

Proof: $(\text{append} (\text{append } xs \ ys) \ zs) == (\text{append } xs \ (\text{append } ys \ zs))$

Case #1: nil case

Given $xs = '()$, $ys = '()$

$(\text{append} (\text{append } xs \ ys) \ zs)$

$= \{ \text{By assumption } xs = '() \}$

$(\text{append} (\text{append } '() \ ys) \ zs)$

$= \{ \text{append-nil law} \}$

$(\text{append } ys \ zs)$

$= \{ \text{append-nil law} \}$

$(\text{append } '() \ (\text{append } ys \ zs))$

$= \{ \text{substitution, } xs = '() \}$

$(\text{append } xs \ (\text{append } ys \ zs))$

Case #2: Inductive case for non-empty lists

Inductive hypothesis: Assume $(\text{append} (\text{append } xs \ ys) \ zs)$

$= (\text{append } xs \ (\text{append } ys \ zs))$ holds true for all xs, ys, zs

$(\text{append} (\text{append } xs \ ys) \ zs))$

$= \{ \text{by assumption, } xs = (\text{cons } a \ as) \}$

$(\text{append} (\text{append } (\text{cons } a \ as) \ ys) \ zs)$

$= \{ \text{append-cons rule} \}$

$(\text{append} (\text{cons } a \ (\text{append } as \ ys) \ zs))$

$= \{ \text{append-cons rule} \}$

$(\text{cons } a \ (\text{append} (\text{append } as \ ys) \ zs))$

$= \{ \text{inductive hypothesis} \}$

$(\text{cons } a \ (\text{append } as \ (\text{append } ys \ zs)))$

$= \{ \text{append-cons rule} \}$

$(\text{append} (\text{cons } a \ as) \ (\text{append } ys \ zs))$

$$= \{ \text{by assumption, cons } a \text{ as } = xs \\ (\text{append } xs (\text{append } ys \text{ } zs)) \}$$

Therefore, by induction, $(\text{append } (\text{append } xs \text{ } ys) \text{ } zs)$
 $= (\text{append } xs (\text{append } ys \text{ } zs))$ for all xs, ys, zs .