CONSTRAINT SATISFACTION PROBLEMS — PART III

Chapter 6

Outline

- Constraint Satisfaction Problems (CSP)
- Backtracking search for CSPs
- Local search for CSPs

Standard search formulation

Let's start with the straightforward approach, then improve it

- Initial state: the empty assignment { }
- Successor function: assign a value to an unassigned variable that does not conflict with current assignment
 - > fail if no legal assignments
- State: partial assignment
- □ Goal test: the current assignment is complete

- This is the <u>same</u> for all CSPs
- Every <u>solution</u> appears at depth n with n variables
 - → use depth-first search

Branching factor: maximum number of

Standard search formulation successors of a node

We could use depth-first search

- But for a CSP with n variables of domain size d
- The branching factor b is
 - at 1st level: nd because <u>any of d values</u> can be assigned to <u>any of n variables</u>
 - \square at 2^{nd} level: (n-1)d
 - **-** ...
 - at n-th level: d
- We generate a tree with n! · dⁿ leaves,
 even if only dⁿ possible complete assignments

Backtracking search

In CSPs, variable assignments are commutative, i.e.,

```
{WA = red then NT = green } \underline{same as}
{NT = green then WA = red }
```

- □ Thus we need only consider a single variable at each level in the search tree \rightarrow b=d \rightarrow the number of leaves is d^n
- Backtracking search = Depth-first search for CSPs with singlevariable assignments

A depth-first search that

- chooses values for SEARCH one variable at a time and
- backtracks when a variable has no legal values left to assign.

Backtracking search

```
function BACKTRACKING-SEARCH(csp) returns a solution, or failure
  return BACKTRACK( { }, csp )
function BACKTRACK(assignment, csp) returns a solution, or failure
  if assignment is complete then return assignment
  var ← SELECT-UNASSIGNED-VARIABLE(csp)
 for each value in ORDER-DOMAIN-VALUES(var, assignment, csp) do
     if value is consistent with assignment then
        add {var = value} to assignment
        result ← BACKTRACK(assignment, csp)
        if result ≠ failure then return result
        remove {var = value} from assignment
 return failure
```

Review: Map-Coloring

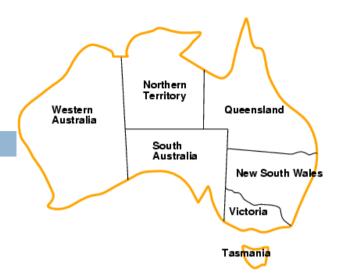
CSP formulation

- □ Variables WA, NT, Q, NSW, V, SA, T
- \square Domains $D_i = \{ \text{red, green, blue} \}$
- Constraints: adjacent regions must have different colors

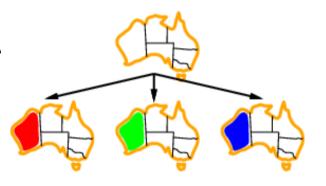


Part of the search tree

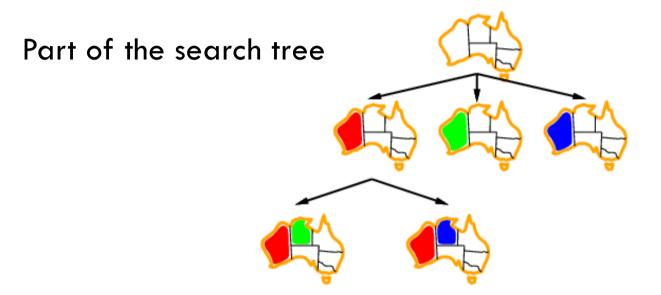


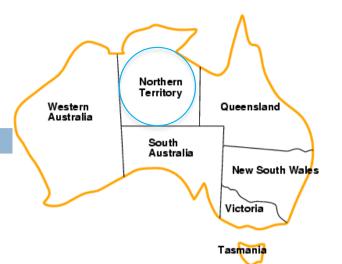


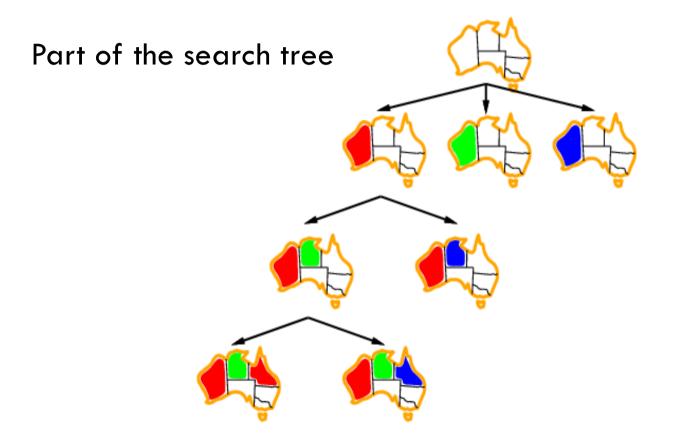
Part of the search tree

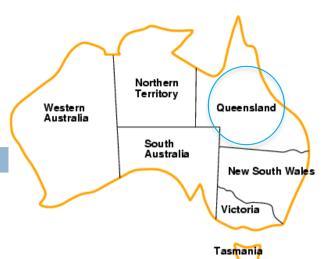


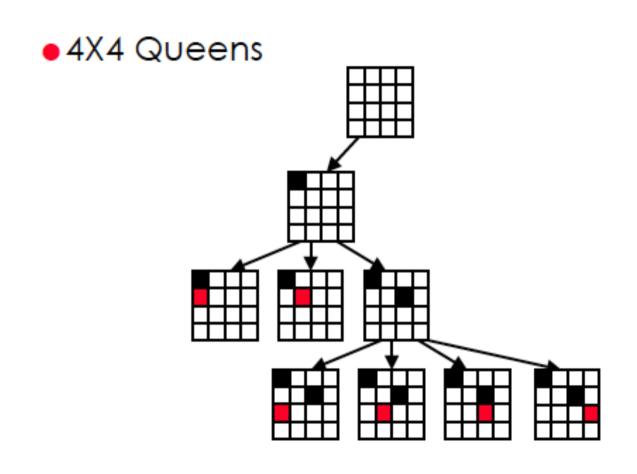


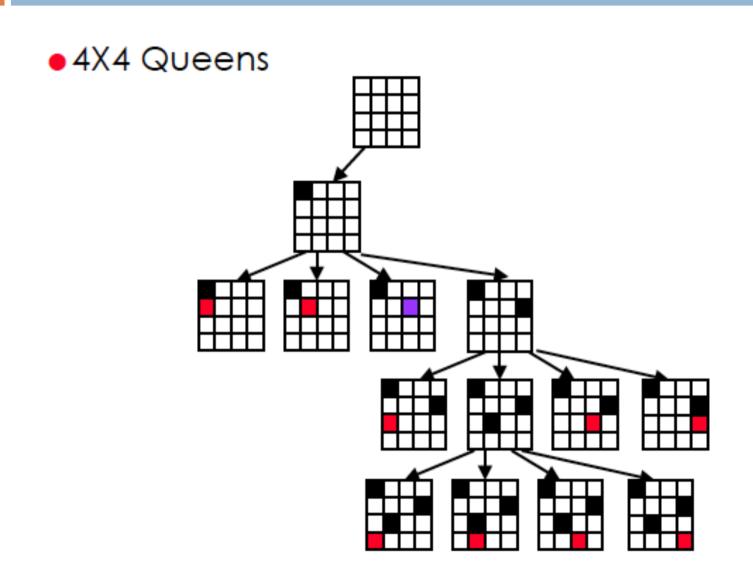




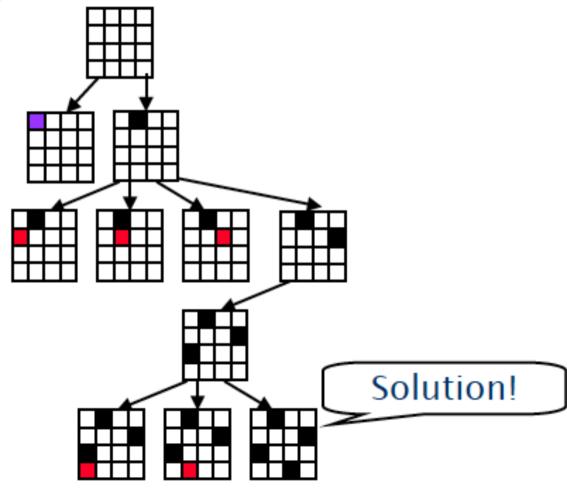








4X4 Queens



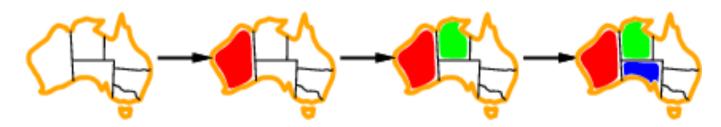
Improving backtracking efficiency

- General-purpose methods can give <u>huge gains</u> in speed:
 - Which variable should be assigned <u>next</u>?
 - In what order should its values be tried?

Most constrained variable

■ Most constrained variable: FIRST FAIL

chooses the variable with the fewest legal values



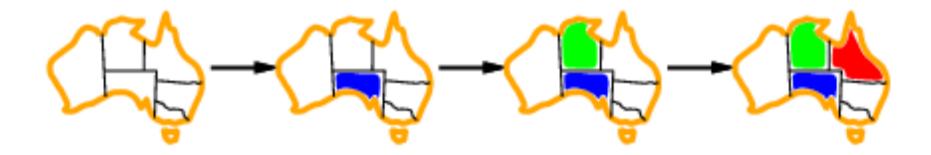
□ a.k.a. minimum remaining values (MRV) heuristic

euristica per scegliere quale variabile selezionare per prima

Most constraining variable

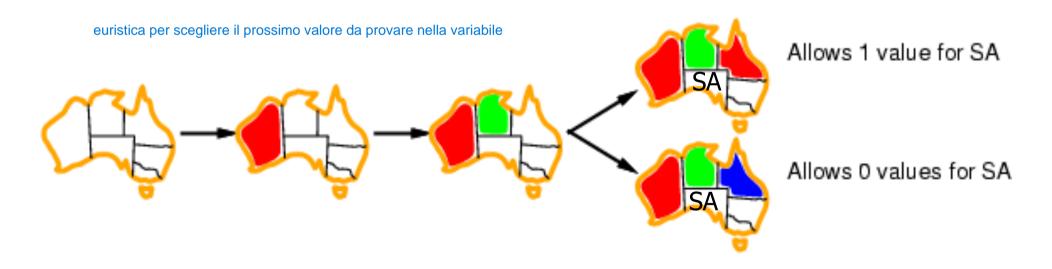


- Tie-breaker among most constrained variables
- Most constraining variable (aka, degree heuristic): chooses the variable involved in the most constraints with unassigned variables



Least constraining value

- Given a variable, chooses the <u>least constraining value</u>:
 - the one that rules out the fewest values in the remaining variables



Value ordering

- □ The ordering of values does not matter if
 - all solutions needed
 - □ no solution

because we have to consider every value