



# Advanced Algorithms

## Assignment 3: Minimum cut

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# 1 Introduction

For this assignment, we implemented and analyzed the performance of two algorithms for the min-cut problem for weighted graphs. The algorithms implemented are:

1. Stoer and Wagner's Deterministic Algorithm;
2. Karger and Stein's Randomized Algorithm.

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## 2 Stoer and Wagner's Deterministic Algorithm

### 2.1 Data Structure

### 2.2 Implementation

### 2.3 Complexity

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### 3 Karger and Stein's Randomized Algorithm

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```
1  KARGER (G,k) :
2  A =  $+\infty$ 
3  for i = 1 to k:
4      t = RECURSIVE_CONTRACT(G)
5      if t < min:
6          min = t
7  return min
8
9  RECURSIVE_CONTRACT(G=(D,W)) :
10 n= number of vertices in G
11 if n<=6:
12     Gp= CONTRACT(G,2)
13     return weight of the only edge (u,v) in Gp
14 t =  $n/\sqrt{2}+1$ 
15 for i = 1 to 2:
16     Gi = CONTRACT(G,t)
17     wi = RECURSIVE_CONTRACT(Gi)
18 return min(w1,w2)
19
20 CONTRACT(G=(D,W),k) :
21 n= number of vertices in G
22 for i = 1 to n-k:
23     (u,v) = EDGE_SELECT(D,W)
24     CONTRACT_EDGE(u,v)
25 return D,W
26
27 CONTRACT_EDGE(u,v) :
28 D[u] = D[u]+D[v]-2W[u,v]
29 D[v] = 0
30 W[u,v] = W[v,u] = 0
31 for each vertex w  $\in$  V: except u and v:
32     W[u,v] = W[u,w] + W[v,w]
33     W[w,u] = W[w,u] + W[w,v]
34     W[v,w] = W[w,v] = 0
35
36 EDGE_SELECT(D,W)
37 1. Choose u with probability proportional to D[u]
38 2. Once u is fixed, choose v with probability proportional to W[u,v]
39 3. return the edge (u,v)
```

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This is a randomized algorithm for the computation of a graph. In the next subsections we explain how we have implemented the data structure and the functions of the algorithm.

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## 3.1 Data Structure

For the implementation of the algorithm we used:

- **k**: is a constant ( $\log^2 n$ ) used by Karger to repeat Recursive-Contract  $k$  times and to obtain an error with probability less or equal to  $1/n$ ;
- **V**: is the list of nodes;
- **D**: is the list of the sum of the weights of each node;
- **W**: is the list of the graph with 3 parameters(node,node,weight). Each node is connected with the others and if they are not connected the third parameter is set to 0 otherwise is set to the correct weight.

## 3.2 Implementation

For the implementation of the algorithm we used these functions:

- **Karger(G,k)**:
  1. This is the main function of the algorithm where we set the timeout to 60 to limit the execution time of large instances;
  2. We start the time and we set the minimum cut to infinite;
  3. We iterate  $k$  times to obtain an error with probability less or equal to  $1/n$ ;
  4. If the time minus the starting time is greater than the timeout we break;
  5. We execute a copy of V,W and D and then we call the function Recursive-Contract;
  6. If we found a value less than our minimum we update it and we set the discovery time;
  7. In the end we print the Minimum Cut, the Total time and the discovery time.
- **Recursive-Contract(V, W, D)**: In this function we follow exactly the function above with our data structure;
- **Contract(s, V, W, D)**: Also in this function we follow the function above. The only difference is that when we select and contract an edge (u,v) then we remove from V the vertex v to respect the contraction;
- **Contract-Edge(u,v, W, D)**: Also in this case we follow the function above (We update D and W with the new values to execute the contraction of the edge) with our data structure;
- **Edge-Select(V1, D, W)**: Edge selecting random select first node and look for connected other node, algorithm use:
  1. **Random-Select(C)**:
    - (a) Build cumulative weights vector by input array of weights;
    - (b) Set random value  $r$  in range (0, max value of weight);

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- (c) Run binary search to return node related to selected edge;
2. **binarySearch(array, x):**
- (a) Special case of binary search according to inequality  $C[i - 1] \leq r < C[i]$ ;
  - (b) Devide array in parts;
  - (c) Check first value of right part, if random value higher, then go right;
  - (d) If value lower check that previous element lower or equal;
  - (e) Return ralated name of node.

### 3.3 Complexity

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## 4 Results

### 4.1 Table with Min-Cut results

### 4.2 Graph of the Time Cost of the two Algorithms

### 4.3 Graph of the Time Cost compared to the Discovery Time of Karger and Stein Algorithm

### 4.4 Graph of the Time Cost compared to the Asymptotic Complexity of the two Algorithms



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## 5 Conclusion