
CPS Activity

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Conversion

$$C[f \text{ arg} = e] \Rightarrow f \text{ arg } k = C[e]_k$$

$$\begin{aligned} C[a]_k &\Rightarrow k \ a \\ C[f \text{ arg}]_k &\Rightarrow f \text{ arg } k \\ C[f \text{ arg}]_k &\Rightarrow C[\text{arg}]_{(\lambda v. f \ v \ k)}, \text{ where } v \text{ is fresh.} \\ C[e_1 + e_2]_k &\Rightarrow k(e_1 + e_2) \\ C[e_1 + e_2]_k &\Rightarrow C[e_1]_{(\lambda v. \rightarrow k(v + e_2))} \text{ where } v \text{ is fresh.} \\ C[e_1 + e_2]_k &\Rightarrow C[e_1]_{(\lambda v_1. \rightarrow C[e_2]_{\lambda v_2. \rightarrow k(v_1 + v_2)})} \text{ where } v_1 \text{ and } v_2 \text{ are fresh.} \end{aligned}$$

Convert To

Convert the following functions to CPS:

```
1 sumList [] = 0
2 sumList (x:xs) = x + sumList xs
```

2. Assume f is written in direct style.

```
1 map f [] = []
2 map f (x:xs) = f x : map f xs
```

3. Assume f is written in CPS and takes one continuation.

```
1 map f [] = []
2 map f (x:xs) = f x : map f xs
```

4. Convert the following code to CPS.

```
1 min a b = if a < b then a else b
2 min4 a b c d = min (min a b) (min c d)
```

More CPS Transforms

5. Write the CPS transform for the `if` expression. You will need two cases.

Reordering Computations

6. Suppose you have a calculator which has an accumulator and a list of instructions. `Add i` adds `i` to the accumulator, and `Sub i` subtracts `i` from the accumulator.

```
1  data Calc = Add Integer
2             | Sub Integer
3  deriving (Eq, Show)
```

The only problem is that our accumulator cannot be negative! Use continuations to fix this.

Here's the original calculator:

```
1  calc xx = aux 0 xx
2  where aux a [] = a
3         aux a ((Add i):xs) = aux (a+i) xs
4         aux a ((Sub i):xs) = aux (a-i) xs
```

Hint: you will need *two* continuations to make this work.