

# CS 477: Homework #7

Due on December 8th, 2016 at 2:30pm

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## Problem 1

[20 points](U & G Required)

Answer the questions below regarding the following graph:



- [5 points] In what order are edges added to the Minimum Spanning Tree (MST) using Kruskal's Algorithm? List the edges by giving their endpoints.
- [5 points] In what order are edges added to the MST using Prim's Algorithm starting from vertex A? List the edges by giving their endpoints.

## Solution

- g-h, A-c, h-i, A-b, c-d, c-g, f-g, c-e
- A-c, A-b, c-d, c-g, g-h, h-i, f-g, c-e

## Problem 2

[30 points](U & G Required) Exercise 22.2-9 (page 602).

- 22.2-9) Let  $G = (V, E)$  be a connected, undirected graph. Give an  $\mathcal{O}(V + E)$ -time algorithm to compute a path in  $G$  that traverses each edge in  $E$  exactly once in each direction. Describe how you can find your way out of a maze if you are given a large supply of pennies.

## Problem 3

[30 points](U & G Required) Exercise 22.5-6 (page 621).

22.5-6) Given a directed graph  $G = (V, E)$ , explain how to create another graph  $G' = (V, E')$  such that

- a.)  $G'$  has the same strongly connected components as  $G$
- b.)  $G'$  has the same component graph as  $G$
- c.)  $E'$  is as small as possible.

Describe a fast algorithm to compute  $G'$ .

## Problem 4

[20 points](U & G Required) Exercise 24.3-2 (page 663).

- 24.3-2) Give a simple example of a directed graph with negative-weight edges for which Dijkstras algorithm produces incorrect answers. Why doesnt the proof of Theorem 24.6 go through when negative-weight edges are allowed?

## Problem 6

[20 points](Extra Credit) Exercise 24.3-6 (page 663).

- 24.3-6) We are given a directed graph  $G = (V, E)$  on which each edge  $(u, v) \in E$  has an associated value  $r(u, v)$ , which is a real number in the range  $0 \leq r(u, v) \leq 1$  that represents the reliability of a communication channel from vertex  $u$  to vertex  $v$ . We interpret  $r(u, v)$  as the probability that the channel from  $u$  to  $v$  will not fail, and we assume that these probabilities are independent. Give an efficient algorithm to find the most reliable path between two given vertices.