CS 421 — Type Semantics Activity (Polytype Version)

Mattox Beckman

The Rules

Constants

 $\frac{}{\Gamma \vdash n : \mathtt{int}}$ Const, when n is an integer.

Similarly for True and False.

Variables

$$\frac{1}{\Gamma \vdash x : \sigma}$$
 VAR, when $x : \sigma \in \Gamma$

Binary Arithmetic

$$\frac{\Gamma \vdash e_1 : \mathtt{int} \qquad \Gamma \vdash e_2 : \mathtt{int}}{\Gamma \vdash e_1 \oplus e_2 : \mathtt{int}} \mathtt{BinOp}$$

Integer Relations

$$rac{\Gamma dash e_1 : \mathtt{int} \qquad \Gamma dash e_2 : \mathtt{int}}{\Gamma dash e_1 \sim e_2 : \mathtt{bool}} \, \mathtt{RelOp}$$

If

$$\frac{\Gamma \vdash e_1 : \mathtt{bool} \qquad \Gamma \vdash e_2 : \tau \qquad \Gamma \vdash e_3 : \tau}{\Gamma \vdash \mathtt{if} \ e_1 \ \mathtt{then} \ e_2 \ \mathtt{else} \ e_3 : \tau} \ \mathtt{If}$$

Application

$$\frac{\Gamma \vdash e_1 : \tau \to \tau' \qquad \Gamma \vdash e_2 : \tau}{\Gamma \vdash e_1 e_2 : \tau'} \mathbf{App}$$

Abstraction

$$\frac{\Gamma \cup \{x:\tau\} \vdash e:\tau'}{\Gamma \vdash \lambda x.e:\tau \to \tau'} \, \mathbf{Abs}$$

Let

$$\frac{\Gamma \vdash e_1 : \sigma \qquad \Gamma \cup [x : \sigma] \vdash e_2 : \tau}{\Gamma \vdash \mathbf{let} \ \mathbf{x} = e_1 \ \mathbf{in} \ e_2 : \tau} \ \mathbf{Let}$$

Gen

$$\frac{\Gamma \vdash e : \sigma}{\Gamma \vdash e : \forall \alpha. \sigma}$$
 GEN, where α is not free in Γ

Inst

$$\frac{\Gamma \vdash e : \sigma'}{\Gamma \vdash e : \sigma}$$
 Inst, where $\sigma' \geq \sigma$

Proofs

Create proofs for the following judgements according to the given rules.

Problem 1)

 $\mbox{\{id:} \forall \alpha.\alpha \rightarrow \alpha \mbox{ , y:} \mbox{Int} \} \vdash (id \ y) : \ Int \label{eq:partial}$

Problem 2)

 $\{y: Int, z: String\} \vdash (\lambda f.(f y, f z)) (\lambda x.x) : (Int, String)$

Problem 3)

 ${x:Int,y:String} \vdash let f = \lambda x.x in (f x, f y) : (Int, String)$