Fast Exponentiation
CS 491 – Competitive Programming

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Spring 2024

Objectives

▶ Use binary representation to compute exponentials in logarithmic time.

▶ Use a similar technique with matrices to compute Fibonacci numbers.

Naïve Exponent

Remember the equation for n^m:

$b^0 = 1$

$b^n = b * b^{n-1}$

Recursive Implementation

```
1  int exponent(int base, int n) {
2    if (n == 0)
3      return 1;
4    return n * exponent(base,n-1);
5  }
```

Naïve Factorial, II

Iterative Implementation

```
1  int fact(int n) {
2    int out = 1;
3    while (n>0) {
4      out *= base;
5      --n;
6    }
7    return out;
8  }
```

▶ What is the time complexity?

Trick: Use a Binary Representation of the Exponent

- Q: What is 3^{22}?
    - Remember $a^x a^y = a^{x+y}$
- A: $3^{22} = 3^{10110_2} = 3^{10000_2} 3^{00100_2} 3^{00010_2} = 3^{16} 3^4 3^2$
  $= 1 \cdot 3^{10000_2} \times 0 \cdot 3^{01000_2} \times 1 \cdot 3^{00100_2} \times 1 \cdot 3^{00010_2} \times 0 \cdot 3^{00001_2}$

Implementation

Call with base=3 and n=22

```
1  int fexp(int b, int n) {
2    int out = 1;
3    int fact = b;
4
5    while (n>0) {
6      if (n & 1)
7        out *= fact;
8      fact *= fact;
9      n >>= 1;
10   }
11
12   return out;
13 }
```

► Initialize: out=1, fact=3, n=22

Implementation

Call with base=3 and n=22

```
int fexp(int b, int n) {
  int out = 1;
  int fact = b;

  while (n>0) {
    if (n & 1)
      out *= fact;
    fact *= fact;
    n >>= 1;
  }

  return out;
}
```

► Initialize: out=1, fact=3, n=22

► n (22) & 1 is 0:
  out=1, fact=9, n=11

Implementation

Call with base=3 and n=22

```
int fexp(int b, int n) {
  int out = 1;
  int fact = b;

  while (n>0) {
    if (n & 1)
      out *= fact;
    fact *= fact;
    n >>= 1;
  }

  return out;
}
```

▶ Initialize: out=1, fact=3, n=22

▶ n (22) & 1 is 0:
  out=1, fact=9, n=11

▶ n (11) & 1 is 1:
  out=9, fact=81, n=5

Implementation

Call with base=3 and n=22

```
int fexp(int b, int n) {
  int out = 1;
  int fact = b;

  while (n>0) {
    if (n & 1)
      out *= fact;
    fact *= fact;
    n >>= 1;
  }

  return out;
}
```

▶ Initialize: out=1, fact=3, n=22

▶ n (22) & 1 is 0:
out=1, fact=9, n=11

▶ n (11) & 1 is 1:
out=9, fact=81, n=5

▶ n (5) & 1 is 1:
out=729, fact=6561, n=2

Implementation

Call with base=3 and n=22

```
1  int fexp(int b, int n) {
2    int out = 1;
3    int fact = b;
4
5    while (n>0) {
6      if (n & 1)
7        out *= fact;
8      fact *= fact;
9      n >>= 1;
10   }
11
12   return out;
13 }
```

▶ Initialize: out=1, fact=3, n=22

▶ n (22) & 1 is 0:
out=1, fact=9, n=11

▶ n (11) & 1 is 1:
out=9, fact=81, n=5

▶ n (5) & 1 is 1:
out=729, fact=6561, n=2

▶ n (2) & 1 is 0:
out=729, fact=43046721, n=1

Implementation

Call with base=3 and n=22

```
  int fexp(int b, int n) {
    int out = 1;
    int fact = b;

    while (n>0) {
      if (n & 1)
        out *= fact;
      fact *= fact;
      n >>= 1;
    }

    return out;
  }
```

- ▶ Initialize: out=1, fact=3, n=22
- ▶ n (22) & 1 is 0:
  out=1, fact=9, n=11
- ▶ n (11) & 1 is 1:
  out=9, fact=81, n=5
- ▶ n (5) & 1 is 1:
  out=729, fact=6561, n=2
- ▶ n (2) & 1 is 0:
  out=729, fact=43046721, n=1
- ▶ n (1) & 1 is 1:
  out=31381059609, n=0

Calculating Fibonacci Numbers

Wrong Way

```
1   int fib(int n) {
2     if (n<2) return 1;
3     return fib(n-1) + fib(n-2)
4   }
```

Reasonable way

```
1   int fib(int n, int a=1, int b=1) {
2     if (n==1)
3       return a;
4     else
5       return fib(n-1,b,a+b);
6   }
```

Fast Fibonacci

▶ We can do even better! Consider this equation:

$$f'_1 = f_1 + f_2$$
$$f'_2 = f_1$$

▶ We can represent this in matrix form.

$$\begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \end{bmatrix} = \begin{bmatrix} 2 \\ 1 \end{bmatrix}$$

▶ Repeating

$$\begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 2 \\ 1 \end{bmatrix} = \begin{bmatrix} 3 \\ 2 \end{bmatrix}$$

▶ Again…!

$$\begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 3 \\ 2 \end{bmatrix} = \begin{bmatrix} 5 \\ 3 \end{bmatrix}$$

Square the Matrix

▶ Square the matrix to do multiple steps:

$$\begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} = \begin{bmatrix} 2 & 1 \\ 1 & 1 \end{bmatrix}$$

▶ Then…

$$\begin{bmatrix} 2 & 1 \\ 1 & 1 \end{bmatrix} \begin{bmatrix} 2 \\ 1 \end{bmatrix} = \begin{bmatrix} 5 \\ 3 \end{bmatrix}$$

▶ You can use the same technique as with fast exponents to "power up" this matrix and compute large Fibonacci numbers in logarithmic time.