

Name: _____

CS 421 — Big Step Semantics Activity

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Reading the Judgments

- $\langle i, \sigma \rangle \Downarrow_e e$ “Expression i in environment σ evaluates to e .”
- $\langle i, \sigma \rangle \Downarrow_b e$ “Boolean Expression i in environment σ evaluates to e .”
- $\langle i, \sigma \rangle \Downarrow \sigma'$ “Executing statement i in environment σ results in environment σ' .”

The Rules

$$\frac{}{\langle i, \sigma \rangle \Downarrow_e i} \text{ if } i \text{ is an integer.}$$

$$\frac{}{\langle u, \sigma \rangle \Downarrow_e v} \text{ if } u := v \in \sigma.$$

$$\frac{\langle e_1, \sigma \rangle \Downarrow_e v_1 \quad \langle e_2, \sigma \rangle \Downarrow_e v_2}{\langle e_1 \oplus e_2, \sigma \rangle \Downarrow_e v_1 \oplus v_2}$$

$$\frac{}{\langle \mathbf{true}, \sigma \rangle \Downarrow_b \mathbf{true}} \quad \frac{}{\langle \mathbf{false}, \sigma \rangle \Downarrow_b \mathbf{false}}$$

$$\frac{}{\langle u, \sigma \rangle \Downarrow_b v} \text{ if } u := v \in \sigma.$$

$$\frac{}{\langle \mathbf{skip}, \sigma \rangle \Downarrow \sigma}$$

$$\frac{\langle e, \sigma \rangle \Downarrow_e v}{\langle x := e, \sigma \rangle \Downarrow \sigma[x := v]}$$

$$\frac{\langle S_1, \sigma \rangle \Downarrow \sigma' \quad \langle S_2, \sigma' \rangle \Downarrow \sigma''}{\langle S_1; S_2, \sigma \rangle \Downarrow \sigma''}$$

$$\frac{\langle e_1, \sigma \rangle \Downarrow_e v_1 \quad \langle e_2, \sigma \rangle \Downarrow_e v_2}{\langle e_1 \sim e_2, \sigma \rangle \Downarrow_b v_1 \sim v_2}$$

$$\frac{\langle B, \sigma \rangle \Downarrow_b \mathbf{true} \quad \langle S_1, \sigma \rangle \Downarrow \sigma'}{\langle \mathbf{if } B \mathbf{ then } S_1 \mathbf{ else } S_2 \mathbf{ fi}, \sigma \rangle \Downarrow \sigma'} \quad \frac{\langle B, \sigma \rangle \Downarrow_b \mathbf{false} \quad \langle S_2, \sigma \rangle \Downarrow \sigma'}{\langle \mathbf{if } B \mathbf{ then } S_1 \mathbf{ else } S_2 \mathbf{ fi}, \sigma \rangle \Downarrow \sigma'}$$

$$\frac{\langle B, \sigma \rangle \Downarrow_b \mathbf{false}}{\langle \mathbf{while } B \mathbf{ do } S \mathbf{ od}, \sigma \rangle \Downarrow \sigma} \quad \frac{\langle B, \sigma \rangle \Downarrow_b \mathbf{true} \quad \langle \mathbf{while } B \mathbf{ do } S \mathbf{ od}, \sigma' \rangle \Downarrow \sigma'' \quad \langle S, \sigma \rangle \Downarrow \sigma'}{\langle \mathbf{while } B \mathbf{ do } S \mathbf{ od}, \sigma \rangle \Downarrow \sigma''}$$

Reductions

Reduce the following programs according to the semantic rules given.

Problem 1)

$\langle x + y + z, \{x:=10, y:=20, z:=12\} \rangle$

Problem 2)

$\langle t:=a, a:=b, b:=t, \{a:=5, b:=10\} \rangle$

Problem 3)

```
<if x>y then m:=x-y else m:=y-x fi, {x:=10,y:=30}>
```

Problem 4)

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<while x > 1 do x:=x/2 od, {x:=2}>
```

Make your own rules!

Problem 5)

Write a rule to explain the `when B S` statement. It executes S only if B is true.

Problem 6)

Write a rule for `do S while B od`. It is like `while`, but executes S at least one time.