

Advanced Algorithms

Assignment 3: Minimum cut

June 13, 2022

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Contents

1 Introduction

For this assignment, we implemented and analyzed the performance of two algorithms for the min-cut problem for weighted graphs. The algorithms implemented are:

- 1. Stoer and Wagner's Deterministic Algorithm;
- 2. Karger and Stein's Randomized Algorithm.

- 2 Stoer and Wagner's Deterministic Algorithm
- 2.1 Data Structure
- 2.2 Implementation
- 2.3 Complexity

3 Karger and Stein's Randomized Algorithm

```
KARGER (G, k):
1
     A = +\infty
2
     for i = 1 to k:
         t = RECURSIVE CONTRACT(G)
            if t < min:
                   \min = t
      return min
     RECURSIVE CONTRACT (G=(D,W)):
     n= number of vertices in G
10
      if n < =6:
11
         Gp = CONTRACT(G, 2)
12
         return weight of the only edge (u,v) in Gp
13
      t = n/\sqrt{2}+1
14
     for i = 1 to 2:
15
         Gi = CONTRACT(G, t)
         wi = RECURSIVE CONTRACT(Gi)
17
      return min(w1,w2)
18
19
     CONTRACT(G=(D,W),k):
20
     n= number of vertices in G
21
     for i = 1 to n-k:
22
         (u, v) = EDGE SELECT(D, W)
23
         CONTRACT EDGE(u, v)
24
      return D,W
25
26
     CONTRACT EDGE(u,v):
27
     D[u] = D[u] + D[v] - 2W[u, v]
     D[v] = 0
29
     W[u, v] = W[v, u] = 0
30
     for each vertex w \in V: except u and v:
31
        W[u, v] = W[u, w] + W[v, w]
32
        W[w, u] = W[w, u] + W[w, v]
33
        W[v, w] = W[w, v] = 0
34
     EDGE SELECT(D,W)
36
      1. Choose u with probability proportional to D[u]
37
      2. Once u is fixed, choose v with probability proportional to W[u,v]
38
      3. return the edge (u, v)
```

This is a randomized algorithm for the computation of a graph. In the next subsections we explain how we have implemented the data structure and the functions of the algorithm.

3.1 Data Structure

For the implementation of the algorithm we used:

- **k**: is a constant (log^2n) used by Karger to repeat Recursive-Contract k times and to obtain an error with probability less or equal to 1/n;
- V: is the list of nodes;

- **D**: is the list of the sum of the weights of each node;
- W: is the list of the graph with 3 parameters(node,node,weight). Each node is connected with the others and if they are not connected the third parameter is set to 0 otherwise is set to the correct weight.

3.2 Implementation

3.3 Complexity

4 Results

- 4.1 Table with Min-Cut results
- 4.2 Graph of the Time Cost of the two Algorithms
- 4.3 Graph of the Time Cost compared to the Discovery Time of Karger and Stein Algorithm
- 4.4 Graph of the Time Cost compared to the Asymptotic Complexity of the two Algorithms

5 Conclusion