

Aperiodic Scheduling

Earliest Deadline Due, Latest Deadline First, Earliest Deadline First

Exercise class 4

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Gliederung

Organisation

Overview Aperiodic Task Scheduling

Task 1

Task 2

Task 3



Task 4

Organisation






Organisation I

► feedback for me: <https://forms.gle/f3YN8EFrZ1vsfPoC6>

Feedback for Introduction to ESE Tutor Jürgen  

Questions Responses Settings

0 responses  

Accepting responses 

Waiting for responses

Overview Aperiodic Task Scheduling



Overview Aperiodic Task Scheduling

	Equal arrival times non preemptive	Arbitrary arrival times preemptive
Independent tasks	EDD (Jackson)	EDF (Horn)
Dependent tasks	LDF (Lawler)	EDF* (Chetto)

Overview Aperiodic Task Scheduling

- ▶ **Lateness:** $L_i = f_i - d_i$
- ▶ **Maximum lateness:** $L_{\max} = \max_i (f_i - d_i)$
 - ▶ this a a metric to compare schedules

Task 1

Task 1

Earliest Deadline Due

Task 1.1:

	J_1	J_2	J_3	J_4
C_i	3	6	2	4
D_i	8	15	3	11

$(\forall J_i \in J : a_i = 0)$

Requirements 1.1:

- ▶ *non-preemptive*
- ▶ tasks have *same arrival times* (synchronous arrivals)
- ▶ tasks are *independent*
- ▶ $\min(D_i)$ for all remaining J_i (d_i if $\forall J_i \in J : a_i = c \wedge c \neq 0$)
- ▶ *minimizing the maximum lateness*

Task 1

Earliest Deadline Due

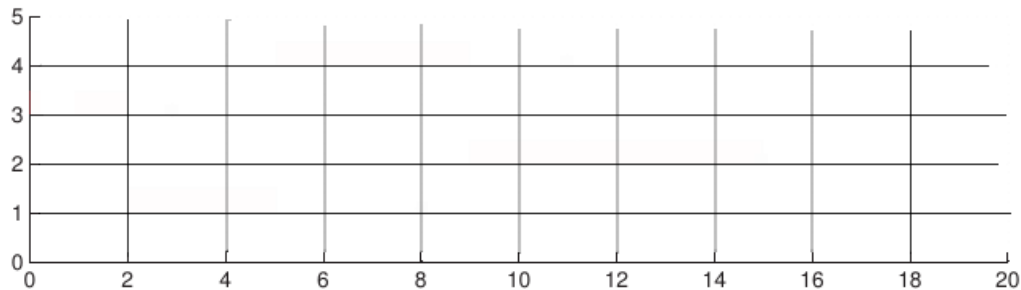


Figure 1: EDD schedule.

Task 1

Earliest Deadline Due

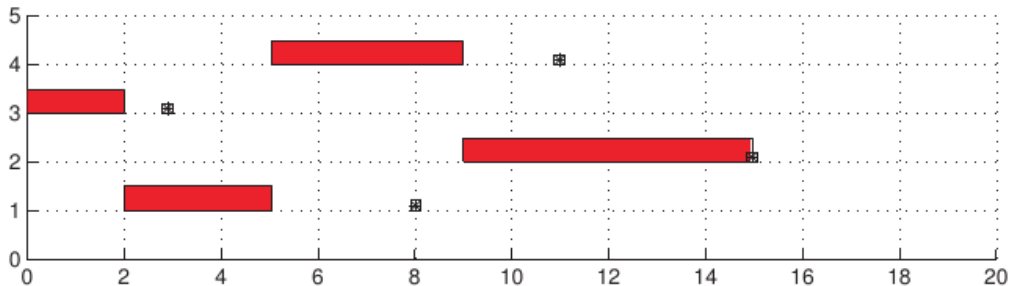


Figure 1: EDD schedule.

Task 2



Task 2

Latest Deadline First

Task 2.1:

	J_1	J_2	J_3	J_4	J_5	J_6	J_7	J_8
C_i	3	4	2	3	3	2	2	1
D_i	5	8	11	15	12	18	19	20

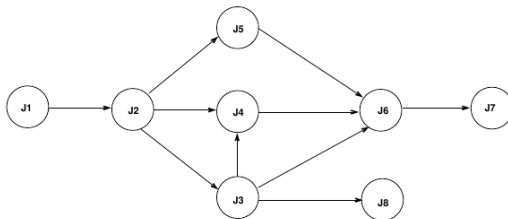


Figure 2: Precedence graph.

Task 2

Latest Deadline First

Requirements 2.1:



- ▶ is *non-preemptive*
- ▶ *synchronous task activations*
- ▶ tasks are *dependent*, use *precedence graph*, going from *tail* to *head*
- ▶ $\max(D_i)$ (d_i if $\forall J_j \in J : a_j = c \wedge c \neq 0$) for all tasks J_i *without successors* or whose *successors* have been all selected in the *precedence graph* inserted into the queue to be *executed last*
- ▶ at runtime, tasks are extracted from the *head of the queue*: the *first task* inserted in the queue will be *executed last*
- ▶ *minimizes the maximum lateness*

Task 2

Latest Deadline First

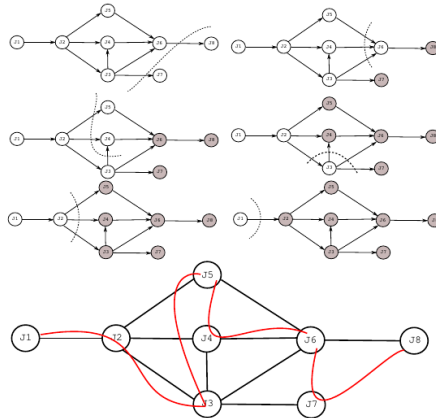


Figure 4: The LDF algorithm proceeds as depicted (figures left to right)

Task 2

Latest Deadline First

► queue of tasks: (, , , , , , ,)

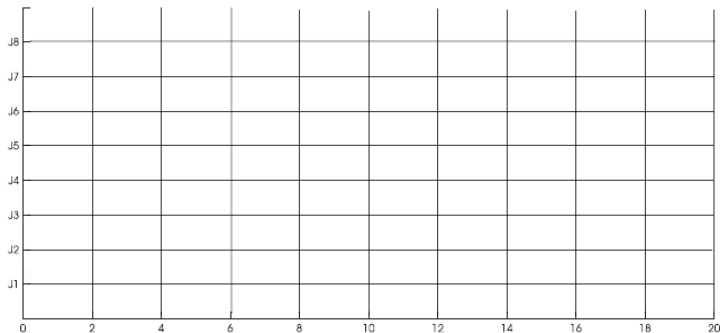


Figure 3: LDF schedule.

Task 2

Latest Deadline First

- queue of tasks: $(J_1, J_2, J_3, J_5, J_4, J_6, J_7, J_8)$

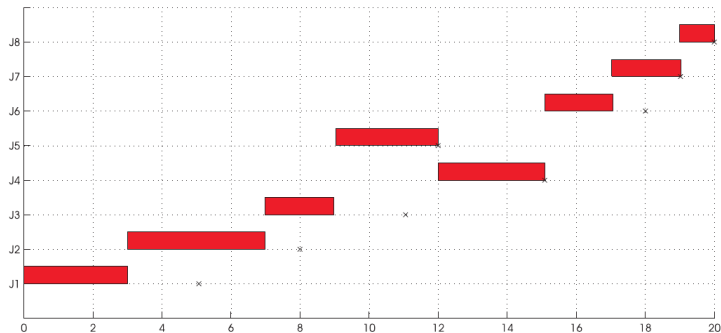


Figure 3: LDF schedule.

Task 3



Task 3

Earliest Deadline First

Task 3.1:

	J_1	J_2	J_3	J_4	J_5
a_i	0	2	0	8	13
C_i	3	1	6	2	3
d_i	16	7	8	11	18

Requirements 3.1:

- ▶ *is preemptive*
- ▶ *arbitrary arrival times*
- ▶ *the tasks are independent*
- ▶ *minimizes the maximum lateness*
- ▶ *$\min(d_i)$ for all remaining tasks J_i that have already arrived (are ready) and not finished every time the arrival time of a task is reached*

Task 3

Earliest Deadline First

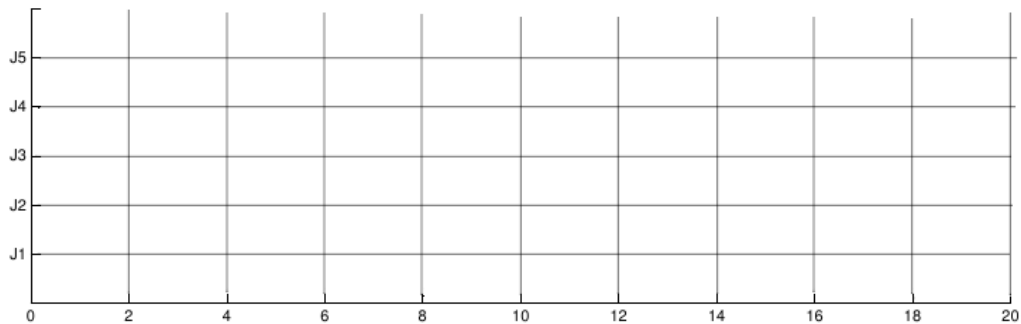


Figure 5: EDF schedule

Task 3

Earliest Deadline First

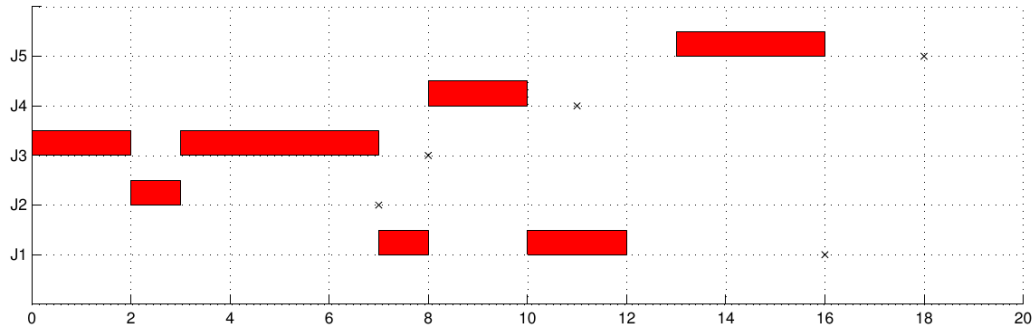


Figure 5: EDF schedule

Task 3

Earliest Deadline First

Task 3.2:

- ▶ at time $t = 3$, a new task J_x arrives with execution time $C_x = 2$ and deadline $d_x = 10$.
- ▶ still **guarantee the schedulability** of task set?

Requirements 3.2:

- ▶ $\forall i = 1, \dots, n \quad t + \sum_{k=1}^i c_k(t) \leq d_i$

Task 4



Task 4 I

EDF* - Example

Task 4.1:

Given tasks A, B, C, D, E, F, G with precedences $A \rightarrow C, B \rightarrow C, C \rightarrow E, D \rightarrow F, B \rightarrow D, C \rightarrow F, D \rightarrow G$.

All tasks arrive at time $t_0 = 0$, have a common deadline $d = 20$ and the following execution times:

	A	B	C	D	E	F	G
C_i	3	2	4	3	2	5	1

We will now prepare the tasks for EDF*

Task 4

EDF*

Requirements 4.1:



- ▶ *preemptive*
- ▶ *arbitrary arrival times*
- ▶ *tasks are dependent*
- ▶ *release time and deadline* of individual tasks are modified such that all the *precedence constraints* are satisfied
- ▶ scheduling problem is *transformed* into a problem *without precedence constraints*, which can then be handled by a "normal" EDF scheduler

Task 4 I

EDF* - Transformation

- ▶ EDF* transforms the arrival time and deadline of every task in the following way:

Deadline:

1. Task must finish the execution time within its deadline: $f_i \leq d_i$
2. Task must not finish the execution time later than the maximum start time of its successor(s): $f_i \leq d_j - C_j$

$$\rightarrow d_i^* = \min(d_i, \min(d_j^* - C_j : J_i \rightarrow J_j))$$

Task 4 II

EDF* - Transformation

Arrival time

1. Task must start the execution not earlier than its release time: $s_j \geq r_j$
2. Task must not start execution earlier than the minimum finishing time of its predecessor(s): $s_j \geq r_i + C_i$

$$\rightarrow r_j^* = \max(r_j, \max(r_i^* + C_i : J_i \rightarrow J_j))$$

Task 4

EDF* - Precedence graph example

Given the precedences $A \rightarrow C$, $B \rightarrow C$, $C \rightarrow E$, $D \rightarrow F$, $B \rightarrow D$, $C \rightarrow F$, $D \rightarrow G$ we first draw the precedence graph:

Task 4

EDF* - Precedence graph example

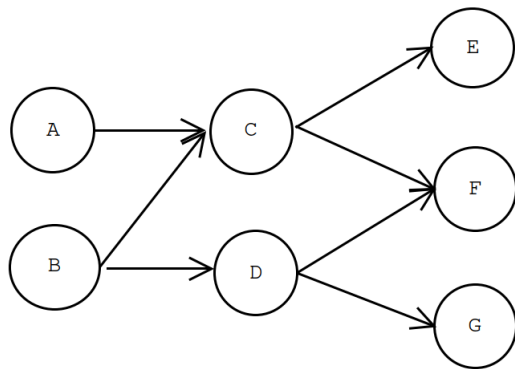


Figure 1: Task 4: precedence graph

Task 4 I

EDF* - Transformation example

- ▶ $r_A^* = r_A, r_B^* = r_B$
- ▶ $r_C^* = \max\{r_C, \max\{r_A^* + C_A, r_B^* + C_B\}\} = \max\{0, \max\{3, 2\}\} = 3$
- ▶ $r_D^* = \max\{r_D, r_B^* + C_B\} = \max\{0, 2\} = 2$
- ▶ $r_F^* = \max\{r_F, \max\{r_C^* + C_C, r_D^* + C_D\}\} = \max\{0, \max\{7, 5\}\} = 7$
- ▶ $r_E^* = \max\{r_E, r_C^* + C_C\} = \max\{0, 7\} = 7$
- ▶ $r_G^* = \max\{r_G, r_D^* + C_D\} = \max\{0, 5\} = 5$

Task 4 II

EDF* - Transformation example

- ▶ $d_E^* = d_F^* = d_G^* = 20$
- ▶ $d_C^* = \min\{d_C, \min\{d_E^* - C_E, d_F^* - C_F\}\} = \min\{20, \min\{18, 15\}\} = 15$
- ▶ $d_D^* = 15$
- ▶ $d_A^* = 11$
- ▶ $d_B^* = 11$

We now successfully have transformed the problem into one without precedence and can simply use EDF!

Task 4

EDF* - Transformation example

- The modified **release times** and **deadlines** are:

	<i>A</i>	<i>B</i>	<i>C</i>	<i>D</i>	<i>E</i>	<i>F</i>	<i>G</i>
r_i^*	0	0	3	2	7	7	5
d_i^*	11	11	15	15	20	20	20

Task 4

EDF* - Schedule

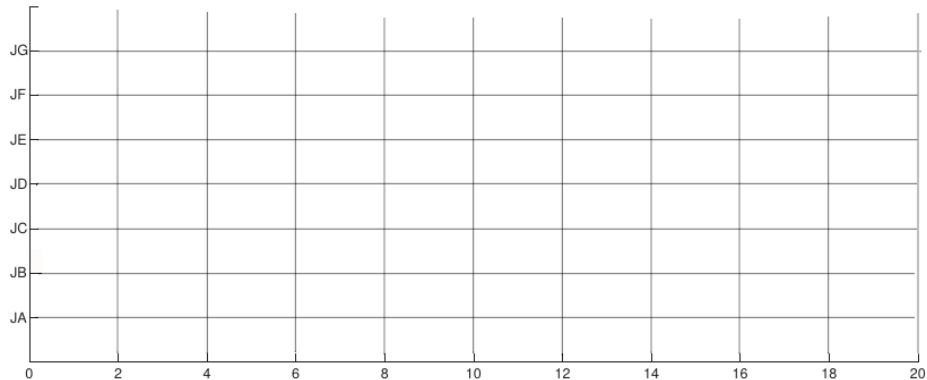


Figure 7: EDF* schedule

Task 4

EDF* - Schedule

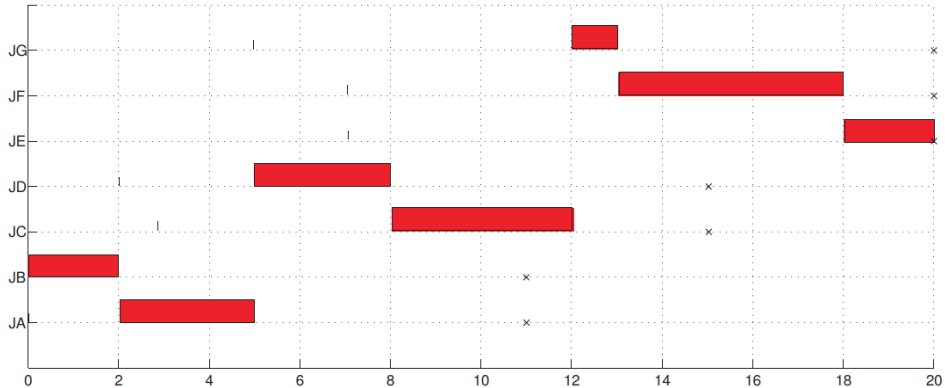


Figure 7: EDF* schedule