CS186 Discussion 11

(ER Diagrams, Functional Dependencies)

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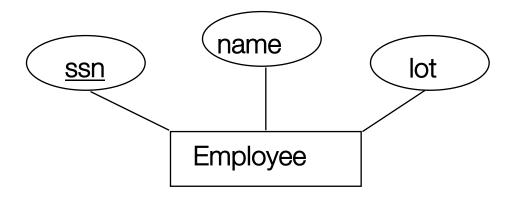
ER Diagrams

ER Diagrams

- Conceptual design of relations
 - Entities
 - Relationships
 - Attributes

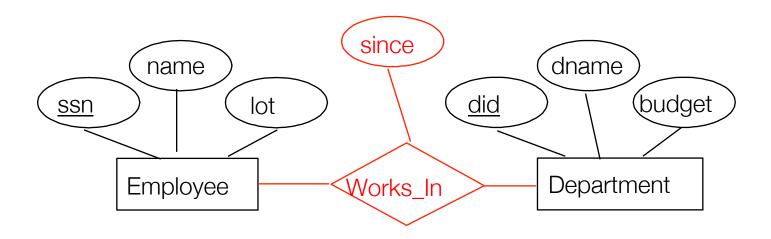
Entities

- Entity: a real world object described by a set of attributes
- Entity set: a collection of similar entities
 - E.g. all employees
 - Entities in entity set have same attributes
 - Has a <u>key</u> attribute
 - Each attribute has a domain



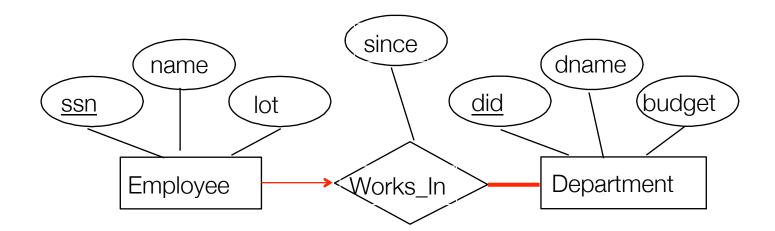
Relationships

- Relationship: association among two or more entities
 - Can have own attributes
- Relationship set: collection of similar relationships



Constraints

- Key constraint: entity participates at most once
 - Key, non-key
 - Represented by
- Participation constraint: entity participates at least once
 - Total, partial
 - Represented by ——



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	Partial Participation		Total Participation	
Non-Key	0 or more		1 or more	
Key	0 or 1	>	1	—

Ternary Relations

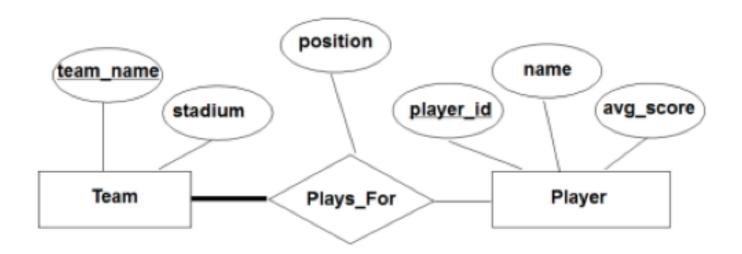
• 3 entities connected to a relationship instead of 2

Weak Entity

- An entity that only makes sense in the context of another entity (its parent)
 - Has a partial key (dashed underline)
- E.g. there can be two songs with the same name
 - Key is (Artists.artist_id, Songs.song_name)

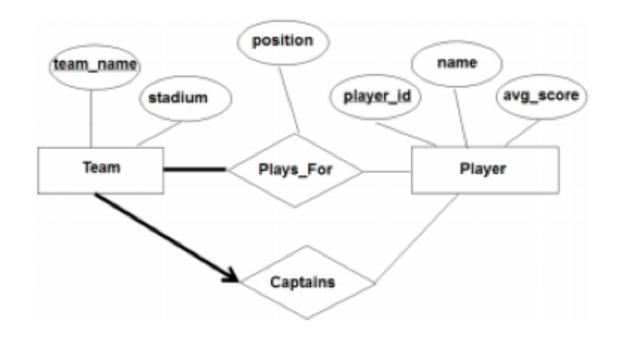
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2. Now let's say we want to also track who is the captain of every team. How will the ER diagram change from the previous case? Note: Every team needs exactly one captain!

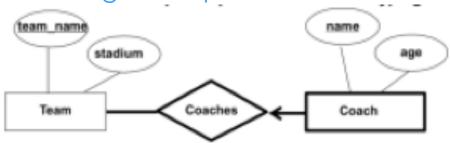
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3. Are there any weak-entity relationships in either of our ER diagrams?

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No. A weak entity can be identified uniquely only by considering the primary key of another (owner) entity. Consider the following example:



A team can have many coaches, but each coach exactly coaches one team. Coach is a weak-entity set and can be identified by its partial key "name".

Functional Dependencies

Functional Dependencies

- Used to identify redundancy in schemas and suggest refinement
- $X \rightarrow Y$
 - X determines Y
- $K \rightarrow \{all \ attributes \ of \ R\}$
 - K is a superkey of R

Armstrong's Axioms

- Reflexivity: If $X \supseteq Y$, then $X \rightarrow Y$
- Augmentation: If X → Y, then XZ → YZ for any Z
 XZ → YZ does not mean X → Y
- Transitivity: If $X \to Y$ and $Y \to Z$, then $X \to Z$
- Union: If $X \to Y$ and $X \to Z$, then $X \to YZ$
- Decomposition: If $X \to YZ$, then $X \to Y$ and $X \to Z$

Closures

 X+: set of all FDs implied by X, including trivial dependencies

```
X+ := X
while not done:
  for U→V in F:
    if U in X+:
    add V to X+
```

Boyce-Codd Normal Form (BCNF)

- R is in BCNF if the only non-trivial FDs over R are key constraints
- R with FDs F is in BCNF if for all X → A in F+:
 - $-A\subseteq X$ (called a trivial FD), or
 - X is a superkey for R

BCNF Decomposition

 For R, if X → A violates BCNF, decompose R into R – A and XA

```
1. Flight schema
Flights(Flight_no, Date, fRom, To, Plane_id),
    ForeignKey(Plane_id)
Planes(Plane_id, tYpe)
Seat(Seat_no, Plane_id, Legroom), ForeignKey(Plane_id)
Find the set of functional dependencies.
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$$FD \rightarrow RTP$$
 $P \rightarrow Y$
 $SP \rightarrow L$

2. Now consider the attribute set R = ABCDE and the FD set $F = \{AB \rightarrow C, A \rightarrow D, D \rightarrow E, AC \rightarrow B\}$. Compute the closure for the following attributes.

A:

AB:

B:

D:

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```
A: {A, D, E}
```

AB: {A, B, C, D, E}

B: {B}

D: {D, E}

3. Decompose R = ABCDEFG into BCNF, given the functional dependency set:

$$F = \{AB \rightarrow CD, C \rightarrow EF, G \rightarrow A, G \rightarrow F, CE \rightarrow F\}$$

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AB→CD => decompose ABCDEFG into ABCD, ABEFG

G→A => decompose ABEFG into AG, BEFG

G→F => decompose BEFG into FG, BEG

Final relations: ABCD, AG, FG, BEG.

4. Does the above decomposition preserve dependencies?

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No, $C \rightarrow EF$ and $CE \rightarrow F$ are not represented in the closure of the union of each subrelation's dependencies