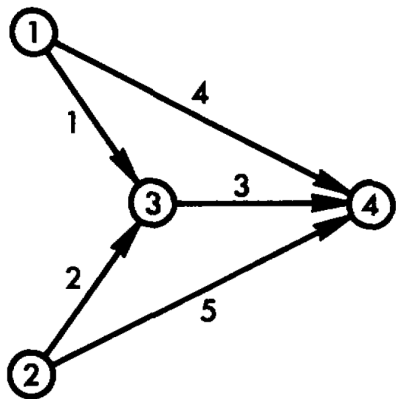


## Multiple origin-destination pairs

Best explained by a simple example — in generality, the indexing notation becomes quite unpleasant!!!



- ▶ Inspired by Sheffi p79
- ▶ No 'unmixing' (a simple case)

▶ Link cost functions:

$$c_1 = 2 + x_1$$

$$c_2 = 2 + x_2$$

$$c_3 = 1 + x_3$$

$$c_4 = 4 + x_4$$

$$c_5 = 4 + x_5$$

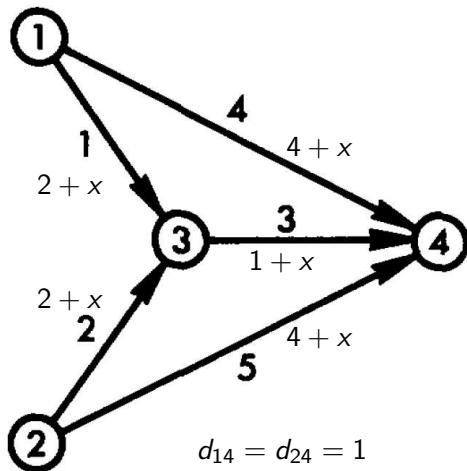
▶ OD demands:

$$d_{14} = 1$$

$$d_{24} = 1 \text{ (say)}$$

▶ Notice the top-down symmetry — to help hand calculations. But point is to formulate general problem (e.g., break symmetry in demand?) and give to Matlab.

## Example — preliminaries



► Find (here, all) routes (here, in links):

route 1:  $\{4\}$

route 2:  $\{1, 3\}$

route 3:  $\{5\}$

route 4:  $\{2, 3\}$

► Route variables  $\mathbf{y} \geq \mathbf{0}$  satisfy demands:

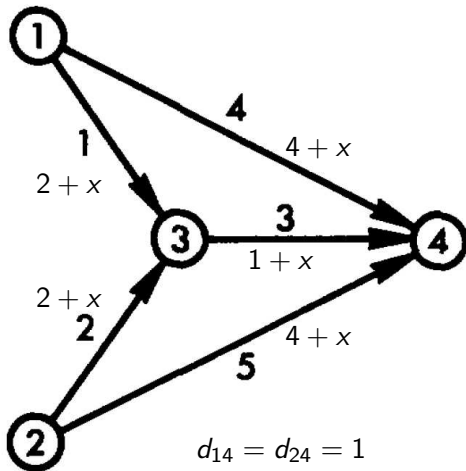
$$y_1 + y_2 = d_{14} \quad \text{and} \quad y_3 + y_4 = d_{24}.$$

► Find link-route incidence matrix:

$$\mathbf{A} = \begin{pmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

(NB unusual — more links than routes.)

## UE solution — complementarity method



R1: {4}. R2: {1, 3}. R3: {5}. R4: {2, 3}

► Compute link costs in route variables — formally  $\mathbf{c}(\mathbf{A}\mathbf{y})$  — then compute route costs in route variables, is  $\mathbf{A}^T \mathbf{c}(\mathbf{A}\mathbf{y})$ :

$$R1: 4 + y_1$$

$$R2: (2 + y_2) + (1 + y_2 + y_4) \\ = 3 + 2y_2 + y_4$$

$$R3: 4 + y_3$$

$$R4: 3 + y_2 + 2y_4$$

► (Complementarity) Solve:

$$4 + y_1 = 3 + 2y_2 + y_4$$

$$4 + y_3 = 3 + y_2 + 2y_4$$

with  $y_1 + y_2 = d_{14} = 1$  and  
 $y_3 + y_4 = d_{24} = 1$  and  $\mathbf{y} \geq \mathbf{0}$ .

► Solution  $y_1 = y_2 = y_3 = y_4 = 1/2$ .

(Symmetry helps a lot!!!)

Usually — need to check for unused routes.

## UE solution — in route variables, by Beckmann form

- ▶ Beckmann function:

$$\begin{aligned}\hat{f}(\mathbf{x}) &= \sum_{\text{links } i} \int_0^{x_i} c_i(\tilde{x}) d\tilde{x}, \\ &= \left(\frac{1}{2}x_1^2 + 2x_1\right) + \left(\frac{1}{2}x_2^2 + 2x_2\right) + \left(\frac{1}{2}x_3^2 + x_3\right) + \left(\frac{1}{2}x_4^2 + 4x_4\right) + \left(\frac{1}{2}x_5^2 + 4x_5\right), \\ &= \frac{1}{2}\mathbf{x}^T \mathbf{H} \mathbf{x} + \mathbf{a}^T \mathbf{x},\end{aligned}$$

where  $\mathbf{H} = \mathbf{I} = \text{diag}(\mathbf{b})$  and  $\mathbf{a} = (2, 2, 1, 4, 4)^T$ .

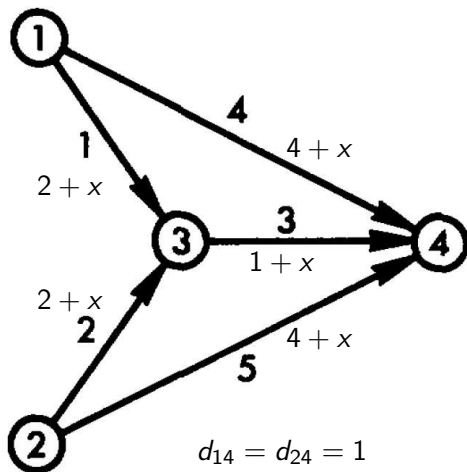
- ▶ Write  $\mathbf{x} = \mathbf{A}\mathbf{y}$ , and minimise subject to  $\mathbf{y} \geq \mathbf{0}$  and demand constraints  $y_1 + y_2 = d_{14}$  and  $y_3 + y_4 = d_{24}$ , or in vector form

$$\begin{pmatrix} 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \end{pmatrix} \mathbf{y} = \begin{pmatrix} d_{14} \\ d_{24} \end{pmatrix},$$

that is  $\mathbf{M}\mathbf{y} = \mathbf{d}$ , where  $\mathbf{M}$  is **OD-route incidence matrix**.

- ▶ SO solution approach is similar: but  $\mathbf{H} = 2 \text{diag}(\mathbf{b})$  (multiply instead of integrate).

## UE solution — in Beckmann form, link variables only



► All the route computations are skipped: saves a lot of work!!!

► Minimise

$$\frac{1}{2} \mathbf{x}^T \mathbf{H} \mathbf{x} + \mathbf{a}^T \mathbf{x}$$

subject to  $\mathbf{x} \geq \mathbf{0}$  and node balances  $\mathbf{B}\mathbf{x} = \mathbf{s}$ , in the form

$$\begin{pmatrix} -1 & 0 & 0 & -1 & 0 \\ 0 & -1 & 0 & 0 & -1 \\ +1 & +1 & -1 & 0 & 0 \\ 0 & 0 & +1 & +1 & +1 \end{pmatrix} \mathbf{x} = \begin{pmatrix} -d_{14} \\ -d_{24} \\ 0 \\ d_{14} + d_{24} \end{pmatrix}$$

► The only catch is for examples with flow 'unmixing' — not needed here.

## With a little help from Matlab ...

---

```
>> help quadprog
```

**quadprog** Quadratic programming.

$X = \text{quadprog}(H,f,A,b)$  attempts to solve the quadratic programming problem:

$$\begin{array}{ll}\min & 0.5*x'*H*x + f'*x \\ & x\end{array} \quad \text{subject to: } A*x \leq b$$

$X = \text{quadprog}(H,f,A,b,Aeq,beq)$  solves the problem above while additionally satisfying the equality constraints  $Aeq*x = beq$ . (Set  $A=[]$  and  $B=[]$  if no inequalities exist.)

$X = \text{quadprog}(H,f,A,b,Aeq,beq,LB,UB)$  defines a set of lower and upper bounds on the design variables,  $X$ , so that the solution is in the range  $LB \leq X \leq UB$ . Use empty matrices for  $LB$  and  $UB$  if no bounds exist. Set  $LB(i) = -\text{Inf}$  if  $X(i)$  is unbounded below; set  $UB(i) = \text{Inf}$  if  $X(i)$  is unbounded above.

$X = \text{quadprog}(H,f,A,b,Aeq,beq,LB,UB,X0)$  sets the starting point to  $X0$ .

$X = \text{quadprog}(H,f,A,b,Aeq,beq,LB,UB,X0,OPTIONS)$  minimizes with the default optimization parameters replaced by values in  $OPTIONS$ , an argument created with the `OPTIMOPTIONS` function. See `OPTIMOPTIONS` for details.