

Scheduling Periodic and Mixed Task Sets

Total Bandwidth Server (TBS), Rate Monotonic (RM), Polling Server

Exercise class 5

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Gliederung

Organisation

Overview

Task 1

Task 2

Task 3

Appendix

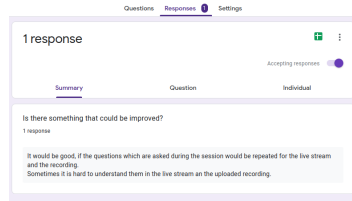
Literature

Organisation



Organisation I

- ▶ **feedback for me:** <https://forms.gle/f3YN8EFrZ1vsfPoC6>



The screenshot shows a Google Forms interface with three tabs: 'Questions', 'Responses' (selected), and 'Settings'. Under the 'Responses' tab, it says '1 response'. There are three sub-tabs: 'Summary', 'Question', and 'Individual'. The 'Summary' sub-tab is selected, showing a question: 'Is there something that could be improved?'. Below the question, it says '1 response' and displays a single response: 'It would be good, if the questions which are asked during the session would be repeated for the live stream and the recording. Sometimes it is hard to understand them in the live stream an the uploaded recording.'

- ▶ **get the slides before the exercise class:** https://github.com/matthejue/Einfuehrung_in_ESE_Tutoratsfolien_out
- ▶ **warning:** the slides often **get changed** just shortly before or after the exercise class. Both the lecture and the exercise classes are pretty **running edge**

Organisation II

- A** *A (offline)* 1:47 PM

Can one task start in one core be interrupted and resume later in another core?
- A** *A (offline)* 1:55 PM

Thanks
- Ne** *Negin* 1:57 PM

so can one specific task be executed in two cores?

Overview

Overview Scheduling I

Aperiodic Task Scheduling		
<ul style="list-style-type: none"> event-driven can arrive at any point in time Optimality: Minimize the maximum lateness of the task set precedence relations between tasks can be described through an acyclic directed graph G where tasks are represented by nodes and precedence relations by arrows. G induces a partial order on the task set 2 types of aperiodic tasks: <ul style="list-style-type: none"> aperiodic task: has a minimum inter-arrival time between consecutive instances of the task first task: maximum inter-arrival time between consecutive instances of the task cannot be bounded 		
	equal arrival times, non-preemptive	arbitrary arrival times, preemptive
independent tasks	Earliest Deadline Due (EDD) <ul style="list-style-type: none"> priority determined by $\min(D_i)$ for all remaining J_i <u>Schedulability Test:</u> $\sum_{k=1}^i C_k \leq d_i$ for each task J_i 	Earliest Deadline First (EDF) <ul style="list-style-type: none"> priority determined by $\min(d_i)$ for all remaining J_i that have already arrived (are ready) and not finished every time the arrival time of a task is reached <u>Schedulability Test:</u> $t + \sum_{k=1}^i c_k(t) \leq d_i$ for all active tasks J_i <ul style="list-style-type: none"> $c_k(t)$ is the remaining worst-case execution time of task J_k
dependent tasks	Latest Deadline First (LDF) <ul style="list-style-type: none"> priority determined by $\max(D_i)$ for all J_i without successors or whose successors have been all selected in the precedence graph inserted into the queue to be executed last at runtime, tasks are extracted from the head of the queue; the first task inserted in the queue will be executed last 	Earliest Deadline First - Star (EDF*) <ul style="list-style-type: none"> release time and deadline of individual tasks are modified such that all the precedence constraints are satisfied <ul style="list-style-type: none"> $r_i^* = \max(r_i, \max(r_i^* + C_i : J_i \rightarrow J_j))$ $d_i^* = \min(d_i, \min(d_j^* - C_j : J_i \rightarrow J_j))$ scheduling problem is transformed into a problem without precedence constraints, which can then be handled by a "normal" EDF scheduler

Overview Scheduling II

Periodic Task Scheduling • Mixed Task Sets (Periodic + Aperiodic)		
	$D_i = T_i$	$D_i \leq T_i$
	Rate Monotonic (RM) <ul style="list-style-type: none"> priority determined by $\text{rate}(T_i)$ for all remaining J_i Schedulability Test 1: $U = \sum_{i=1}^n \frac{C_i}{T_i} \leq n \left(2^{1/n} - 1 \right)$ (sufficient but not necessary) Schedulability Test 2: same as Schedulability Test 1 for DM 	Deadline Monotonic (DM) <ul style="list-style-type: none"> priority determined by $\text{rate}(D_i)$ for all remaining J_i Schedulability Test 1: $\sum_{i=1}^n \frac{C_i}{D_i} \leq n \left(2^{1/n} - 1 \right)$ (sufficient but not necessary) Schedulability Test 2: for all tasks τ_i $R_i = C_i + \sum_{j=1}^{i-1} \left\lceil \frac{B_j}{T_j} \right\rceil C_j$ that satisfies: $R_i \leq C_i + \sum_{j=1}^{i-1} \left\lceil \frac{B_j}{T_j} \right\rceil C_j$ and $R_i \leq D_i$ (necessary and sufficient)
	Background Scheduling (BS) <ul style="list-style-type: none"> aperiodic tasks scheduled after FCFS when no periodic task ready to execute 	
	Putting Server (PS): <ul style="list-style-type: none"> PS scheduled as periodic task, it serves any pending aperiodic requests until its capacity (execution time) C_p is exhausted if no aperiodic requests are pending, PS suspends itself until the beginning of its next period, and the budget originally allocated for aperiodic service is freed up and assigned to periodic tasks Schedulability Test 1: $\sum_{i=1}^n \frac{C_i}{T_i} + \frac{C_p}{T_p} \leq (n+1) \left(2^{1/(n+1)} - 1 \right)$ (sufficient but not necessary) Schedulability Test 2: $C_p \leq C_s$ and $2T_p \leq D_p$ (necessary and sufficient, using completion time) Schedulability Test 3: $T_p + \left\lceil \frac{C_p}{C_s} \right\rceil T_s \leq D_p$ (necessary and sufficient, not using completion time) 	
	Harvest Deadline First (HDF) <ul style="list-style-type: none"> Schedulability Test: $\sum_{i=1}^n \frac{C_i}{T_i} = U \leq 1$ 	
	Total Bandwidth Server (TBS) <ul style="list-style-type: none"> $d_k = \max \{ \tau_k, d_{k-1} \} + \frac{C_k}{U_s}$ Schedulability Test: $U_s + U_a \leq 1$ (necessary and sufficient) 	Harvest Deadline First (HDF) <ul style="list-style-type: none"> Schedulability Test: Necessary. Have the first computing system: pre-emptive scheduling
	Dynamic priority	

Overview Scheduling III

III

Important Parameters

Periodic Tasks

- Γ : set of periodic tasks
- τ_i : periodic task
- $\tau_{i,j}$: j -th instance of task τ_i
- $r_{i,j}, s_{i,j}, f_{i,j}, d_{i,j}$: release time, start time, finishing time, and absolute deadline of the j -th instance of task τ_i
- Φ_i : phase of task τ_i (release time of first instance)
- C_i : worst-case execution time of task τ_i
- T_i : period of task τ_i
- D_i : relative deadline of task τ_i

Real-Time Tasks

- Arrival time** a_i or **release time** r_i is the time at which a task becomes ready for execution.
- Start time** s_i is the time at which a task starts its execution.
- Execution time** C_i is the time needed by the processor to execute a task without interruption.
- Finishing time** f_i is the time at which a task finishes its execution.
- Absolute deadline** d_i is the time by which a task should be completed.

- Relative deadline** D_i is the difference between the absolute deadline and the arrival time of a task, that is, $D_i = d_i - a_i$.
- Response time** R_i is the difference between the finishing time and the arrival time of a task, that is, $R_i = f_i - a_i$.
- Lateness** $L_i = f_i - d_i$ represents the delay of a task completion with respect to its deadline, that is, $L_i \leq 0$ if a task completes within its deadline.
- Forfeiture** $F_i = \max(0, L_i)$ is the time a task pays after its deadline.
- Slack** $X_i = d_i - a_i - C_i$ is the maximum time a task can be delayed on its execution in order to complete within its deadline.

Critical Instant

- Definition:** A critical instant of a task is the release time r_i that produces the largest response time R_i , that is, the largest difference between release time r_i and finishing time f_i .
- Lemma:** For any task, a critical instant occurs whenever the task is released simultaneously with all higher-priority tasks.

Overview Scheduling IV

Important Formulas	
<ul style="list-style-type: none"> release time: $r_{i,j} = \Phi_i + (j-1)T_i$ absolute deadline: $d_{i,j} = r_{i,j} + D_i = \Phi_i + (j-1)T_i + D_i$ <p>and if</p> $D_i = T_i$ <p>then</p> $d_{i,j} = \Phi_i + jT_i$ rate: $\frac{1}{T_i}$ processor utilization: $U_p = \sum_{i=1}^n \frac{C_i}{T_i}$ 	
Metrics to compare schedules <ul style="list-style-type: none"> Average response time: $\bar{r}_r = \frac{1}{n} \sum_{i=1}^n (f_i - a_i)$ Total completion time: $t_c = \max_i(f_i) - \min_i(a_i)$ Weighted sum of response times: $t_w = \frac{\sum_{i=1}^n w_i(f_i - a_i)}{\sum_{i=1}^n w_i}$ Maximum lateness: $L_{\max} = \max_i(f_i - d_i)$ Number of late tasks: $N_{\text{late}} = \sum_{i=1}^n \text{miss}(f_i)$ where $\text{miss}(f_i) = \begin{cases} 0 & \text{if } f_i \leq d_i \\ 1 & \text{otherwise} \end{cases}$ 	<p>Only these metrics are useful to evaluate real-time schedules, as they involve task deadlines.</p>

Task 1

Task 1 I

Earliest Deadline First (EDF) and Total Bandwidth Server (TBS)

Task 1.1:

	τ_1	τ_2	τ_3
C_i	1	1	2
T_i	3	5	13

- ▶ what can be the maximum value of U_s such that the whole set (i.e. periodic tasks and the TBS) is schedulable with EDF?

Task 1 II

Earliest Deadline First (EDF) and Total Bandwidth Server (TBS)

Requirements 1.1:

Schedulability test:

Given a set of n periodic tasks with processor utilization U_p and a total bandwidth server with utilization U_s , the whole set is schedulable by EDF if and only if

$$U_p + U_s \leq 1$$

► *processor utilization factor:* $U = \sum_{i=1}^n \frac{C_i}{T_i}$

Task 1 III

Earliest Deadline First (EDF) and Total Bandwidth Server (TBS)

Solution 1.1:

- ▶ *Maximum utilization of the Total Bandwidth Server:*

$$U_{s,\max} = 1 - U_p = 1 - \left(\frac{1}{3} + \frac{1}{5} + \frac{2}{13}\right) = \frac{61}{195} \approx 0.3128$$

Task 1 I

Earliest Deadline First (EDF) and Total Bandwidth Server (TBS)

Task 1.2:

- ▶ construct EDF-Schedule
- ▶ assume $U_s = 0.25$
- ▶ three aperiodic requests served by TBS:

	J_4	J_5	J_6
r_i	0	15	10
C_i	2	1	1

- ▶ arrival time of first instance is 0:

	τ_1	τ_2	τ_3
a_i	0	0	0
C_i	1	1	2
T_i	3	5	13

Task 1 II

Earliest Deadline First (EDF) and Total Bandwidth Server (TBS)

Requirements 1.2:

- ▶ $d_i = \max(r_i, d_{k-1}) + \frac{C_k}{U_s}$
 - ▶ d_{k-1} denotes the previously calculated deadline ($k - 1$ means the predecessor in the ordering according to the release time)

Solution 1.2:

- ▶ order the aperiodic tasks by **increasing release time** r_i : J_4, J_6, J_5
- ▶ calculate the deadlines:
 - ▶ $d_4 = \max(r_4, d_0) + \frac{2}{0.25} = \max(0, 0) + 8 = 8$
 - ▶ $d_6 = \max(r_6, d_4) + \frac{1}{0.25} = \max(10, 8) + 4 = 14$
 - ▶ $d_5 = \max(r_5, d_6) + \frac{1}{0.25} = \max(15, 14) + 4 = 19$
- ▶ periodic tasks already ordered by **increasing period**: t_i : τ_1, τ_2, τ_3

Task 1 III

Earliest Deadline First (EDF) and Total Bandwidth Server (TBS)

Solution 1.2:

	τ_1	τ_2	τ_3
a_i	0	0	0
C_i	1	1	2
T_i	3	5	13
	J_4	J_5	J_6
r_i	0	15	10
d_i	8	19	14
C_i	2	1	1

Task 1 IV

Earliest Deadline First (EDF) and Total Bandwidth Server (TBS)

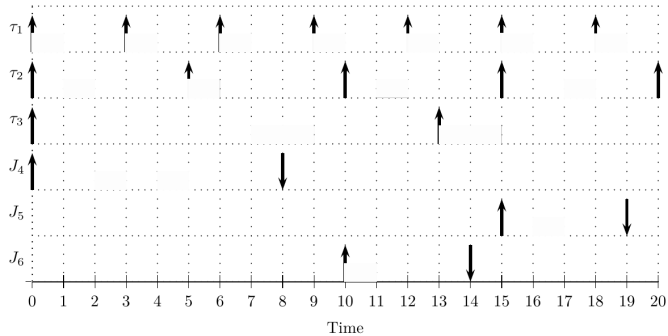


Figure 1: EDF schedule solution for Task 1

Task 1 V

Earliest Deadline First (EDF) and Total Bandwidth Server (TBS)

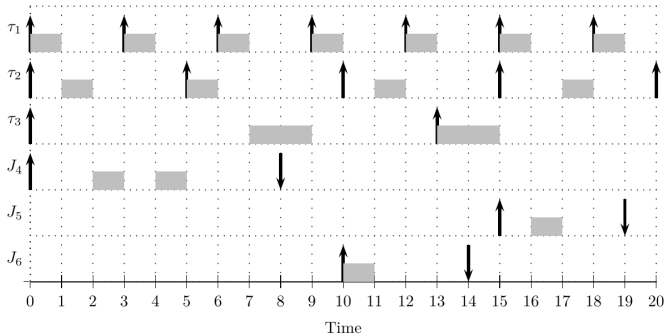


Figure 1: EDF schedule solution for Task 1

Task 2

Task 2 I

Schedulability Test for Fixed Priorities – Rate Monotonic (RM)

Task 2.1:

- ▶ task-set schedulable with **RM**
- ▶ using **sufficient** test

	τ_1	τ_2	τ_3
C_i	1	3	2
T_i	3	8	9

Requirements 2.1:

- ▶ $U = \sum_{i=1}^n \frac{C_i}{T_i} \leq n \left(2^{1/n} - 1 \right)$, U is the **fraction** of the **processor time** spent on **executing task set**

Task 2 II

Schedulability Test for Fixed Priorities – Rate Monotonic (RM)

Solution 2.1:

- ▶ $\frac{1}{3} + \frac{3}{8} + \frac{2}{9} = 0.93 \leq 3(2^{\frac{1}{3}} - 1) = 0.78 \quad \times$
- ▶ condition is *not necessary*, hence we *don't know* whether the task set is schedulable with *RM* or not

Task 2 III

Schedulability Test for Fixed Priorities – Rate Monotonic (RM)

Task 2.2:

- ▶ task-set schedulable with RM
- ▶ using necessary test

	τ_1	τ_2	τ_3
C_i	1	3	2
T_i	3	8	9

Task 2 IV

Schedulability Test for Fixed Priorities – Rate Monotonic (RM)

Requirements 2.2:

- ▶ guarantee that *all* the tasks can be scheduled in *any possible instance*
- ▶ in particular, if a task can be scheduled in its *critical instances*, then the schedulability guarantee condition holds
 - ▶ a *critical instance* of a task occurs whenever the task is *released simultaneously* with all higher priority tasks
- ▶ **Schedulability Test:** For all tasks τ_i smallest R_i that satisfies

$$R_i = C_i + \sum_{j=1}^{i-1} \left\lceil \frac{R_i}{T_j} \right\rceil C_j \text{ and } R_i \leq D_i \text{ (necessary and sufficient)}$$

Task 2 V

Schedulability Test for Fixed Priorities – Rate Monotonic (RM)

Requirements 2.2:

```

DM_guarantee ( $\Gamma$ ) {
  for (each  $\tau_i \in \Gamma$ ) {
     $I_i = \sum_{k=1}^{i-1} C_k$ ;
    do {
       $R_i = I_i + C_i$ ;
      if ( $R_i > D_i$ ) return(UNSCHEDULABLE);
       $I_i = \sum_{k=1}^{i-1} \left\lceil \frac{R_i}{T_k} \right\rceil C_k$ ;
    } while ( $I_i + C_i > R_i$ );
  }
  return(SCHEDULABLE);
}
    
```

Task 2 VI

Schedulability Test for Fixed Priorities – Rate Monotonic (RM)

Solution 2.2:

- ▶ The tasks are first ordered by their priorities: τ_1 , τ_2 and τ_3
 - ▶ In this case the tasks are already ordered

▶ τ_3 :

- ▶ $R_3^0 = C_3 = 2 \quad I_3^0 = \left\lceil \frac{2}{3} \right\rceil 1 + \left\lceil \frac{2}{8} \right\rceil 3 = 1 + 3 = 4 \quad 4 + 2 \neq 2$
- ▶ $R_3^1 = 4 + 2 = 6 \quad I_3^1 = \left\lceil \frac{6}{3} \right\rceil 1 + \left\lceil \frac{6}{8} \right\rceil 3 = 2 + 3 = 5 \quad 5 + 2 \neq 6$
- ▶ $R_3^2 = 5 + 2 = 7 \quad I_3^2 = \left\lceil \frac{7}{3} \right\rceil 1 + \left\lceil \frac{7}{8} \right\rceil 3 = 3 + 3 = 6 \quad 6 + 2 \neq 7$
- ▶ $R_3^3 = 6 + 2 = 8 \quad I_3^3 = \left\lceil \frac{8}{3} \right\rceil 1 + \left\lceil \frac{8}{8} \right\rceil 3 = 3 + 3 = 6 \quad 6 + 2 = 8 \dots \checkmark$
 (since $R_3 = 8 \leq T_3 = 9$)

Task 2 VII

Schedulability Test for Fixed Priorities – Rate Monotonic (RM)

Solution 2.2:



▶ τ_2 :

▶ $R_2^0 = C_2 = 3 \quad I_2^0 = \lceil \frac{3}{3} \rceil 1 = 1 \quad 1 + 3 \neq 3$

▶ $R_2^1 = 1 + 3 = 4 \quad I_2^1 = \lceil \frac{4}{3} \rceil 1 = 2 \quad 2 + 3 \neq 4$

▶ $R_2^2 = 2 + 3 = 5 \quad I_2^2 = \lceil \frac{5}{3} \rceil 1 = 2 \quad 2 + 3 = 5 \dots \checkmark$

(since $R_2 = 5 \leq T_2 = 8$)

▶ τ_1 :

▶ $R_1^0 = C_1 = 1 \quad I_1^0 = 0 \quad 0 + 1 = 1 \dots \checkmark$

(since $R_1 = 1 \leq T_1 = 3$)

▶ *The necessary and sufficient test succeeds. This means that the task set is schedulable with RM.*

Task 2 VIII

Schedulability Test for Fixed Priorities – Rate Monotonic (RM)

Task 2.3:

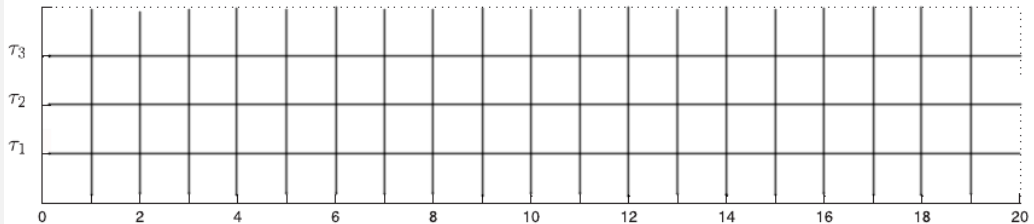
- ▶ first job of each task arrives at time 0
- ▶ construct schedule for interval $[0, 20]$

	τ_1	τ_2	τ_3
r_i	0	0	0
C_i	1	3	2
T_i	3	8	9

Task 2 IX

Schedulability Test for Fixed Priorities – Rate Monotonic (RM)

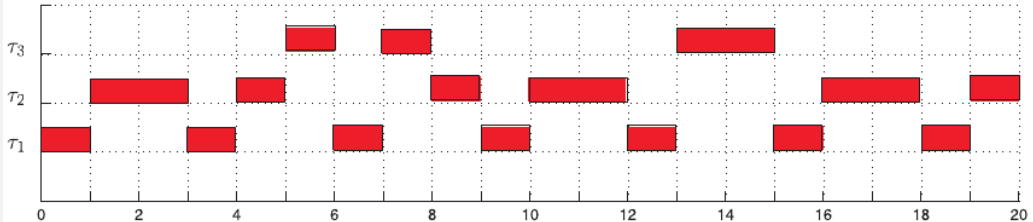
Solution 2.3:



Task 2 X

Schedulability Test for Fixed Priorities – Rate Monotonic (RM)

Solution 2.3:



Task 3

Task 3

Scheduling with Polling Server

Task 3.1:



	τ_1	τ_2	τ_3
C_i	2	2	2
D_i	6	8	16
T_i	6	8	16

- ▶ In addition to the above periodic tasks, we have an aperiodic job J_a with $C_a = 1$ and relative deadline D_a . Let $T_s = 25$ and $C_s = 1$ respectively, where T_s denotes the period and C_s the computing time (or capacity) of the polling server (PS)
- ▶ Compute the minimum relative deadline of J_a which is guaranteed not to be missed, that is, its aperiodic guarantee.

Task 3 I

Scheduling with Polling Server

Solution 3.1:

- ▶ *As reminder: The polling server itself acts like a **periodic** task, that uses its capacity to serve **aperiodic** tasks. If the polling server has the current highest priority it begins to serve any pending aperiodic requests within the limits of its capacity.*
- ▶ *If there are no pending aperiodic tasks at that time it **suspends** its entire capacity! (until the beginning of the next period)*
- ▶ *Therefore, the worst possible case occurs when J_a appears slightly later than the polling server checks for pending aperiodic tasks. Meaning, the polling server suspends its capacity and J_a has to wait $T_s + \lceil \frac{C_a}{C_s} \rceil T_s = (1 + \lceil \frac{C_a}{C_s} \rceil) T_s$*

Task 3 II

Scheduling with Polling Server

Solution 3.1:



- ▶ To guarantee to not miss the deadline D_a the condition $(1 + \lceil \frac{C_a}{C_s} \rceil) T_s \leq D_a$ needs to hold.
- ▶ Entering the values for the exercise gives us $D_a = (1 + \frac{1}{1}) \cdot 25 = 50$

Sidenote 🔍

Note that the above computation of course only holds if the RM schedule meets all the deadlines.

Task 3

Scheduling with Polling Server

Task 3.2:

Using the sufficient test of RM, test if the polling server of task 3.1 is schedulable along with the periodic task-set.

Requirements 3.2:

► *Schedulability Test:*
$$\sum_{i=1}^n \frac{C_i}{T_i} + \frac{C_s}{T_s} \leq (n+1) \left[2^{1/(n+1)} - 1 \right]$$

Task 3

Scheduling with Polling Server

Solution 3.2:



- ▶ We have already seen the sufficient but not necessary condition of $\sum_{i=1}^n \frac{C_i}{T_i} \leq n(2^{1/n} - 1)$ for rate monotonic scheduling.
- ▶ The convenient part about our polling server: we can just treat it as an additional *periodic* task.
- ▶ Therefore, the same test offers us a sufficient condition for rate monotonic scheduling with a polling server! Simply increase from n tasks to $(n+1)$
- ▶ Check: $\sum_{i=1}^n \frac{C_i}{T_i} + \underbrace{\frac{C_s}{T_s}}_{\text{server task}} \leq (n+1)(2^{1/(n+1)} - 1)$

Task 3

Scheduling with Polling Server

Solution 3.2:

► *Putting in values we obtain:*

$$\frac{2}{6} + \frac{2}{8} + \frac{2}{16} + \frac{1}{25} \leq (3 + 1)(2^{\frac{1}{3+1}} - 1) \Leftrightarrow \frac{449}{600} \leq \sqrt[4]{256} \cdot \sqrt[4]{2} - 4 \Leftrightarrow 0.75 \leq 0.76.$$

Since this is true, we know (as this is a sufficient condition) that the RM schedule meets all deadlines!

Appendix

Overview periodic Task Scheduling

	Deadline equals period	Deadline smaller than period
static priority	RM (rate-monotonic)	DM (deadline-monotonic)
dynamic priority	EDF	EDF

Overview Aperiodic Task Scheduling

Schedulability test

	Deadline equals period ($D_i = T_i$)	Deadline smaller than period ($D_i \leq T_i$)
static priority	<p>(1) $\sum_{i=1}^n \frac{C_i}{T_i} \leq n \left(2^{1/n} - 1 \right)$ (sufficient but not necessary) (2) same as Schedulability Test 2 for DM</p>	<p>(1) $\sum_{i=1}^n \frac{C_i}{D_i} \leq n \left(2^{1/n} - 1 \right)$ (sufficient but not necessary) (2) smallest R_i that satisfies $R_i = C_i + \sum_{j=1}^{i-1} \left\lceil \frac{R_j}{T_j} \right\rceil C_j$ and $R_i \leq D_i$ for all tasks τ_i (necessary and sufficient)</p>
dynamic priority	<p>$\sum_{i=1}^n \frac{C_i}{T_i} = U \leq 1$ (necessary and sufficient)</p>	<p>→ Buttazzo, <i>Hard real-time computing systems: predictable scheduling algorithms and applications</i></p>

Mixed Task Sets

- ▶ So far: we differentiated between **periodic** and **aperiodic** tasks.
- ▶ Now: Consider a **mixed** task set!
- ▶ We want to be able to find a schedule when there's both **periodic** and **aperiodic** tasks.

Schedulability tests

Sufficient? Necessary?

- ▶ We're interested in whether a given problem can be scheduled by algorithms.
- ▶ Depending on the algorithm we can derive sufficient and necessary conditions.

Sufficient: If $A \implies B$ then A is a sufficient condition for B.

Necessary: If $B \implies A$ then A is a necessary condition for B.

- ▶ A necessary and sufficient condition means, both statements are logically equivalent.

Schedulability tests

Utilization

Different kind of utilizations also play a big role in our analysis. We introduced the **processor utilization factor** $U = \sum_{i=1}^n \frac{C_i}{T_i}$ and later on U_s as the server utilization.

(More about servers later)

RM - Rate Monotonic Scheduling

Schedulability

- ▶ RM is optimal among all fixed-priority assignments in the sense that no other fixed-priority algorithm can schedule a task set that cannot be scheduled by RM.
- ▶ As in the lecture, we have $\sum_{i=1}^n \frac{C_i}{T_i} \leq n(2^{1/n} - 1)$ as a **sufficient** but not **necessary** condition.

RM(PS) - Rate Monotonic Polling Server

- ▶ One way to handle both periodic and aperiodic tasks is to use a so called server.
- ▶ This PS (Polling Server) acts as a periodic task (meaning it is instantiated at regular intervals T_s) whose job it is to, once it has the highest priority, serve any pending aperiodic requests within the limits of a server capacity C_s .
- ▶ Since we introduce yet another periodic task, the schedulability analysis simply is the same as normal *RM* with one additional task. Again, we have the **sufficient** but not **necessary** condition:
$$\frac{C_s}{T_s} + \sum_{i=1}^n \frac{C_i}{T_i} \leq (n+1)(2^{1/(n+1)} - 1)$$

Literature

Bücher



[Buttazzo, Giorgio C.](#) *Hard real-time computing systems: predictable scheduling algorithms and applications*. Vol. 24. Springer Science & Business Media, 2011.