# Edit Distance CS 491 CAP

Dr. Mattox Beckman

UNIVERSITY OF ILLINOIS AT URBANA-CHAMPAIGN
DEPARTMENT OF COMPUTER SCIENCE

Fall 2024

# **Objectives**

▶ Use DP to determine the edit distance between two strings.

### The Problem

- Given two strings s and t, how many "edits" does it take to tranform one to another?
  - Edit = insert, delete, or change.
  - Usually each of these "costs" one unit.
- Usually called the Levinstein Distance
- Examples:
  - changing DATA to BETA needs 2 steps.
  - changing ETA to BETA needs 1 step.
  - changing GRETA to BETA needs 2 steps.

### The Näive Algorithm

```
Base Cases
1 // Thanks, Wikipedia!
   int LD(string s, int len_s, string t, int len_t) {
     int cost;
4
     /* base case: empty strings */
     if (len_s == 0) return len_t;
6
     if (len t == 0) return len s;
7
8
     /* test if last characters of the strings match */
9
     if (s[len s-1] == t[len t-1])
10
         cost = 0:
11
     else
12
         cost = 1;
13
```

# The Näive Algorithm, ctd

#### Recursive Case

### How can you convert this to DP?

You have to decide what is the state being remembered....

### Dynamic Programming using Memoization

### Base Cases

```
int LD(const char *s, int len_s, const char *t, int len_t)
2
     vvi dp = vvi(len_s + 1, vi(len t +1));
3
     int cost;
4
5
     for(int i=0; i<=len s; ++i)
6
        dp[i][0] = i;
7
8
     for(int i=0; i<=len t; ++i)
9
        dp[0][i] = i;
10
```

# Dynamic Programming using Memoization, ctd

#### Memoized Part

```
for(int i=1; i<=len s; ++i)
11
        for(j=1; j<=len_t; ++j) {
12
           cost = s[i] == t[j] ? 0 : 1;
13
14
           dp[i][j] = minimum(dp[i-1][j] + 1,
15
                                dp[i][j-1] + 1,
16
                                dp[i-1][j-1] + cost);
17
           }
18
      return dp[len s][len t];
19
   }
20
```