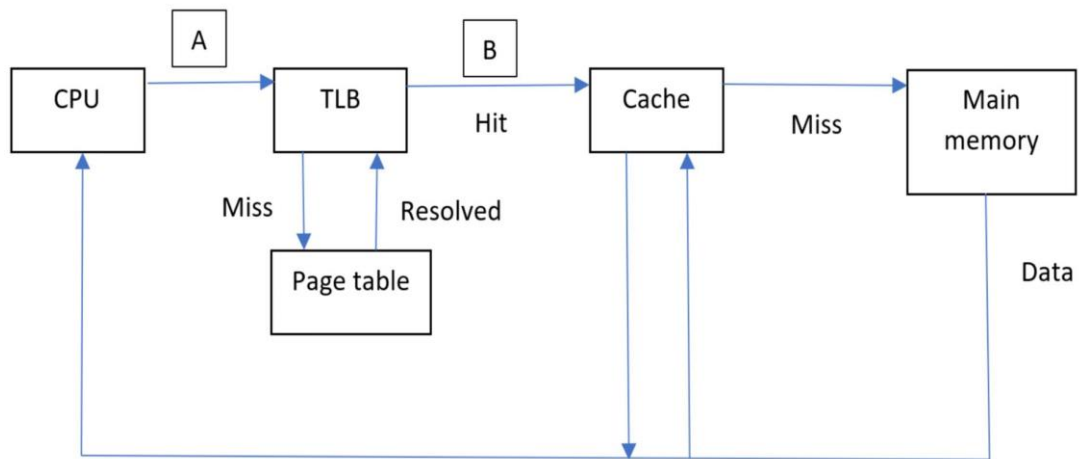


Group:

組員簽名 : _____

[group 14] (對抗賽)

1. Consider the following figure



- (A) What is the name of the TLB's input (Symbol A in the figure)?
- (B) What is the name of the TLB's output (Symbol B in the figure)?
- (C) What is the name of this type of cache (with B as its input)?
- (D) "We could have a hit in the cache, and get a TLB miss and a page table miss." Is the statement true (Yes or No)? Explain your answer.
- (E) Consider the processor operating at 1 GHz. The processor stalls during a cache miss and has the following properties: (i) a cache access time of 2 clock cycles for a hit, (ii) a miss penalty of 100 clock cycles, and (iii) a miss rate of 0.03 misses per reference. Please compute the average memory access time.

Ans:

- (A) Virtual address
- (B) Physical address
- (C) Physically addressed cache
- (D) No page table miss → data are not in memory → data are not in cache
- (E) Average memory access time = $(2 + 0.03 \times 100) \times 1 \text{ ns} = 5 \text{ ns}$

[group 9] (對抗賽)

2. Please fill the block below

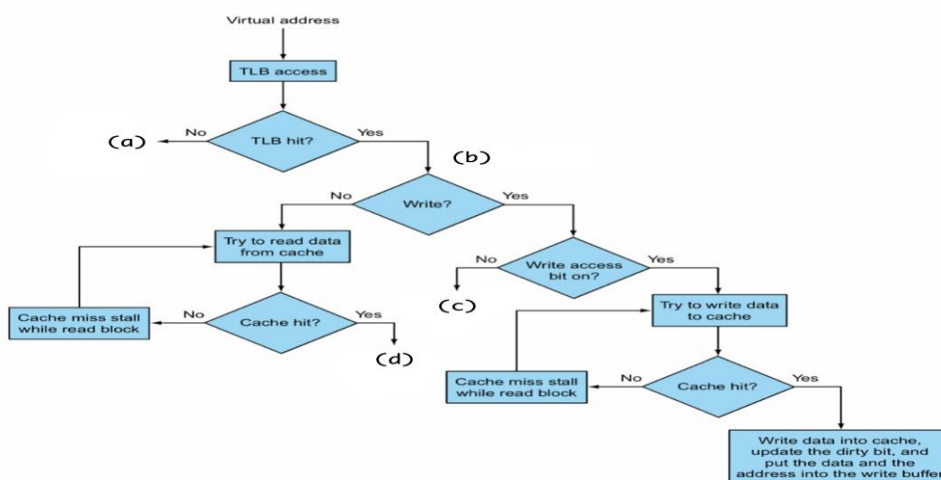
Cache	TLB	Page table	Possible?
Hit	Miss	Miss	
Miss	Hit	Miss	
Miss	Miss	Hit	
Hit	Hit	Miss	
Miss	Hit	Hit	
Hit	Miss	Hit	

Ans:

Cache	TLB	Page table	Possible?
Hit	Miss	Miss	Impossible
Miss	Hit	Miss	Impossible
Miss	Miss	Hit	Possible
Hit	Hit	Miss	Impossible
Miss	Hit	Hit	Possible
Hit	Miss	Hit	Possible

[group 6] (對抗賽)

3. Fill in the following blanks by: **TLB miss exception**, **Physical address**, **Write protection exception**, and **Deliver data to the CPU**.



Ans:

(a) TLB miss exception

(b) Physical address

(c) Write protection exception

(d) Deliver data to the CPU

[group 10] (對抗賽)

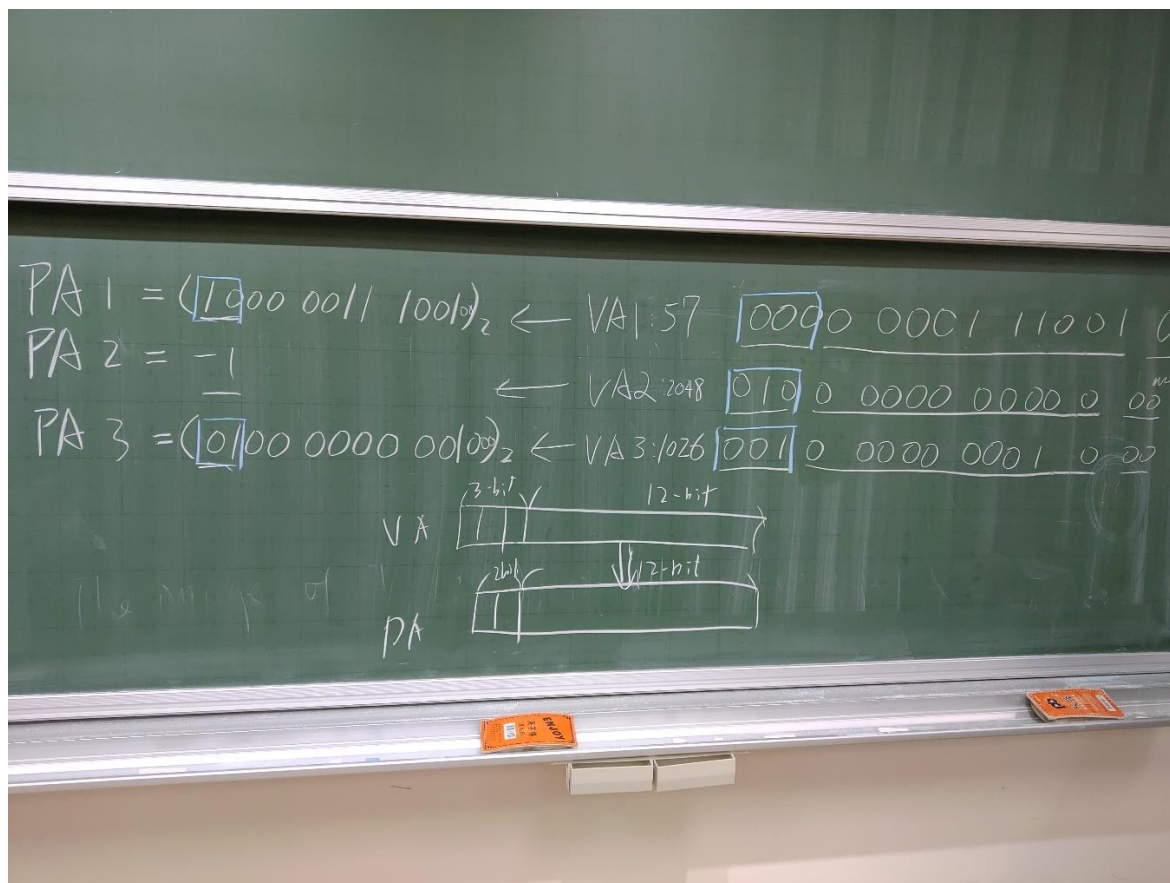
4. The information of a virtual memory system is assumed to be:

- Page size: 1024 words
- Number of virtual pages: 8
- Number of physical pages: 4
- The current page table is

VPN	0	1	2	3	4	5	6	7
PPN	2	1	NULL	NULL	3	NULL	0	NULL

For the following (decimal) virtual word addresses: VA1: 57, VA2: 2048, VA3: 1026, VA4:7749, VA5 6150, their corresponding physical addresses are PA1, PA2, PA3, PA4, PA5. And PA = -1 if it is a page fault. Please write down PA1~PA5, separately.

Ans:



[group12]

5. Answer the following questions:

- (1) What design changes can be used to improve any one of the 3Cs?
- (2) What possible negative effect may occur for each of the design changes?

	Compulsory	Capacity	Conflict
(1)			
(2)			

Hints: Increase cache size; Increase block size; Increase access time; Increase associativity; Increase miss penalty.

ANS:

	Compulsory	Capacity	Conflict
(1)	Increase block size	Increase cache size	Increase associativity
(2)	Increase miss penalty	Increase access time	Increase access time

[group8] (對抗賽)

6. Which of the following statements are true?
- (a) All addresses generated by the program are physical addresses
 - (b) One memory reference only needs one memory access.
 - (c) We use hardware to detect page fault and handle the faults in software.
 - (d) Array of page table entries, indexed by page offset.
 - (e) PTE can refer to location in swap space on disk.

ANS: C, E

- (a) all addresses generated by the program are **virtual** addresses
- (b) One memory reference only needs **two** memory access: access page table and access physical memory.
- (d) Array of page table entries, indexed by **virtual page number**.

[group13] (對抗賽)

7. Which statements are true? If it's false, please explain why.
- (a) When you want to access data in the memory without using TLB strategy, you can directly access it.
 - (b) Page fault means that page is not resident in memory.
 - (c) Hardware can handle the page faults by itself.
 - (d) Following (c), we can infer that OS doesn't need to know where to find the page.
 - (e) As TLB hit on write happens, dirty bit should be toggled in addition to write back to page table on replacement.
 - (f) TLB hit and Page table miss can happen simultaneously.

ANS:

- (a) False, you should access page table first so that you can access data stored in the memory.
- (b) True.

- (c) Hardware should trap to the OS so that it can remedy page faults.
- (d) False, OS should need to know where to find the page.
- (e) True.
- (f) False, PTE haven't been moved to memory so that TLB hit is impossible.

[group 5] (對抗賽)

8. True or False

- a. Every process on OS has its own page table.
- ~~b. 刪除~~
- c. CPU doesn't need to visit TLB before fetching instruction.
- d. Without virtual memory, memory access can be faster.
- e. In terms of general-purpose OS, when a user process does something like:

```
int x = 1;
printf("%p", &x);
```

the standard output shows the physical address of variable x.

Ans:

- (a) T
- ~~(b) 刪除~~
- (c) F, TLB is needed.
- (d) T
- (e) F, Virtual address

[group 4]

9. Which of the following statements are true? if not please elaborate the reasons

- (A) Reducing the page fault is important because the miss penalty is huge up to millions of cycles.
- (B) As the TLB miss occurs, if the page is not in the memory, it can be fixed by hardware method.
- (C) TLB is like a cache for a disk.
- (D) It is possible that miss happens in TLB and page table in the same time while cache is hit.
- (E) When TLB hit and read it will trigger the dirty bit.

Ans:

- (A) True
- (B) False, it can only be solved by software method which OS will handle fetching the page and updating the page table.
- (C) True
- (D) False, it is not possible since it will not be cache if it is not in the memory.
- (E) False, as TLB hit and write, it will trigger the dirty bit.