

## Option 3: Risk-aware balanced recursive partition into axis-aligned rectangles (2D)

**Goal.** Given calibration data

$$\mathcal{D} = \{(u_{1,k}, u_{2,k}, E_k)\}_{k=1}^n, \quad (u_{1,k}, u_{2,k}) \in [0, 1]^2, \quad E_k \in \{0, 1\},$$

construct  $B$  axis-aligned rectangles (leaf regions) with *near-uniform mass* (roughly  $n/B$  points per leaf), where splits are chosen to *separate error rates* between children.

**Node representation.** A node (region)  $\nu$  stores:

- an index set  $S_\nu \subseteq [n]$  of points currently in the node,
- a bounding box  $R_\nu = [\ell_{1,\nu}, r_{1,\nu}] \times [\ell_{2,\nu}, r_{2,\nu}] \subseteq [0, 1]^2$ ,
- the empirical error rate  $\hat{\eta}(\nu) \frac{1}{|S_\nu|} \sum_{k \in S_\nu} E_k$ .

The root node is  $\nu_0$  with  $S_{\nu_0} = [n]$  and  $R_{\nu_0} = [0, 1] \times [0, 1]$ .

**Size constraint (near-uniform mass).** Fix

$$n_{\min} \left\lfloor \frac{n}{B} \right\rfloor \quad (\text{e.g., } n = 4000, B = 30 \Rightarrow n_{\min} \approx 133).$$

Every split must produce children with at least  $n_{\min}$  points.

**Candidate thresholds.** Fix a small set of quantile levels

$$\mathcal{Q} = \{q_1, \dots, q_M\} \subset (0, 1) \quad (\text{e.g., } \{0.1, 0.2, \dots, 0.9\}).$$

For a node  $\nu$  and dimension  $d \in \{1, 2\}$ , define candidate thresholds

$$\tau \in \left\{ \text{Quantile}_q(\{u_{d,k} : k \in S_\nu\}) : q \in \mathcal{Q} \right\},$$

optionally discarding duplicates (ties).

**Split and score.** For node  $\nu$ , a candidate split  $(d, \tau)$  induces children:

$$S_L = \{k \in S_\nu : u_{d,k} \leq \tau\}, \quad S_R = S_\nu \setminus S_L,$$

with corresponding rectangles (updating only the split coordinate bounds):

$$R_L = R_\nu \cap \{u_d \leq \tau\}, \quad R_R = R_\nu \cap \{u_d > \tau\}.$$

The split is *admissible* if  $|S_L| \geq n_{\min}$  and  $|S_R| \geq n_{\min}$ .

Define empirical child error rates

$$\hat{\eta}_L = \frac{1}{|S_L|} \sum_{k \in S_L} E_k, \quad \hat{\eta}_R = \frac{1}{|S_R|} \sum_{k \in S_R} E_k.$$

A simple, effective split score is the absolute separation:

$$\Delta(d, \tau; \nu) |\hat{\eta}_L - \hat{\eta}_R|.$$

Optionally (to discourage very unbalanced splits even when admissible), use

$$\Delta'(d, \tau; \nu) \Delta(d, \tau; \nu) \cdot \sqrt{\frac{|S_L||S_R|}{|S_\nu|^2}}.$$

(Use either  $\Delta$  or  $\Delta'$  consistently.)

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**Algorithm 1** Risk-aware balanced rectangle partition (2D)

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**Require:** Data  $\mathcal{D} = \{(u_{1,k}, u_{2,k}, E_k)\}_{k=1}^n$ , target leaves  $B$ , quantiles  $\mathcal{Q}$ , minimum leaf size  $n_{\min}$ .  
**Ensure:** A binary tree; leaves define rectangles  $\{R_\nu\}_{\nu \in \mathcal{L}}$ , with associated  $S_\nu$  and  $\widehat{\eta}(\nu)$ .

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1: Initialize root leaf  $\nu_0$  with  $S_{\nu_0} = [n]$ ,  $R_{\nu_0} = [0, 1] \times [0, 1]$ .
2:  $\mathcal{L} \leftarrow \{\nu_0\}$ .
3: while  $|\mathcal{L}| < B$  do
4:    $bestLeaf \leftarrow \text{None}$ ;  $bestSplit \leftarrow \text{None}$ ;  $bestScore \leftarrow -\infty$ .
5:   for all leaves  $\nu \in \mathcal{L}$  do
6:     Compute the best admissible split  $(d^*, \tau^*)$  for  $\nu$ :
7:      $score^*(\nu) \leftarrow \max\{\Delta(d, \tau; \nu) : d \in \{1, 2\}, \tau \in \mathcal{T}_{d,\nu}, \text{admissible}\}$ 
    where  $\mathcal{T}_{d,\nu}$  are candidate thresholds from quantiles in  $\mathcal{Q}$ .
8:     if no admissible split exists for  $\nu$  then
9:       continue
10:      end if
11:      if  $score^*(\nu) > bestScore$  then
12:         $bestScore \leftarrow score^*(\nu)$ 
13:         $bestLeaf \leftarrow \nu$ 
14:         $bestSplit \leftarrow (d^*, \tau^*)$ 
15:      end if
16:    end for
17:    if  $bestLeaf$  is None then ▷ no admissible splits anywhere
18:      break
19:    end if
20:    Split  $bestLeaf$  using  $bestSplit$  into children  $\nu_L, \nu_R$ :
     $S_{\nu_L} = \{k \in S_{bestLeaf} : u_{d^*, k} \leq \tau^*\}, S_{\nu_R} = S_{bestLeaf} \setminus S_{\nu_L}$ 
     $R_{\nu_L} = R_{bestLeaf} \cap \{u_{d^*} \leq \tau^*\}, R_{\nu_R} = R_{bestLeaf} \cap \{u_{d^*} > \tau^*\}$ 
21:  end while

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**Tree growth strategy (exactly  $B$  leaves).** Maintain a set of current leaves  $\mathcal{L}$  (initially  $\{\nu_0\}$ ). Repeatedly split *one* current leaf at a time until  $|\mathcal{L}| = B$ . At each iteration, among all leaves, select the leaf whose *best admissible split* has the largest score.

**Output.** The final leaves  $\mathcal{L}$  define  $|\mathcal{L}| \leq B$  rectangles  $\{R_\nu\}$ . For each leaf, store:

$$R_\nu = [\ell_{1,\nu}, r_{1,\nu}) \times [\ell_{2,\nu}, r_{2,\nu}), \quad \widehat{\eta}(\nu) = \frac{1}{|S_\nu|} \sum_{k \in S_\nu} E_k, \quad |S_\nu|.$$

In downstream calibration/certification, treat  $Q(u_1, u_2) = \nu$  if  $(u_1, u_2) \in R_\nu$ .

**Inference-time assignment.** Given a new point  $(u_1, u_2)$ , traverse the tree from the root: at each internal node with split  $(d, \tau)$  go left if  $u_d \leq \tau$  and right otherwise, until a leaf  $\nu$  is reached. Return  $Q(u_1, u_2) = \nu$ .

**Implementation notes (important for correctness).**

- **Ties at thresholds.** Use the convention “left if  $u_d \leq \tau$ ” consistently in both training and inference.
- **Quantile computation.** Quantiles are empirical over  $\{u_{d,k} : k \in S_\nu\}$ ; remove duplicates to avoid empty children.
- **Exact  $B$  leaves.** The loop increases leaf count by +1 each split. If no admissible split exists before reaching  $B$ , you stop early.
- **Uniformity.** This enforces a *lower bound*  $n_{\min}$ ; if you want tighter near-uniformity, also enforce an upper bound  $|S_{\text{child}}| \leq n_{\max}$  (e.g.  $n_{\max} \approx \lceil 1.3n/B \rceil$ ) by discarding splits that create too-large leaves.