Using and Understanding the Real-Time Cyclictest Benchmark

Cyclictest results are the most frequently cited real-time Linux metric. The core concept of Cyclictest is very simple. However the test options are very extensive. The meaning of Cyclictest results appear simple but are actually quite complex. This talk will explore and explain the complexities of Cyclictest. At the end of the talk, the audience will understand how Cyclictest results describe the potential real-time performance of a system.

What Cyclictest Measures

Latency of response to a stimulus.

external interrupt triggers (clock expires)

- possible delay until IRQs enabled
- IRQ handling
- cyclictest is woken
- possible delay until preemption enabled
- possible delay until cyclictest is highest priority
- possible delay until other process is preempted
- scheduler overhead

transfer control to cyclictest

What Cyclictest Measures

Latency of response to a stimulus.

Causes of delay list on previous slide is simplified:

- order will vary
- may occur multiple times
- there are additional causes of delay

Many factors can increase latency

- additional external interrupts
- SMI
- processor emerging from sleep states
- cache migration of data used by woken process
- block on sleeping lock
 - lock owner gets priority boost
 - lock owner schedules
 - lock owner completes scheduled work
 - lock owner releases lock, loses priority boost

How Cyclictest Measures Latency

(Cyclictest Pseudocode)

The source code is nearly 3000 lines, but the algorithm is trivial

Test Loop

```
clock_gettime((&now))
next = now + par->interval
while (!shutdown) {
   clock_nanosleep((&next))
   clock_gettime((&now))
   diff = calcdiff(now, next)
   # update stat-> min, max, total latency, cycles
   # update the histogram data
   next += interval
```

The Magic of Simple

This trivial algorithm captures all of the factors that contribute to latency.

Mostly. Caveats will follow soon.

Cyclictest Program

```
main() {
   for (i = 0; i < num_threads; i++) {
      pthread_create((timerthread))
   while (!shutdown) {
      for (i = 0; i < num\_threads; i++)
         print_stat((stats[i]), i))
      usleep(10000)
   }
   if (histogram)
      print_hist(parameters, num_threads)
}
```

timerthread()

```
*timerthread(void *par) {
    # thread set up
    # test loop
}
```

Thread Set Up

```
stat = par->stats;
pthread_setaffinity_np((pthread_self()))
setscheduler(({par->policy, par->priority))
sigprocmask((SIG_BLOCK))
```

Test Loop (as shown earlier)

```
clock_gettime((&now))
next = now + par->interval
while (!shutdown) {
   clock_nanosleep((&next))
   clock_gettime((&now))
   diff = calcdiff(now, next)
   # update stat-> min, max, avg, cycles
   # Update the histogram
   next += interval
```

Why show set up pseudocode?

The timer threads are not in lockstep from time zero.

Multiple threads will probably not directly impact each other.

September 2013 update

```
linux-rt-users
[rt-tests][PATCH] align thread wakeup times
Nicholas Mc Guire
2013-09-09 7:29:48
And replies
```

"This patch provides and additional -A/--align flag to cyclictest to align thread wakeup times of all threads as closly defined as possible."

"... we need both.

same period + "random" start time
same period + synced start time

it makes a difference on some boxes that is significant."

The Magic of Simple

This trivial algorithm captures all of the factors that contribute to latency.

Mostly. Caveats, as promised.

Caveats

Measured maximum latency is a floor of the possible maximum latency

- Causes of delay may be partially completed when timer IRQ occurs
- Cyclictest wakeup is on a regular cadence, may miss delay sources that occur outside the cadence slots

Caveats

Does not measure the IRQ handling path of the real RT application

- timer IRQ handling typically fully in IRQ context
- normal interrupt source IRQ handling:
 - irq context, small handler, wakes IRQ thread
 - IRQ thread eventually executes, wakes RT process

Caveats

Cyclictest may not exercise latency paths that are triggered by the RT application, or even non-RT applications

- SMI to fixup instruction errata
- stop_machine()
 - module load / unload
 - hotplug

Do not use cyclictest. :-)

Instrument the RT application to measure latency

Run the normal RT application and non-RT applications as the system load

Run cyclictest with a higher priority than the RT application to measure latency

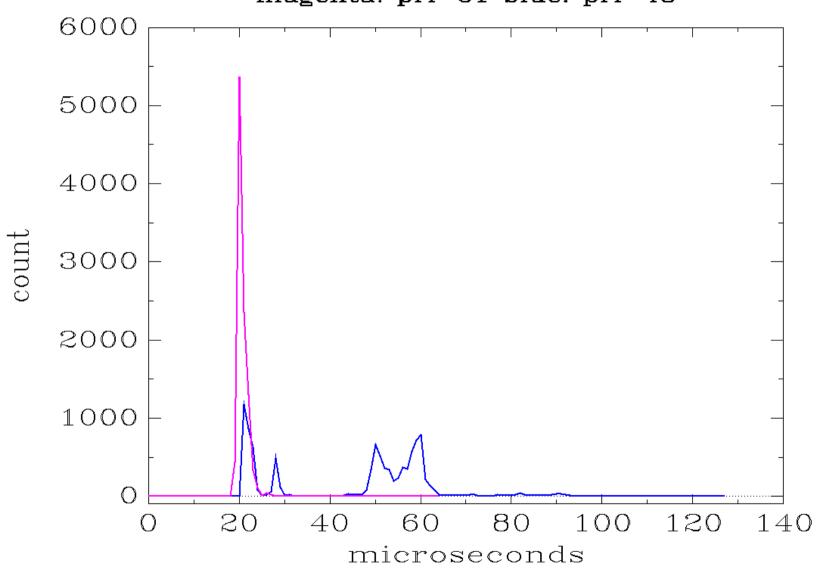
Typical real time application will consist of multiple threads, with differing priorities and latency requirements

To capture latencies of each of the threads, run separate tests, varying the cyclictest priority

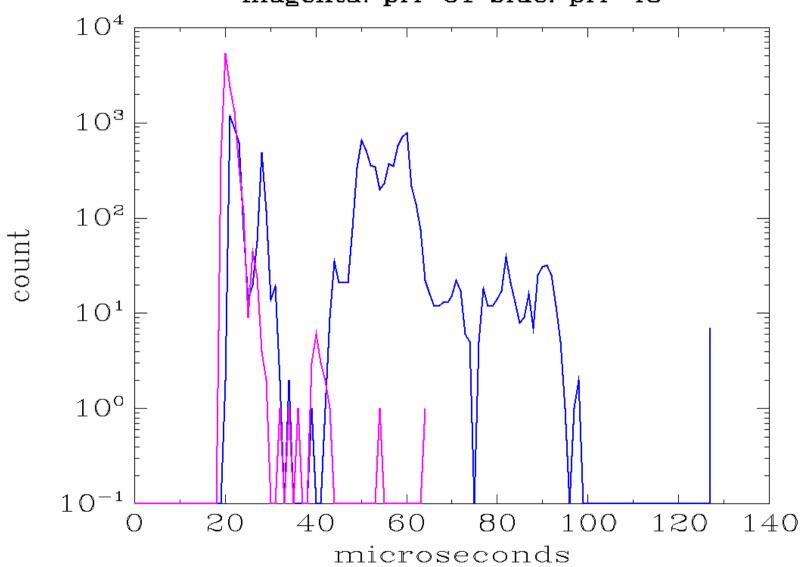
Example

	deadline <u>constraint</u>	latency <u>constraint</u>		cyclictest priority	
Α	critical	80 usec	50	51	
В	0.1% miss	100 usec	47	48	

Cyclictest Latency magenta: pri=51 blue: pri=48



Cyclictest Latency magenta: pri=51 blue: pri=48



Aside:

Cyclictest output in these slides is edited to fit on the slides

Original:

Example of edit:

```
$ cyclictest_0.85 -l100000 -q -p80 -S
T:0 I:1000 Min: 37 Avg: 45 Max: 68
T:1 I:1500 Min: 37 Avg: 42 Max: 72
```

Cyclictest Command Line Options

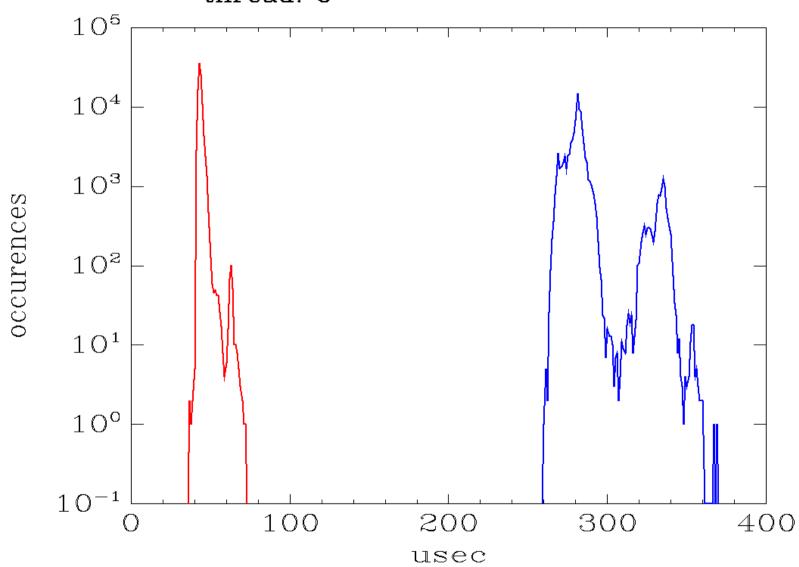
Do I really care???

Can I just run it with the default options???

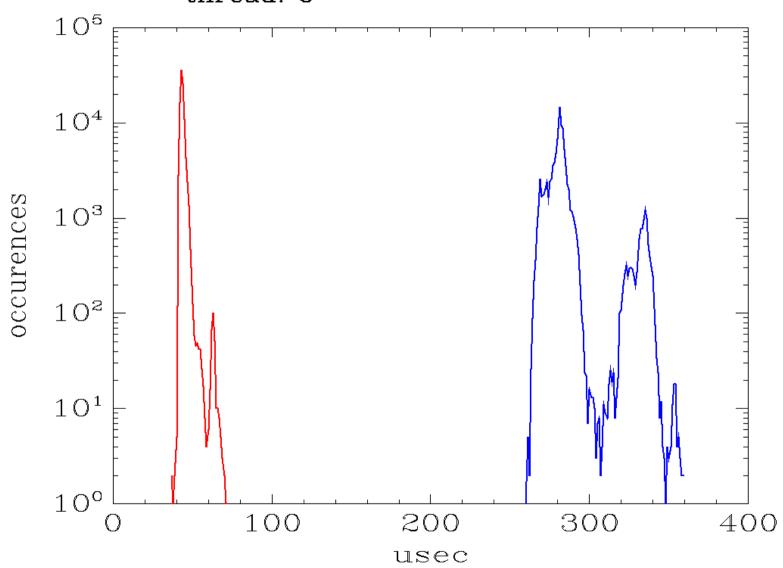
Do I really care???

```
$ cyclictest_0.85 -l100000 -q -p80
    T:0 Min: 262 Avg: 281 Max: 337
    $ cyclictest_0.85 -l100000 -q -p80 -n
    T:0 Min: 35 Avg: 43 Max:
          stop after 100000 loops
-1100000
          quiet
- Q
-p80
          priority 80, SCHED_FIF0
          use clock_nanosleep()
-n
          instead of nanosleep()
```

red: clock_nanosleep() blue: nanosleep()
thread: 0



red: clock_nanosleep() blue: nanosleep()
thread: 0



Impact of Options

More examples

Be somewhat skeptical of maximum latencies due to the short test duration.

Examples are:

100,000 loops

1,000,000 loops

Arbitrary choice of loop count. Need large values to properly measure maximum latency!!!

Priority of Real Time kernel threads for next two slides

```
PPID S RTPRIO CLS
PID
                         CMD
                   1 FF
                          [ksoftirqd/0]
  3
         2 S
         2 S
                          [posixcputmr/0]
                  70 FF
         2 S
                          [migration/0]
                  99 FF
                          [posixcputmr/1]
         2 S
                  70 FF
         2 S
                  99 FF
                          [migration/1]
                          [ksoftirqd/1]
         2 S
                   1 FF
 11
                          [irq/41-eth%d]
353
         2 S
                  50 FF
         2 S
                          [irq/46-mmci-pl1]
374
                  50 FF
                          [irq/47-mmci-pl1]
375
         2 S
                  50 FF
                          [irq/36-uart-pl0]
394
                  50 FF
```

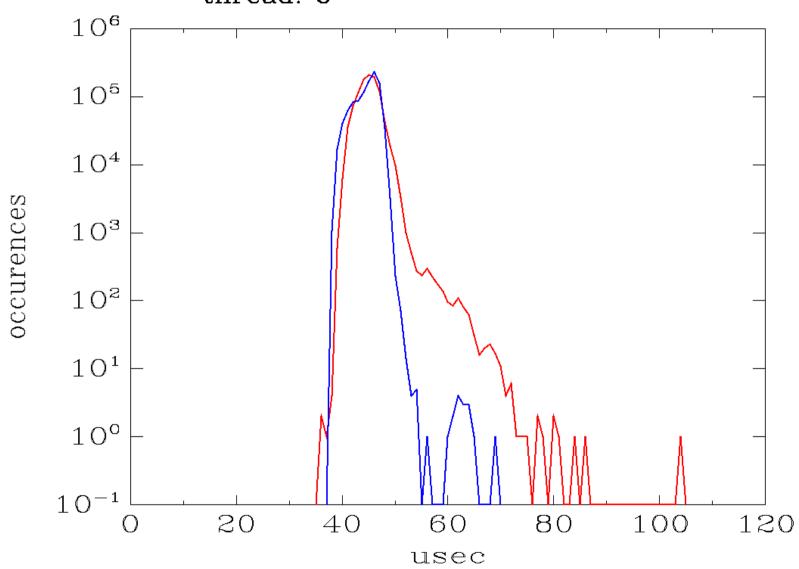
-1100000

T:0	Min:	128	Avg:	189	Max:	2699		<u>live update</u>
T:0	Min:	125	Avg:	140	Max:	472	- q	no live update
T:0	Min:	262	Avg:	281	Max:	337	- p80	SCHED_FIFO 80
T:0	Min:	88	Avg:	96	Max:	200	-n	<u>clock_nanosleep</u>
T:0 T:1	Min: Min:		•		Max: Max:	496 509	-q -p80 -a -t	pinned
T:0	Min:	35	Avg:	43	Max:	68	-q -p80 -n	SCHED_FIFO, c_n
T:0	Min:	34	Avg:	44	Max:	71	-q -p80 -a -n	pinned
T:0	Min:	38	Avg:	43	Max:	119	-q -p80 -a -n -m	mem locked
T:0 T:1	Min: Min:	36 37	Avg: Avg:	43 45	Max: Max:	65 78	-q -p80 -t -n	<u>not pinned</u>
T:0 T:1	Min: Min:	36 37	Avg: Avg:	44 45	Max: Max:	91 111	-q -p80 -a -t -n	pinned
T:0 T:1	Min: Min:	34 34	Avg: Avg:	44 43	Max: Max:	94 104	-q -p80 -S	<u>=> -a -t -n</u>

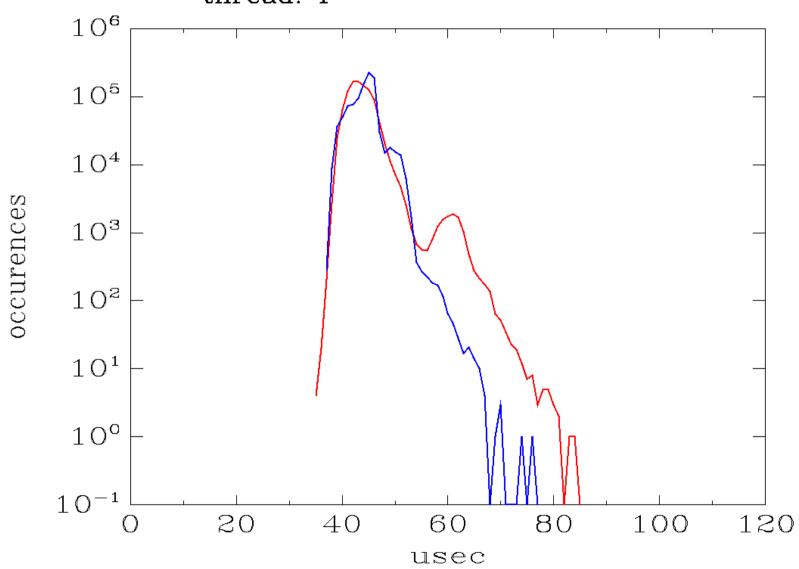
-11000000

T:0	Min:	123	Avg:	184	Max:	3814				<u>live update</u>
T:0	Min:	125	Avg:	150	Max:	860	- q			no live update
T:0	Min:	257	Avg:	281	Max:	371	-q -p	080		SCHED_FIFO 80
T:0	Min:	84	Avg:	94	Max:	319	-q -n	າ		<u>clock nanosleep</u>
T:0 T:1	Min: Min:		Avg: Avg:			682 506	-q -p	080 -a	-t	pinned
T:0	Min:	38	Avg:	44	Max:	72	-q -p	080 -n		SCHED FIFO, c_n
T:0	Min:	33	Avg:	42	Max:	95	-q -p	080 -a	-n	pinned
T:0	Min:	36	Avg:	42	Max:	144	-q -p)80 -a	-n -m	mem locked
T:0 T:1	Min: Min:	36 37	Avg: Avg:	44 45	Max: Max:	84 94	-q -p	080 -t	- n	<u>not pinned</u>
T:0 T:1	Min: Min:	36 36	Avg: Avg:	43 43	Max: Max:	87 91	-q -p)80 -a	-t -n	pinned
T:0 T:1	Min: Min:	36 34	Avg: Avg:	43 42	Max: Max:	141 88	-q -p	080 -S		<u>=> -a -t -n</u>

blue: cpu pinned red: cpu not pinned thread: 0



blue: cpu pinned red: cpu not pinned thread: 1



Simple Demo -- sched_normal

- single thread
- clock_nanosleep(), one thread per cpu, pinned
- clock_nanosleep(), one thread per cpu
- clock_nanosleep(), one thread per cpu, memory locked
- clock_nanosleep(), one thread per cpu, memory locked, non-interactive

What Are Normal Results?

What should I expect the data to look like for my system?

Examples of Maximum Latency

```
https://rt.wiki.kernel.org/index.php/CONFIG_PREEMPT_RT_Patch #Platforms_Tested_and_in_Use_with_CONFIG_PREEMPT_RT Platforms Tested and in Use with CONFIG_PREEMPT_RT Comments sometimes include avg and max latency table is usually stale
```

linux-rt-users email list archives http://vger.kernel.org/vger-lists.html#linux-rt-users

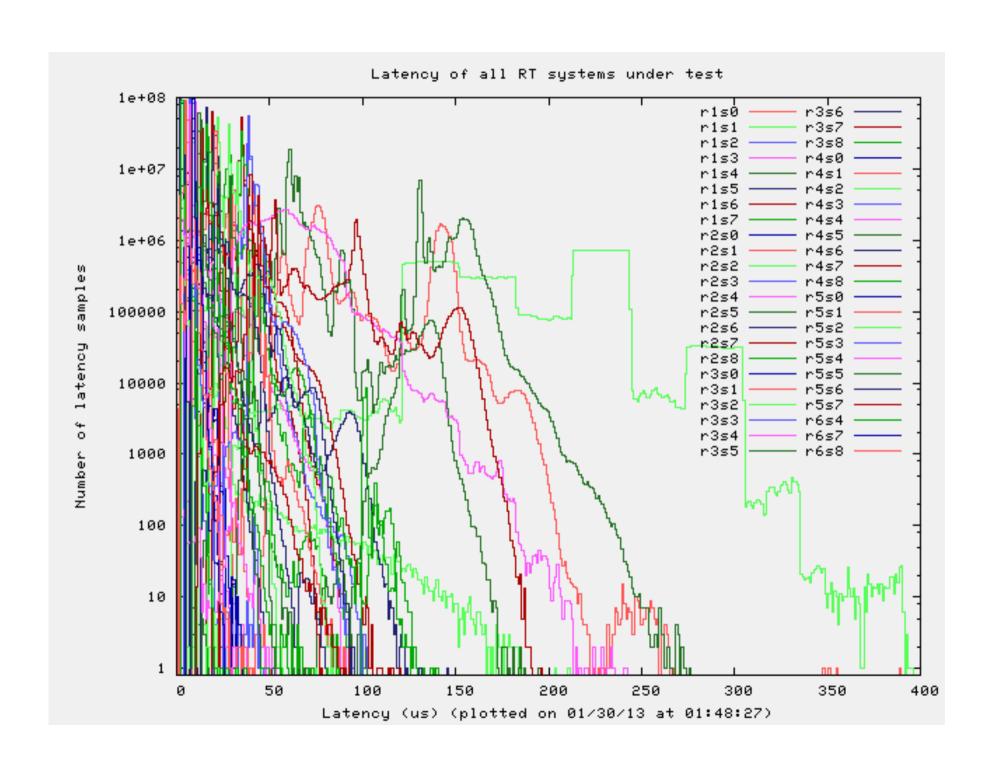
Graphs of Maximum Latency

OSADL.org

Graphs for a wide variety of machines

List of test systems:

https://www.osadl.org/Individual-system-data.qa-farm-data.0.html



Full URL of previous graph

https://www.osadl.org/Combined-latency-plot-of-all-RT-systems.qa-latencyplot-allrt.0.html?latencies=&showno=

Typical command:

cyclictest -l100000000 -m -Sp99 -i200 -h400 -q

OSADL Realtime QA Farm:

https://www.osadl.org/QA-Farm-Realtime.qa-farm-about.0.html

OSADL Latency plots:

https://www.osadl.org/Latency-plots.latency-plots.0.html

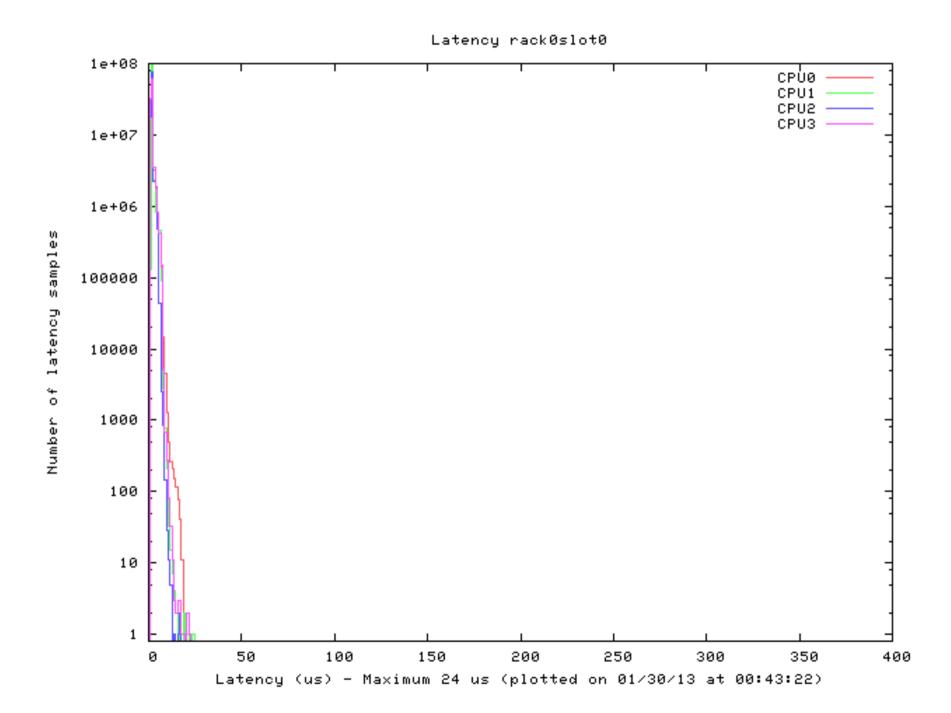
Additional OSADL Data

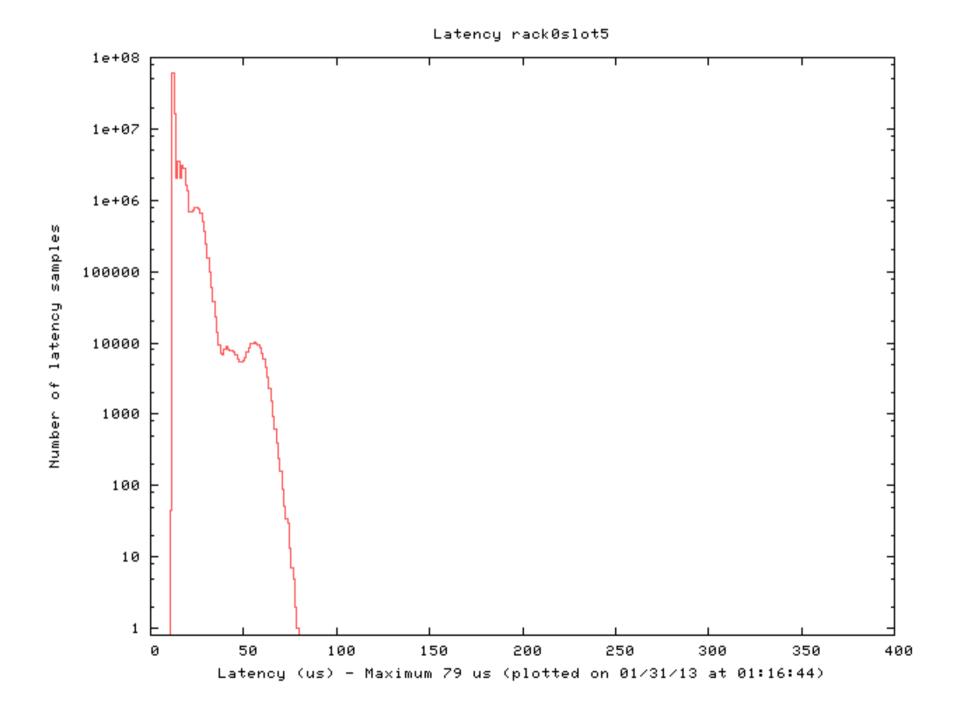
OSADL members have access to additional data, such as

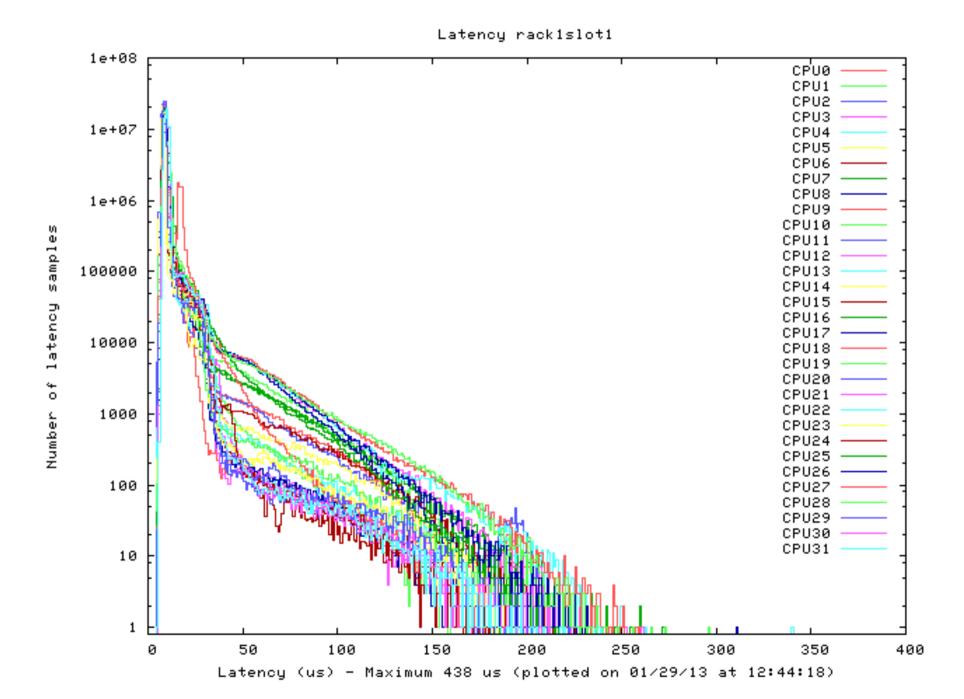
- the data used to create the graphs
- the latency graphs extended in a third dimension, showing all test runs

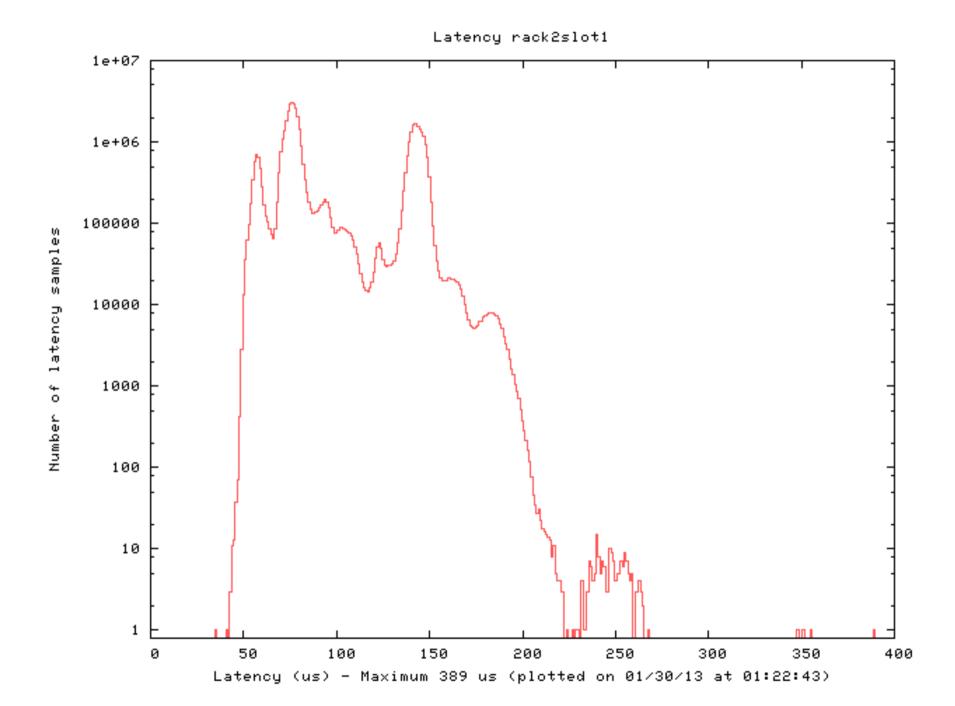
Some Random Individual Systems

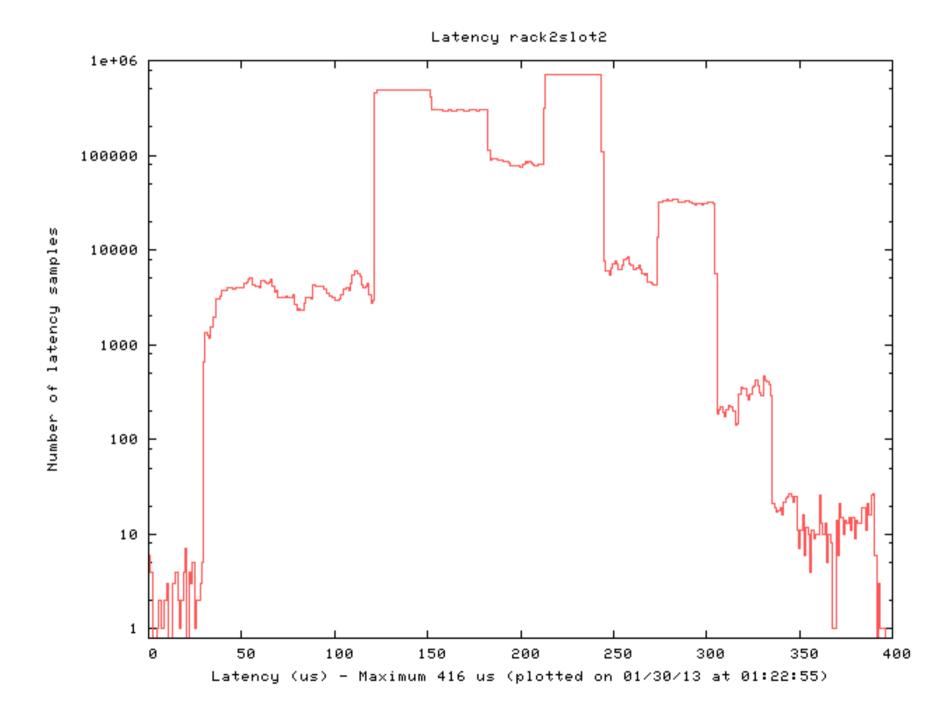
Picked from the OSADL spaghetti graph

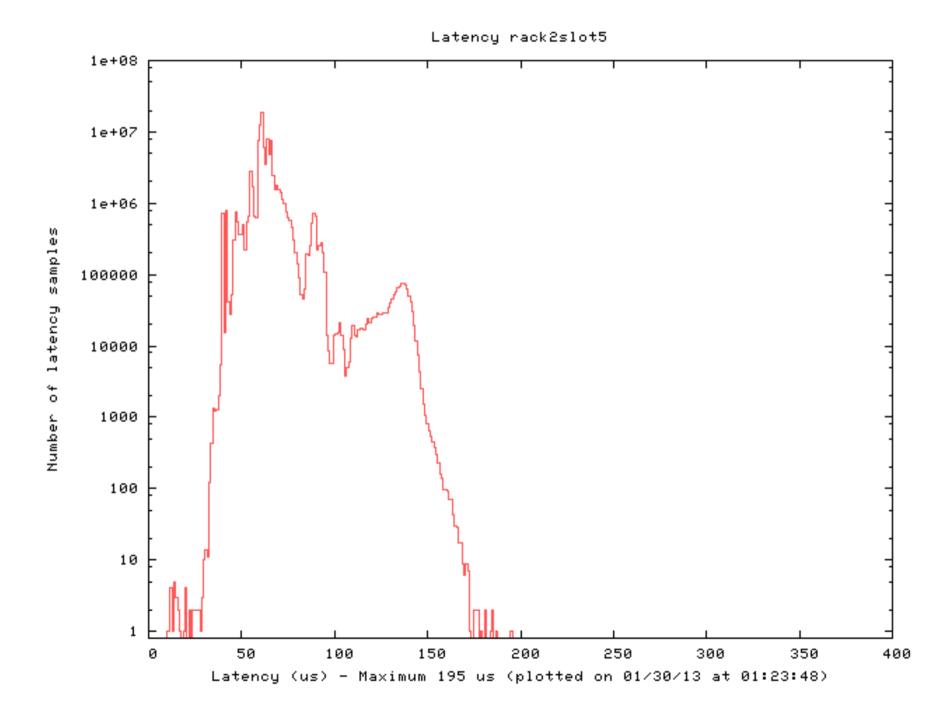


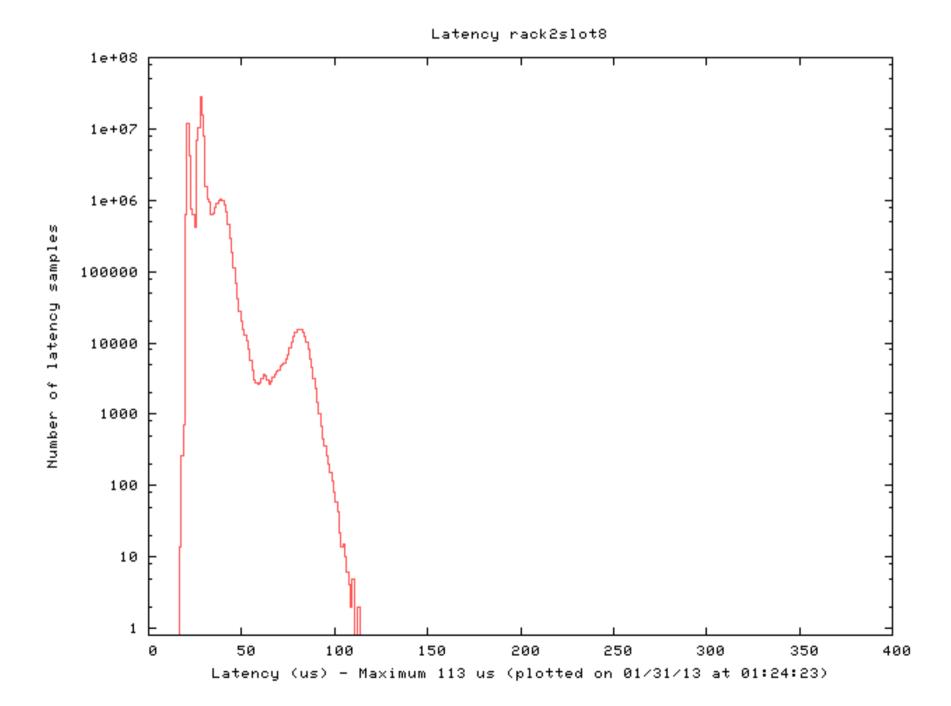


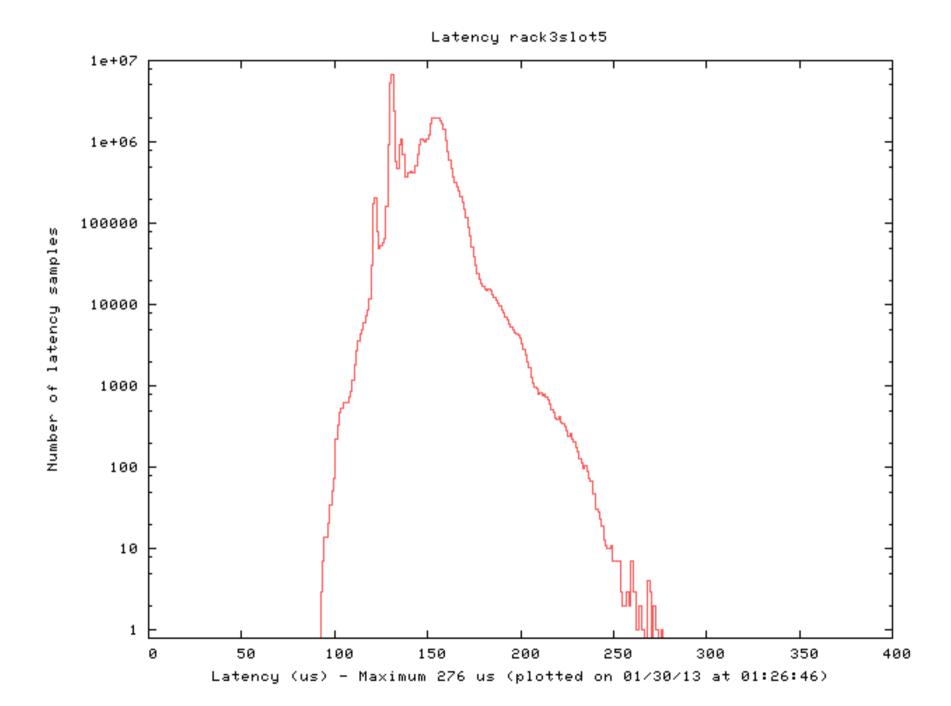


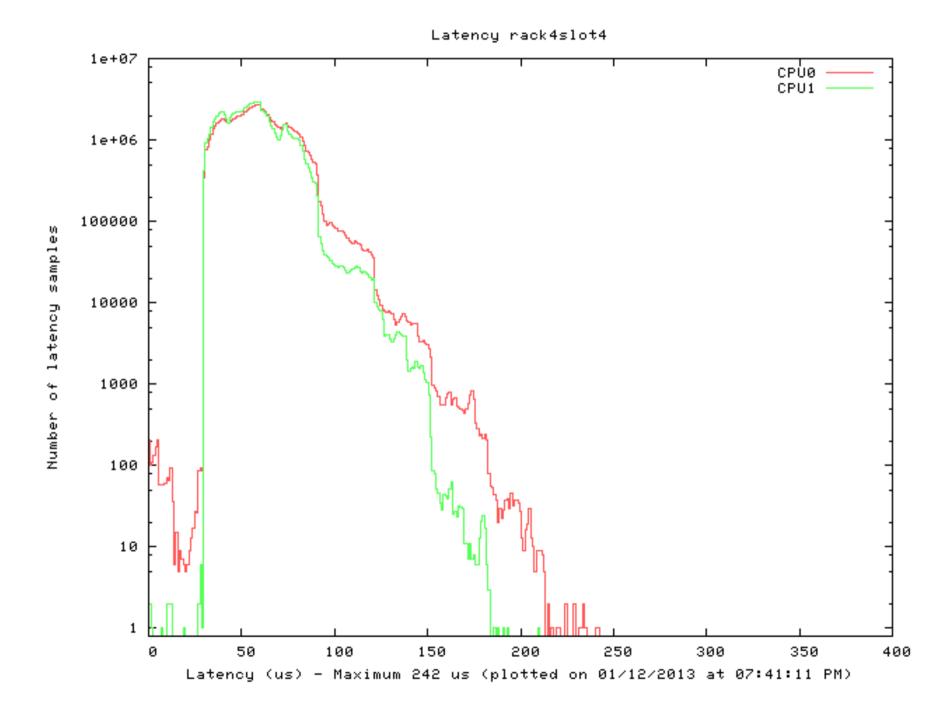


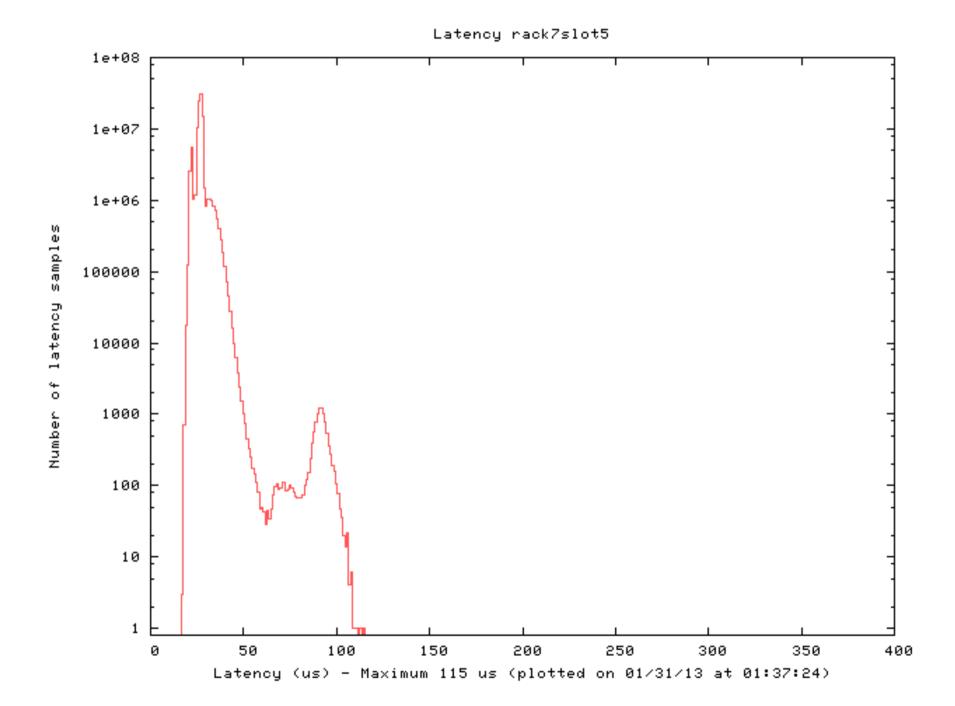




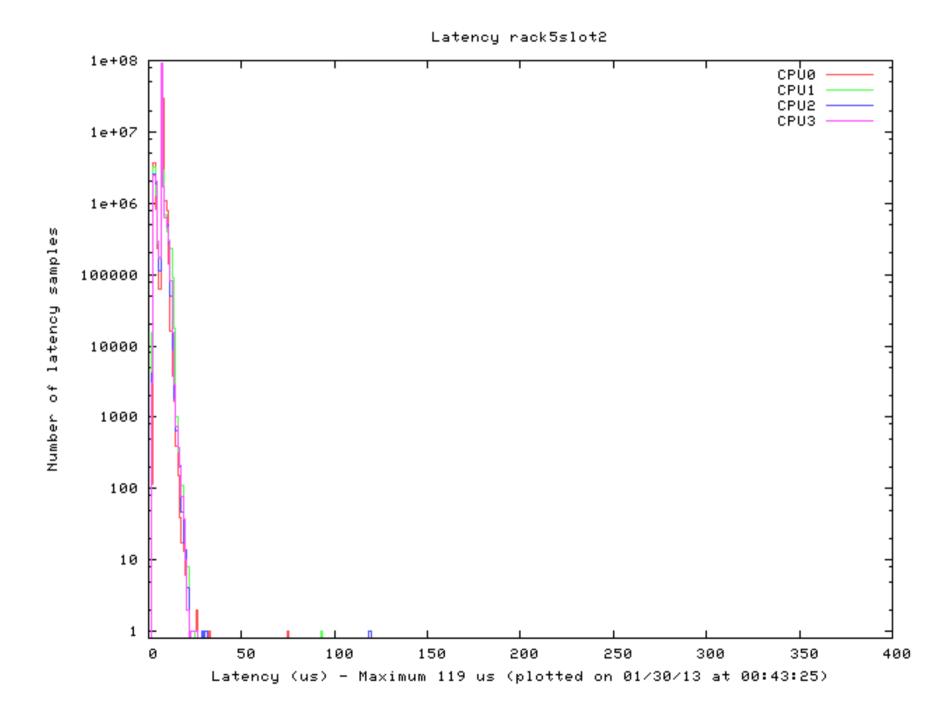


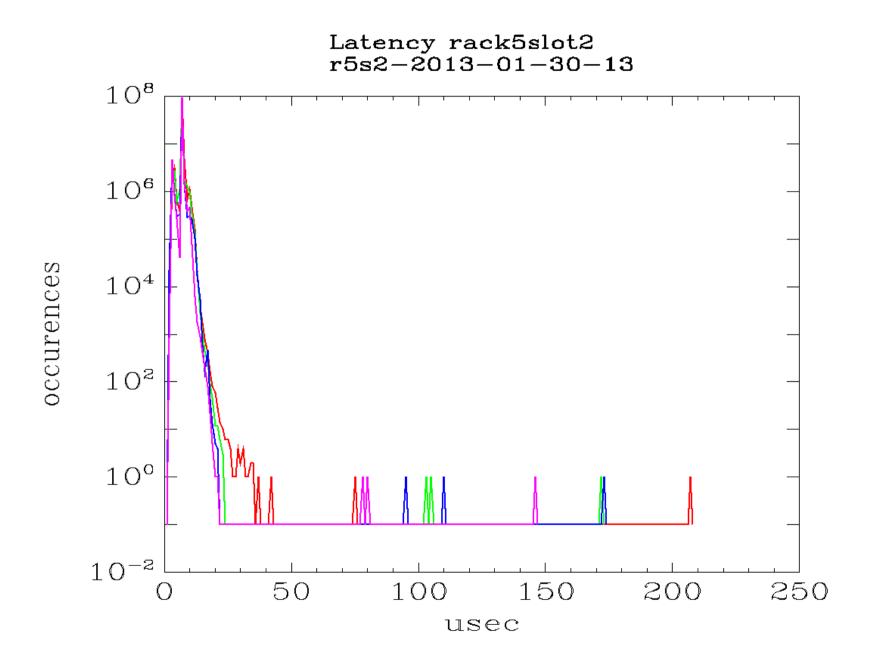


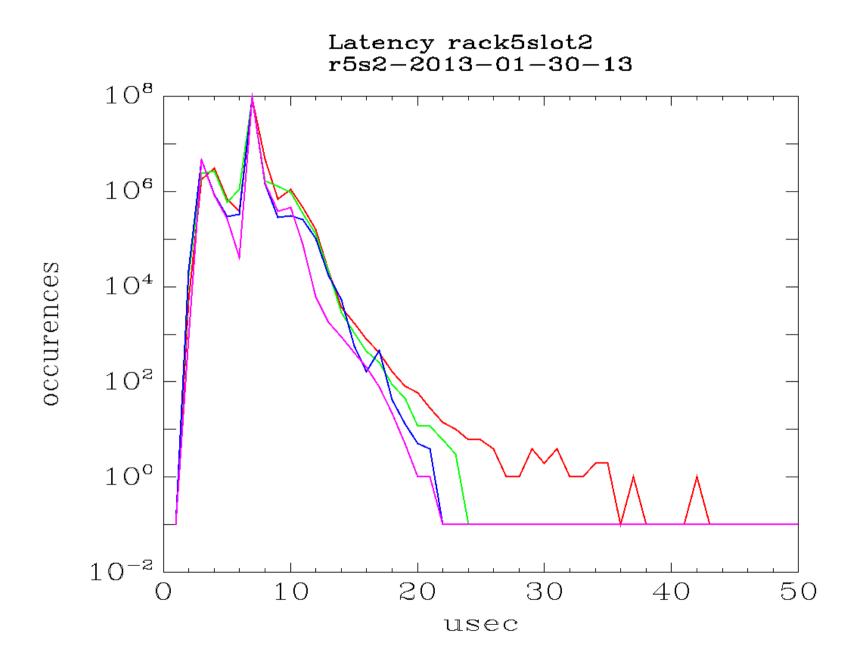




Even "boring" graphs may contain interesting details







Command Line Options

An unruly, out of control, set of control knobs

cvclictest V 0.85 Usage: cyclictest <options> run thread #N on processor #N, if possible -a [NUM] --affinity with NUM pin all threads to the processor NUM -b USEC --breaktrace=USEC send break trace command when latency > USEC --preemptirgs both preempt and irgsoff tracing (used with -b) -c CLOCK --clock=CLOCK select clock 0 = CLOCK MONOTONIC (default) 1 = CLOCK REALTIME - C --context context switch tracing (used with -b) -d DIST --distance=DIST distance of thread intervals in us default=500 --duration=t specify a length for the test run default is in seconds, but 'm', 'h', or 'd' maybe added to modify value to minutes, hours or days --latency=PM 00S write PM OOS to /dev/cpu dma latency - e - E --event event tracing (used with -b) -f --ftrace function trace (when -b is active) -q MAX Report time in ms (up to MAX) for histogram overflows --of max=MAX - h --histogram=US dump a latency histogram to stdout after the run (with same priority about many threads) US is the max time to be be tracked in microseconds - H --histofall=US same as -h except with an additional summary column base interval of thread in us default=1000 -i INTV --interval=INTV - I --irqsoff Irqsoff tracing (used with -b) -1 LOOPS --loops=LOOPS number of loops: default=0(endless) --mlockall lock current and future memory allocations - m delay updating the screen until a new max latency is hit - M --refresh on max --nanosleep use clock nanosleep - n print results in ns instead of us (default us) - N --nsecs oscilloscope mode, reduce verbose output by RED -o RED --oscope=RED -0 TOPT --traceopt=TOPT trace option -p PRIO --prio=PRIO priority of highest prio thread -P --preemptoff Preempt off tracing (used with -b) print only a summary on exit - q --auiet spread priority levels starting at specified value -0 --priospread -r --relative use relative timer instead of absolute -R --resolution check clock resolution, calling clock gettime() many times. list of clock_gettime() values will be reported with -X - S --system use svs nanosleep and svs setitimer -S --smp Standard SMP testing: options -a -t -n and same priority of all threads -t --threads one thread per available processor -t [NUM] --threads=NUM number of threads: without NUM, threads = max cpus without -t default = 1 -T TRACE --tracer=TRACER set tracing function configured tracers: blk function graph wakeup rt wakeup function nop force unbuffered output for live processing - u --unbuffered -U --numa Standard NUMA testing (similar to SMP option) thread data structures allocated from local node output values on stdout for statistics - V --verbose format: n:c:v n=tasknum c=count v=value in us --wakeup task wakeup tracing (used with -b) -W rt task wakeup tracing (used with -b) -W --wakeuprt --dbg_cyclictest print info useful for debugging cyclictest - X policy of realtime thread, POLI may be fifo(default) or rr -y POLI --policy=POLI format: --policy=fifo(default) or --policy=rr

\$ cyclictest --help

Thread Behavior Options

```
-a [NUM] --affinity
                          run thread #N on processor #N, if possible
                          with NUM pin all threads to the processor NUM
-c CLOCK --clock=CLOCK
                          select clock
                          0 = CLOCK_MONOTONIC (default)
                          1 = CLOCK REALTIME
-d DIST --distance=DIST
                          distance of thread intervals in us default=500
-i INTV --interval=INTV
                          base interval of thread in us default=1000
        --mlockall
                          lock current and future memory allocations
– m
       --nanosleep
                          use clock nanosleep
- n
                          priority of highest prio thread
-p PRIO --prio=PRIO
        --priospread
                           spread priority levels starting at specified value
- Q
                          use relative timer instead of absolute
        --relative
- r
- S
        --system
                          use sys nanosleep and sys setitimer
-S
        --smp
                          Standard SMP testing: options -a -t -n and
                          same priority of all threads
-t --threads
                          one thread per available processor
-t [NUM] --threads=NUM
                          number of threads:
                          without NUM, threads = max_cpus
                          without -t default = 1
                          Standard NUMA testing (similar to SMP option)
-U
       --numa
                          thread data structures allocated from local node
-y POLI --policy=POLI
                          policy of realtime thread, POLI may be fifo(default) or rr
                          format: --policy=fifo(default) or --policy=rr
```

side effect, sets -d0

```
-h --histogram=US dump a latency histogram to stdout after the run (with same priority about many threads)
US is the max time to be be tracked in microseconds same as -h except with an additional summary column
```

Benchmark and System Options

```
-D --duration=t specify a length for the test run default is in seconds, but 'm', 'h', or 'd' maybe added to modify value to minutes, hours or days number of loops: default=0(endless)

-e --latency=PM_QOS write PM_QOS to /dev/cpu_dma_latency
```

Display Options

-g MAX	of_max=MAX	Report time in ms (up to MAX) for histogram overflows
- h	histogram=US	dump a latency histogram to stdout after the run
		(with same priority about many threads)
		US is the max time to be be tracked in microseconds
- H	histofall=US	same as -h except with an additional summary column
- M	refresh_on_max	delay updating the screen until a new max latency is hit
- N	nsecs	print results in ns instead of us (default us)
-o RED	oscope=RED	oscilloscope mode, reduce verbose output by RED
- q	quiet	print only a summary on exit
- u	unbuffered	force unbuffered output for live processing
- V	verbose	output values on stdout for statistics
		format: n:c:v n=tasknum c=count v=value in us

Debug Options

```
-b USEC --breaktrace=USEC send break trace command when latency > USEC
                          both preempt and irgsoff tracing (used with -b)
- B
        --preemptirgs
- C
                          context switch tracing (used with -b)
        --context
                          event tracing (used with -b)
- E
        --event
-f
        --ftrace
                          function trace (when -b is active)
        --irgsoff
                          Irgsoff tracing (used with -b)
- T
-0 TOPT --traceopt=TOPT
                          trace option
                          Preempt off tracing (used with -b)
        --preemptoff
- P
        --resolution
                          check clock resolution, calling clock gettime() many
- R
                          times. list of clock_gettime() values will be
                          reported with -X
                          set tracing function
-T TRACE --tracer=TRACER
   configured tracers: blk function graph wakeup rt wakeup function nop
        --wakeup
                          task wakeup tracing (used with -b)
– W
        --wakeuprt rt task wakeup tracing (used with -b)
-W
        --dbg cyclictest print info useful for debugging cyclictest
- X
```

Debug Options

No time to describe in this talk

Hooks to invoke various tools that can capture the cause of large latencies

Options Trivia

Options parsing is not robust - example 1

```
# affinity will be 0
$ cyclictest -t -l100 -a0
$ cyclictest -t -l100 -a 0
$ cyclictest -t -l100 -a7 -a0
# affinity will be 7, with no error message
$ cyclictest -t -l100 -a7 -a 0
-a cpu affinity
```

Options Trivia

Options parsing is not robust - example 2

-t one thread per cpu

```
$ cyclictest -ant
T: 0 (26978) P: 0 I:1000 C: 2091 Min:
$ cyclictest -an -t
T: 0 (26980) P: 0 I:1000 C:
                              1928 Min:
T: 1 (26981) P: 0 I:1500 C:
                              1285 Min:
   cpu affinity
- a
-n clock_nanosleep()
```

Options Trivia

Options parsing is not robust

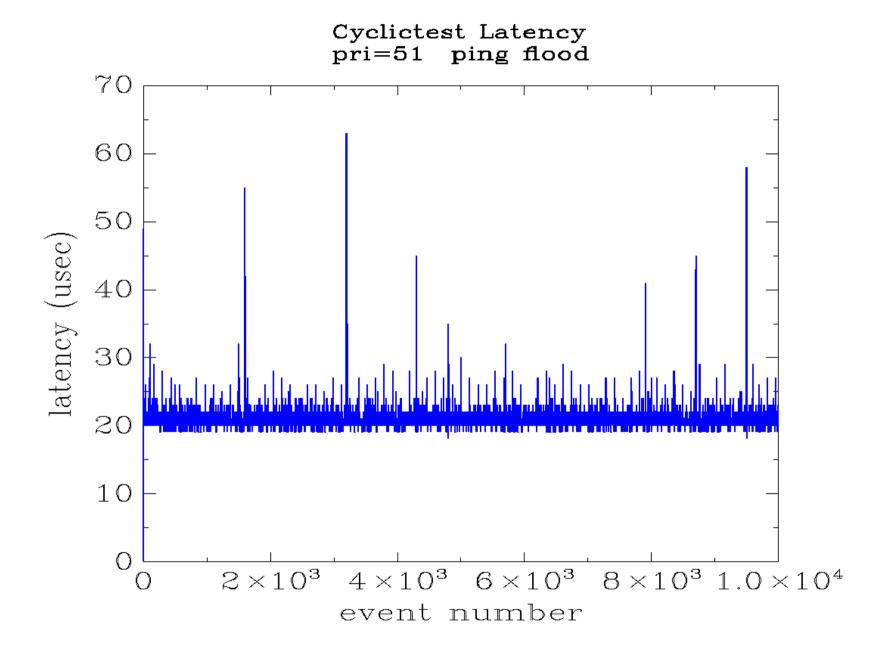
Best Practice:

- do not combine options
- specify each separately with a leading "-"

Third Data Format

Report each latency

```
$ cyclictest -q -n -t1 -p 48 -l 10000 -v
```



Hitting the RT sched throttle

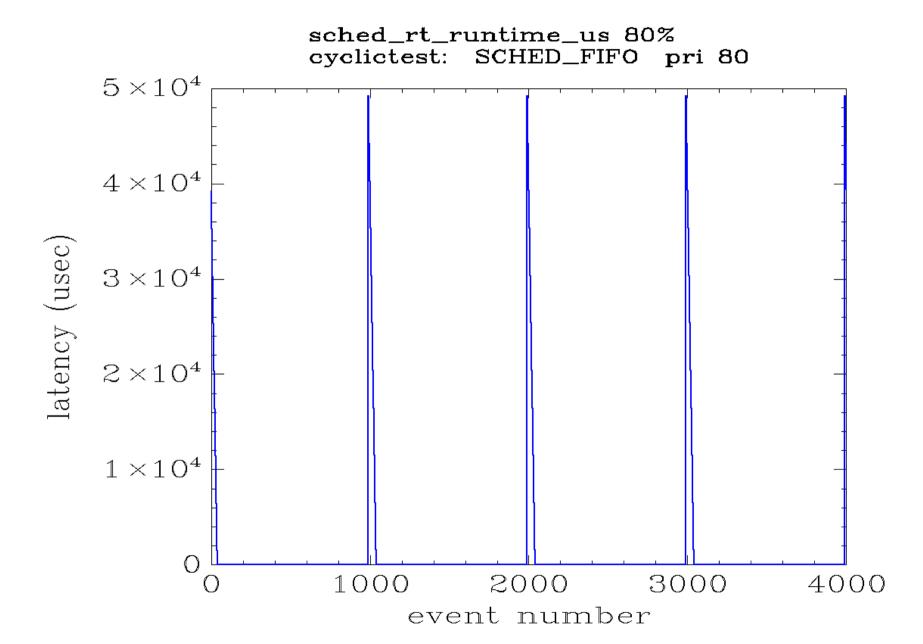
/proc/sys/kernel/sched_rt_runtime_us

/proc/sys/kernel/sched_rt_period_us

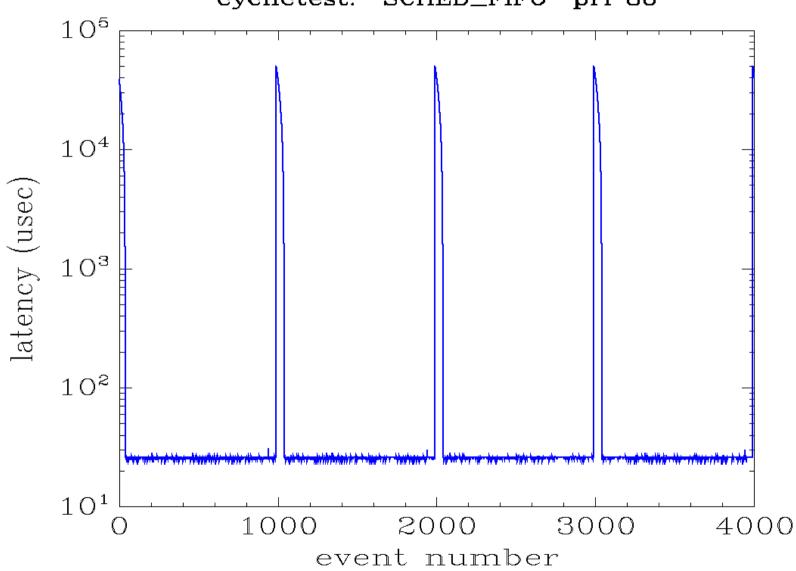
cyclictest: SCHED_FIFO priority=80

background load:

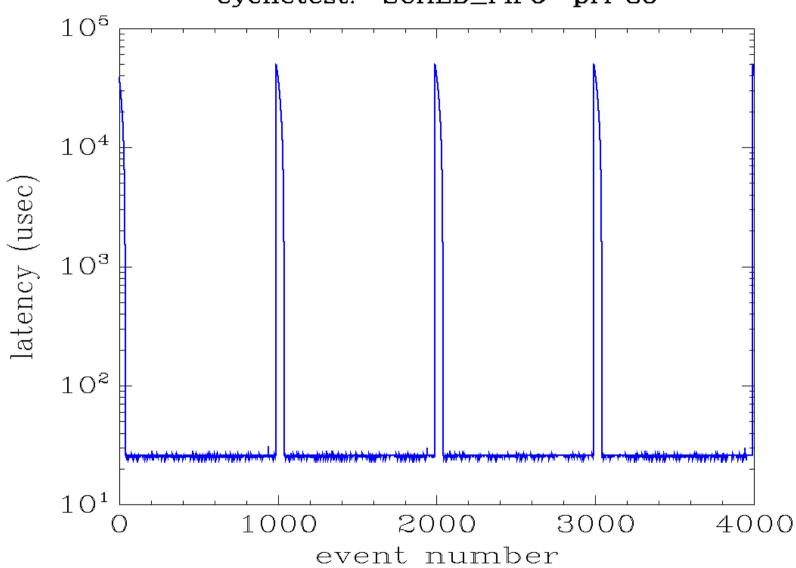
- continuous
- SCHED_FIFO priority=40

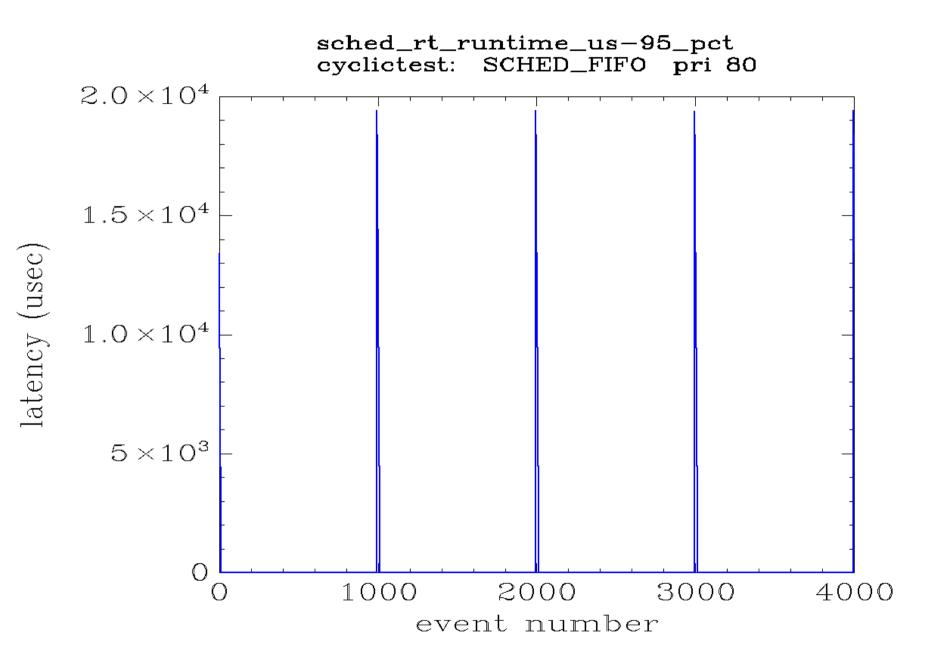


sched_rt_runtime_us 80% cyclictest: SCHED_FIFO pri 80

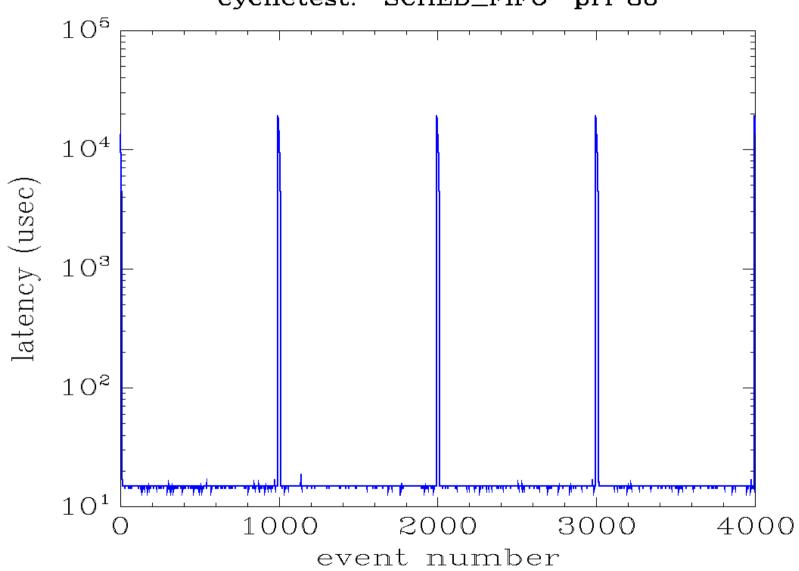


sched_rt_runtime_us 80% cyclictest: SCHED_FIFO pri 80





sched_rt_runtime_us-95_pct cyclictest: SCHED_FIFO pri 80



Hitting the RT sched throttle

```
/proc/sys/kernel/sched_rt_runtime_us
```

/proc/sys/kernel/sched_rt_period_us

cyclictest: SCHED_NORMAL

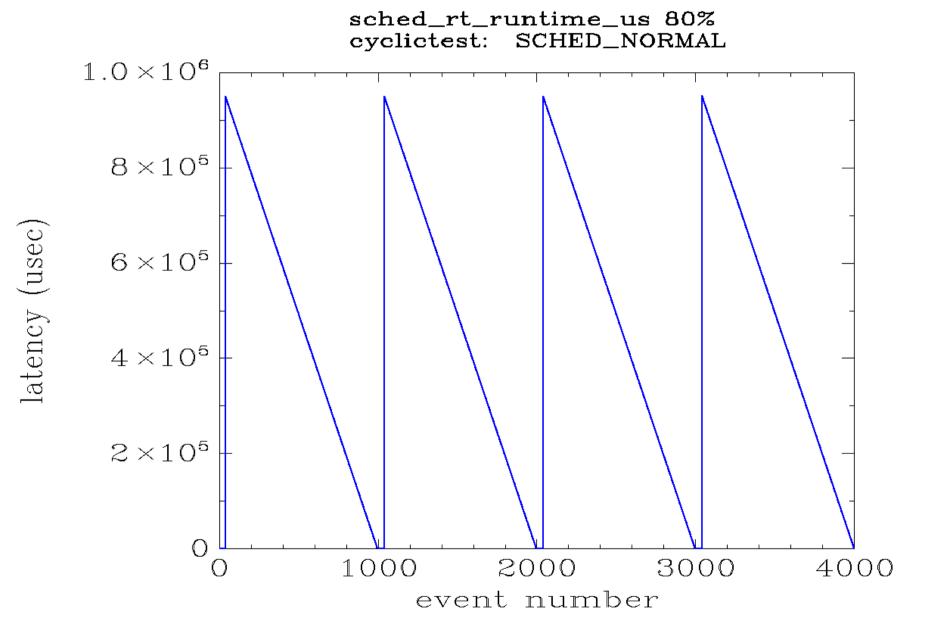
background load:

- continuous
- SCHED_FIFO priority=40

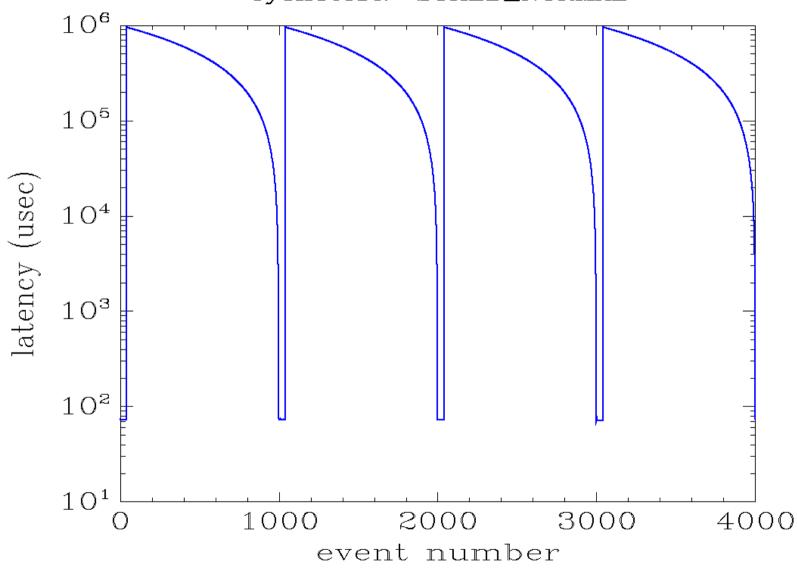
Hitting the RT sched throttle

Why is this measurement interesting???

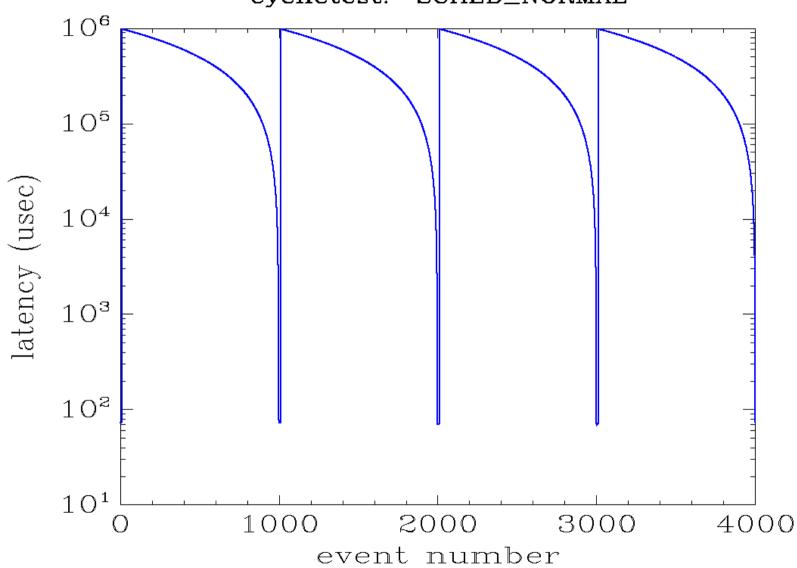
Gives a picture of how much cpu is NOT used by the real time tasks



sched_rt_runtime_us 80% cyclictest: SCHED_NORMAL



sched_rt_runtime_us-95_pct cyclictest: SCHED_NORMAL



Unusual Uses of Cyclictest

Rough measurement of response time of a real time application, without instrumenting task.

cyclictest latency =~ task latency + task work duration

This is not an accurate measurement, but it does provide a rough picture.

Response Time of a Task

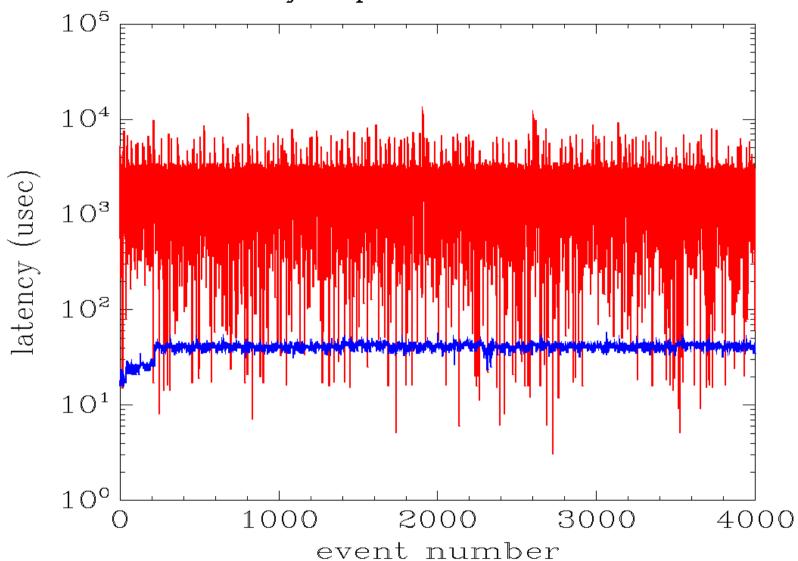
Cyclictest:

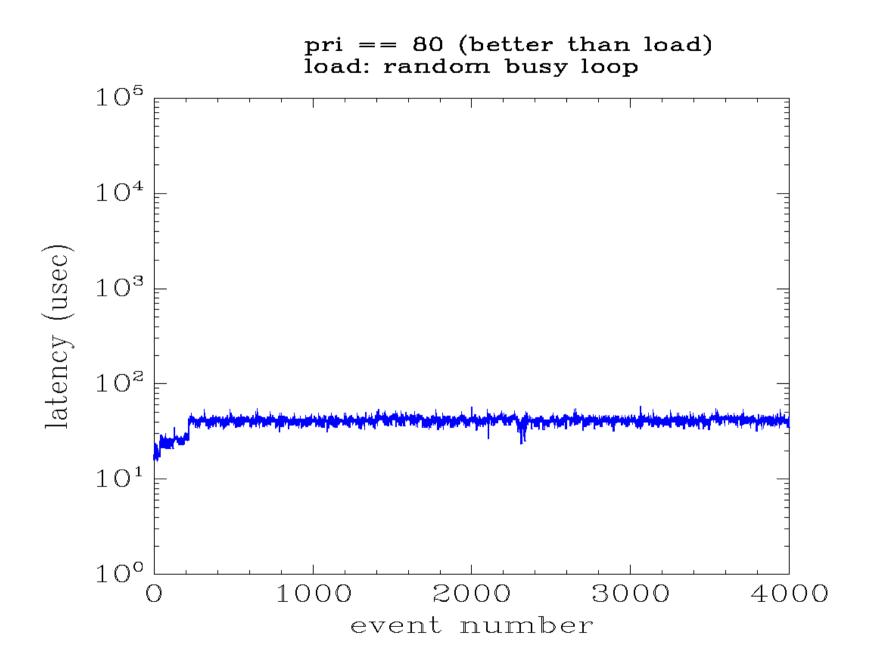
- (1) SCHED_FIFO priority=80 baseline latency
- (2) SCHED_FIFO priority=30 approximation of task response time

Real time application:

- busy loop (random number of iterations), followed by a short sleep
- SCHED_FIFO priority=40

blue: pri better than load red: pri worse than load load: random busy loop





pri == 30 (worse than load) load: random busy loop 105 10^4 latency (usec) 10³ 10² 101 10° 1000 2000 3000 4000 event number

Response Time of a Task

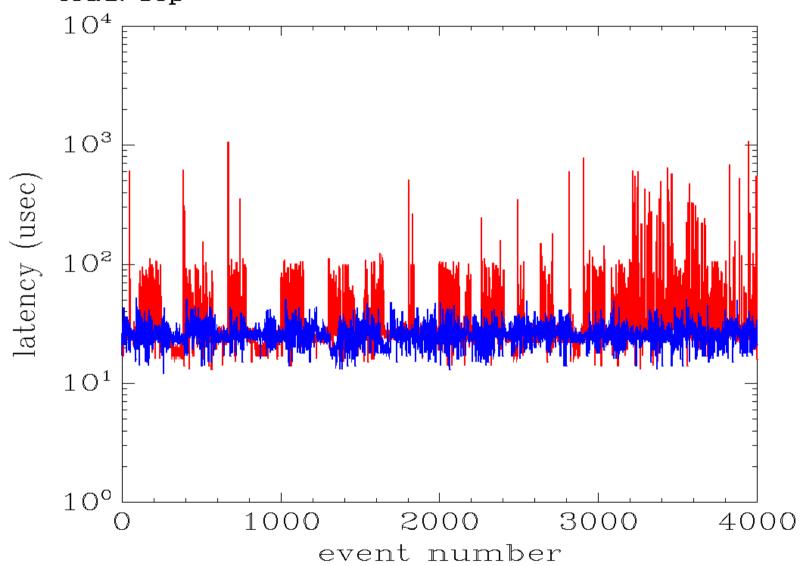
Cyclictest:

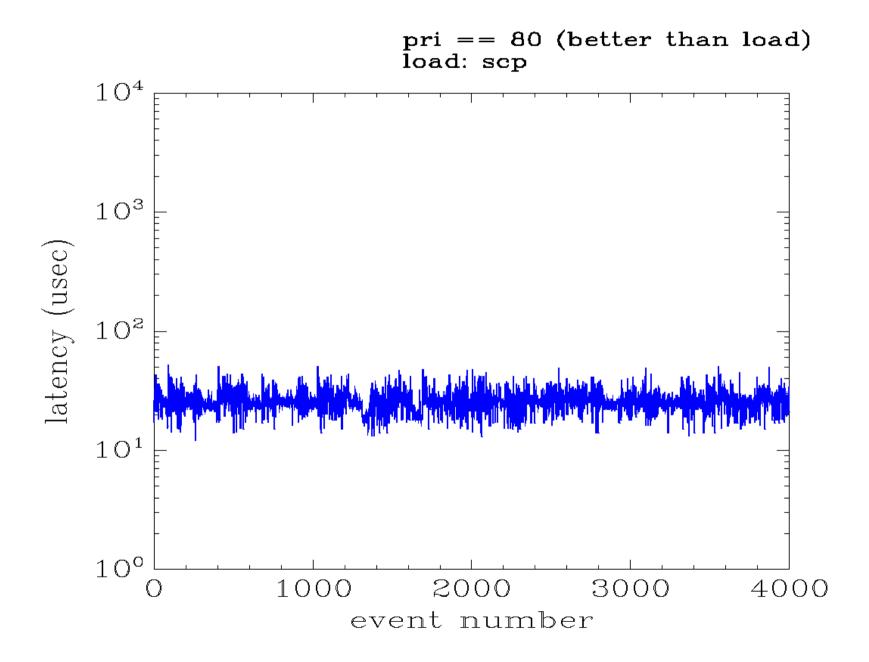
- (1) SCHED_FIFO priority=80 baseline latency
- (2) SCHED_FIFO priority=30 approximation of task response time

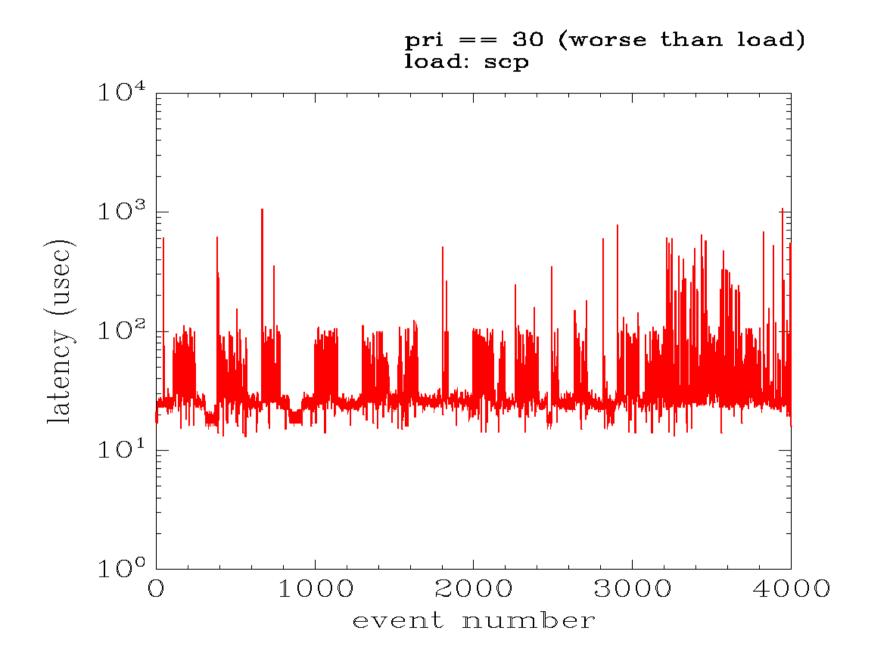
Real time application:

- recursive scp of a panda root file system
- SCHED_FIFO priority=40
- no guarantee of sleep between scp "transactions" - response time may include multiple transactions

blue: pri better than load red: pri worse than load load: scp

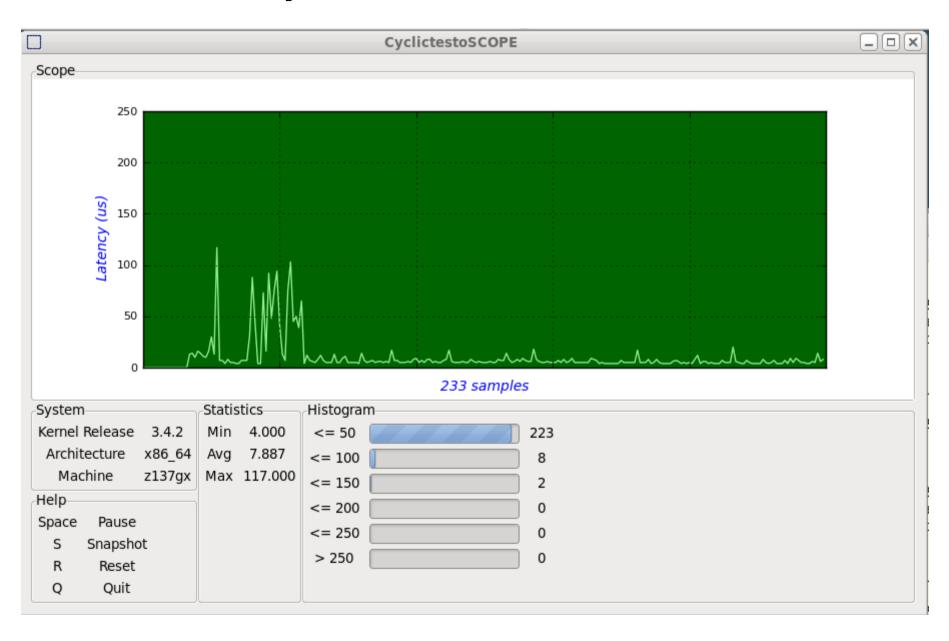






Demo - oscilloscope

oscilloscope screen shot



Fourth Data Format

Report time of each histogram overflow

Should be in next version of cyclictest (0.86?)

\$ cyclictest -q -h 400 -g 1000

The same information can be extracted from the third data format (-v), but this method is lower overhead.

Finding and Building

```
git clone \
 git://git.kernel.org/pub/scm/linux/kernel/git/clrkwllms/rt-tests.git
source: src/cyclictest/cyclictest.c
self-hosted:
  make
self-hosted without NUMA:
  make NUMA=0
cross-build without NUMA:
  make NUMA=0 CC="${CROSS_COMPILE}gcc"
```

Review

- Simple methodology captures all sources of latency fairly well
- Options must be used with care
- Options are powerful
- Different data formats are each useful
- Debug features can capture the cause of large latencies

THE END

Thank you for your attention...

Questions?

How to get a copy of the slides

1) leave a business card with me

2) frank.rowand@sonymobile.com