

Embedded Systems Exercise 2 - HS 2020 21./23.10.2020

Stefan Draskovic

Exercise structure

- Goal of today's exercise:
 - Cyclic Executive Scheduling
- Agenda:
 - Wednesday 16:15 17:00 Introduction and solving a sample question (recorded)
 - Friday 16:15 17:00 Solutions (recorded)
- Available assistants:
 - Stefan Drašković
 - Julian Keller



Exercise structure

• Interactions:

- Exercise Zoom: Questions can be asked throughout the exercise in this room by raising your hand. Please feel free to write in the chat in case we oversee your question.
- Help Zoom: Student assistants are available after the session for questions and 1-on-1 meetings.
- Matrix Chatroom: Questions that are relevant for everyone can be asked in the Matrix chatroom where the responsible assistants can answer as quickly as possible.





Content of today's exercise

Cyclic Executive Scheduling: Analyze one schedule and construct another



Definitions

Γ: Task-set, or the set of all tasks

 τ_i : Task

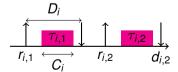
 $\tau_{i,j}$: Job, or the j^{th} instance of task τ_i

 $r_{i,j}$: Release time of job $\tau_{i,j}$

 $d_{i,j}$: Absolute deadline of job $\tau_{i,j}$

 D_i : Relative deadline of task τ_i , $D_i = d_{i,j} - r_{i,j}$

 C_i : Worst case execution time of task τ_i



Time-triggered Cyclic-executive Scheduling

• **Assumption:** Tasks are periodic, but may have different periods. Task τ_i has period T_i , and phase Φ_i .

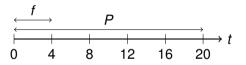
$$r_{i,j} = \Phi_i + (j-1)T_i$$

 $d_{i,j} = \Phi_i + (j-1)T_i + D_i$

Objective: Schedule the task-set using a simple scheduling scheme

Time-triggered Cyclic-executive Scheduling

• The period *P* of the system is divided into frames *f*



- Assignment of jobs to frames is made off-line
- Timer interrupts regularly every frame start, and releases the jobs for this frame

Conditions for P and f

1. A task executes at most once within frame

$$\forall \tau_i: f \leq T_i$$

- 2. P is a multiple of f
- 3. Tasks start and complete within a single frame

$$\forall \tau_i: f \geq C_i$$

4. Between the release time and deadline of every task there is at least one full frame

$$\forall \tau_i$$
: $2f - \gcd(T_i, f) \leq D_i$

Γ	Ti	Фі	Di	Ci	frame
$ au_{ extsf{1}}$	12	2	8	2.8	
$ au_2$	12	3	9	3	
$ au_3$	4	0	4	1	



Γ	T_i	Фі	Di	Ci	frame
$ au_{ extsf{1}}$	12	2	8	2.8	
$ au_2$	12	3	9	3	
$ au_3$	4	0	4	1	

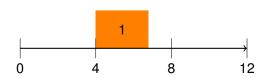




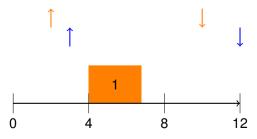


Γ	T_i	Фі	Di	Ci	frame
$ au_{ extsf{1}}$	12	2	8	2.8	2
$ au_2$	12	3	9	3	
$ au_3$	4	0	4	1	

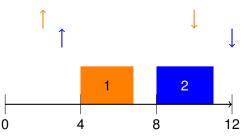




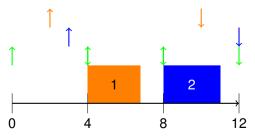
Γ	T_i	Фі	Di	Ci	frame
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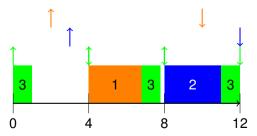
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Γ	T_i	Фі	Di	Ci	frame
$ au_{ extsf{1}}$	12	2	8	2.8	2
$ au_2$	12	3	9	3	3
$ au_3$	4	0	4	1	1, 2, 3



Correctness of Schedule

let f_{ij} note a frame in which that job $\tau_{i,j}$ executes

- Is *P* a common multiple of all periods *T_i*? Is *P* a multiple of *f*?
- Is the frame sufficiently long?

$$\sum_{\{i|f_{ij}=k\}} C_i \leq f \qquad \forall 1 \leq k \leq \frac{P}{f}$$

Are release times respected? or
 Determine offsets such that instances start after release time

$$\forall \tau_i: \qquad \Phi_i = \min_{1 \leq j \leq \frac{P}{T_i}} \left\{ (f_{ij} - 1)f - (j-1)T_i \right\}$$

Are deadlines respected?

$$\forall au_i, \ 1 \leq j \leq \frac{P}{T_i}: \qquad (j-1)T_i + \Phi_i + D_i \geq f_{ij}f_{\text{Stefan Draskovic}}$$
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Task 1: Check Schedule Correctness!

let f_{ii} note a frame in which that job $\tau_{i,i}$ executes

- Is P a common multiple of all periods T_i ? Is P a multiple of f?
- Is the frame sufficiently long?

$$\sum_{\{i|f_{ij}=k\}} C_i \le f \qquad \forall 1 \le k \le \frac{P}{f}$$

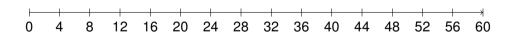
Are release times respected? or
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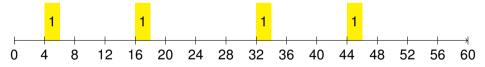
Are deadlines respected?

$$\forall \tau_i, \ 1 \leq j \leq \frac{P}{T_i}: \qquad (j-1)T_i + \Phi_i + D_i \geq f_{ij}f_{\text{Stefan Draskovic}}$$

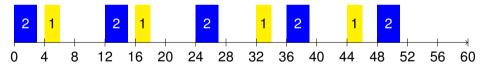
Is P a common multiple of all periods T_i? Is P a multiple of f?
 Yes!



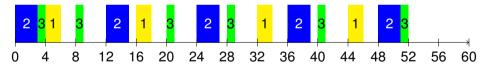
Is P a common multiple of all periods T_i? Is P a multiple of f?
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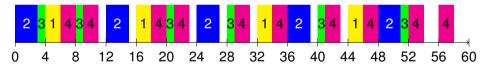
Is P a common multiple of all periods T_i? Is P a multiple of f?
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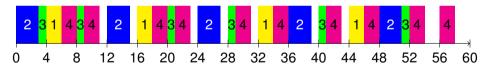


Is P a common multiple of all periods T_i? Is P a multiple of f?
 Yes!



Is P a common multiple of all periods T_i? Is P a multiple of f?
 Yes!

Is the frame sufficiently long?



Yes!

• Determine offsets such that instances start after release time.

$$\Phi_1 = \min \begin{cases} (2-1)4 - (1-1)15 \\ (5-1)4 - (2-1)15 \\ (9-1)4 - (3-1)15 \\ (12-1)4 - (4-1)15 \end{cases} = \min \begin{cases} 4 \\ 1 \\ 2 \\ -1 \end{cases}$$

Determine offsets such that instances start after release time.

$$\Phi_{1} = \min \begin{cases} (2-1)4 - (1-1)15 \\ (5-1)4 - (2-1)15 \\ (9-1)4 - (3-1)15 \\ (12-1)4 - (4-1)15 \end{cases} = \min \begin{cases} 4 \\ 1 \\ 2 \\ -1 \end{cases}$$

$$\Phi_{2} = 0 \quad \Phi_{3} = -2 \quad \Phi_{4} = 2$$

• Are deadlines respected?

Yes! For τ_1 :

$$\begin{cases} (1-1)15-1+9=8\geq 8=2\cdot 4\\ (2-1)15-1+9=23\geq 20=5\cdot 4\\ (3-1)15-1+9=38\geq 36=9\cdot 4\\ (4-1)15-1+9=53\geq 48=12\cdot 4 \end{cases}$$

• Are deadlines respected?

Yes! For τ_1 :

$$\begin{cases} (1-1)15-1+9=8\geq 8=2\cdot 4\\ (2-1)15-1+9=23\geq 20=5\cdot 4\\ (3-1)15-1+9=38\geq 36=9\cdot 4\\ (4-1)15-1+9=53\geq 48=12\cdot 4 \end{cases}$$

... deadlines are also respected for τ_2 , τ_3 and τ_4

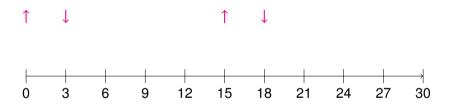
Task 2: Find Schedule

Task	Period	Deadline	Execution Time
$ au_1$	15	3	3
$ au_2$	10	5	3
$ au_3$	6	6	3

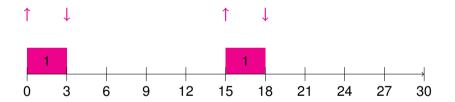
Task 2: Possible Solution



Task 2: Possible Solution

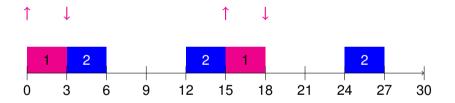


Task 2: Possible Solution



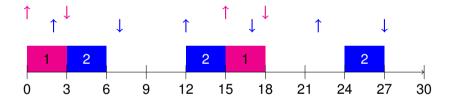


Task 2: Possible Solution



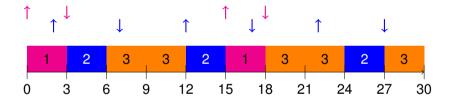


Task 2: Possible Solution



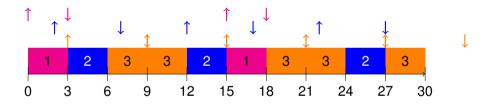


Task 2: Possible Solution





Task 2: Possible Solution







Questions?

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