

APPLIED LINEAR ALGEBRA: PROGRAMMING ASSIGNMENT

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ABSTRACT. This is a recommended programming assignment. It deals with reading a matrix from a file, storing it in two formats, and performing matrix times a vector operation.

1. EFFICIENT WAY OF STORING SPARSE MATRICES

Now we describe an efficient way to store matrices, taking advantage that in a row, we may have only few non-zero entries. We describe next the CSR format (CSR stands for “compressed sparse row”). The CSR format is a popular way to store sparse matrices. For an $n \times m$ sparse matrix $A = (a_{ij})$, the CSR format exploits two one-dimensional integer arrays I and J and if the matrix is not Boolean (as the relation tables discussed in class) a real array “Data” is needed in addition to store the values/the actual entries a_{ij} of A .

Let A have at row i , $m_i \geq 1$ non-zero entries at positions $(i, j_1^{(i)})$, \dots , $(i, j_{m_i}^{(i)})$.

The one-dimensional array I has length $n + 1$. With $I[0] = 0$, we set

$$I[i] = I[i - 1] + m_i \text{ for } i \geq 1.$$

The array J has length $I[n]$, which is the total number of all nonzero entries of A . Similarly, the data array, Data, has the same length $I[n]$.

For each row $i = 1, \dots, n$ of A , we list consecutively in the one-dimensional array J the indices $j_s^{(i)}$, $s = 1, \dots, m_i$ starting at position $I[i - 1]$ till position $I[i] - 1$, that is

$$J[I[i - 1] + s - 1] = j_s^{(i)}, \text{ for } s = 1, \dots, m_i.$$

The data array is filled-in similarly, i.e., we let

$$\text{Data}[I[i - 1] + s - 1] = a_{i, j_s^{(i)}} \text{ for } s = 1, \dots, m_i.$$

Having sparse matrices stored in CSR format in practice it is useful to have algorithms that implement matrix operations such as A^T , matrix-matrix multiply $C = AB$. I.e., if A is stored in CSR format we need to store A^T in CSR format using only $\mathcal{O}(n)$ operations. Similarly, if the sparse matrices A and B are represented in CSR format with $\mathcal{O}(n)$ non-zero entries, we want to find an algorithm that computes and stores C in CSR format for $\mathcal{O}(n)$ storage and operations. All this is feasible for matrices corresponding to sparse relation tables.

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2. THE CURRENT ASSIGNMENT

- Read a matrix from a file. The file will have for each row two integers, i and j and a real number, which will be the (i, j) -the entry a_{ij} of the matrix A . The file may contain only the two integers i and j , then we assume that $a_{ij} = 1$. In both cases, all other entries of A that are not read from the file are assumed zero.
- Choose a way to store the matrix A . Possible options:
 - (1) Use two dimensional array $A[.,.]$. Initialize all entries of A with zeros. Then, for each pair of indices (i, j) that you have read, you set $A[i, j] = a_{ij}$ (the value which you have read, or if it is missing in the file, you set it to 1).
 - (2) Alternatively, you can use the more efficient way of storing A (the CSR format described above) as triplet of three one-dimensional arrays; two integer arrays $I[.]$, $J[.]$ and one array of reals $data[.]$.
 - (3) Use two one-dimensional arrays, $v[.]$ and $w[.]$, of corresponding length. Initialize $w[.]$ with some values (choose random values or unit values).
 - (4) Compute the product $v = Aw$.
 - (5) Compare the timings for the two formats of A (two-dimensional and the CSR) for fairly large matrices. Document your observation(s) with some conclusions.
 - (6) A simple file of any length $n > 1$ is (assuming that indices run from 0)

0	0	2.0
0	1	-1.0
1	0	-1.0
1	1	2.0
1	2	-1.0
\vdots	\vdots	\vdots
i	$i - 1$	-1.0
i	i	2.0
i	$i + 1$	-1.0
\vdots	\vdots	\vdots
n	$n - 1$	-1.0
n	n	2.0

- (7) A large set of matrix files is found at
http://www.cise.ufl.edu/research/sparse/matrices/list_by_id.html

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