## **Purpose and Efficacy of Test Cases**

For my testing file, I copied the testing file for project 4 and then altered or added tests when I felt necessary. So that I do not repeat myself from the previous writeup, I will only explain these changes and additions that I made, but I will still make note of all of the tests for completeness.

First, I kept the preliminary empty tree test that included testing the string and get\_height() methods, except I also included a to\_list() test as well. This would make sure that all three methods handle the empty tree case properly.

### **Get and Set Height**

In this section, I included all of the same tests with the addition of one more test aimed towards verifying \_\_balance(). This test was performed on a tree with seven elements where each of the elements that is inserted is lower than the first value. This way, \_\_balance() must keep calling \_\_rotation() to eventually reach a tree of height 3.

Additionally, the heights of some of the trees in this section needed to be changed, as \_\_balance() caused them to perform one or more rotations. For example, there were left and right leaning trees with three elements and a height of 3 that changed to a height of 2 after rotating. By including these tests, I am making sure that the height is always updated whenever necessary. This also verifies the height aspect of other methods, including \_\_rins(), insert\_element(), \_\_balance(), and \_\_rotate().

#### **Insert**

With my insert tests, I did not need to add any tests. In my previous tests, there were already checks for both single and double rotations on the left and right side, so there was no need to repeat anything. However, since this is only checking insertions, I only verified that the insertions themselves worked, and must wait until the traversal section to verify that the rotations did in fact happen.

In terms of changes, when checking situations that raise ValueErrors, since I check the height afterwards, I also had to change some of the heights.

#### Remove

Unlike my insert section, I did have to add six tests to my removal section. This is because I had no removal tests that resulted in single or double rotations, so I made two for each, one for the left side and one for the right side. Additionally, two of the tests that I added had floaters that were not None to make it easier to visualize any mistake that may happen with the floaters. Again, this only checks to make sure that removal works in a situation where a rotation happens and does not actually check if the rotations happened, so that will be done in the traversal section as well. I did not need to change any heights for the ValueError subsection.

#### **Traversals**

This section is where I made the most changes. In addition to the three traversals, I also added the to\_list traversal with all of the same tests as the others. The only difference these tests had was the assertEqual statement.

With the other traversals, I did decide to add several tests, however, as I felt that maybe my testing was incomplete from project 4. For each traversal, I added eight tests. These all focused on removal and included removing a leaf from a tree with three nodes, removing a leaf from a tree with four nodes to cause a single rotation, removing a leaf from a tree with four nodes to cause a double rotation, and removing a leaf from a tree with four nodes to cause a single rotation with a non-None floater. Each of these tests had a left variant and a right variant as well.

The removal from a tree with three leaves was the test I felt could have been included in project 4, as it would uniquely identify a tree after a removal, which was not included previously. The remainder of the tests were to ensure that the rotations did happen as they were expected to, and the floaters relocated to the correct place. Since the combination of the traversals can uniquely identify a specific tree, I used this to confirm that that was the case. Of course, I did include these tests for to\_list as well, verifying my testing for all four of the traversals.

# **Analysis of the Worst-Case Asymptotic Performance**

BST_Node:
init(): This method is unchanged from project 4.
Binary_Search_Tree:
init(): This method is unchanged from project 4.

set_hei	ht(): This method is unchanged from project 4.	
	e_height(): This method checks if the node is None and then her `0' or the node's height, all of which can be done in constant	t
of t. Since constant t	(): This method relies on accessing the children and grandchildr we store the children as attributes, we can access them in me. Additionally, this method relies onset_height(), however also a constant time operation, so is the entirerotation()	
conditiona get_no	(): This method involves usingget_node_height() with s to then determine the necessary rotations. Since e_height() androtation() are both constant time operations,balance().	
returns se ensures th it's perfor	This method is almost the exact same as project 4, except it fbalance(node) instead of node. However, sincebalance() at the tree stays within a height difference of 1 between childre nance changes from linear to logarithmic time. Additionally, () can be done in constant time, so it does not make the ce worse.	
time, inse	nent(): Just like howrins() changed from linear to logarithmi t_element() does too, as it is only dependent onrins() and me assignment.	C
	Sincerrem() still has the same worst case asrins(), also changes to logarithmic time because ofbalance().	
	ement(): Again, like how insert_element() relies onrins(), ement() relies on rrem(), so it also performs in logarithmic time	€.
rin_ord	er(): This method is unchanged from project 4.	
in_order()	This method is unchanged from project 4.	
traversals value to the a linear time because it	):rto_list() is implemented slightly differently from the other Instead of using string concatenation, it appends each node's e given list. Appending is a constant time operation as opposed ne operation, sorto_list() has linear time performance. This i must still traverse through the tree and visit each node, but sin e stored as attributes, visiting each node is only linear time.	l to

```
to_list(): Since to_list() creates an empty list and then calls __rto_list(), it has the same performance as __rto_list(), which is linear time.

__rpre_order(): This method is unchanged from project 4.

pre_order(): This method is unchanged from project 4.

__rpost_order(): This method is unchanged from project 4.

post_order(): This method is unchanged from project 4.

get_height(): This method is unchanged from project 4.

__str__(): This method is unchanged from project 4.
```

### **Fraction:**

For my fraction main section, in my first test, I created twenty-five Fraction objects to illustrate my sorting. In order to sort them, I had to first insert each fraction into the tree and then call to\_list. Since insert\_element() has logarithmic performance and I must insert each fraction one at a time, this step has a performance of O(nlog(n)).

Then, I called to\_list() and assigned the return value to a variable, in\_order\_rep. Since to\_list() has linear time performance and it was the only thing I did in this step, then as a whole this step was also done in linear time. Of course, I then printed out my results, showing first the original fraction order and then their sorted order.

In order to demonstrate  $\underline{\phantom{a}}$  eq $\underline{\phantom{a}}$ , I had a second test that had two fraction objects:  $\frac{2}{3}$  and  $\frac{4}{6}$ . Since these two fractions are different, yet equivalent, I felt they would be good for demonstration. I then had a try statement where I used the same code as my first test with the except statement catching a ValueError. This test correctly printed that the ValueError happened, which meant that my private  $\underline{\phantom{a}}$  lt $\underline{\phantom{a}}$ , and  $\underline{\phantom{a}}$  methods all functioned properly.