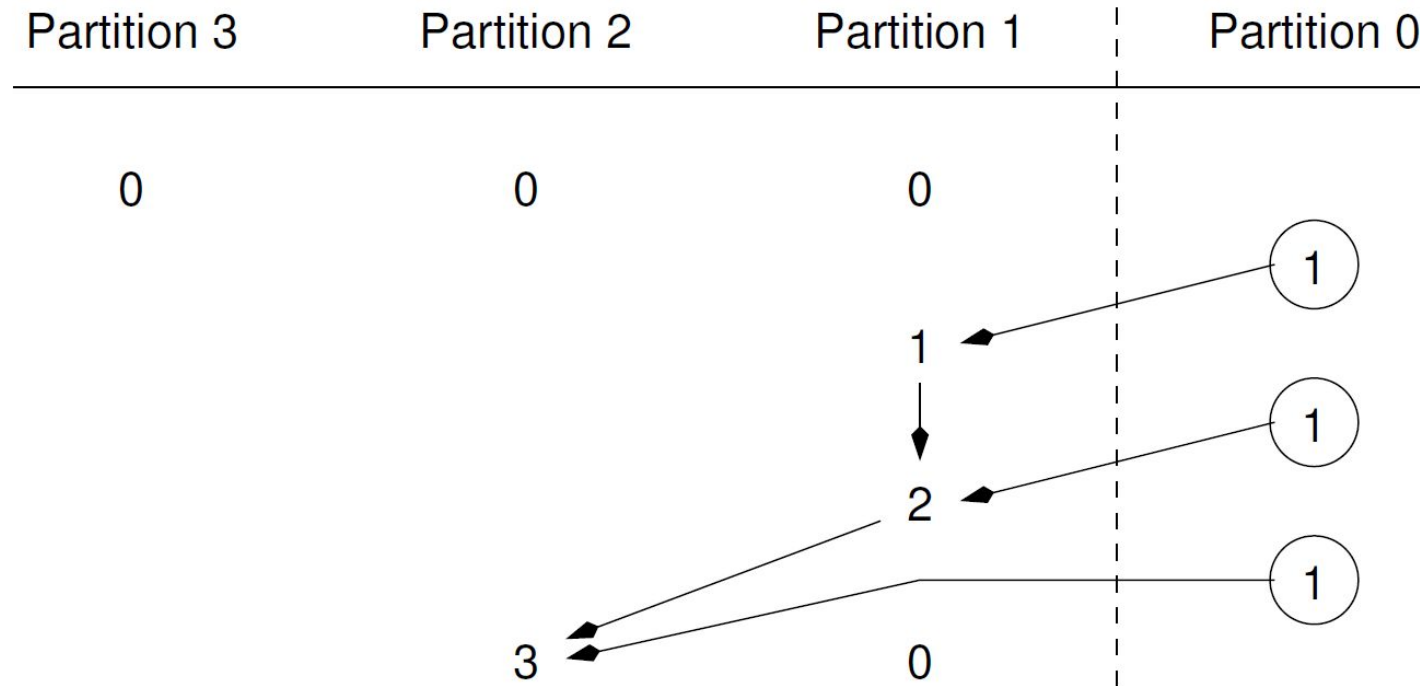
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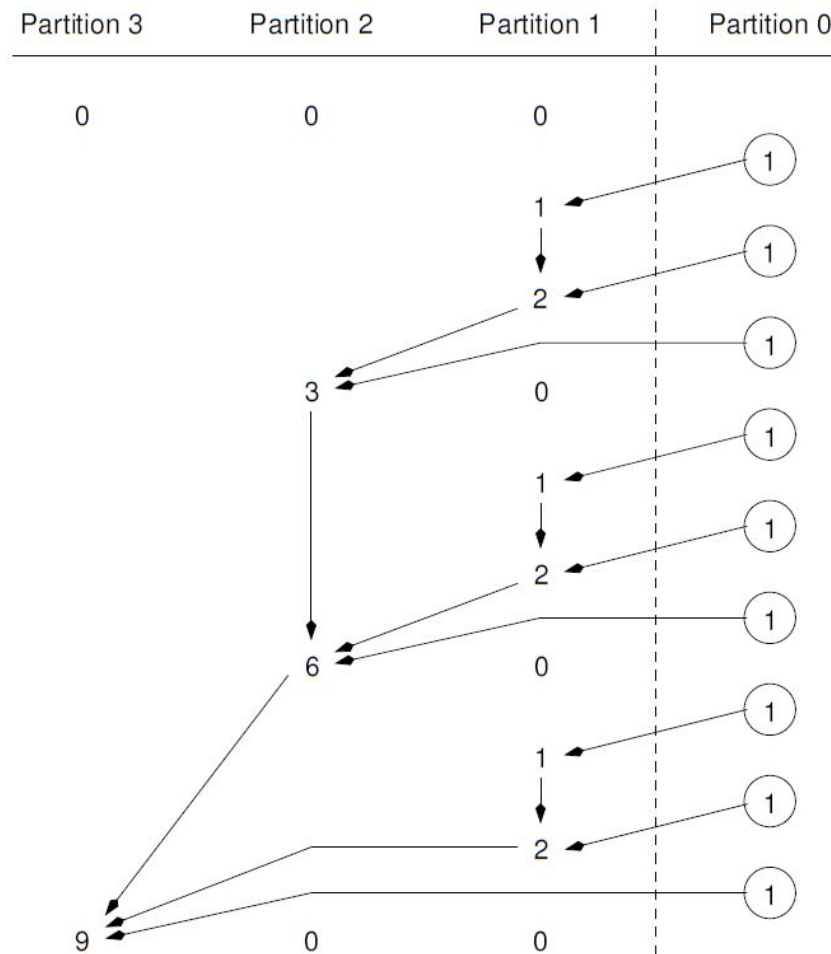


Hierarchical Merging, $r = 3$



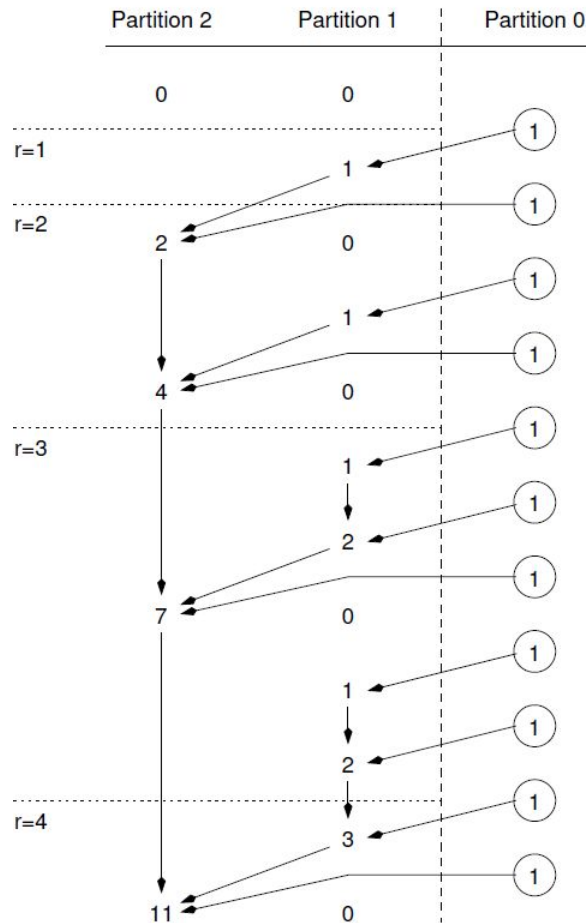
Merge with fixed r

Figure 2. Merge with fixed $r = 3$.
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Merge with fixed p

Figure 3. Merge with fixed $p = 3$.
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Analysis - Construction vs Access

r	Build cost	Access cost
2	0.29	3.6
3	0.31	3.1
4	0.32	2.9
6	0.34	2.6
8	0.35	2.4
12	0.39	2.2
16	0.42	2.0

Table 1. Construction cost versus access cost.

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Analysis - Index Construction Time

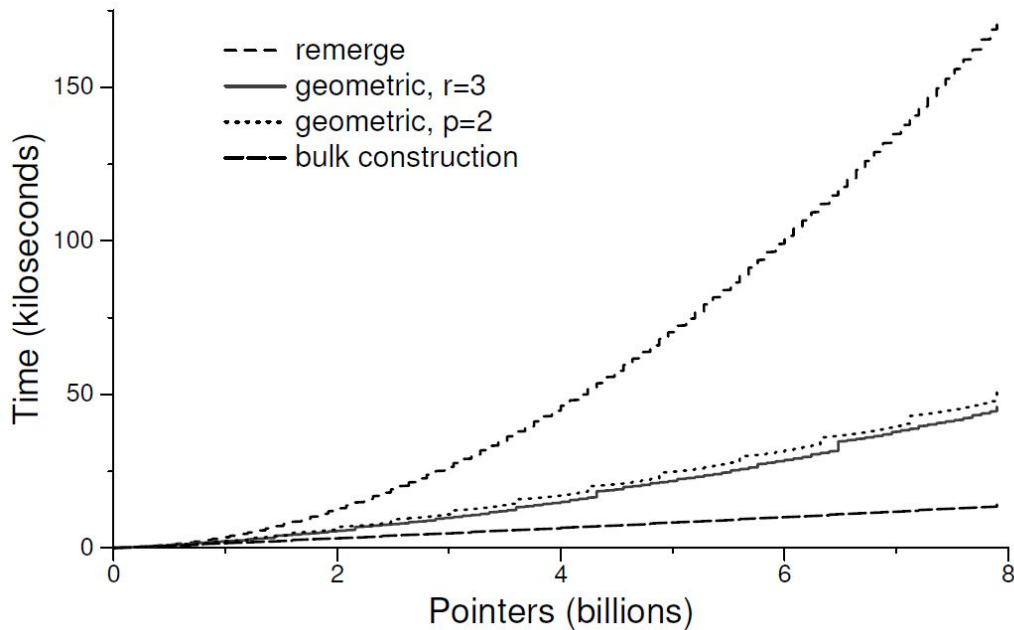


Figure 4. Build time vs number of pointers.

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Analysis - Fixed r vs Fixed p

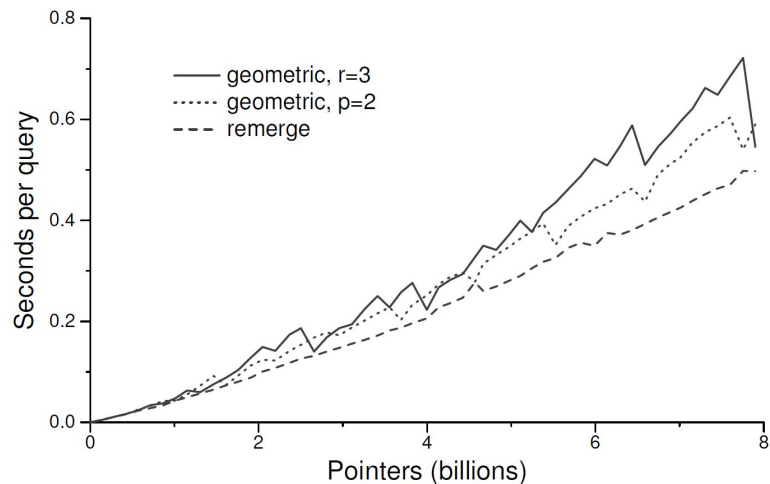
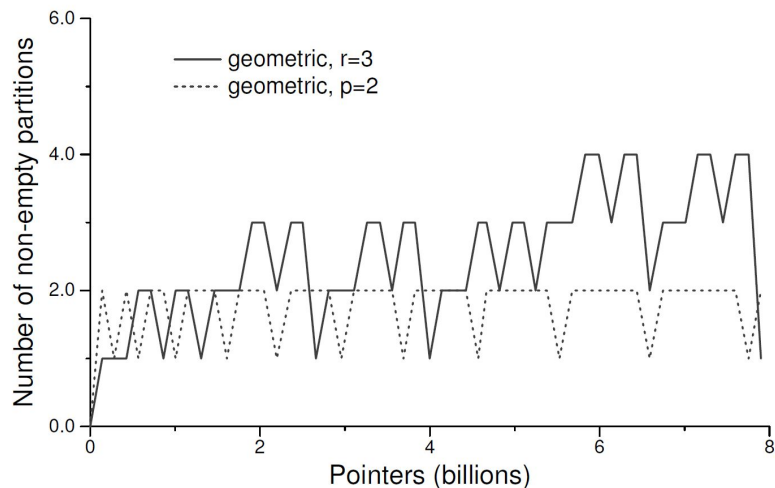


Figure 5. Query time versus number of pointers.

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Conclusions

- Index construction cost significantly reduced versus standard contiguous inverted indexes
- 100GB collection index construction 4x faster than Rmerge
- As index grows, speed gain increases
 - Adding last Gig of collection takes
 - 1 hour with Rmerge
 - 9 minutes with Geometric Partitioning
- Query time only 20% slower than Rmerge

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In-Place Index Maintenance

- In-memory data appended to on-disk lists without merging
- Generally requires relocating on-disk lists
 - Can be costly with large collections
- Over-allocating can reduce cost

Merge-Based Index Maintenance

- **Immediate Merge** - Rebuild/Remerge
 - Maintains one in-memory and one on-disk index
 - When in-memory full → merge with on-disk
 - Re-processes entire on-disk image
 - Still performs better than In-Place
- **Sqrt Merge**
 - Geometric Partitioning with fixed $p = 2$
- **Logarithmic Merge**
 - Geometric Partitioning with fixed r

Motivation

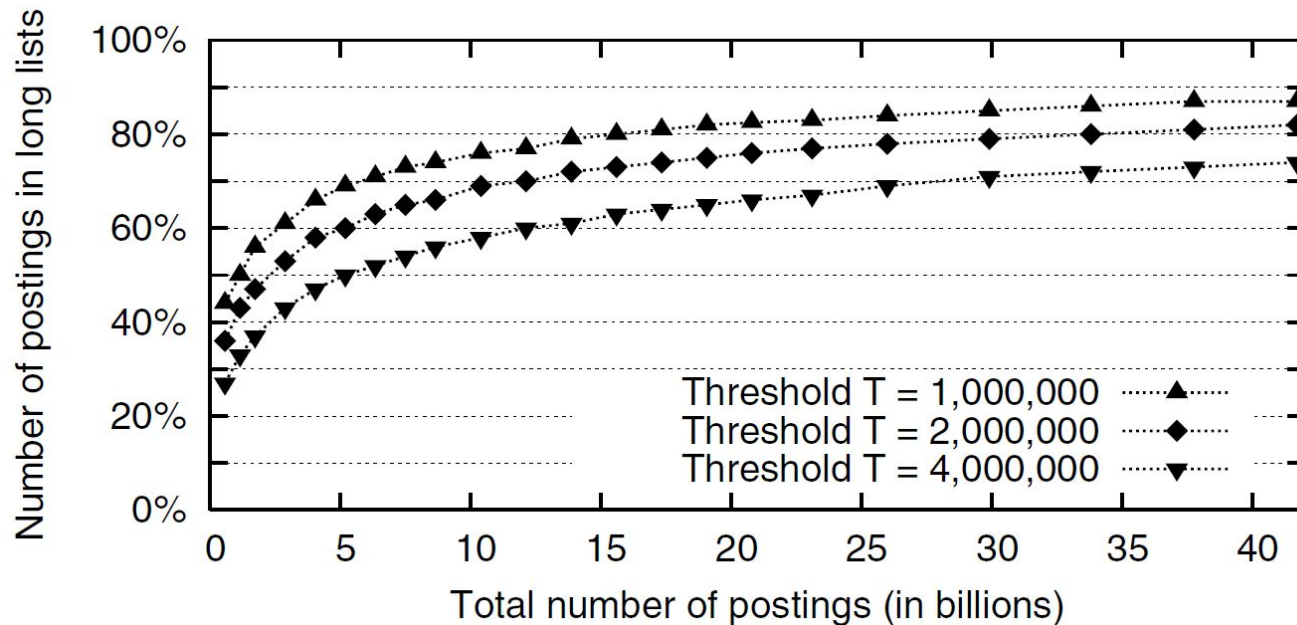


Figure 2. Total number of postings in the index vs. number of postings in long lists.

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First Hybrid Strategy

- Combines Merging and In-Place approaches
- Quantity T designates *short* or *long* term list length
 - Short lists managed with merge strategy
 - Long lists managed with in-place using filesystem, stored contiguously
- Avoids costly on-disk relocations of long posting lists
- Combined hybrid versions (“C” indicates contiguous lists)
 - Hybrid Immediate Merge $\rightarrow \text{HIM}_C$
 - Hybrid Sqrt Merge $\rightarrow \text{HSM}_C$
 - Hybrid Logarithmic Merge $\rightarrow \text{HLM}_C$

First Hybrid Strategy Flaws

- File system in-place update hard to analyze
 - File system operation differs system to system
 - Cannot be sure if posting lists are actually contiguous
- Static per-term T value inflexible for ever growing index
 - Highly sensitive to the amount of main memory available

New Hybrid Strategy

- Terms lists split into short and long, non-contiguous portions
 - Query time suffers, but disk reads can be amortized over longer read op
 - Relative query time increases only 5% over old hybrid
- New Hybrid (“NC” indicates non-contiguous lists)
 - Hybrid Immediate Merge $\rightarrow \text{HIM}_{\text{NC}}$
 - Hybrid Sqrt Merge $\rightarrow \text{HSM}_{\text{NC}}$
 - Hybrid Logarithmic Merge $\rightarrow \text{HLM}_{\text{NC}}$

New Hybrid Strategy Illustrated

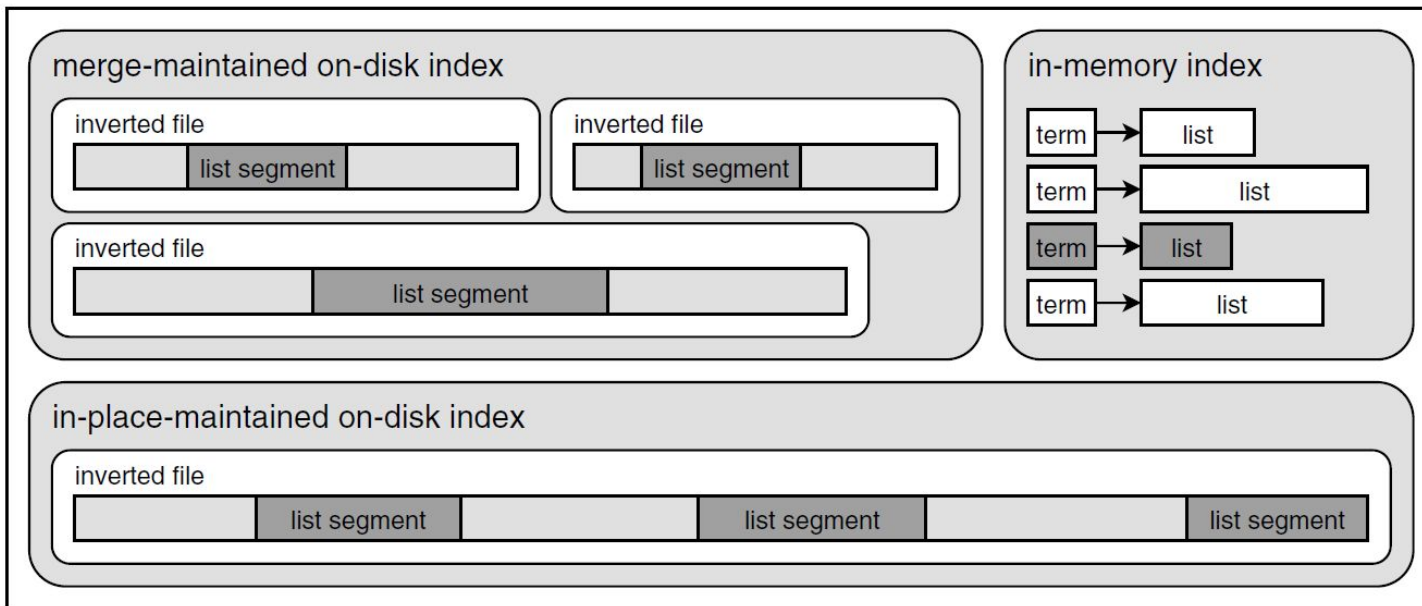


Figure 3. Index layout for a hybrid maintenance strategy with non-contiguous posting lists.

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Indexing Time and Query Processing Performance

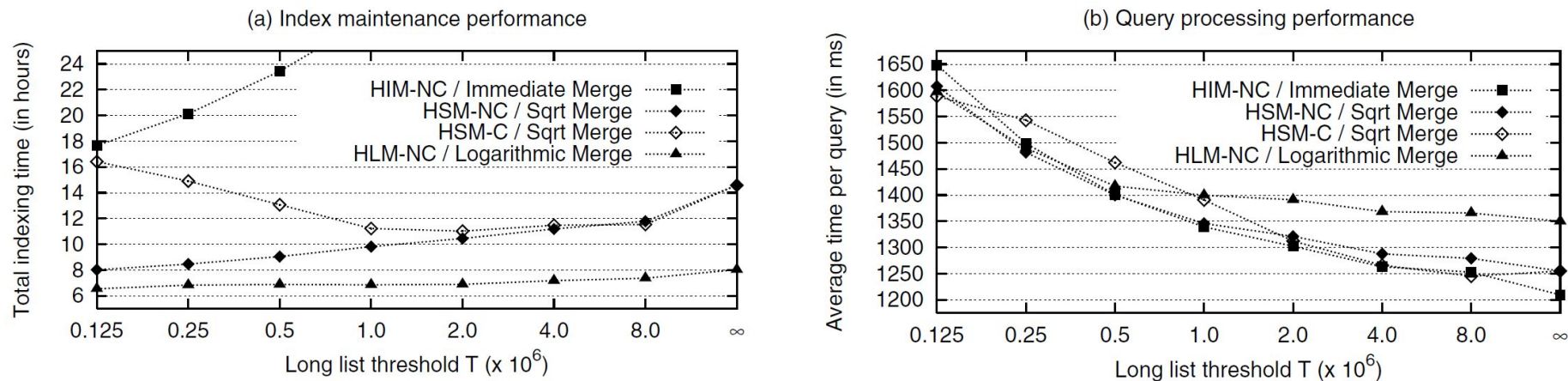


Figure 5. Index maintenance and query processing performance for different strategies with various parameter settings.

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Memory Size versus Maintenance Speed

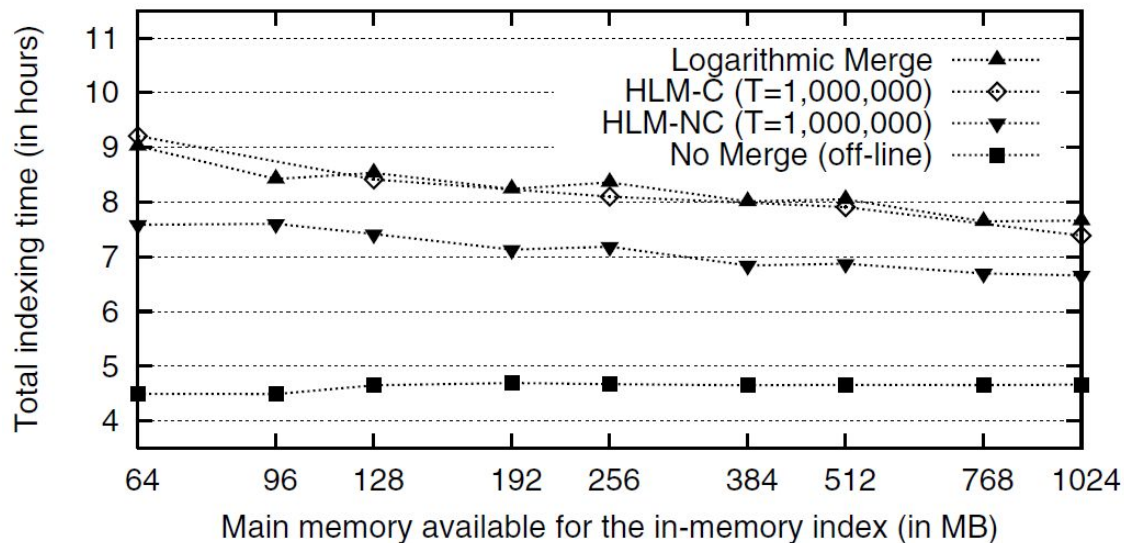


Figure 6. Impact of memory size on index maintenance performance.

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Conclusions

- New Hybrid approach outperforms previous index maintenance strategies including previous hybrid approach
- HLM_{NC} require only $\Theta(N)$ disk operations for text collection of size N
- Main shortfall is increased query times due to multiple partition fragments, future work could improve this metric with over-allocating disk space

Questions?