Aperiodic Scheduling

Earliest Deadline Due, Latest Deadline First, Earliest Deadline First

Exercise class 4

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In cooperation with:
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Based on the lecture of: Marco Zimmerling

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University of Freiburg, Chair for Embedded Systems

Gliederung

Organisation

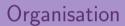
Overview Aperiodic Task Scheduling

Task 1

Task 2

Task 3

Task 4





Organisation I

► feedback for me: https://forms.gle/f3YN8EFrZ1vsfPoC6

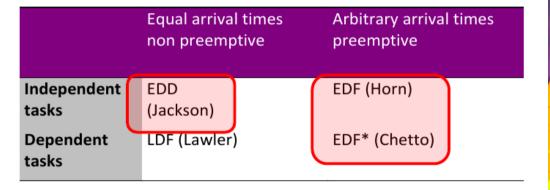
Feedback for Introduction to ESE Tutor Jürgen 🗀 🛕			
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Overview Aperiodic Task Scheduling

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Overview Aperiodic Task Scheduling



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Earliest Deadline Due

Task 1.1:

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	J_1	J_2	J_3	J_4
C_i	3	6	2	4
D_i	8	15	3	11

$$(\forall J_i: r_i=0)$$

Requirements 1.1:



8/33

- non-preemptive
- tasks have same arrival times (synchronous arrivals)
- tasks are independent
- $ightharpoonup min(D_i)$ for all remaining J_i
- minimizing the maximum lateness

Earliest Deadline Due

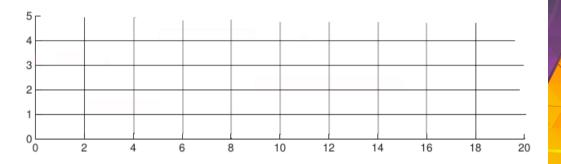


Figure 1: EDD schedule.

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Earliest Deadline Due

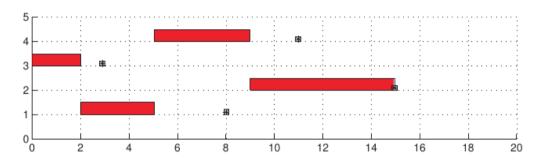


Figure 1: EDD schedule.



Latest Deadline First

Task 2.1:

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12/33

	J_1	J_2	J_3	J_4	J_5	J_6	J_7	J_8
C_i	3	4	2	3	3	2	2	1
D_i	5	8	11	15	12	18	19	20

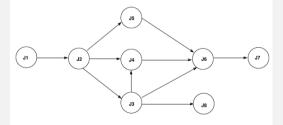


Figure 2: Precedence graph.

Latest Deadline First

Requirements 2.1:

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- ▶ is non-preemptive
- synchronous task activations
- tasks are dependent, use precedence graph, going from tail to head
- ► max(D_i) for all tasks J_i without successors or whose successors have been all selected in the precedence graph inserted into the queue to be executed last
- ▶ at runtime, tasks are extracted from the head of the queue: the first task inserted in the queue will be executed last
- minimizes the maximum lateness

Latest Deadline First

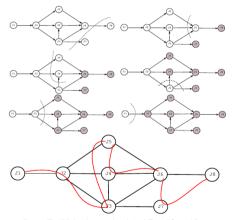


Figure 4: The LDF algorithm proceeds as depicted (figures left to right)

Latest Deadline First

queue of tasks: (, , , , , , , , , , ,)

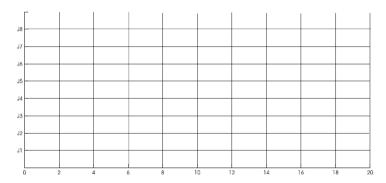


Figure 3: LDF schedule.

Latest Deadline First

• queue of tasks: $(J_1, J_2, J_3, J_5, J_4, J_6, J_7, J_8)$

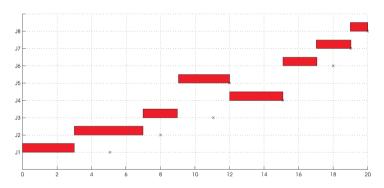


Figure 3: LDF schedule.

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Earliest Deadline First

Task 3.1:



		J_1	J_2	J_3	J_4	J_5
a_i		0	2	0	8	13
C_i	į	3	1	6	2	3
d_i		16	7	8	11	18

Requirements 3.1:



18 / 33

- ▶ is preemptive
- ► arbitrary arrival times
- the tasks are independent
- minimizes the maximum lateness.
- $ightharpoonup min(D_i)$ for all remaning tasks J_i that have already arrived (are ready) and not finished every time the arrival time of a task is reached

Earliest Deadline First

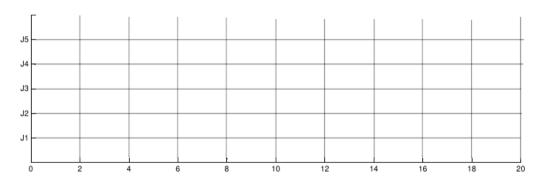


Figure 5: EDF schedule

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Earliest Deadline First

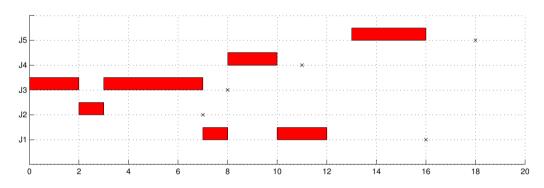


Figure 5: EDF schedule

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Earliest Deadline First

Task 3.2:

7

- ▶ at time t = 3, a new task J_x arrives with execution time $C_x = 2$ and deadline $d_x = 10$.
- ▶ still guarantee the schedulability of task set?

Requirements 3.2:



$$\forall i=1,\ldots,n \quad t+\sum_{k=1}^{i}c_k(t)\leq d_i$$





Task 4 I

EDF* - Example

Task 4.1:

1

Given tasks A, B, C, D, E, F, G with precedences $A \rightarrow C, B \rightarrow C, C \rightarrow E, D \rightarrow F, B \rightarrow D, C \rightarrow F, D \rightarrow G.$

All tasks arrive at time $t_0 = 0$, have a common deadline d = 20 and the following execution times:

	Α	В	С	D	Е	F	G
C_i	3	2	4	3	2	5	1

We will now prepare the tasks for EDF*

EDF*

Requirements 4.1:



24 / 33

- non-preemptive
- ► arbitrary arrival times
- ► tasks are dependent
- release time and deadline of individual tasks are modified such that all the precedence constraints are satisfied
- scheduling problem is transformed into a problem without precedence constraints, which can then be handled by a "normal" EDF scheduler

Task 4 I

EDF* - Transformation

▶ EDF* transforms the arrival time and deadline of every task in the following way:

Deadline:

- 1. Task must finish the execution time within its deadline: $f_i \leq d_i$
- 2. Task must not finish the execution time later than the maximum start time of its successor(s): $f_i \leq d_j C_j$
- $\rightarrow d_i^* = min(d_i, min(d_j^* C_j : J_i \rightarrow J_j))$

Task 4 II

EDF* - Transformation

Arrival time

- 1. Task must start the execution not earlier than its release time: $s_j \geq r_j$
- 2. Task must not start execution earlier than the minimum finishing time of its predecessor(s): $s_j \ge r_i + C_i$
- $\rightarrow r_j^* = max(r_j, max(r_i^* + C_i : J_i \rightarrow J_j))$

EDF* - Precedence graph example

Given the precedences $A \to C$, $B \to C$, $C \to E$, $D \to F$, $B \to D$, $C \to F$, $D \to G$ we first draw the precedence graph:

EDF* - Precedence graph example

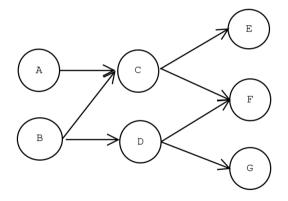


Figure 1: Task 4: precedence graph

Task 4 I

EDF* - Transformation example

- $r_A^* = r_A, r_B^* = r_B$
- $r_C^* = \max\{r_C, \max\{r_A^* + C_A, r_B^* + C_B\}\} = \max\{0, \max\{3, 2\}\} = 3$
- $r_D^* = \max\{r_D, r_B^* + C_B\} = \max\{0, 2\} = 2$
- $r_F^* = \max\{r_F, \max\{r_C^* + C_C, r_D^* + C_D\}\} = \max\{0, \max\{7, 5\}\} = 7$
- $r_E^* = \max\{r_E, r_C^* + C_C\} = \max\{0, 7\} = 7$
- $r_G^* = \max\{r_G, r_D^* + C_D\} = \max\{0, 5\} = 5$

Task 4 II

EDF* - Transformation example

$$d_E^* = d_F^* = d_G^* = 20$$

$$d_C^* = \min\{d_C, \min\{d_E^* - C_E, d_F^* - C_F\}\} = \min\{20, \min\{18, 15\}\} = 15$$

$$d_D^* = 15$$

$$d_A^* = 11$$

$$d_B^* = 11$$

We now successfully have transformed the problem into one without precedence and can simply use EDF!

EDF* - Transformation example

► The modified release times and deadlines are:

	A	B	C	D	E	F	G
r_i^*	0	0	3	2	7	7	5
d_i^*	11	11	15	15	20	20	20

EDF* - Schedule

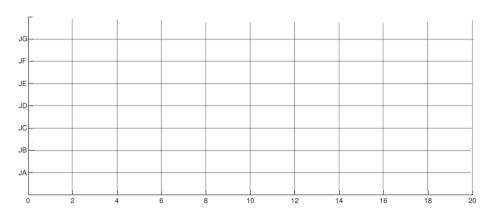


Figure 7: EDF* schedule

EDF* - Schedule

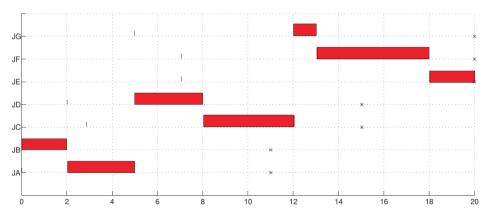


Figure 7: EDF* schedule

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