Lecture Months Tuesday, March 17.

Today

- · Dual feasible solutions as optimality "certificates"
- · Complementary Slackness.
- · KKT conditions (Karush-Kuhn-Tucker)
- · Water filling example.

Primal
$$p^* = \min_{s \leftarrow f_i(x) \leq 0} f_{s(x)} = \min_{s \leftarrow f_i(x) \leq 0} \max_{s \leftarrow f_i(x) \leq 0} L(x, \lambda, \nu)$$

Dual:
$$dt = \max_{s:t: 1 \ge 0} g(\lambda_1 z) = \max_{s:t: 1 \ge 0} \min_{s:t: 1 \ge 0} L(x_1, \lambda_2 z)$$

$$L(\vec{x},\vec{\lambda}) = c^{T}\vec{x} + \lambda^{T}(A\vec{x} - \vec{b})$$

$$= -b^{T}\lambda + (A^{T}\lambda + c)^{T}\vec{x}$$

$$bT\lambda + (A^T\lambda + c)^T z^2$$

$$g(\lambda) = \min_{\chi} L(\chi, \lambda) = \begin{cases} -b^T\lambda & \text{if } (A^T\lambda + c) = 0 \\ -\infty & \text{otherwise} \end{cases}$$

d= maximize -bT)

s.t. ATA+C=0

720

So why do we care about the dual? "Sertificate" function. (BEV. 5.5.1) He g can find (1,2,) some 1, and 2, that are dual feasible, i.e. salisfy the constraints of the dual, then 9 can say $p^* \ge g(\lambda_1, 2)$. It so, is a primal feasible point, then I can say: $f_0(x_1) - p^{+} \leq f_0(x_1) - g(x_1, y_1)$ any any dual
prime feasible feasible points.

print

for the points. If f(2,) is E close to any dual feasible point, then it must be E close to the primal optimal pt itself. -> This is very useful when thinking of stopping criteria for different algorithms. So far we have seen gradient Descent as one example of such a search

afgorithm - if you also have dual information you algorithm - if you also have dual information you are.

Can know how close to the optimal you are.

Now instead of what happen at any primal feasible and dual feasible point, let us consider the primal optimal (1*, 2*) prints.

Assume strong duality holds, so pt = dt $p^* = f_0(\vec{z})$ $d^* = g(\vec{x}^*, \vec{z}^*)$

 $p^{*}=d^{*}=d^{*}=0$ $f_{0}(\vec{x},\vec{x},\vec{y}^{*})$ $f_{0}(\vec{x},\vec{y}^{*})$ $f_{0}(\vec{x},\vec{y}^$

 $L(\vec{x}^*, \vec{x}, \vec{z}^*) \leq L(\vec{x}, \vec{x}^*, \vec{z}^*) \stackrel{(3)}{\leq} f_0(x^*) + \sum_{i=1}^{m} \lambda_i^* f_i(x^*) + \sum_{i=1}^{p} \nu_i^* f_i(x^*)$

L(2*, 7*, 2)*) (4) $\leq f_0(x^+) + 0 + 0$ $f_i(\vec{x}^+) \leq 0$ $h_i(\vec{x}^+) = 0$

But this means inequalities (3) and (4) must in fact be equalities!

Most importantly:

$$\sum_{i=1}^{m} \lambda_{i}^{*} f_{i}(\vec{x}^{*}) = 0$$
 kecause (3) = (4).

But neall λ_i^{+2} 0 and $f_i(\hat{x}^{+}) \leq 0$ $\Rightarrow \lambda; f; (\vec{x}^+) \leq 0.$

The only way a sum of non-negative terms is O is of each turn is zero. Therefor:

$$\lambda : f_i(\overrightarrow{x}^*) = 0 \quad \forall i = 1, 2, \dots m.$$

This is called "complementary slactross."

Non-convex publishers. Strong duality holds.

(i.e. convexity is not needed).

KKT Necessary Conditions i.e.

Like 24 be primal optimal, (1*, 2)*) be dual op

Let 2t be primal optimal, (1*,2)*) be dual opt.
Then, their optimality implies:

 $f_i(x^*) \le 0$, i=1, ..., m Primal feasible $h_i(x^*) = 0$, i=1..., p. Dual feasible. $\lambda_i^* \ge 0$, i=1..., m. Dual feasible. $\lambda_i^* f_i(x^*) = 0$, i=1..., m. Complementary Slackness

 $\nabla f_0(x^*) + \sum_{i=1}^{m} \lambda_i^* \nabla f_i(x^*) + \sum_{i=1}^{p} \nu_i^* \nabla h_i(x^*) = 0$

 $\Rightarrow \chi \neq \text{ minimizes } L(\chi, \lambda^*, \nu^*)$

Hence, graient wirtiz must le zero at optimum point xx.

It is impodant to note that for non-convex problems, these one necessary conditions, i.e. IF xt, It, 21t are optimal they must satisfy the conditions, but there can be points that satisfy the conditions, but one not optimal.

KKT Conditions for convex problems. + diffrentiable

"Sufficient Condition"

1-1- Now if fi(21) (fo(21), f,(21) ... fm(21)) are all Convex, and hi(2) are affine, and \(\tilde{\pi}. \), \(\tilde{\pi} \) are points that satisfy the KKT conditions . ie. if.

 $f_i(\tilde{x}) \leq 0$, j=1,2,...m $h_i(\widetilde{x}) = 0$, $i = 1, 2 \dots P$

, i= 1, ... m $\tilde{\lambda}_{i} \geq 0$

, i=1, ... m $\widetilde{\lambda}_{i} \cdot f_{i}(\widetilde{x}) = 0$

 $\nabla f_0(\tilde{\alpha}) + \sum_{i=1}^{m} \tilde{\lambda}_i \nabla f_i(\tilde{\alpha}) + \sum_{i=1}^{g} \tilde{\nu}_i \nabla \tilde{h}_i(\tilde{\alpha}) = 0$

then: $\tilde{\chi}$ is primal optimal and $(\tilde{\chi}, \tilde{\chi})$ are duel optimal.

200 2, 7, 2 satisfy KKT (=) 2,5,24 are primal and the problem is convex dual optimal.

+ hi are affine

KXT conditions are necessary AND sufficient.

Proof:

Consider: $L(x, \tilde{\lambda}, \tilde{z})$ function of x.

Since $\chi \geq 0$, this is convex in se.

$$L(x, \tilde{\lambda}, \tilde{\omega}) = f_0(x) + \sum \tilde{\lambda}; f_i(x) + \sum \omega; h_i(x)$$

affine

: If $\nabla L(x, \tilde{\lambda}, \tilde{\omega}) \Big|_{x=\tilde{x}} = 0$, then \tilde{x} is minimum point.

$$g(\tilde{\lambda}, \tilde{\Sigma}) = \min_{\chi} L(\chi, \tilde{\lambda}, \tilde{\Sigma}) = L(\tilde{\chi}, \tilde{\lambda}, \tilde{\Sigma})$$

= $f_0(\tilde{\alpha}) + \sum_i \tilde{\lambda}_i f_i(\tilde{\alpha}) + \sum_i \tilde{\lambda}_i f_i(\tilde{\alpha})$

=
$$f_0(\tilde{x}) + 0 + 0$$

Complementary

This is a second se

: (%) and (5, 2) have zero duality gap.

- $g(\lambda, 2)$ is a lower bound for every primal value force). So if $g(\tilde{J}, \tilde{Z}) = fo(\tilde{Z})$ for some value force). So if $g(\tilde{J}, \tilde{Z}) = fo(\tilde{Z})$ this must be the optimum.