v-FORTH 1.5

ZX Spectrum Next version

1990-2021 Matteo Vitturi

Introduction &
Technical Info

Build 20210425

1. Forewords

This document introduces a Forth implementation suitable to run on Sinclair ZX Spectrum Next.

This is in essence a FIG-Forth ported to the new **Sinclair ZX Spectrum Next** based on my previous work **v-Forth 1.413** available at https://sites.google.com/view/vforth/vforth1413 and at https://github.com/mattsteeldue/vforth.

This version **v-Forth 1.5** is available at https://github.com/mattsteeldue/vforth-next. The main difference from the previous version is that it uses a dedicated file on SD instead of on ZX Microdrive cartridges to provide a Block/Screen facility. Even if this is a "working" piece of software, the porting is still a work-in-progress, there are many things to do.

Disclaimer

Copying, modifying and distributing this software is allowed provided this copyright notice is kept.

This work is available "as-is" with no whatsoever warranty.

Changes may occur at any time without notification.

The author – me – is not a native English speaking and, for certain, you will find grammatical errors. In case, it would be very appreciated if you could drop me a line with any suggestion and/or correction at matteo _underscore_ vitturi@yahoo.com. I am not able to write a longer disclaimer than the above.

2. Getting started

The most recent version of this software can be downloaded from GitHub repository as .zip file at

https://github.com/mattsteeldue/vforth-next/tree/master/download

In alternative, the same executable programs are available in the same repository:

https://github.com/mattsteeldue/vforth-next/tree/master/SD/tools/vforth

Unzip or copy the software to "C:/tools/vForth" directory inside your Next's SD card so it appears as follow:

| 100 | | 114 | 8.4 T T1.T1 T1.T1 |
|-----------------|-----|-------------|---------------------|
| € [.] | | <dir></dir> | 30/10/2020 23:28 |
| [doc] | | <dir></dir> | 30/10/2020 23:28 |
| [src] | | <dir></dir> | 30/10/2020 23:28 |
| Blocks-64 | bin | 16.776.704 | 30/10/2020 23:32-a- |
| Forth 15 | bas | 655 | 30/10/2020 23:32-a- |
| forth 15e | bin | 9.344 | 30/10/2020 23:32-a |
| Forth 15_loader | bas | 573 | 07/09/2020 22:16-a |

If you wish to use a different directory instead of C:/forth, you need to modify the paths in the two Basic programs.

The Forth System is activated by running a Basic program **C:/forth/forth15_loader.bas**. This can be done using the Browser and selecting it, then clicking ENTER.



The Basic loader **forth15_loader.bas** frees upper memory setting RAMTOP to address 25345; it loads **forth15e.bin** (the Forth core) and then it loads a smaller Basic launcher **Forth15.bas** you can customize for your purposes.

```
now LOADing code...
"forth15e.bin" CODE 25600
sleep 5
LOADing wrapper...
"Forth15.bas"
```

A Splash screen displays "Version number" and "Build date" followed by some technical system information. Within a few seconds the system will ask if you would like to "Run Scr# 11 autoexec": the only way to refuse is using [N] key. It is anyway a good idea allowing Forth to continue and LOAD Screen # 11 that in turn loads a few useful Screens which make available, among the others, two pecuiliar words: EDIT the "Screen Editor" and SEE the "Debugger Inspector". This phase is executed only at *first* startup, but you can run it again using AUTOEXEC word.

```
v-Forth 1.5 NextZXO5 version
build 20210407
1990-2021 Matteo Vitturi
28.0 MHz Z80n CPU Speed.
19142 bytes free in Dictionary.
65533 bytes free in Heap.
Autoexec says :
Do you wish to load Scr# 11 ? (Y/n) ■
```

The Basic launcher **Forth15.bas** usually auto-starts the first time at LINE 20, so you won't care, but in case you STOP it or the Forth system encounters an Error that forces it to return to Basic, you have two main choices:

- a. give RUN: This does a WARM start, preserving your previous work and buffer status.
- b. give RUN 20: This does a COLD start, restoring all as you just loaded from SD card.

Before entering Forth, the Basic launcher does an <code>OPEN# 13</code>, "o>output.txt" that can be later chosen from Forth via <code>13 SELECT</code> to collect any output you send to this output channel. To restore sending output to video there is a easy <code>VIDEO</code> definition that simply does <code>2 SELECT</code>.

You can modify the Basic launcher by adding commands to OPEN# any other file *for read* so that it can be fed to Forth as a text source; for example you can add the following Basic line:

```
92 OPEN # 12, "src/Z80N-asm.f"
```

Later, this allows Forth to load such a source file using the following:

-12 LOAD

In this case, a negative number such as -12 says LOAD Forth definition to start reading text from input stream #12 instead of loading from Screen # "-12", that doesn't exists. This feature, i.e. passing a negative "screen" number to LOAD, is not Forth standard, but an idea of mine.

Anyway, there is no more need to OPEN# a stream from Basic, since two new specific definitions allow you to include source from any file: INCLUDE and NEEDS. See chapter "5. The dictionary" for more details.

In this Forth implementation I preferred LAYER 1,2 display mode to allow 64 character per line: this is quite necessary to be able to display a whole 1024 characters in a single screen.

If you prefer LAYER 1,1 you can add a line 61 in Forth15.bas wrapper as follow

61 LAYER 1,1: PRINT CHR\$ 30; CHR\$ 4;

to switch to LAYER 1,1 and condensed character set. The result is quite poor in my opinion:

v-Forth 1.5 NextZXOS version
build 20210407
1990-2021 Matteo Vitturi
28.0 MHz Z8On CPU Speed.
19142 bytes free in Dictionary.
65533 bytes free in Heap.
Autoexec says :
Do you wish to load Scr# 11 ? (Y/n) ■

The Full Screen Editor Utility

The EDIT definition is available after you give: NEEDS EDIT or in the old way 190 LOAD.

On this Forth system, as in many others, a Screen has 1.024 bytes of text-data spread in 16 lines, 64 bytes each.

This "Full Screen Editor Utility" is invoked using EDIT definition. This enters a simple full-screen editor that can modify current Screen, one screen at a time. While using EDIT, you are allowed move to next Screen or to previous Screen using the command explained below.

Remember: to quit EDIT editor, you have to use [Edit] key followed by [Q] key, in a way similar as Unix vi editor

This editor works only while the display-mode allow 64 character per line at least.

EDIT ---

For example, to select, show and edit Screen # 196 you can give:

```
DECIMAL 196 LIST (to set 196 the "current screen")

EDIT (to enter the editor on "current screen")
```

It shows a header reporting the Screen number and a line-ruler followed by 16 lines that compose the Screen itself.

A flashing cursor is visible at home position. The cursor has two flashing mode to distinguish CAPS-LOCK enabled or disabled.

The cursor keys, i.e. [Shift] key + 5 / 6 / 7 / 8 keys, allow the flashing cursor to be moved across the screen to point the current position inside the Screen, so text can be typed at any position in the Screen.

Current cursor positions (**row** number and **col**umn number) is shown at bottom status bar along with current character, **dec**imal ASCII-code and **hex**adecimal code of it.

Pad line shows the current PAD content. Line oriented commands handle or work with PAD. See the "Line Editor" chapter.

After [Edit] key (i.e. Shift + 1) the Editor recognizes the following single key-stroke commands:

[Edit] + Q : Quit EDIT Utility

[Edit] + U: Undo, that is re-read current screen from disk ignoring any modification done since last FLUSH.

[Edit] + N : go to Next screen

[Edit] + B : go Back to previous screen

[Edit] + S: make Space at current row shifting down lower lines; last line will be lost.

[Edit] + D : Delete current line shifting up lower lines. Also, copy current line to PAD

[Edit] + I: Insert PAD at current line: it does S and R together.

[Edit] + R: Replace current line with the current PAD content.

[Edit] + H : Hold or take current line content and keep it in PAD

[Edit] + P: accepts two hexadecimal digits representing a byte and Put it at cursor position.

any other key has no meaning and returns the flashing cursor back to its position.

[Delete] removes a character at current cursor position, shifting left the rest of the line.

[Break] (that is Caps-Shift + SPACE) inserts a space at current cursor position, shifting right the rest of the line.

Beware, any modification immediately affects the underlying Buffers, so if you mess things too much so that [Edit] + U is not enough, there is only a way to recover it: using EMPTY-BUFFERS to erase all buffers without flushing to disk.

This "Full Screen Editor" is a work-in-progress and can be improved if needed.

Search and Locate Utility

The following definition are available after you give alternatively:

```
NEEDS LOCATE or
NEEDS GREP or
NEEDS BSEARCH or
NEEDS COMPARE
```

Using the old way 70 LOAD will compile them all reading from Screens# 70-75.

LOCATE ---

Used in the form

LOCATE cccc

this word examines all Screens (between 1 and 1000) looking for the definition of **cccc** and shows the Screen where it found the first occurrence, and makes it the "current screen", just like LIST for example:

LOCATE COMPARE

takes a few seconds to search in which Screen COMPARE is defined, and if found it shows the Screen using LIST.

```
Scr# 70

0 .( Compare Utility. ) CR

1 \ Compare two strings and return 0 if they're equal

2 \ or 1 if s1 \ s2 or -1 if s1 \ s2

3 : COMPARE (a1 c1 a2 c2 -- -1001)

4 ROT 2DUP SWAP - \R \ a1 a2 c2 c1 \ c1-c2

5 MIN \ a1 a2 min(c2,c1) \ c1-c2

6 (COMPARE) \ b \ c1-c2

7 R \ SWAP ?DUP \ c1-c2 b b b \ 0

8 IF \ c1-c2 b that is not zero

9 SWAP DROP \ b that is 1 or -1

10 ELSE \ c1-c2

11 1 SWAP * \ sign(c1-c2) or zero

12 THEN;

13 -->

14 CREATE S1 ," Hello world!"

15 CREATE S2 ," Hello world?"
```

GREP ---

Used in the form

GREP cccc

this word examines all Screens looking for any occurrence of word cccc showing them in a table form; for example

GREP COMPARE

will take some more time to complete and gives something like the following

BSEARCH n1 n2 ---

Used in the form

```
n1 n2 BSEARCH cccc
```

this word examines all Screens between n1 and n2 looking for any occurrence of word \mathbf{ccc} showing them in a table form. This definition is used by GREP that in fact is defined as 1 1000 BSEARCH .

COMPARE al cl a2 c2 ---

Given two string descriptors, that is address and length, (a1, c1) and (a2, c2), this definition compares the two strings and returns

- 1 if they're equal
- 2 if String1 > String2
- -1 if String1 < String2

For example:

```
CREATE S1 ," Hello world!"
CREATE S2 ," Hello world?"
S1 COUNT S2 COUNT COMPARE .
```

will print -1 since the two strings differs only for the last character and ! is before ? in the ASCII table, so the string comparison S1 < S2 is true. Compare the result of the following two rows:

```
S2 COUNT S1 COUNT COMPARE . S1 COUNT S1 COUNT COMPARE .
```

Debugger Utility

The following definitions are available after loading if you give <code>NEEDS SEE</code> or via 20 <code>LOAD</code> or after a regular <code>AUTOEXEC</code>.

SEE ---

Used in the form

SEE cccc

it will print how the word cccc is defined along with its NFA, CFA, PFA data.

If cccc is a regular colon-definition the result will show something close to the original source the word was defined from.

For example, the word **TYPE** is a colon-definition defined as follow:

```
: type ( a n -- )
    over + swap
    ?Do
        i c@ emit
    Loop
;
```

If you give

SEE TYPE

it will print (depending on which build you're running) something like:

```
Nfa: 727F 84
Lfa: 7284 COUNT
Cfa: 7286 6BCF
OVER + SWAP (?DO) 12 I C@ EMIT (LOOP) -8 ok
```

The first line shows **TYPE** NFA (\$727F in this case) followed by \$84 that is the counter byte of a 4-bytes length name: the counter byte always has the most significant bit set, so \$80 is added to \$04 giving \$84.

The second line is **TYPE** LFA (\$7190) which address holds a pointer to COUNT NFA i.e. the previous definition in dictionary.

The third line is **TYPE** CFA which address holds a pointer to the machine-code part of a regular colon-definition (that's the ENTER routine of every colon-definition).

The fourth line represents the PFA and, in this case, is in some way a definition "de-compilation" where literals and offsets are shown in "inverse video" mode.

Another example, the word **NIP** isn't a colon-definition, but it is coded directly in machine-code as follow:

```
CODE nip ( n1 n2 -- n2 )
POP HL|
EX(SP)HL
Next
C;
```

and if you give

SEE NIP

it will print

```
Nfa: 6AD5 83

Lfa: 6AD9 DROP

Cfa: 6ADB 6ADD

6ADD E1 E3 DD E9 84 54 55 43 ac]I.TUC

6AOD CB D5 6A EA 6A E1 D1 E5 KUjjjaQe
```

In this case, since **NIP** is not a colon-definition, its PFA part is just a **DUMP** (you can stop at any time).

Again, the first line shows **NIP**'s CFA (\$6AD9 in this case) followed by \$83, the counter byte, that indicates a 3-bytes length word name.

The second line is NIP's LFA (\$6AD9) which has a pointer to DROP's NFA, that is the previous definition in dictionary.

The third line is **NIP**'s CFA (\$6ADB) which cell points to the following address (\$6ADD), that is **NIP**'s PFA where the small piece of machine-code lies. We should be able to see **E1** for POP HL, **E3** for EX (SP), HL and **DDE9** for JP (IX) to "Next" inner interpreter address **\$6434**.

The following bytes 84 54 55 43 are the beginning of the following definition in dictionary (TUCK in this case).

This utility is not perfect, but is a good way to debug and understand Forth.

DUMP a ---

Performs a "dump" of a memory area from address a for 128 bytes or until [Break] is pressed. Visualization is always in hexadecimal, current base is maintained. For example:

```
DECIMAL 448 DUMP
```

will print the Standard ROM content starting from address 448 (\$01C0):

```
01C0
     4C 49 53 D4
                  4C 45 D4 50 LISTLETP
01C8
     41 55 53 C5
                  4E 45 58 D4
                              AUSENEXT
01D0
     50 4F 4B C5
                  50 52 49 4E POKEPRIN
01D8
     D4 50 4C 4F D4 52 55 CE
                              TPLOTRUN
     53 41 56 C5
                  52 41 4E 44
01E0
                               SAVERAND
01E8
     4F 4D 49 5A
                  C5 49 C6 43
                              OMIZEIFC
01F0
     4C D3 44 52
                  41 D7 43 4C
                              LSDRAWCL
01F8 45 41 D2 52 45 54 55 52 EARRETUR
```

.WORD a ---

Given a CFA, this word prints the ID. It is used by SEE to perform some word "decompilation"

.s ---

Prints the current status of the Calculator Stack. For example, supposing to start with an empty stack,

0 1 2 3 .s

will print

0 1 2 3 ok

DEPTH --- n

It leaves the depth of the Calculator Stack before it was executed. For example, supposing to start with an empty stack,

0 1 2 DEPTH .

will print

3 ok

3. Technical specifications

CPU Registers

Registers are used in the in the following way:

- AF Used for normal operations.
- BC Forth Instruction Pointer: should be preserved on enter-exit a definition and during ROM/OS calls.
- DE Free (Low part when used for 32-bit manipulations)
- HL Work Register (High part when used for 32-bit manipulations)
- AF'- Not used, somewhere used for backup purpose
- BC'- Used only in I/O operations: available in fast Interrupt via EXX
- DE'- Used only in I/O operations: available in fast Interrupt via EXX
- HL'- Used only in I/O operations: available in fast Interrupt via EXX (saved at startup from Basic)
- SP Calculator Stack Pointer
- IX Used to point to the Forth "inner-interpreter" (this saves 2 T-States compared to a normal Jump). See (NEXT) word.
- IY Used by ZX System, must be preserved to let keyboard to be served

Much care has been taken to avoid any use of alternate registers (at least with interrupts enabled). This should allow users to create their own fast-response interrupt routine with EXX instead of pushing away all registers.

Single Cell 16 bits Integer Number Encoding

A 16 bits *integer* represents an integer number between –32768 and +32767 inclusive. The sign is kept in the most significant bit. Alternatively, the it represents an *unsigned integer* between 0 and +65535.

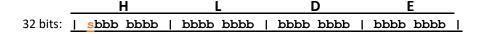
In the CPU registers, an *integer* is kept in H and L where H is the most significant part.

In memory, an *integer* is stored in two contiguous bytes in "little-endian" way, that is, the lower address has the least significant part, the in L register. The byte at higher address has the most significant part, the one in H register, as usual for Zilog Z80.

Double cell 32 bits Integer Number Encoding

The second integer format requires two *integers* that form a 32 bits number said *double* or *long* that allows an integer between -2.147.483.648 and +2.147.483.647, and the sign is kept on the most significant bit of the first *integer*.

Imagine a *double integer* kept in CPU register in the in this way:



using register H, L, D and E, with the most significant part in H, and the least in E.

Then, on Calculator Stack the *double integer* requires four contiguous bytes split in the two *integers* that forms it with the most significant integer (HL) on top of Calculator Stack (i.e. in the lower addresses), and the least significant integer (DE) the second element from top is in the higher address, that is the second element from top. so it appears as L H E D,

| CPU | Calculator Stack | |
|-----|-----------------------|--|
| D | SP + 3 | |
| E | SP + 2 | |
| Н | SP + 1 | |
| L | SP + 0 (Top Of Stack) | |

More confusingly, in RAM it is kept as E D L H. See how 2VARIABLE is defined to understand this fact.

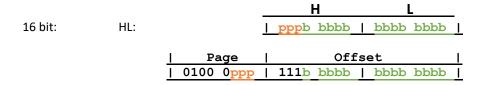
| CPU | 2VARIABLE |
|-----|-------------|
| Н | Address + 3 |
| L | Address + 2 |
| D | Address + 1 |
| E | Address + 0 |

Double Cell Floating-Point Number Encoding

There is another optional format that use 32 bits as a *double integer*, but all bits are used in a different way to allows to represent a *floating point number* approximately between $-1.7 * 10^38$ and $+1.7 * 10^38$ with 6-7 precision digits. The sign is kept in the most significant bit, the same way as a *double integer*; then eight bits follow as the exponential part, then 23 bits of mantissa. The sign in this position allows (IMO) using most of the same semantics of *double integers* as per the sign of the number.

Single Cell 16 bits Heap Pointer Address Encoding

There is Spectrum Next's peculiar 16 bits Heap Pointer Address Encoding that leverage on MMU7 i.e. Z80 memory space addresses between 0E000h and 0FFFFh. The three most significant bits represent an 8K-page between 64 and 71, lower bits are taken as offset from 0E000h. A specific definition FAR takes care of converting an heap-pointer address to an E000 offset and paging to MMU7 the correct 8K of physical RAM. Any NextZXOS call and most of I/O operations restore page 1.



4. Error messages.

Error messages strings are stored at Screens from 4 to 6 that are therefore reserved.

| | 0 0 |
|------|-----------------------------|
| Code | Message |
| #0 | is undefined. |
| #1 | Stack is empty. |
| #2 | Dictionary full. |
| #3 | No such line. |
| #4 | has already been defined. |
| #5 | Invalid stream. |
| #6 | No such block. |
| #7 | Stack is full! |
| #8 | Old dictionary is full. |
| #9 | Tape error. |
| #10 | Wrong array index. |
| #11 | Invalid floating point. |
| #17 | Can't be executed. |
| #18 | Can't be compiled. |
| #19 | Syntax error. |
| #20 | Bad definition end. |
| #21 | is a protected word. |
| #22 | Aren't loading now. |
| #23 | Forget across vocabularies. |
| #24 | RS loading error. |
| #25 | Cannot open stream. |
| | |

#28 Unexpected fixup/commaer.

Error at postit time. Inconsistent fixup.

- #29 Commaer data error.
- #30 Commaer wrong order.
- #31 Programming error.
- #33 Programming error.
- #43 File not found.

#26

#27

- #44 NexZXOS doscall error.
- #45 NextZXOS pos error.
- #46 NextZXOS read error.
- #47 NextZXOS write error.

5. The Dictionary

Legenda

In this list:

| a | memory address | 16 bits |
|------|---|---|
| b | byte, small unsigned integer | 8 bits |
| С | character | 8 bits, but often only lower 7 are significant. |
| d | signed double integer | 32 bits |
| fp | floating point number | 32 bits |
| ha | heap-pointer address (see >FAR) | 16 bits. |
| n | signed integer | 16 bits |
| u | unsigned integer | 16 bits |
| ud | unsigned double integer | 32 bits |
| | | |
| f | flag: a number evaluated as a boolean | 16 bits |
| ff | false flag: zero | 16 bits |
| tf | true flag: non-zero | 16 bits |
| | | |
| nfa | name field address | 16 bits |
| lfa | link field address | 16 bits |
| cfa | code field address | 16 bits |
| pfa | parameters field address | 16 bits |
| xt | execution token – same as cfa | 16 bits |
| | | |
| cccc | character string or word name available in the vocabulary | |
| | a list of words | |
| TOS | top of calculator stack | |

Core dictionary

'null' --- (immediate)

This is a "ghost" word executed by INTERPRET to go back to the caller once the text to be interpreted ends. This word allows you to use a **0x00** (NULL ASCII) as the end-of-text indicator in the input text stream.

! n a ---

stores an integer $\, n \,$ in the memory cell at address $\, a \,$ and $\, a \,$ + $\,$ 1. Pronounced "store".

Zilog Z80 microprocessor is a little-endian CPU that holds lower byte at lower address and higher byte in the higher address.

!CSP ---

saves the value of SP register in CSP user variable. It is used by : and ; for syntax checking. Also, CASE use it for the same purpose.

d1 --- d2

From a double number d1 it produces the next ASCII character to be put in an output string using HOLD. The number d2 is d1 divided by BASE and is kept for subsequent elaborations. This word is used between <# and #>. See also #S.

#> d --- a b

terminates a numeric conversion started by <# . This word removes d and leaves the values suitable for TYPE.

#BUFF --- n

Constant, the number of available buffers. This build has 3 buffers located at address between FIRST @ and LIMIT @.

#S d1 --- d2

This word is equivalent of a series of # that is repeated until d2 becomes zero. It is used between <# and #>.

#SEC --- n

This is a constant that gives the number of available Screens/blocks.

' --- cfa

Pronounced "tick". Used in the form

' cccc

this definitions leaves the **cfa** of word <code>cccc</code>, that is its xt or value to be compiled or passed to <code>EXECUTE</code>. If the word <code>cccc</code> is not found after the <code>CURRENT</code> and <code>CONTEXT</code> search phases, then an error #0 is raised, that is the message "cccc is undefined". In a previous version of this Forth, this word returned **pfa**: we changed this previous standard to return **cfa**.

```
( immediate)
```

Enclose a comment. Used in the form

(cccc)

ignores what is between brackets. The space after (is not considered in ccc. The comment must be delimited in the same row with a closing) followed by a space or the end of line.

(+LOOP) n ---

This is the primitive definition compiled by +LOOP.

This is the primitive definition compiled by . " and . (. It executes \mathtt{TYPE} .

(;CODE) ---

This is the primitive definition compiled by ; CODE. It rewrites the **cfa** of LATEST word to make it point to the machine code starting from the following address.

(?DO) ---

This is the primitive definition compiled by ?DO.

At compile-time it compiles the **cfa** of (?DO) followed by an offset as for BRANCH used to jump after the whole ?DO ... LOOP structure in case the limit equals the initial index, otherwise it is equivalent to (DO).

(?EMIT) c1 --- c2

Decodes the character cl using the following table. It is used internally by EMIT.

HEX 06 → print-comma

HEX 07 → bell rings

HEX 08 → back-space

HEX 09 → tabulator

HEX 0D → carriage return

HEX $0A \rightarrow$ new line (emitted as a 0D on the fly)

For not listed character, c2 is equal to c1.

(ABORT) ---

Definition executed in case of error issued by ERROR when WARNING contains a negative number. This word usually executes ABORT but can be patched with some user defined word at the **pfa** of (ABORT).

(COMPARE) al a2 n -- b

This word performs a lexicographic compare of $\,n\,$ bytes of text at address all with $\,n\,$ bytes of text address all. The compare is case-sensitive or case-insensitive based on the last execution of CASEON and CASEOFF.

When executed, this word returns a numeric value

- 0: if strings are equal
- +1: if string at a1 greater than string at a2
- -1: if string at al less than string at a2

See also CASEON and CASEOFF.

(DO) ---

This is the primitive compiled by DO.

Searches in the dictionary starting from address a2 a word which text name is kept at address a1; returns a cfa, the first byte b of nfa and a tf on a successful search; elsewhere a ff only.

The search is case-sensitive or case-insensitive based on the last execution of CASEON and CASEOFF.

Address a 2 must be the nfa of the first word involved in the search in the vocabulary.

In previous versions of this Forth, it returned a pfa, we change our mind.

Byte b keeps the length of the found word in the least significant 5 bits, bit 6 is the IMMEDIATE flag. Bit 5 is the SMUDGE bit. Bit 7 is always set to mark the beginning or end of the nfa.

See also CASEON and CASEOFF.

(LINE) n1 n2 --- a b

Retrieves line n1 of block n2 and send it to buffer. It returns the address a within the buffer and a counter b that is C/L (=64) meaning a whole line.

(LOOP) ---

This is the primitive compiled by LOOP. See also DO and +LOOP.

(MAP) a2 a1 n c1 --- c2

Translate...***

(NEXT) --- a

Constant. It is the address of "next" entry point for the **Inner Interpreter**. When creating word using machine code, the last op-code should be an unconditional jump to this address. If the created word wants to return an *intger* value on TOS, it should jump to the previous address; and if it wants to return a *double integer* value, it should jump to the next previous one. For example, to create two definitions that disable and enable interrupts, without an ASSEMBLER, you could use the following snippet:

```
CODE
       INT-DI
                HEX
  F3 C,
                 \
                    di
  C3 C, (NEXT)
                  ١
                    jp (NEXT)
  SMUDGE
                    now a dictionary search will find this word
CODE
       INT-EI
                HEX
  FB C,
                  ١
                    еi
                     jp (NEXT)
  C3 C, (NEXT)
                 ١
                    now a dictionary search will find this word
  SMUDGE
```

This Forth implementation always keeps (NEXT) value in **IX register**, so that the previous snipped should be written as:

```
CODE
                 HEX
       INT-DI
  F3 C.
                     di
  DD C, E9 C,
                     jp (ix)
CODE
       INT-EI
                 HEX
  FB C,
                     еi
                  ١
  DD C, E9 C,
                     jp (ix)
  SMUDGE
                     now a dictionary search will find this word
```

(NUMBER) d = -- d2 = 2

Converts the ASCII text at address a+1 in a double integer using the current BASE. Number d2 is left on top of stack for any subsequent elaborations, a2 is the address of the first non-converted character.

In the CPU registers a double integer is kept as HLDE, on the stack is treated as two distinct integers where HL is on TOS and DE is the second from top, so that in memory it appears as LHED. Instead, in a variable declared with 2VARIABLE is stored as EDHL.

Used by NUMBER.

Determines if the character at address $\,a$ is a sign (+ o -) and if found increments $\,a$. The flag f indicates the sign: ff for a positive sign + or no sign at all, tf for a negative sign - . If a is incremented then variable DPL is incremented as well. Used by da NUMBER and (EXP) in the floating-point option.

Computes the product of two integers.

Compute $(n1 \cdot n2) / n3$ using a double integer for the intermediate value to avoid precision loss.

Leaves the quotient n5 and the reminder n4 of the operation ($n1 \cdot n2$) / n3 using a double integer for the intermediate to avoid precision loss.

Leaves the sum of two integer.

Adds to the cell at address a the number n. It is the same as the sequence $a \in n + a$!

Computes n3 as n1 with the sign of n2. If n2 is zero, it means positive.

+BUF a1 --- a2 f

Advances the address of the buffer from a1 to a2, that is the next buffer. The flag f is false if a2 is the buffer pointed by PREV.

+LOOP n1 --- (run time)
a n2 --- (compile time)

Used in colon definition in the form

DO ... n1 +LOOP

At run-time + LOOP checks the return to the corresponding DO, n1 is added to the index and the total compared with the limit. The jump back happens:

- a) while index < limit if n1 > 0;
- b) while index > limit if n1 < 0.

Otherwise the execution leaves the loop. On leaving the loop, the parameters are discarded and the execution continues with the following word.

At compile-time + LOOP compiles (+LOOP) and a jump is calculated from HERE to a which is the address left on the stack by DO. The value n2 is used internally for syntax checking.

+ORIGIN n --- a

Returns the address n bytes after the "origin". In this build the origin is 6400h. Used rarely to modify the boot-up parameters in the origin area.

, n ---

It puts n in the following cell of the dictionary and increments DP (dictionary pointer) of two locations.

, "I

Compile a "Counted-ZString". It calls WORD to read characters from the current input stream up to a delimiter " and stores such a string at HERE. It stores the length of the string as the first byte, then it ends everything with a NUL character (0x00).

- n1 n2 --- n3

Computes n3 = n1 - n2 as the difference from the penultimate and the last number on the stack.

-->

Continues the interpretation in the next Screen during a ${\tt LOAD}$.

-1 --- n

This is the constant value -1 that in this implementation is OFFFFh. Compiling a constant result in a faster execution than a literal.

```
-DUP n --- n n (non zero)
n --- n (zero)
```

Duplicates n if it is non zero.

Used in the form -FIND cccc.

It accepts a word (delimited by spaces) from the current input stream, storing it at address <code>HERE</code>. Then, it run a search in the <code>CONTEXT</code> vocabulary first, then in the <code>CURRENT</code> vocabulary. If the word is found, it leaves the <code>cfa</code> of the word, its length-byte <code>b</code> and a <code>tf</code>. Otherwise only a <code>ff</code>.

```
-TRAILING a1 n1 --- a2 n2
```

This definition assumes that a string n1 characters long is already stored at address a1 containing a space right-delimited word. It determines n2 as the position of the first delimiter after the word.

```
. n ---
```

Prints the integer $\, n \,$ followed by a space.

```
." --- (immediate)
Used in the form
." cccc "
```

At compile-time, within a colon-definition, compiles the primitive to output the text followed by the string ccc (delimited by "). The text ccc is prepended by a length-counter that TYPE will use at run-time. When interpreted, i.e. outside a colon-definition, immediately sends the text to output.

```
.( immediate)
```

Used in the form

. (cccc)

acting as . " cccc " but the string is delimited in a different way

```
.C c --- (immediate)
```

Used in the form

c .C xxxx C

acting as ."xxxx" but the string is delimited by character c. It is a more generic form of ." that, in fact, use this word as their primitive.

```
.LINE n1 n2 ---
```

Sends line n1 of block n2 to the current peripheral ignoring the trailing spaces.

.R n1 n2 ---

Prints a number n1 right aligned in a field n2 character long, with no following spaces. If the number needs more than n2 characters, the excess protrudes to the right.

/ n1 n2 --- n3

Computes n3 = n1/n2, the quotient of the integer division.

/MOD n1 n2 --- n3 n4

Computes the quotioent n4 and the reminder n3 of the integer division n1/n2. The reminder has the sign of n1.

0 --- n

This is a constant value zero. Compiling a constant results in a faster execution than a literal.

0< n --- f

Leaves a tf if n is less than zero, ff otherwise.

0= n --- f

Leaves a tf if n is not zero, ff otherwise. It is like a NOT n.

0> n --- f

OBRANCH f ---

Direct procedure that executes a conditional jump. If $\,f\,$ is zero the offset in the cell following <code>OBRANCH</code> is added to the Instruction Pointer to jump forward of backward.

It is compiled by IF, UNTIL and WHILE.

1 --- n

Constant value 1. Compiling a constant results in a faster execution than a literal.

1+ n1 --- n2

Increments by one the number on TOS.

1- n1 --- n2

Decrements by one the number on TOS.

2 --- n

Constant value 2. Compiling a constant results in a faster execution than a literal.

2! d a ---

n-lo n-hi a ---

Stores the double integer held on TOS to address a.

2* n1 --- n2

Doubles the number on TOS.

2+ n1 --- n2

Increments by two the number on TOS.

2/ n1 --- n2

Halves the number on TOS.

20 a --- d

a --- n-lo n-hi

Fetches the double integer at address a. to TOS.

2CONSTANT d --- (immediate) (compile time)
--- d (run time)

Defining word that creates a double constant. Used in the form

d 2CONSTANT cccc

it creates the word \mbox{ccc} and \mbox{pfa} holds the number d. When \mbox{ccc} is later executed it put d on TOS. This definition is not available at startup, it has to be loaded via NEEDS $\mbox{2CONSTANT}$.

2VARIABLE d --- (immediate) (compile time)
--- a (run time)

Defining word used in the form:

d 2VARIABLE cccc

creates the word cocc with the pfa containing the initial value d. When cocc is executed, it puts on TOS the pfa of cocc that is the address that holds the value d.

When used in the form

cccc @

the content of the double-variable cccc is left on TOS.

When used in the form

d cccc !

the double-value on TOS is stored to the double-variable ccc.

This definition is not available at startup, it has to be loaded via NEEDS 2CONSTANT.

2DROP d --n1 n2 ---

Discards a double integer from the TOS, i.e. discards the top two integer.

2DUP d --- d d

Duplicates the double integer on TOS, i.e. duplicates in order the two top integer.

20VER d1 d2 --- d1 d2 d1

n1 n2 n3 n4 --- n1 n2 n3 n4 n1 n2

Copies to TOS the second double integer from top.

2ROT d1 d2 d3 --- d2 d3 d1

n1 n2 n3 n4 n5 n6 --- n3 n4 n5 n6 n1 n2

Rotates the three top double integers, taking the third an putting it on top. The other two double integer are pushed down from top by one place.

2SWAP d1 d2 --- d2 d1

Swaps the two double integers on TOS.

3 --- n

Constant value 3. Compiling a constant results in a faster execution than a literal.

: --- (immediate)

This is a defining word that creates and begins a colon-definition. Used in the form

: cccc ... ;

creates in the dictionary a new word cccc so that it executes the sequence of already existing words '...'.

The CONTEXT vocabulary is set to be the CURRENT and compilation continues while STATE is not zero. Words having the bit 6 of its length-byte set are immediately executed instead of being compilated.

; --- (immediate)

Ends a colon definition and stops compilation.. It compiles; S and execute SMUDGE to make the word findable.

;CODE --- (immediate)

Used in the form

terminates a colon definition stoppin copilation of word ccc and compiling (; CODE). Usually ; CODE is followed by suitable machine code sequence..

;S --- (immediate)

This is usually the last word compiled in a colon definition by ; it does the action of returning to the calling word. It is used to force the immediate end of a loading session started by LOAD.

Leaves a tf if n1 is less than n2, ff otherwise.

<# ---

Sets <code>HLD</code> to the value of <code>PAD</code>. It is used to format numbers using #, #S, <code>SIGN</code> and #>. The conversion is performed using a double integer, and the formatted text is kept in <code>PAD</code>.

<BUILDS ---

Used in a colon definition in the form

Subsequent execution of ccc in the form

cccc nnnn

creates a new word nnnn with an high-level procedure that at run-time calls the DOES> part of cccc. When nnnn is executed, the pfa of nnnn is put on TOS and the executed the following DOES>.

<BUILD and DOES> allow writing high-level procedures instead of using machine code as ; CODE would require.

<NAME cfa --- nfa

Converts a cfa in its nfa. It is the same as the sequence >BODY NFA. See also: CFA, LFA, NFA, PFA, >BODY.

Leaves a tf if n1 equals to n2, ff otherwise.

Leaves a tf if n1 is greater than n2, ff otherwise.

>BODY cfa --- pfa

Converts a cfa in its pfa.

See also: CFA, LFA, NFA, PFA, <NAME.

>IN --- a

User variable that keeps track of text position within an input buffer. WORD uses and modifies the value of IN that is incremented when consuming input buffer.

>R n ---

Takes an integer from TOS and puts it on top of the Return Stack. It should be used only within a colon definition and the use of >R should be balanced with a corresponding R>.

? a ---

Prints the content of cell at address a. It is the same as the sequence: a @.

?COMP ---

Raises an error message #17 if the current STATE is not compiling state.

?CSP ---

Raises an error message #20 if the value of CSP is different from the current value of SP register. It is used to check the compilation in a colon definition.

?DO n1 n2 --- (immediate) (run time) --- a n (compile time)

Used in a colon definition in the form

?DO ... LOOP ?DO ... n3 +LOOP

It is used as DO to put in place a loop structure, but at run-time it first checks if n1 = n2 and in that case the loop is skipped. At run-time ?DO starts a sequence of words that will be repeated under control of an initial-index n2 and a limit n1. ?DO consumes these two value from stack and the corresponding LOOP increments the index. If the index is less than the limit, the executions returns to the corresponding ?DO, otherwise the two parameters are discarded and the execution continues after the LOOP.

The limt n1 and the initial value n2 are determined during the execution and can be the result of other previous operations. Inside a loop the word I copies to TOS the current value of the index.

Se also: I, DO, LOOP, +LOOP, LEAVE. In particular LEAVE allows leaving the loop at the first opportunity. At compile-time ?DO compiles (?DO) followed by an offset like BRANCH and leaves the address of the following location and the number $\,$ n to syntax-check

?DO- [a1 n1] a n ---

Computes and compiles a relative offset from a to HERE and in case it completes the BRANCH part previously compiled by ?DO that left al and nl. It is used by LOOP, +LOOP. If the loop begin with DO then al and nl aren't there.

?DUP n n (non zero) (zero) n n

Duplicates the value on TOS if it is not qual to zero. This is the same as <code>-DUP</code>.

?ERROR f n

Raises an error message #n if f is true.

?EXEC

Raises an error message #18 if we aren't compiling.

?LOADING

Raises an error message #22 if we aren't loading. It show the illegal use of -->.

?PAIRS n1 n2

Raises an error message #19 if n1 is different from n2. It is used for syntax checking by the words that completes the construction of structures DO, BEGIN, IF, CASE.

?STACK

Raises an error message #1 if the stack is empty and we tried to consume an element from the calculator stack. On the other hand, an error message #7 if the stack is full.

?TERMINAL f

Tests the keyboard. Leaves a tf if the [BREAK] key is pressed, ff oherwise.

(a а n

Reads cell at address a and put an integer on TOS.

ABORT

Clears the stack and pass to prompt command, prints the copyright message and returns the control to the human operator executing QUIT.

ABS n u

Leaves the absolute value of n.

ACCEPT n1 n2 a

Transfers characters from the input terminal to the address a for n1 location or until receiving a 0x13 "CR" character. A 0x00 "null" character is added. It leaves on TOS $\,\mathrm{n2}\,$ as the actual length of the received string. More, $\,\mathrm{n2}\,$ is also copied in SPAN user variable. See also QUERY.

ACCEPT- a n1 --- n2

As for ACCEPT, but it reads at most n1 characters text from current channel/stream via INKEY one character at a time, It stores the text at address a. Not so efficient, but it allows to compile an external souce-file attached to a Basic's OPEN# stream. It does not modify SPAN.

AGAIN --- (immediate) (run time)
a n --- (compile time)

Used in colon definition in the form

BEGIN ... AGAIN

At run-time AGAIN forces the jump to the corresponding BEGIN and has no effect on the calculator stack. The execution cannot leave the loop (at least until a R> is executed at a lower level).

At compile-time AGAIN compiles BRANCH with an offset from HERE to a. The number n is used for syntax-check.

ALLOT n ---

This definition is used to reserve some space in the dictionary or to free memory. It adds the signed integer n to DP (Dictionary Pointer) user variable.

It executes an AND binary operation between the two integers. The operation is performed bit by bit.

AUTOEXEC ---

This word is executed the first time the Forth system boots and **loads Screen# 1**. Once called it patches ABORT to prevent any further executions at startup.

B/BUF --- n

Constant. Number of bytes per buffer. In this implementation is 512.

B/SCR --- n

Constant that indicates the number of Blocks per Screen. In Next version is 2. In Microdrive version it was 1.

BACK a ---

Computes and compiles a relative offset from a to HERE. Used by AGAIN, UNTIL, LOOP, +LOOP.

BASE --- a

User variable that indicates the current numbering base used in input/output conversions. It is changed by $\mathtt{DECIMAL}$ that put ten, HEX that put sixteen, and with some extensions \mathtt{BINARY} that put two and \mathtt{OCTAL} that put eight.

BASIC u ---

Quits Forth and returns to Basic returning to the caller USR the unsigned integer on TOS.

BEGIN --- (immediate) (run time)
--- a n (compile time)

Used in colon definition in one of the following forms

BEGIN ... F UNTIL or
BEGIN ... F WHILE ... REPEAT or
BEGIN ... F END

At compile-time, it starts one of these structures.

At run-time BEGIN marks the beginning of a words sequence to be repeatedly executed and indicates the jump point for the corresponding AGAIN, REPEAT, UNTIL or END.

With UNTIL, the jump to the corresponding BEGIN happens if on TOS there is a ff, otherwise it quits the loop.

With AGAIN and REPEAT, the jump to the corresponding BEGIN always happens.

The WHILE part is executed if and only if on TOS there is a tf, otherwise it quits the loop.

BL --- c

Constant for "Blank". This implementation uses ASCII and BL is 32.

BLANKS a n ---

Fills with "Blanks" n location starting from address a.

BLK --- a

User variable that indicates the current block to be interpreted. If zero then the input is taken from the terminal buffer TIB.

BLK-FH --- a

Variable containing file-handle to Block's file !Blocks-64.bin.

BLK-FNAME --- a

Variable containing the counted-zstring "!Block-64.bin". as produced by ," definition.

BLK-INIT ---

Initialize BLOCK system. It opens for update (read/write) file "!Block-64.bin".

BLK-READ a n ---

Read block n to address a. See also F READ.

BLK-SEEK n ---

Seek block n within blocks!.bin file. See also ${\mathbb F}$ SEEK.

BLK-WRITE a n ---

Take text content at address a to disk block n. See also ${\tt F}$ WRITE.

BLOCK n --- a

Leaves the address of the buffer that contains the block n. If the block isn't already there, it is fetched from disk. If in the buffer there was another buffer and it was modified, then it is re-written to disk before reading the block n. See also BUFFER, R/W, UPDATE, FLUSH.

BRANCH ---

Direct procedure that executes an unconditional jump. The memory cell following BRANCH has the offset to be relatively added to the Instruction Pointer to jump forward or backward. It is compiled by AGAIN, ELSE, REPEAT.

BUFFER n --- a

Makes the next buffer available assigning it the block number n. If the buffer was marked as modified (by UPDATE), such buffer is re-written to disk. The block is not read from disk. The address point to the first character of the buffer.

BYE ---

Executes FLUSH and EMPTY-BUFFERS, then quits Forth and returns to Basic returning to the caller USR the value of 0 +ORIGIN. See also BASIC.

C! b a ---

Stores a byte b to address a.

C, b ---

Puts a byte b in the next location available in the dictionary and increments DP (dictionary pointer) by 1.

C/L --- c

Constant that indicate the number of characters per screen line. In this implementation it is 32.

C@ a --- b

Puts on TOS the byte at address a.

CASEOFF ---

Sets case-sensitive search OFF. changes the system behavior so that (FIND) can search the dictionary ignoring case, and (COMPARE) compares two strings ignoring case.

CASEON ---

Sets case-sensitive search ON. It changes the system behavior so that (FIND) will search the dictionary case sensitive, and (COMPARE) will compare the two strings case sensitive.

CELL+ n1 --- n2

Increments n1 by 1 "cell", that is two units. In this implementation a cell is two bytes.

CELL- n1 --- n2

Decrements n1 by 1 "cell", that is two units. In this implementation a cell is two bytes.

CELLS n1 --- n2

Doubles the number n1 on TOS giving the number of bytes equialent to n1 "cells". In this implementation a cell is two bytes.

CFA pfa --- cfa

Converts a pfa in its cfa. See also LFA, NFA, PFA, >BODY, <NAME.

CHAR --- c

Used in the form

CHAR c

determines the first character of the next word in the input stream.

CLS ---

Clears the screen using the ZX Spectrum ROM routine 0DAFh.

CMOVE a1 a2 n ---

Copies the content of memory starting at address all for n bytes, storing them from address all is moved first. See also CMOVE>.

CMOVE> a1 a2 n ---

The same as CMOVE but the copy process starts from location al + n - 1 proceding backward to the location al.

CODE ---

Defining word used in the form

CODE ccc

it creates a new dictionary entry for the definition <code>cccc</code> with the cfa of such a definition pointing to its pfa that is empty for the moment, <code>HERE</code> points that location; then some machine-code instruction should be added using <code>C</code>, that will be compiled at <code>HERE</code>. The new word is created in the <code>CURRENT</code> vocabulary but won't be found by (<code>FIND</code>) because it has the <code>SMUDGE</code> bit set. Once the word construction is complete, it is a programmer responsibility to execute <code>SMUDGE</code>. This word is overridden by <code>ASSEMBLER</code> vocabulary available after <code>LOADing</code> Screens 100-165, this allows the programmer

to use a pseudo-standard Z80 notation to create a new low-level definition using assembler directly.

Here is an example that creates a definition SYNC-FRAME to wait for the next maskable interrupt:

CODE SYNC-FRAME HEX

76 C, \ halt ; wait for interrupt or reset

DD C, E9 C, \ jp (ix) ; jump to the inner interpreter

SMUDGE

COLD ---

This word executes the Cold Start procedure that restore the system at its startup state. It sets DP to the minimum standard and executes ABORT.

COMPILE ---

At compile-time, it determines the cfa of the word that follows COMPILE and compile it in the next dictionary cell.

CONSTANT n --- (immediate) (compile time)
--- n (run time)

Defining word that creates a constant. Used in the form

n CONSTANT cccc

it creates the word cccc and pfa holds the number n. When cccc is later executed it put n on TOS.

CONTEXT --- a

User variable that points to the vocabulary address where a word search begins.

COUNT a1 --- a2 b

Leaves the address of text a2 and a length b. It expects that the byte at address a1 to be the length-counter and the text begins to the next location.

CR ---

Transmits a 0x0D to the current output peripheral.

CREATE --- a (compile time)

Defining word used in the form

CREATE cccc

it creates a new dictionary entry for the definition cccc with the **pfa** still empty.

When cccc is executed, it puts on TOS the pfa of cccc

Often used with ALLOT to reserve space in the dictionary to be later used, for instance as an array.

See also VARIABLE.

CSP --- a

User variable that temporarily holds the value of SP register during a compilation syntax error check.

CURRENT --- a

User variable that points to the address in the Forth vocabulary where a search continues after a failing search executed in the CONTEXT vocabulary. See also LATEST.

D+ d1 d2 --- d3

Computes d3 as the sum of d1 and d2. This is a 32 bits sum.

Computes d that is ud with the sign of n.

Prints a double integer followed by a space. The double integer is kept on stack in the format n-lo n-hi and the integer on TOS is the most significant.

D.R d n ---

Prints a double integer rigth aligned in a field n character wide. No space follows. If the field is not large enough, then the excess protrudes to the right.

DABS d --- ud

Leaves the absolute value of a double integer.

DECIMAL ---

Sets BASE to 10, that is the decimal base.

DEFINITIONS ---

To be used in the form

cccc DEFINITIONS

it sets the CURRENT vocabulary to be the CONTEXT vocabulary and this allows adding new definitions to cccc vocabulary.

For example: FORTH DEFINITIONS or ASSEMBLER DEFINITIONS.

In this implementation a Forth oriented ASSEMBLER vocabulary is available as an extra-option that can be LOADed from Screens 100 -160.

DEVICE --- a

Variable that holds the number of current channel: 2 for video, 3 for printer, and any number between 4 and 15 to refer to an Basic OPEN# channel.

DIGIT c n --- u tf (ok)
c n --- ff (ko)

Converts the ASCII character c in the equivalent number using the base n, followed by a a tf. If the conversion fails it leaves a ff only.

DLITERAL d --- d (immediate) (run time)
d --- (compile time)

Same as LITERAL but a 32 bits number is compiled. DLITERAL is an immediate word that is executed and not compiled.

DMINUS d1 --- d2

Computes the opposite double number.

DO n1 n2 --- (immediate) (run time)
--- a n (compile time)

Used in colon definition in the form

DO ... LOOP or DO ... n +LOOP

It is used to put in place a loop structure: The execution of DO starts a sequence of words that will be repeated, under control of an initial-index n2 and a limit n1. DO drops these two value from stack and the corresponding DO, otherwise the two parameters are discarded and the execution continues after the DOP.

The limt n1 and the initial value n2 are determined during the execution and can be the result of other previous operations. Inside a loop the word I copies to TOS the current value of the index.

See also: I, DO, LOOP, +LOOP, LEAVE. In particular LEAVE allows leaving the loop at the first opportunity.

At compile-time DO compiles (DO) and leaves the address of the following location and the number $\,n\,$ to syntax-check.

DOES> ---

Word that defines the execution action of a high-level defining word. DOES> changes the pfa of the word being defined to point the words sequence compiled after DOES>. It is used in conjunction with <BUILDS. When the machine-code part of DOES> is executed, it leaves on TOS the pfa of the new word, this allows the interpreter to use this area. Obvious uses are new vocabularies (Assembler), multidimensional array and other compiling operations.

DP --- a

User variable (Dictionary Pointer) that holds the address of next available memory location in the dictionary. It is read by HERE and modified by ALLOT.

DPL --- a

User variable that holds the number of digits after the decimal point during the interpretation of double integer. It can be used to keep track of the column of the decimal point during a number format output. For 16 bit integer it defaults to -1. It takes into account the exponential part and its sign for floating point numbers.

DROP n ---

Drops the value on TOS. See also OVER, NIP, TUCK, SWAP, DUP, ROT.

DUP n --- n n

Duplicates the value on TOS. See also OVER, DROP, NIP, TUCK, SWAP, ROT.

Used in colon definition in the form

At run-time ELSE forces the execution of the false part of an IF-ELSE-ENDIF structure. It has no effects on the stack. At compile-time ELSE compiles BRANCH and prepares the following cell for the relative offset, stores at a1 the previous offset from HERE; then it leaves a2 and n2 for syntax checking.

EMIT c ---

Sends a printable ASCII character to the current output peripheal. OUT is incremented. 7 EMIT activates an acoustic signal. The 'null' OxOO ASCII character is not transmitted.

EMITC b ---

Sends a byte b character to the current output peripheal selected with SELECT. See also DEVICE.

EMPTY-BUFFERS ---

Erases all buffers. Any data stored to buffers after the previous FLUSH is lost.

ENCLOSE a c --- a n1 n2 n3

Starting from address a, and using a delimiter character c, it determines the offset n1 of the first non-delimiter character, n2 of the first delimiter after the text, n3 of first character non enclosed.

This word doesn't go beyond a 'null' ASCII that represent a unconditional delimiter. For example:

1:
$$c c x x x c x \rightarrow 2 5 6$$

2: $c c x x x x' null' \rightarrow 2 5 5$
3: $c c' null' \rightarrow 2 3 2$

Synonym of UNTIL.

ENDIF a n --- (immediate) (compile time)

At run-time, ENDIF indicates the destination of the forward jump from IF or ELSE. It marks the end of a conditional structure. THEN is a synonym of ENDIF.

At compile-time ENDIF calculates the forward jump offset from a to HERE and store it at a. The number n is used for syntax checking.

ERASE a n ---

Erases n memory location starting from a, filling them with 0x00 'null' characters.

ERROR b --- n1 n2 --- ff

Notifies an error b and resets the system to command prompt. First of all, the user variable WARNING is examined.

If WARNING is 0 then the offending word is printed followd by a "?" character and a short message "MSG#n".

If WARNING is 1, instead of the short message, the text available on line b of block 4 (of drive 0) is displayed. Such a number can be positive or negative and lay beyond block 4.

If WARNING is -1 then ABORT is executed, which resets the system to command prompt. The user can (with care) modify this behavior of that by altering (ABORT).

If BLK is non zero, then ERROR leaves on the stack n1 that is the value of IN and n2 that is the value of BLK at the error moment. These numbers can then be used by WHERE to determine and show the exact error position.

If BLK is zero, then only a ff is left on TOS.

In all cases, the final action is \mathtt{QUIT} .

EXECUTE cfa ---

Executes the word which cfa is held on TOS.

EXP --- a

User variable that holds the exponent in a floating-point conversion.

EXPECT a n ---

Transfers characters from the input terminal to the address a for n location or until receiving a 0x13 "CR" character. A 0x00 "null" character is added in the following location. The actual length of the received string is kept in SPAN user variable. See also ACCEPT.

FENCE --- a

User variable that holds the (minimum) address to where FORGET can act.

FILL anb ---

Fills n memory location starting from address a with the value of b.

FIRST --- a

User variable that holds the address of the first buffer. See also LIMIT.

FLD --- a

User variable that holds the width of output field.

FLUSH ---

Executes SAVE-BUFFERS. It saves to disk the buffers marked "modified" by UPDATE.

FORGET ---

Used in the form

FORGET cccc

removes from the dictionary the word <code>cccc</code> and all the preceding definitions. Care must be put when more than one vocabulary is involved. See <code>MARKER</code>.

FORTH --- (immediate)

This is the name of the first vocabulary. Executing FORTH sets this to be the CONTEXT vocabulary. As soon as no new vocabulary is defined, all new colon definitions became part of FORTH vocabulary. FORTH is immediate, so it is executed during the creation of a colon definition to select the needed vocabulary. See also ASSEMBLER (optional vocabulary).

F CLOSE n --- f

NextZXOS option: it closes a file handle $\,n\,$ previously opened with F_OPEN. Flag $\,f\,$ is 0 for OK. It uses an RST 8 call followed by \$9B service number.

F FGETPOS n --- d f

NextZXOS option: given an open file handle n returns the position d. Flag f is 0 for OK.

F GETLINE a fh --- a n

Given a filehandle read next line (terminated with \$0D or \$0A. Address a is left for subsequent process and n as the number of byte read, i.e. the length of line just read.

F INCLUDE n ---

Given a file-handle n, it includes the source from file

F OPEN a1 a2 n1 --- n2 f

NextZXOS option: it opens a file using filespec given at address al and returns filehandle number n, nl is "mode" as specified in "NextZXOS and esxDOS APIs" standard documentation. Filespec is a NUL-termianted string. Flag f is 0 for OK. It uses an RST 8 call followed by \$9A service number. See F CLOSE.

F READ a n1 n2 --- n3 f

NextZXOS option: it reads at most n1 bytes from file handle n2 and stores them at address a. Returns n3 as the actual bytes read. Flag $\, f \,$ is 0 for OK. It uses RST 8 call followed by \$9D service number.

F SEEK d n ---

NextZXOS option: it seeks position $\,d\,$ at open file given by filehandle $\,n.$ It uses an RST 8 call followed by \$9F service number. Flag $\,f\,$ is 0 for OK.

F SYNC n --- f

NextZXOS option: it syncs to disk open file given by filehandle n. It uses an RST 8 call followed by \$9C service number. Flag f is 0 for OK.

F WRITE a n1 n2 --- n3 f

NextZXOS option: it takes n1 bytes at address a and writes them to filehandle n2. It uses an RST 8 call followed by \$9F service number. Returns n3 as the actual bytes written. Flag f is 0 for OK.

HERE --- a

Leaves the address of next location available on the dictionary.

HEX --- a

Changes the base to hexadecimal, setting BASE to 16.

HLD --- a

User variable that holds the address of last character used in a numeric conversion output.

HOLD c ---

Used between <# and #> to put a ASCII character during a numeric format.

T --- n

Used between DO and LOOP (or DO and +LOOP, ?DO and LOOP, ?DO and +LOOP) to put on TOS the current value of the loop index.

ID. nfa ---

It prints the definition name whose nfa is on TOS.

Used in colon definition in the form

IF ... ELSE ... ENDIF

At run-time IF selects which words sequence to execute based on the flag on TOS:

If f is true, the execution continues with the instruction that follows IF ("true" part).

If f is false, the execution continues after the ELSE ("false" part).

At the end of the two parts, the executions always continues after ENDIF.

ELSE and its "false" part are optional and if omitted no "false part" will be executed and execution continues after ENDIF. At compile time IF compiles 0BRANCH reserving a cell for an offset to the point after the corresponding ELSE or ENDIF. The integer n is used for syntax checking.

IMMEDIATE

Marks the latest defined word such that at compile-time it is always executed instead of being compiled. The bit 6 of the length byte of the definition is set. This allows such definitions to handle complex compilation situation instead of burdening the main compiler.

The user can force the compilation of an immediate definition prepending a [COMPILE] to it.

INCLUDE

It is used in in the form:

INCLUDE cccc

starts interpretation of text read from file $\ensuremath{\mathsf{cccc}}$.

See also LOAD

INDEX n1 n2 ---

Prints the first line of all screens between n1 and n2. Useful to quick check the content of a series of screens.

INKEY --- b

Reads the next character available from current stream and previously selected with SELECT leving it on TOS. It is the opposite of EMITC.

INTERPRET ---

This is the text interpreter. It executes or compiles, depending on the value of STATE, text from input buffer a word at a time. It first searches on CONTEXT and CURRENT vocabularies; if they fail, the text is interpreted as a numeric value, converted using the current BASE, and put on TOS. If that numeric conversion fails too, an error is notified with the symbol "?" followed by the word that caused the error. INTERPRET executes NUMBER and the presence of a decimal point "." indicates that the number is assumed as double integer instead of a simple integer.

After execution of the word found, the control is given back to the caller procedure.

INVV ---

"Inverse video". It enables Inverse-Video attribute mode. See also TRUV.

KEY --- b

Shows a (flashing) cursor on current video position and waits for a keypress. It leaves the ASCII code b of the character read from keyboard without printing it to video. In this implementation some SYMBOL-SHIFT key combinations are decoded as follow:

```
E2
     STOP
                       7E
С3
                \rightarrow
     NOT
                       7C
CD
     STEP
                       5C
CC
     TO
                       7в
СВ
                       7 D
     THEN
С6
     AND
                        5В
C5
                       5D
     OR
                             ]
AC
     AΤ
                       7 F
C7
                       20
     <=
                             space
С8
     >=
                       20
                             space
С9
      <>
                       06
                             as CAPS-SHIFT + 2 and toggles CAPS-SHIFT On and Off,
```

L/SCR --- n

Constant that indicates the number of lines per Screen. In this implementation is 16.

LATEST --- nfa

Leaves the nfa of the latest word defined in CURRENT vocabulary.

LEAVE ---

Forces the conclusion of a DO \dots LOOP setting the limit at the current index I, inducing an exit at the first occasion. The index remains unaltered and the execution continues normally up to the following LOOP or +LOOP.

LFA pfa --- lfa

Converts a pfa in its lfa. See also CFA, NFA, PFA, >BODY, <NAME.

LIMIT --- a

User variable that points to the first location above the last buffer. Normally it is the top of RAM, but not always. In this implementation, it can be set at E000h to allow MMU7 as a general purpose 8K RAM bank. See also: FIRST.

LIST n ---

Prints screen number n. Sets SCR to n.

LIT --- n

Puts on TOS the value hold in the following location. It is automatically compiled a before each literal number.

LITERAL n --- n (immediate) (run time)
n --- (compile time)

Compile-time, LITERAL compiles LIT followed by the value n in the following cell. This is an immediate word and, a

colon definition, it will be executed.

It is used in the form

```
: cccc ... [ calculations ] LITERAL ... ;
```

the compilation is suspended during the calculations and, when compilation resumes, LITERAL compiles the value put on TOS during the previous calculations.

LOAD n ---

Start interpretation of Screen $\,n$. The loading phase ends at the end of the screen or at the first occurrence of ; S. If $\,n$ is negative, instead of loading from Screen# $\,n$, it loads text directly from stream $\,n$ as previously OPEN# from Basic. See also -->

LOAD+ n ---

Start interpretation of screen $\,$ n. The loading phase ends at the end of the screen or at the first occurrence of ; S. See also --> and LOAD.

LOAD- n ---

Start interpretation of text read directly from stream n as from Basic's OPEN# n . It uses ACCEPT-. See also --> and LOAD.

LOOP a n --- (immediate) (run time)
n --- (compile time)

Used in colon defintion in the form

DO ... LOOP ?DO ... LOOP

At run-time LOOP checks the jump to the corresponding DO. The index is incremented and the total compared with the limit; the jump back happens while the index is less than the limit. Otherwise the execution leaves the loop. On loop leaving, the parameters are discarded and the execution continues with the following word.

At compile-time LOOP compiles (LOOP) and the jump is calculated from HERE to a which is the address left by DO on the stack. The value n2 is used internally for syntax checking.

LP --- a

User variable for printer purposed. Not used.

LSHIFT n1 u --- n2

Shifts left an integer n1 by u bit.

M* n1 n2 --- d

Mixed operation. It leaves the product of n1 and n2 ad a double integer.

M+ d u --- d2

Mixed operation. It leaves the sum of d and unsigned u as a double integer d2.

This definition is available after NEESS M+

M/ d n1 --- n2 n3

Mixed operation. It leaves the remainder n2 and the quotient n3 of the integer division of a double integer d by the divisor n1. The sign of the reminder is the same as d.

M/MOD ud1 u1 --- u2 ud3

Mixed operation. Leaves the remainder u2 and the quotient ud3 of the unsigned integer division of a double integer d by the divisor n1.

MARK a n ---

TYPE in inverse video. This word is not available at startup, it has to be loaded via NEEDS MARK.

MARKER --- (immediate) (run time)

Used outside a colon defintion in the form

MARKER cccc

this creates a new definition <code>cccc</code> that once executed restores the dictionary to the status before <code>cccc</code> was created. This removes <code>cccc</code> and all subsequent definitions. This word allows forgetting across vocabularies since it keeps track of <code>VOC-LINK</code>, <code>CURRENT</code>, <code>CONTEXT</code> values.

MAX n1 n2 --- n3

Leaves the maximum between n1 and n2.

MESSAGE n ---

Prints to the current device the error message identified by n. If WARNING is zero, a short MSG#n is printed. If WARNING is non zero 1, line n of screen 4 (of drive 0) is displayed. Such a number can be positive or negative and lay beyond block 4. See also ERROR.

MIN n1 n2 --- n3

Leaves the minimum between n1 and n2.

MINUS n1 --- n2

Changes the sing of n1

MMU7! n ---

This word accepts n between 0 and 223 and map the corresponding 8K-page at E000-FFFh addresses. It is coded in Assembler and uses NEXTREG A,n Next's peculiar op-code (ED 92). See MMU7@.

MMU7@ --- r

This word returns a number n between 0 and 223 by asking the hardware which 8K-page is currently fitted in MMU7. See $\,^{MMU7}$!

MOD n1 n2 --- n3

Divides n1 by n2 and leaves the remainder n3. The sign is the same as n1.

M P3DOS n1 n2 n3 n4 a --- n4 n5 n6 n7 f

This is the NZXOS call wrapper. Parameters passed on stack are used as follow:

- n1 = input parameter value for **hI** registers pair
- n2 = input parameter value for **de** registers pair
- n3 = input parameter value for **bc** registers pair
- n4 = **a** register input parameter value
- a = service routine address
- n5 = hl returned value
- n6 = **de** returned value
- n7 = **bc** returned value
- n8 = a register
- f = 0 for OK, non zero for KO.

This word calls uses RST 08 followed by \$94 to call the specified routine.

Value returned on register IX is also stored at HEX 2A +ORIGIN before IX is restored to its fixed value.

NEEDS ---

It is used in in the form:

NEEDS cccc

if the definition cccc is not present in dictionary, then it starts interpretation of text read from file **inc/cccc.f**. This definitions differs from INCLUDE because NEEDS cccc refers to a dictionary entry whilst INCLUDE cccc refers to a full-path filename with explicit extension.

Some characters are illegal for filename: noticeably the "double-quote" character (") is among them. In such case, these characters are converted into "tilde" (~) and this file is then searched for. For example:

At the moment we are writing, this definitions has a flaw: in case of interpretation/compilation error, the handle to the file remains open and you must close it manually using something like 3 F CLOSE. [...to be fixed...].

NFA pfa --- nfa

Converts a word's pfa into its nfa. See also CFA, LFA, PFA, >BODY, <NAME.

NIP n1 n2 --- n2

Removes the second element from TOS. See also: OVER, DROP, TUCK, SWAP, DUP, ROT.

NMODE --- a

User variable that indicates how double numbers are interpreted. During the input, numbers can be read as double integer numbers or floating-point numbers. This variable is modified by the optional words INTEGER that sets it to 0 and FLOATING that sets it to 1.

NOOP ---

This token does nothing. Useful as a placeholder or to prevent crashes in INTERPRET.

NUMBER a --- d a --- fp (compile time)

Converts a counted string at address a with a in a double number. If NMODE is 0, the string is converted to double integer. Position of the last decimal point encountered is kept in DPL.

If NMODE is 1, a floating-point number conversion is tried.

If no conversion can be done, and error #0 is raised.

OFFSET --- a

User variable that states the beginning of "blocks area". The content of OFFSET is added by BLOCK to the number on TOS to determine the right offset to read from file open to "!Blocks.bin". Messages issued by MESSAGE are independent from OFFSET.

OR n1 n2 --- n3

Executes an OR binary operation between the two integers. The operation is performed bit by bit.

OUT --- a

User variable incremented by EMIT. The user can examine and alter OUT to control the video formatting.

Copies the second number from TOS and put it on the top. See also DROP, NIP, TUCK, SWAP, DUP, ROT.

P! u b ---

Sends to port u a byte b. Note: u is a 16 bit port address and an OUT (C) op-code is internally executed.

P@ n --- b

Accepts the byte b from port u. Note: u is a 16 bit port address and an IN(C) op-code is internally executed.

PAD ---

Leaves on TOS the address of text output buffer. It is at a fixed distance of 68 byte over HERE.

PFA nfa --- pfa

Converts a word's nfa to its pfa. See also CFA, LFA, NFA, >BODY, <NAME.

PICK n --- pfa

Picks the **n-th** element from TOS. This means:

- O PICK is the same as DUP
- 1 PICK is the same as OVER

PLACE --- a

User variable that holds the number of places after the decimal point to be shown during a numeric output conversion. See also PLACES.

PREV --- a

User variable that points to the last referred buffer. UPDATE marks that buffer so that it is later written to disk.

QUERY ---

Awaits from terminal up to 80 characters or until a CR is received. The text is stored in TIB. User variable IN is set to zero.

OUIT ---

Clears the Return-Stack, stops any compilations and return the control to the operator terminal. No message is issued.

R --- n

Copies to TOS the value on top of Return Stack without alter it.

R# --- a

User variable that holds the position of the editing cursor or other function relative to files.

R/W anf ---

Standard FIG-FORTH read-write facility. Address a specifies the buffer used as source or destination; n is the sequential number of the block; f is a flag, 0 to Write, 1 to Read. \mathbb{R}/\mathbb{W} determines the location on mass storage, performs the transfer and error checking.

R0 --- a

User variable that holds the initial value of the Return Stack Pointer. See also RP! and RP@.

R> --- n

Removes the top value from Return Stack and put it on TOS. See also >R, R and RP!.

RECURSE ---

Used only at compile-time inside a colon-definitions, it compiles the definition being created to put in place a recursion call. This word is available after a <code>NEED RECURSE</code>.

REG! b n ---

Writes value b to Next REGister n.

REG@ n --- b

Reads Next REGister n giving byte b.

RENAME ---

Used in the form:

RENAME cccc xxxx

it searches the word cocc in the CONTEXT vocabulary and changes its name to xxxx. The two word names cocc and xxxx must have the same length. This definition is available after NEEDS RENAME.

REPEAT a1 n1 a2 n2 --- (immediate) (compile time)
--- (run time)

Used in colon defintion in the form:

BEGIN ... WHILE ... REPEAT

At run-time REPEAT does an inconditional jumt to the corresponding BEGIN.

At compile-time REPEAT compiles BRANCH and the offset from HERE to all and resolves the offset from all to the location after the loop; n1 and n1 are used for sysntax check.

Rotates the three top integers, taking the third an putting it on top. The other two integer are pushed down from top by one place. See also OVER, DROP, NIP, TUCK, SWAP, DUP.

RP! a ---

System procedure to initialize the Return Stack Pointer to the value passed on TOS that should be the address held in R0 user variable.

RP@ --- a

Leaves the current value of Return Stack Pointer.

RSHIFT n1 u --- n2

Shifts right an integer n1 by u bit.

S->D n --- d

Converts a 16 bit integer into a 32 bit double integer, sign is preserved.

so --- a

User variable that holds the initial value of che SP register. See also: SP! and SP@.

SCR --- a

User variable that hold the number of the last screen retrieved with LIST.

SELECT n ---

Selects the current channel. As usual for ZX Spectrum, n is 0 and 1 for lower part of screen, 2 for the upper part, 3 for printer, 4 for "!Blocks.bin" stream. Note: KEY always select chanle 2 to display the (flashing) cursor.

SIGN n d --- n

If n is negative, it puts an ASCII "-" at the beginning of the numeric string converted in the text buffer. Then, n is discarded while d is kept. Used between <# and #>.

SMUDGE ---

Used by the creation word: during the definition of a new word; it toggles the smudge-bit of the first byte in the nfa of the LATEST defined word. When a word's smudge-bit is set, it prevents the compiler to find it. This is typical for uncomplete or not correctly defined words.

It is also used to remove malformed incomplete words via

SMUDGE FORGET ccc

SOURCE-ID --- a

User variable that keeps the file-handle used during INCLUDE or NEEDS.

SP! a ---

System procedure to initialize the SP register to the address a that should be the address hold in S0 user variable.

SP@ --- a

Returns the content of SP register before SP@ was executed.

SPACE ---

Emits a space to the current output peripheal, usually the video. See also SELECT.

SPACES n --Emits n spaces.

SPAN --- a

User variable that holds the number of characters got from the last ${\tt EXPECT}$.

SPLASH --- a

Shows splash screen build date-number.

STATE --- a

User variable that holds the compilator status. A non-zero value indicates a compilation in progress.

STRM --- a

Variable containing the stream number used by the Screens/Blocks facility. Used by NextZXOS calls.

See also NXTDRV, NXTSTP, NXTRD, NXTWR.

SWAP n1 n2 --- n2 n1

Swaps the two top element at the TOS. See also OVER, DROP, NIP, TUCK, DUP, ROT.

THEN a n --- (immediate)
--- (compile time)

Synonym of ENDIF.

TIB --- a

User variable that holds the address of the Terminal Input Buffer.

TO n ---

Used in the form:

TO cccc

It assigns the value n to the variable ccc previously defined via VALUE. This definition available after NEEDS TO.

TOGGLE a b ---

The byte at location address a is XOR-ed with the model b.

TRAVERSE a1 n --- a2

Spans through the name-field of a definition depending on the value of n.

If n = 1, then all must be the beginning of the name-field, i.e. nfa itself; all is the address of the last byte of the name field.

If n = -1, then a1 must be the last byte of name-field and a2 will be the nfa. Used by da NFA and PFA.

TRUV ---

"True video". It disables Inverse-Video attribute mode. See also INVV.

Questa definizione è disponibile solo dopo il caricamento del "Line Editor" tramite 10 LOAD.

TUCK n1 n2 --- n2 n1 n2

Takes the top element of calculator stack and copies after the second. See also OVER, DROP, NIP, SWAP, DUP, ROT.

TYPE a n ---

Sends to the current output peripheal n characters starting from address a.

U. u ---

Prints an unsigned integer followed by a space.

U< u1 u2 --- f

Leaves a tf if u1 is less than u2, a ff otherwise.

UM* u1 u2 --- ud

Unsigned product of the two integers u1 and u2. The result is a double integer.

UM/MOD ud u1 --- u2 u3

Leaves the quotient u3 and the reminder u2 of the integer division of ud / u1.

UNTIL a n --- (immediate) (compile time)
f --- (run time)

Used in colon definition in the forms

BEGIN ... UNTIL

At run-time UNTIL controls a conditional jump to the corresponding BEGIN when f is false; the exit from the loop happens if f is true.

At compile-time UNTIL compiles <code>OBRANCH</code> and an offset from <code>HERE</code> to a; n is used for syntax checking.

UPDATE ---

Marks as modified the most recent used buffer, the one pointed by PREV. The block contained in the buffer will be transferred to disk when that buffer is requested for another block.

UPPER c1 --- c2

This word converts a character to upper-case. If c1 is not between "a" and "z", then c1 is left unchanged.

USE --- a

User variable that holds the buffer address of the block to be read from disk or that has just been written to.

USER n ---

Defining word used in the form

n USER cccc

creates an user variable 'ccc'. The first byte of pfa of ccc is a fixed offset for the User Pointer, that is the pointer for the user area. In this implementation there is only one User Area and a fixed User Pointer.

When ccc is later executed, it put on TOS the sum of offset and User Pointer, sum to be used ad the address for that specific user variable. The user variable are: TIB, WIDTH, WARNING, FENCE, DP, VOC-LINK, FIRST, LIMIT, EXP, NMODE, BLK, IN, OUT, SCR, OFFSET, CONTEXT, CURRENT, STATE, BASE, DPL, FLD, CSP, R#, HLD, USE, PREV, LP, PLACE, DL.

VALUE n ---

Defining word used in the form:

n VALUE cccc

Creates the word ccc that acts as a variable. To store a value in such a variable you have to use TO.

When ccc is later executed it directly returns the value of the variable without the need to access its address using @.

This definition is available after NEEDS VALUE.

VARIABLE n ---

Defining word used in the form:

n VARIABLE cccc

creates the word cocc with the pfa containing the initial value n. When cocc is executed, it puts on TOS the pfa of cocc that is the address that holds the value n.

When used in the form

cccc @

the content of the variable cccc is left on TOS.

When used in the form

n cccc !

the value on TOS is stored to the variable cccc.

VIDEO ---

Sets DEVICE 2 to select the video as current output peripheral. See SELECT and DEVICE.

VOC-LINK --- a

User variable that holds the address of a field in the definition of the last vocabulary. Each vocabulary is part of a linked-list that uses that field, in each vocabulary definition, as pointer-chain.

VOCABULARY ---

Defining word used in the form

VOCABULARY cccc

creates the word $\, {\tt ccc} \,$ that gives the name of a new vocabulary.

Later execution of

cccc

makes such vocabulary the CONTEXT vocabulary, so that it is possible to search for words defined in this vocabulary first and execute them.

Used in the form

ccc DEFINITIONS

makes such vocabulary the CURRENT vocabulary, so that it is possible to insert new definitions in it.

WARM ---

Executes a warm system restart. It closes and reopen Block/Screen file then does ABORT. It does not EMPTY-BUFFERS.

WARNING --- a

User variable that determines the way an error message is reported. If zero, only a short "MSG#n" is reported. If non zero, a long message is reported. See also ERROR.

WHILE f --- (immediate) (run time)
a n --- al nl a2 n2 (compile time)

Used in colon defintion in the form:

BEGIN ... WHILE ... REPEAT

At run-time WHILE does a conditional execution based on f. If f is true, the execution continues to a REPEAT which will jump to the corresponding BEGIN. If f is false, the execution continues after the REPEAT quitting the loop.

At compile-time WHILE compiles <code>OBRANCH</code> leaving a2 for the offset; a2 will be comsumed by a <code>REPEAT</code>. The address a1 and the number n1 was left by a <code>BEGIN</code>.

WIDTH --- a

User variable that indicates the maximum number of significant characters of the words during compilation of a definition. It must be between 1 and 31.

WORD c --- a

Reads one or more characters from the current input stream up to a delimiter c and stores such string at HERE that is left on TOS. WORD leaves at HERE the length of the string as the first byte and ends everything with at least two spaces. Further occurrences of c will be ignored.

If BLK is zero, the text is taken from the terminal input buffer TIB. Otherwise the text is taken from the disk block held in BLK. User variable >IN is added with the number of character read, the number ENCLOSE returns.

WORDS ---

Shows a list of words of CONTEXT vocabulary. Pressing Break stops.

Executes a XOR binary operation between the two integers. The operation is performed bit by bit.

[--- (immediate)

Used in colon defintion in the form:

: cccc [...] ... ;

it suspends compilation. The words that follows [will be executed instead of being compiled. This allows to perform some calculations or start other compilers before resuming the original compilation with]. See also LITERAL.

[CHAR] --- (immediate) (compile time)

It is the same as the sequence [CHAR c] LITERAL.

It is used in colon defintion in the form:

: cccc ... [CHAR] c ...;

At compile time, [CHAR] compiles LIT and the numeric value of ASCII character c in the following cell.

[COMPILE] --- (immediate)

Used in colon defintion in the form:

```
: cccc ... [COMPILE] wwww ...;
```

[COMPILE] forces the compilation of a definition wwww that is immediate. Normally immediate words aren't compiled but executed and to compile an immediate word it is not possible to use the sequence COMPILE wwww but it is necessary using the sequence [COMPILE] wwww.

Used in the from:

١ . . .

Any character that follow \ until the end of line are treated as a comment.

] ---

Resumes the compilation suspended by $\ [$ so it is possible to complete the definition.

Line Editor

The following definitions are available after you give 90 LOAD.

There is no single definition that needs all the following words, but EDIT is the best candidate: so instead of using the old fashion 90 LOAD you may use NEEDS EDIT.

The Line Editor has a dozen words that can operate on a single line of a given Screen and helps inspect things around.

An edit session normally starts with a LIST on the desired Screen, this sets SCR user variable to the passed Screen number. LIST is a word already available in the "core" dictionary. To clear a Screen I foreseen a BCLEAR word, but I left it commented in Screen# 13 for now, deeming it too dangerous for my tastes; instead I usually use BCOPY from an actually empty Screen. You may give NEEDS BCOPY.

The word FLUSH flushes to disk any modification you've done on any Screen. Beware, a Screen is re-written to disk as soon as the BUFFERS containing it are modified. To save space, this implementation has only three BUFFERS.

EMPTY-BUFFERS is another vital word: it empties all buffers. It is very useful if you mistakenly overwrite or spoil a Screen during an edit operation, with it, you have the chance to "rollback" the things before the anything is written to disk.

To write a line from scratch or to overwrite line, you can use P to "put" the following text to the given line on current screen. For example:

```
1000 LIST
0 P \ One thousand screens
L
```

This sequence selects Screen#1000 and put a text "One thousand screens" on the first line of it. The word L repeats the LIST of current screen.

To move or copy a line around, you can use <code>H</code> to "hold in <code>PAD</code>" a given line on current screen, you can change Screen if you wish, then you can complete this **copy-and-paste** operation with <code>INS</code> to "insert" or <code>RE</code> to "replace" the line you copied in advance with <code>H</code>. None of above words, but <code>H</code>, modify PAD content, so you can repeat the operation. There is also a way to **cut-and-paste** a line using <code>D</code> to "delete and copy to <code>PAD</code>" instead of <code>H</code>.

See also BLOCK, BUFFER, INDEX, L/SCR, LIST, LOAD, MESSAGE, PAD, SCR, STRM., TIB.

This is a quick reference of involved memory areas and words that work on them.

| Text Input Buffer | Parsing | | Edit | One | Blanking |
|-------------------|-----------|-----|------------|--------|------------|
| (keyboard) | Operation | | Operations | | Operations |
| | | | | BLOCK | |
| TIB | | PAD | | BUFFER | |
| | TEXT > | | ← H RE → | | ← E |
| | | | ← D INS → | | ← s |
| | | | ₽ → | | |

| -MOVE "Line move". It moves a li Current screen is the one | | gth, from address $$ a $$ to the line $$ n $$ of current screen, then it does an <code>UPDATE</code> . |
|---|-----------------------------------|--|
| . PAD "Show PAD". It prints the | current PAD cont | ent. |
| B "Back" one Screen. This w | ord set to previou | us Screen by decreasing SCR and prints it using LIST. |
| | | screen (the one indicated by ${\tt SCR}$), the following lines are moved up and the at it can be followed by an ${\tt INS}$ to perform a line move. |
| BCOPY "Block-Copy" utility that of | n1 n2 copies Screen n1 t | co Screen n2. SCR will contain n2. |
| E "Erase" a row. This word | n fills line n with spa | aces. It does UPDATE. |
| H "Hold" a row in PAD. This the one kept in SCR. | n word put line n o | f current Screen to PAD without altering the block on disk. Current Screen is |
| INS "Insert" from PAD. This wand the last is lost. | n vord inserts line n | $$ using text in PAD. The original line ${\tt n}$ and the following ones are moved down |
| L "List" current Screen. This | s word does SCR | @ LIST. |
| LINE | n | a |

Leaves the address $\ a$ of line $\ n$ of current screen, the one kept in SCR. Such a screen is currently held in a buffer.

"Next" Screen. This word sets to next Screen by increasing ${\tt SCR}$ and prints it using ${\tt LIST}.$

N

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| P | n | |
|---|----|--|
| P | 11 | |

"Put" a line. This word accepts the following text (delimited by a tilde character \sim) as the text of line n of current Screen. Text is taken from TIB and sent to the current Screen

RE n ---

"Replace". This word takes text currently in ${\tt PAD}\,$ and put it to line n.

S n ---

"Space" one row. This word frees line n moving the following lines down by one. The last line is lost

SAVE ---

Does <code>UPDATE</code> and <code>FLUSH</code> saving this Screen and all previously modified Screens back to disk.

ROOM ---

This word shows the room available in the dictionary, that is the difference between SP@ and PAD addresses.

TEXT C ---

This word accepts the following text and stores it to PAD. c is a text delimiter. TEXT does not go beyond a 0x00 [null] ASCII.

UNUSED --- n

It returns the number of byte available in dictionary.

WHERE n1 n2 ---

Usually executed after an error has been reported during a LOAD session. Maybe, this word should be included in "core" dictionary. n1 is the value of IN and n2 the value of BLK as were left by ERROR.

WHERE shows on screen the block number, the line number, the very same line highlighting in "inverse video" the word that caused the error.

Case -Of structure

The following definitions are available after you give 17 LOAD.

CASE n0(immediate) (run time) n (compile time) а

Used in colon definition in the form

```
n0 CASE
            ... ENDOF
    n1 OF
            ... ENDOF
    nz OF
    ... ( else )
ENDCASE
```

The word CASE marks the beginning of Case-Of structure i.e. a set of branches where only one is performed based on the value of n0. If none of the "OF clause" values matches, the \mathtt{ELSE} part is performed.

At compile time CASE leaves previous CSP address a and a number n for syntax checking. CASE has to be balanced by a corresponding ENDCASE.

OF n0 nk (immediate) (run time) n1 a n2 (compile time)

This word is used in colon-definition within a Case-Of structure.

At run-time it compares the value now on TOS nk with the value n0 that was on TOS just before the beginning of the Case-Of structure.

At compile-time, it compiles (OF) and OBRANCH using n1 and n2 for syntax checking and leaving a to be used by ENDCASE to resolve OBRANCH.

See also CASE.

(immediate) **ENDOF** (run time) a1 n1 a n2 (compile time)

This word ends an "Of-EndOf" clause started with OF.

At compile-time it acts like a THEN, first compiling a BRANCH that will be resolved by ENDCASE to skip any subsequent "Of-End-Of" clauses and resolving the <code>OBRANCH</code> compiled by the corresponding previous <code>OF</code> to continue the Case-Of structure.

See also CASE.

ENDCASE (immediate) (run time) a a1 ... az (compile time)

This word ends a Case-Of structure started with CASE.

At compile-time it compiles a DROP to discard the value n0 put on TOS before CASE and resolves all OF-ENDOF clauses to jump after the ENDCASE. Finally, it restores previous content of CSP. See also CASE.

n0 nk (run time) (OF)

This word represents the run-time semantic compiled by OF word. At run-time, it compares the value now on TOS nk

with the value n0 that was on TOS just before the beginning of the Case-Of structure and leave a flag to be used by the following OBRANCH (that was compiled by OF). When n0 equals nk, the definitions between OF and ENDOF will be executed, otherwise a jump to the word after ENDOF is performed.

Heap Memory Facility

The definitions that handle the Heap are available after loading via 80 LOAD or NEEDS HEAP .

Among ZX Spectrum Next new features is the huge amount of RAM. The question is how to leverage all that memory in Forth.

Comparing the current version to how previous Forth's system areas are sorted out, the first challenge is to move them down to free the top 8K CPU's addressable memory between 0E000h and 0FFFFh allowing MMU7 to map to any physical 8K RAM page. Strings are dictionary expensive, so it would be useful storing them in heap as constant-strings and fetch them at need. More, 8K of room is a good place to store an array of strings, or even numeric array and implement some matrices algebra.

There are some peculiar addresses that identify the following Forth system areas:

OF840h: Calculator Stack (SP) grows downward, Text Input Buffer (TIB) upward.,

OF8E0h: Return Stack (RP) grows downward, User Variables Area upward.

0F94Ch: the FIRST disk buffer starts here and buffers area ends just before LIMIT 0FF58h.

I coded this "move" in a few words (available in Screens #220-#223) summarized in the definition DOWN that moves these pointers "down" as follow:

0FF58h \rightarrow 0E000h : LIMIT 0F9C4h \rightarrow 0D9F4h : FIRST

0F8E0h → 0D9A0h : Return Stack and User Variables Area

0F840h → 0D900h : Stack Pointer and TIB

Heap Pointer encoding and decoding

The second issue arises when we need a way to encode both **page number** and **address offset** in a usual Z80 16-bits pointer variable.

Two definitions are made available to perform these coding and decoding operations: >FAR and <FAR.

Given an address a (to be intended as an offset of addresses between E000h and FFFFh) and a page number in the definition >FAR encodes the page number in the most significant bits of ha and an offset in the remaining less significant bits. The inverse function is performed by <FAR. Splitting a 16-bit "heap-pointer-number" into the page part and the offset part again.

In the following paragraphs a couple of possible implementations are described in detail.

Heap Pointer description for 64 KBytes space

The following solution opens to 64K of physical RAM: Since an 8K offset requires 13 bits, the remaining 3 bits can be used to encode, say, from page 64 (\$40) to page 71 (\$47). For example:

The encoding/decoding definitions would be something like the following:

```
CODE
        >FAR ( ha --- a n )
        pop
                de
        ld
                a,d
        and
                $E0
        rlca
        rlca
        rlca
                $40
        add
        ld
                1,a
        ld
                h,0
                        ; hl = page number between 64 and 71
        ld
                a,d
                $E0
        or
        ld
                d,a
                        ; de = offset at $E000
                de
        push
                hl
        push
                (ix)
        jр
```

```
CODE
       <FAR ( a n --- ha )
       pop
               hl
                        ; hl = page number between 64 and 71
                de
                        ; de = offset at $E000
       pop
       ld
                a,l
       and
                $07
       rrca
       rrca
       rrca
       ld
               h,a
       ld
                a,d
                $1F
       and
        or
                h
       ld
                d,a
       push
                de
                        ; de = heap-pointer
       jр
                (ix)
```

Heap structure

The Heap is a "linked-list" starting at page \$40 offset \$0002. The User variable HP keeps the "heap-pointer" to the next available location on Heap. So, at startup, HP is \$0002 that correspond to page \$40 offset \$0002.

Any Heap new memory allocation reserves the requested number of bytes and advances HP to point to the next available location on Heap. The value of ${\tt HP}$ is then stored at the location that was available before the memory allocation was requested to put in place a "linked-list".

At startup, HP is \$0002 and the Heap memory looks like the following

```
Location Content

$40:$0000 $0000 (this marks the "end" of Heap )
$40:$0002 free memory pointed by HP
```

Suppose we want to reserve 5 bytes of Heap, we can give 5 HEAP that will return \$0004 as the "Pointer" the Heap area just allocated, while HP User variable will be advanced to \$0009. After the execution the memory will look like:

| \$40:\$0000 \$0000 (this marks the "end" of Heap) \$40:\$0002 \$0009 (final value of HP) \$40:\$0004 00 00 00 00 (5 bytes just allocated) \$40:\$0009 free memory pointed by HP | Location | Content |
|--|----------|---|
| | | |
| | | 00 00 00 00 00 (5 bytes just allocated) free memory pointed by HP |

Then, we want to reserve another 7 bytes, we can give 7 HEAP that will return \$000B as the "Pointer" the this new chunk of memory and HP will be advanced to \$0012. After the execution the memory will look like:

Now, you should be able to see the Linked-List starting at \$0002 that points to \$0009 that points to \$0012. You can follow all these Pointers using the following procedure:

```
2 .S \ Stack is 0002 as Heap-Pointer, that is $40:$0002, the beginning of Heap Memory. FAR .S \ Stack is E002 as real Address (and page $40 is fitted in MMU7)
```

@ .S \ Stack is 0009 as Heap-Pointer

FAR .S \ Stack is E009 as real Address (and page \$40 is fitted in MMU7)

@ .S \ Stack is 0012 as real Address

Some low-level definitions are available to allow to store on heap and retrieve from it and how to avoid that my string is "paged away" in the middle of processing, that is how to let a page to stay in place across Standard-ROM calls or I/O disk operations, that use addresses C000-FFFF for their purposes:

MMU7! to fit a given 8K page number at E000h (i.e. MMU7).

>FAR to quick decode a "16 bit pointer" splitting it into "page & offset" as shown in the post above.

User variable HP has been introduced to keep track of room in Heap: it is the "pointer" to the next available space on Heap.

The following definitions are available after loading via 80 LOAD. or via NEEDS

+" ha --- ha

Appends a string to the last string, return an heap-address pointer to a counted string.

>FAR ha --- a n

Given a heap-encoded pointer ha this definition decodes the top bits as one of the 8K-page available and the lower bits as the offset from E000h. It does not change the MMU7 page. See <FAR, MMU7!

<FAR a n --- ha

Given an address a (to be intended as a physical address between E000h and FFFFh) and a page number $\,\mathrm{n}\,$ for an 8K-page between 64 and 71 (or 40h – 47h) this definition encodes the page number in the most significant bits of $\,\mathrm{ha}\,$ and an offset in the remaining bits. It does not change MMU7 page. See >FAR, MMU7!

FAR ha --- a

This definition converts a heap-pointer ha into an offset a (at E000h) and perform the correct 8K paging on MMU7. It simply calls >FAR and MMU7!

H" --- ha

Accepts a string and store it to Heap, and return an heap-address pointer to a counted string.

HEAP n --- ha

This definition reserves n bytes on Heap

POINTER ha --- a

It works like CONSTANT but it returns a "FAR-resolved" pointer.

A possible use is: S" ccc" POINTER P1

SKIP-PAGE n ---

Check if n bytes are available on current 8K-page in Heap, otherwise advance HP to skip to the beginning of next 8K-page.

S" --- a n

Accepts a string and store it to Heap: at compile time it compiles (S") and the heap-pointer just after it, during direct interpret (and at runtime) it returns a real-address at MMU7 and a counter.

This is the run-time counterpart of S" that use the Heap-Pointer that follows to fit the correct 8K-Page in MMU7 using FAR definition and leaving the real-address a and the length of the string n.

+C ha c --- ha

Consume a character c and append the string being created in Heap at ha. The heap pointer ha is returned unchanged.

+" ha --- ha

Accept a string and append it to the string being created in Heap at ha. The pointer ha is returned unchanged.

HEAP-INIT ---

Ask NEXTZXOS to use pages \$40-\$47 for Heap. From this point Heap command can be used safely.

HEAP-DONE ---

Release to NEXTZXOS pages \$40-\$47. Heap commands should not be used after.

Heap Pointer description for 128 KBytes space

To allow **twice** the number of pages we need one more bit for the page number at the expense of the remaining offset bits; for example, 4 bits for page number allows encoding from page 64 (\$40) to page 79 (\$4F), that is 128K, and leaves the remaining 12 bits for offset to addresses. We notice that **only even** addresses can be directly referenced.

Encoding routine

```
CODE
        >FAR ( ha --- a n )
        pop
                hl
        ld
                a,h
                $F0
        and
        rlca
        rlca
        rlca
        rlca
        add
                $40
        ld
                e,a
        ld
                d,0
                        ; de = page number between 64 and 71
        add
                hl,hl
                        ; shift
        ld
                a,h
                $E0
        or
                        ; hl = offset at $E000
        ld
                h,a
        push
                hl
                de
        push
        jр
                (ix)
```

Heap Pointer description for 256 KBytes space

With 5 bits for page number and 11 bits for offset, we'll have 32 pages between 64 (\$40) and 95 (\$5F), that is 256K and only addresses divisible by 4 can be directly referenced.

```
16 bit pointer

| pppp pbbb | bbbb bbbb |
| | | | |
| Page | Offset |
| 010p pppp | 111b bbbb | bbbb bb00 |
```

Encoding routine

```
CODE
       >FAR ( ha --- a n )
       pop
               hl
       ld
                a,h
       and
                $F8
       rrca
       rrca
       rrca
                $40
       add
       ld
                e,a
       ld
                d,0
                        ; de = page number between 64 and 71
       add
               hl,hl
                        ; shift
       add
               hl,hl
                        ; shift
       ld
                a,h
                $E0
       or
       ld
               h,a
                        ; hl = offset at $E000
       push
               hl
       push
                de
                (ix)
       jр
```

Heap Pointer description for 512 KBytes space

With 6 bits for page number and 10 bits for offset, we'll have 64 pages between 64 (\$40) and 127 (\$7F), that is 512 and only addresses divisible by 8 can be directly referenced.

```
16 bit pointer

| pppp ppbb | bbbb bbbb |
| | | | | |
| Page | Offset |
| 01pp pppp | | 111b bbbb | bbbb b000 |
```

Encoding routine

```
CODE
       >FAR ( ha --- a n )
               hl
       pop
       ld
               a,h
       and
               $FC
       rrca
       rrca
               $40
       add
       ld
               e,a
               d,0
       ld
                       ; de = page number between 64 and 71
       add
               hl,hl
                       ; shift
                      ; shift
       add
               hl,hl
                       ; shift
       add
               hl,hl
       ld
               a,h
       or
               $E0
       ld
               h,a
                       ; hl = offset at $E000
       push
               hl
       push
               de
                (ix)
       jр
```

Heap Pointer description for 1024 KBytes space

Pursuing this path to the limit we can use 7 bits for page number we can pick 128 distinct pages, for example from page 72 (\$48) to page 199 (\$C7) that leads to **1024K** of physical RAM, at the downside to be able to reference only physical addresses divisible by 16.

```
16 bit pointer

| pppp pppb | bbbb bbbb |
| | | | |
| Page | Offset |
| xppp pppp | | 111b bbbb | bbbb 0000 |
```

Coding/decoding routines

```
CODE
        >FAR ( ha --- a n )
        pop
                hl
        ld
                a,h
        srl
                а
                $48
        add
        ld
                e,a
                        ; de = page number between 72 and 199
        ld
                d,0
        add
                hl,hl
        add
                hl,hl
        add
                hl,hl
                        ; shift hl 4 bits left
        add
                hl,hl
        ld
                a,h
        or
                $E0
        ld
                        ; de = offset at $E000
                h,a
                hl
        push
        push
                de
                (ix)
        jр
```

```
CODE
        <FAR ( a n --- ha )
                de
                        ; de = page number between 72 and 199
       pop
       pop
                hl
                        ; hl = offset at $E000
       ld
                a,e
                $48
        sub
       add
               hl,hl
       add
               hl,hl
               hl,hl
        add
                        ; tricky shift h 4 bits right
        add
               hl,hl
        rla
                        ; A receives HL msb
       ld
                1,h
       ld
                h,a
       push
                hl
                (ix)
        jр
```

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| (COMPARE) a1 a2 n b | |
| (DO) | |
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| (EIND) at az Cla D CL | |

| (LINE) n1 n2 a b | |
|--|----|
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| ?EXEC | |
| ?LOADING | |
| ?PAIRS n1 n2 | |
| ?STACK | |
| ?TERMINAL f | |
| @ a n | |
| ABORT | |
| ABS n u | _ |
| ACCEPT a n1 n2 | |
| ACCEPT- a n1 n2 | |
| AGAIN (immediate) (run time) | |
| ALLOT n | |
| AND n1 n2 n3 | |
| AUTOEXEC | _ |
| B/BUF n | |
| B/SCR n | _ |
| BACK a n | |
| BASE a | |
| | |
| | |
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| EMITC b | |
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| ERASE a n | |
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