Overview
Difference to BFS
Input graphs
Implementations
Architecture
Results

PARALLEL TOPOLOGICAL SORTING

Design of High Performance Computing, Fall 2015

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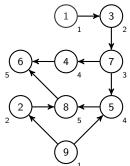
ETH Zürich

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OVERVIEW
DIFFERENCE TO BFS
INPUT GRAPHS
IMPLEMENTATIONS
ARCHITECTURE
RESULTS

OVERVIEW

- DAG defines partial order
- Topological sorting defines one total order on a DAG
- Parallel algorithm: finds one topological sorting of a given DAG



DIFFERENCE TO BFS

- BFS visits every node
- Topological sorting algorithm needs to visit every edge

Example:

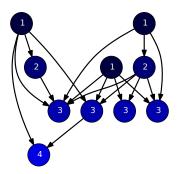


Consider order A,C,B \rightarrow valid in BFS, invalid in topological sorting

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RANDOM GRAPH

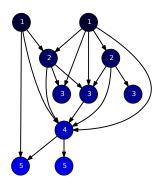
Parameter: Average node degree



SOFTWARE DEPENDENCY GRAPH

Nodes	Edges	Degree (Median)
100'000	266'680	2

10'000 27'416 2



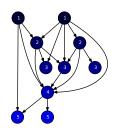
 ${\hbox{Musco, V. et al. (2014) "A Generative Model of Software Dependency Graphs to Better Understand Software} \\$

Evolution."



PARALLEL ALGORITHM (SHARED MEMORY)

- As a preparing step, initialize a counter for every child with the number of its parents.
- ② Distribute parent nodes over threads and process them in parallel.
- ③ For every parent node, get a list of all child nodes and append the parent itself to solution (→ lock).
- **4** For every child of the parent, decrement its counter $(\rightarrow \mathbf{lock})$. Once the counter is zero, we can move on
- When all parents are processed (→ barrier), distribute new nodes and repeat



LOCAL LISTS APPROACH

Idea: perfect load balancing by redistributing nodes at every step

- Parent nodes stored in a global list.
- Distribution of parent nodes: <u>scatter</u> the list among the threads. Each thread has now its local list.
- Add new nodes to the end of local list.
- When all parents were processed, gather all local lists into the global list.
- Repeat until there are no parents left in the global list.

BOOLEAN ARRAY APPROACH

Idea: minimizing memory access by using lookup table

- Array of length N. 1 if node i is a parent node, 0 otherwise.
- Distribution of parent nodes: Parallel for-loop through the array.
- Mark new parents by setting a 1 in a second array.
- When all parents were processed, swap arrays.
- Loop through the array until there are no new parents.

- Decrement shared counter
- Return true if counter is zero

```
inline bool counterCheck() {
    bool lastone;
    #pragma omp critical
    {
        --parcount_;
        lastone = (parcount_ == 0);
    }
    return lastone;
}
```

- Decrement shared counter
- Return true if counter is zero

```
inline bool counterCheck() {
    #pragma omp atomic
    --parcount_;
    return (parcount_ == 0);
}
```

 Multiple threads could return true, although only one thread should do so.

```
inline bool counterCheck() {
    #pragma omp atomic
    --parcount_;
    return (parcount_ == 0);
}
```

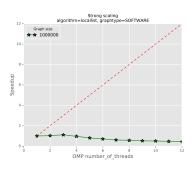
- List based approach: Child is inserted multiple times to solution ⇒ Wrong.
- Array approach: Multiple threads write 1 to the array ⇒ Ok, doesn't matter.

EULER

- Intel Xeon E5 on Euler cluster
- 2 processors per node
- 12 cores
- 30 MB shared last-level cache
- Software graph with around 1M nodes should fit into cache

STRONG SCALING SOFTWARE GRAPH

Local lists



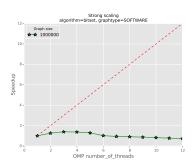
Absolute runtimes on 1 core

serial 0.45 s

bool array 0.58 s

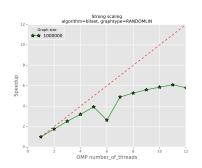
local list 0.48 s





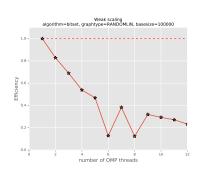
BOOL ARRAY & RANDOM GRAPH (NODE DEGREE 30)

Strong scaling



Absolute runtimes on 1 core

Weak scaling

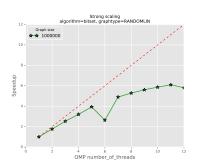


e serial

bool array 3.51 s

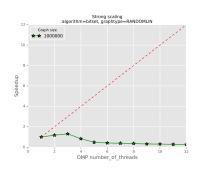
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Optimistic



Absolute runtimes on 1 core

Conventional



ore

conventional 3.52 s

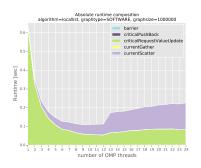
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optimistic

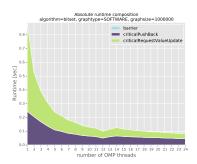
3.51 s

ABSOLUTE RUNTIME COMPOSITION

Local lists



Bool array



REMAINING ISSUES

- Pinpoint the reasons for bad scaling
- Work stealing could help to improve local list approach

REFERENCES



M. C. Er.

A Parallel Computation Approach to Topological Sorting The Computer Journal, Vol 28, 1983, Wiley Heyden Ltd



V. Musco, M. Monperrus, P. Preux

A Generative Model of Software Dependency Graphs to Better Understand Software Evolution

arXiv 1410.7921