

REAL ANALYSIS

MATH 131, HARVEY MUDD COLLEGE

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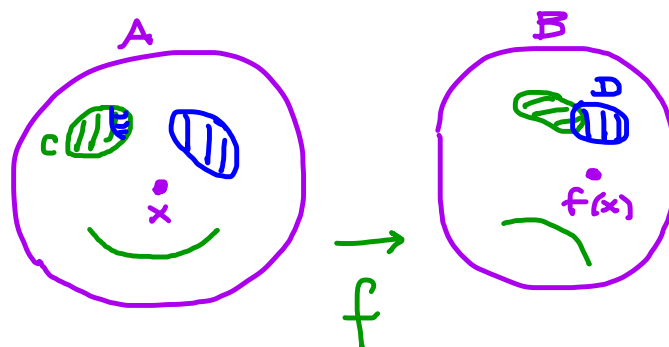
TODAY: COUNTABLE SETS

Q: How do we count?

A: We construct a function.

Recall: $f: A \rightarrow B$

associates $x \mapsto f(x)$
↑
"maps"



Write: $f(C) = \{ f(x) : x \in C \}$ the image of C
 $f^{-1}(D) = \{ x : f(x) \in D \}$ the pre-image of D

- When $f(A) = \text{all } B$, say f is onto (surjective) \rightarrow
- When $f(x) = f(y)$ implies $x = y$, say f is 1-1 (injective) \leftarrow
- When f is 1-1 and onto, say f is bijective. \leftrightarrow

A bijection puts A, B into "1-1 correspondence"

Write: $A \sim B$.

Ex. $\overbrace{\{1, \text{☺}, \infty\}}^A \sim \{1, 2, 3\}$. Say A has 3 elts.

Def'n. Call A finite if $A \sim \overbrace{\{1, 2, \dots, n\}}^{[n]}$ for some $n \in \mathbb{N}$.
or if A is empty.

Else A is infinite.

an infinite

Defn. Call A countable if $A \sim \mathbb{N}$.

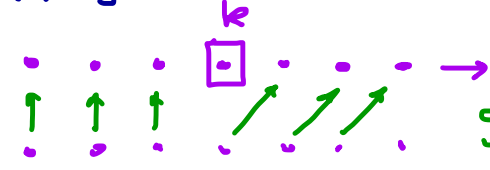
Else A is uncountable.

Ex. \mathbb{N} is ctble: use bijection $f(n) = n$.

Ex. A sequence x_1, x_2, x_3, \dots of distinct terms is ctble:
 $\begin{matrix} \updownarrow & \updownarrow & \updownarrow \\ 1 & 2 & 3 \end{matrix}$ use $f(n) = x_n$.

NOTICE: Any ctble set can be "listed" in a seq.

Ex. $\{2, 3, 4, \dots\}$ is ctble: use $f: \mathbb{N} \rightarrow \mathbb{N} - \{1\}$
 $\begin{matrix} \uparrow & \uparrow & \uparrow \\ 1 & 2 & 3 \end{matrix}$ mapping $n \mapsto n+1$.

Ex. $\mathbb{N} - \{k\}$ is ctble:
use $g: \mathbb{N} \rightarrow \mathbb{N} - \{k\}$
def'd by $\begin{cases} g(n) = n & \text{if } n < k \\ g(n) = n+1 & \text{if } n \geq k \end{cases}$


Thm. \mathbb{N} is infinite.

proof. We'll show \nexists bijection $[n] \leftrightarrow \mathbb{N}$ by induction on n .

base case: If $\exists f_1: [1] \leftrightarrow \mathbb{N}$, consider $\mathbb{N} - f_1(1)$.

It's not empty. So f_1 not bijective.

ind. step: we'll show $[k] \not\leftrightarrow \mathbb{N} \Rightarrow [k+1] \not\leftrightarrow \mathbb{N}$.

equiv: show if $\exists f_{k+1}: [k+1] \leftrightarrow \mathbb{N} \Rightarrow \exists f_k: [k] \leftrightarrow \mathbb{N}$.

idea: $[k+1] \xrightarrow{f_{k+1}} \mathbb{N}$

restricts to $[k] \xrightarrow{f_k} \mathbb{N} - f_{k+1}(k+1) \xrightarrow{g^{-1}} \mathbb{N}$ is desired bijection.

Ex. $2\mathbb{N}$ (even #'s) is ctble: use $f(n)=2n$.

Ex. $\mathbb{Z} = \{\dots, -2, -1, 0, 1, 2, \dots\}$ is ctble.



Thm. Every infinite subset of ctble set is ctble.

pf idea:

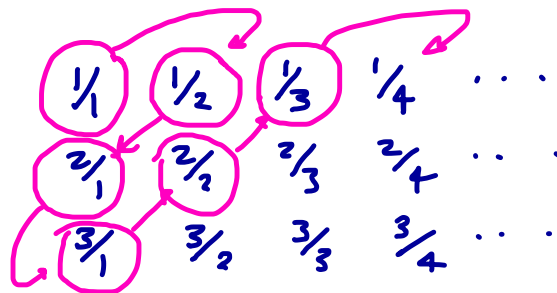
subset $E = \{x_1, x_2, \dots\}$
 $x_{n_1} \quad x_{n_2} \quad x_{n_3} \rightarrow$
 $x_1 \quad x_2 \quad x_3 \quad x_4$
 set $A = \{x_1, x_2, \dots\}$.

Let $n_1 = \inf \{i : x_i \in E\} \leftarrow \exists \text{ least elt by WOP}$
 $\dots \quad n_k = \inf \{i : x_i \in E \text{ \& } i > n_{k-1}\}.$

Then $E = \{x_{n_1}, x_{n_2}, \dots\}$. Or use
 $f(k) = x_{n_k}$

Thm. \mathbb{Q} is ctble.

Idea: \mathbb{Q}_+ is ctble:

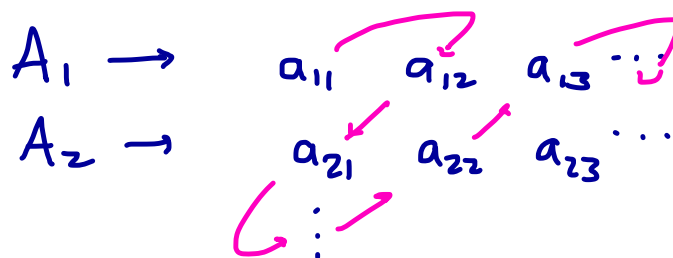


See elts in array are ctble.

Subset (w. one rep of each
class of $\frac{m}{n}$) is ctble!

Thm. A ^{at most} ctble union of ^{at most} ctble sets is ^{at most} ctble.

pf. same idea: if A_1, A_2, \dots ctble sets



Appl. # computer progs is ctble.

Q. Are there uncountble sets?

Q. Is \mathbb{R} unctble?

If so there are ^{"computable" #'s} #'s whose decimal exp
is not the output of
comp. prog!