

Solution of Homework #1

Mattia Lecci, Federico Mason

October 29, 2017

1 Problem 3.2

Given a Z channel with probabilities $p_{Y|X}(0|0) = 1$, $p_{Y|X}(1|0) = 0$, $p_{Y|X}(1|1) = p_{Y|X}(0|1) = \frac{1}{2}$, we want to find the channel capacity C . We know that the channel capacity for a discrete memory channel is defined as follows.

$$C = \max_{p_X} I(X; Y) = H(Y) - H(Y|X) \quad (1)$$

First we compute $H(Y)$ in function of $p = p_X(1)$.

$$\begin{aligned} H(Y) &= -p_Y(1) \cdot \log(p_Y(1)) - p_Y(0) \cdot \log(p_Y(0)) \\ &= -(p_{Y|X}(1|0) \cdot p_X(0) + p_{Y|X}(1|1) \cdot p_X(1)) \cdot \log(p_{Y|X}(1|0) \cdot p_X(0) + p_{Y|X}(1|1) \cdot p_X(1)) \\ &\quad - (p_{Y|X}(0|0) \cdot p_X(0) + p_{Y|X}(0|1) \cdot p_X(1)) \cdot \log(p_{Y|X}(0|0) \cdot p_X(0) + p_{Y|X}(0|1) \cdot p_X(1)) \\ &= -\frac{1}{2} \cdot p_X(1) \log\left(\frac{1}{2} \cdot p_X(1)\right) - \left(p_X(0) + \frac{1}{2} \cdot p_X(1)\right) \cdot \log\left(p_X(0) + \frac{1}{2} \cdot p_X(1)\right) \\ &= -\frac{1}{2} \cdot p \cdot \log\left(\frac{1}{2} \cdot p\right) - \left(1 - \frac{1}{2} \cdot p\right) \cdot \log\left(1 - \frac{1}{2} \cdot p\right) \end{aligned}$$

Then we compute $H(Y|X)$ always in function of $p = p_X(1)$.

$$\begin{aligned} H(Y|X) &= -p_{YX}(0,0) \cdot \log(p_{Y|X}(0|0)) - p_{YX}(1,0) \cdot \log(p_{Y|X}(1|0)) \\ &\quad - p_{YX}(0,1) \cdot \log(p_{Y|X}(0|1)) - p_{YX}(1,1) \cdot \log(p_{Y|X}(1|1)) \\ &= -p_{Y|X}(0|0) \cdot p_X(0) \cdot \log(p_{Y|X}(0|0)) - p_{Y|X}(1|0) \cdot p_X(0) \cdot \log(p_{Y|X}(1|0)) \\ &\quad - p_{Y|X}(0|1) \cdot p_X(1) \cdot \log(p_{Y|X}(0|1)) - p_{Y|X}(1|1) \cdot p_X(1) \cdot \log(p_{Y|X}(1|1)) \\ &= 0 + 0 + \frac{1}{2} \cdot p + \frac{1}{2} \cdot p = p \end{aligned}$$

Now we are able to write $I(X; Y)$ in function of p as we can see from (2).

$$I(X; Y) = -\frac{1}{2} \cdot p \cdot \log\left(\frac{1}{2} \cdot p\right) - \left(1 - \frac{1}{2} \cdot p\right) \cdot \log\left(1 - \frac{1}{2} \cdot p\right) - p \quad (2)$$

In 1 we show the function's trend for all possible values of p .

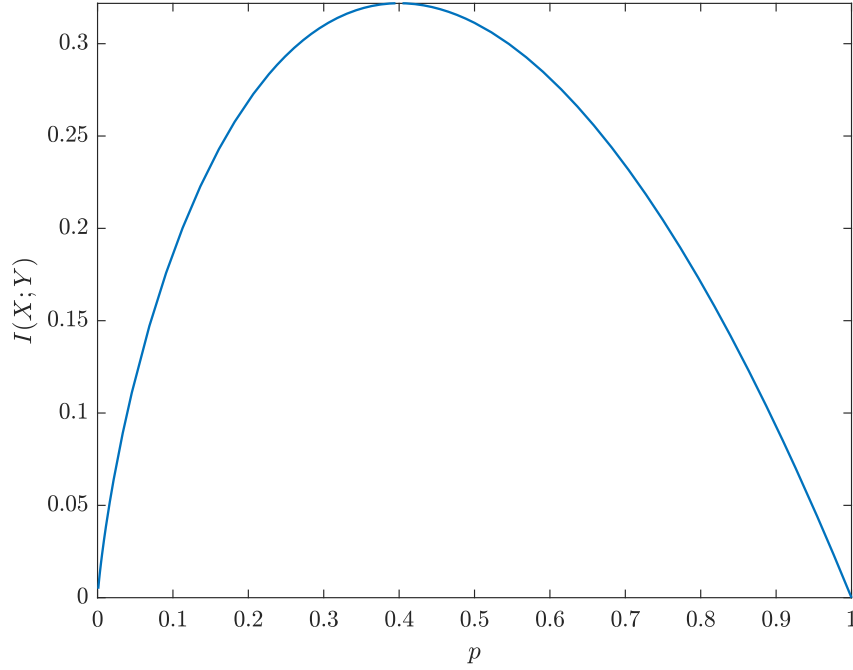


Figure 1: Mutual information as a function of $p \triangleq p_X(1)$

Now we want to find the value of p for which (2) is maximize. With this aim we compute the first and the second derivates of $I(X; Y)$.

$$\begin{aligned} \frac{dI(X; Y)}{dp} &= \frac{1}{2} \cdot \log\left(1 - \frac{p}{2}\right) + \frac{1}{2} \cdot \frac{1 - \frac{p}{2}}{1 - \frac{p}{2}} - \frac{1}{2} \cdot \log\left(\frac{p}{2}\right) - \frac{p}{2} \cdot \frac{2}{p} \cdot \frac{1}{2} - 1 \\ &= \frac{1}{2} \cdot \log\left(\frac{2}{p} - 1\right) - 1 \end{aligned}$$

$$\frac{d^2 I(X; Y)}{dp^2} = \frac{1}{2} \cdot \frac{p}{2 - p} \cdot \frac{1}{2} = \frac{1}{4} \cdot \frac{p}{2 - p}$$

We notice that $\frac{d^2 I(X; Y)}{dp^2}$ is always equal or greater than zero for every possibli value of p . This means that $I(X; Y)$ is a convex function and then the value of p for which $\frac{dI(X; Y)}{dp}$ is zero coincides with the maximum point of $I(X; Y)$. Therefore we compute the values for which we have $\frac{dI(X; Y)}{dp} = 0$.

$$\begin{aligned} \frac{1}{2} \cdot \log\left(\frac{2}{p} - 1\right) - 1 &= 0 \\ \log\left(\frac{2}{p} - 1\right) &= 2 \\ \frac{2}{p} - 1 &= 2^2 \\ p &= \frac{2}{5} \end{aligned}$$

So for $p = \frac{2}{5}$ we maximize $I(X; Y)$ and therefore we obtain the value of the channel capacity C .

$$\begin{aligned} C &= I(X; Y)|_{p=\frac{2}{5}} \\ &= -\frac{4}{5} \cdot \log\left(\frac{4}{5}\right) - \frac{1}{5} \cdot \log\left(\frac{1}{5}\right) - \frac{2}{5} \\ &= \frac{1}{5} \cdot \log\left(\frac{3125}{256}\right) - \frac{2}{5} \end{aligned}$$