# CONSTRAINT SATISFACTION PROBLEMS

Chapter 6

### Outline

- Constraint Satisfaction Problems (CSP)
- Backtracking search for CSPs
- Local search for CSPs

#### Constraint satisfaction problems (CSPs)

- In standard <u>search problems</u>, <u>states</u>:
  - Atomic ("black box" with no internal structure)
  - Evaluated by domain-specific heuristics
  - ☐ <u>Tested</u> to see whether they are goal states
- In CSPs: a factored representation for each state
  - $\square$  State is defined by variables  $X_i$  with values from domain  $D_i$
  - Goal test is a set of constraints specifying allowable combinations of values for subsets of variables
  - Solution: one value for each variable that satisfies all the constraints

#### Constraint satisfaction problems (CSPs)

- Allows useful general-purpose algorithms to solve complex problems
  - with more power than standard search algorithms
  - general-purpose heuristics rather than problem-specific heuristics
- □ The main idea of algorithms for solving CSPs
  - □ To eliminate large portions of the search space all at once
  - by identifying variable/value combinations that violate the constraints

## Constraint satisfaction problem

- □ Set of variables  $X = \{X_1, X_2, ..., X_n\}$
- $\square$  Set of domains  $D = \{D_1, D_2, ..., D_n\}$ 
  - □ Each domain **Di** consists of a set of **allowable values** for variable **X**<sub>i</sub>.
  - In many cases the domain is assumed to be the same for all variables
- □ Set of constraints  $C = \{ c_i = (scope_i, rel_i) \mid i=1,...,h \}$ 
  - scope; subset of X, the variables that are constrained by c;
  - rel<sub>i</sub>: is a relation and tells us which simultaneous assignments of values to variables in scope; are allowed

## Constraint satisfaction problem

- State: defined by an assignment of values to some or all of the variables,  $\{Xi = vi, Xj = vj, ...\}$
- Assignment can be:
  - Consistent: it does not violate any constraints
  - Complete: every variable is assigned
  - Partial: only some of the variables are assigned
- Solution: a consistent and complete assignment

## **Example: Map-Coloring**

Coloring each region either red, green, or blue in such a way that no neighboring regions have the same color.



#### **CSP** formulation

- $\square$  Set of variables  $X = \{WA, NT, Q, NSW, V, SA, T\}$
- $\square$  Domain of each variabile  $D_i = \{\text{red, green, blue}\}$
- Constraints: adjacent regions must have different colors

```
C = \{SA \neq WA, SA \neq NT, SA \neq Q, SA \neq NSW, SA \neq V, WA \neq NT, NT \neq Q, Q \neq NSW, NSW \neq V \}
WA \neq NT, or
(WA,NT) in
\{(red,green),(red,blue),(green,red),(green,blue),(blue,red),(blue,green)\}
```

## **Example: Map-Coloring**



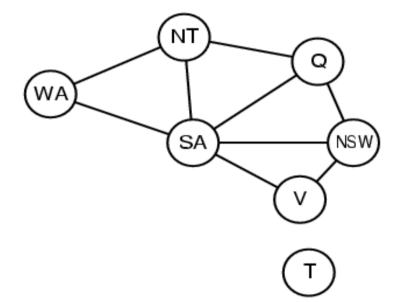
#### Solutions are complete and consistent assignments

e.g.,  $\{WA = red, NT = green, Q = red, NSW = green, V = red, SA = blue, T = green\}$ 

## Constraint graph



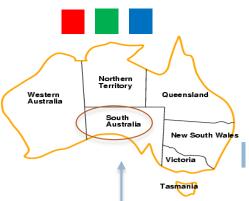
- □ Constraint graph
  - nodes are variables
  - arcs show constraints



#### Why formulate a problem as a CSP?

- A natural representation for many problems
- We already have a CSP-solving system
  - it is often easier to solve a problem using it
  - than to design a custom solution using another search technique
- □ CSP solvers are faster than state-space searchers because the CSP solver can quickly eliminate large parts of the search space

## Why <u>formulate</u> a problem as a CSP?



#### CSP solvers are faster than state-space searchers

Eg., if we choose  $\{SA = blue\}$  in the Australia problem

- None of the five neighboring variables can take blue value
- Without constraint propagation
  - a search procedure should consider 3<sup>5</sup> = 243 assignments for the five neighboring variables
- With constraint propagation
  - we never have to consider blue as a value
  - $\blacksquare$  so we have only  $2^5 = 32$  assignments to look

#### Varieties of CSPs

The simplest kind of CSP involves variables with discrete finite domains

#### Discrete variables

#### **□** Finite domains:

- n variables, domain size d
- e.g., variables WA, NT, Q, NSW, V, SA, T in the map coloring problem and each variable has the domain  $Di = \{red, blue, green\}$

#### Infinite domains:

- integers, strings, etc.
- e.g., job scheduling, variables are start/end days for each job
- $\blacksquare$  constraints: StartJob<sub>1</sub> + 5  $\le$  StartJob<sub>3</sub>

#### Continuous variables

 common in the real world problems, studied in the field of operations research

#### Varieties of constraints

- □ Unary constraints involve a single variable
  - □ e.g., SA ≠ green
- Binary constraints involve pairs of variables
  - e.g., SA ≠ WA
- ☐ **Higher-order** constraints involve 3 or more variables
- Global constraints involve an <u>arbitrary</u> number of variables
  - e.g., Alldiff, which says that all of the variables involved in the constraint must have different values

## Conversion to binary

Binary CSP: CSP where each constraint relates two variables

Any CSP can be converted into a CSP with only binary constraints