

PARALLEL TOPOLOGICAL SORTING

DESIGN OF HIGH PERFORMANCE COMPUTING, FALL 2015

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PROBLEM DESCRIPTION

- Directed acyclic graph (DAG) describes a partial order
- A topological ordering is a total order on a DAG
- Any DAG has at least one topological ordering

"EFFICIENT" PARALLEL AND DISTRIBUTED TOPOLOGICAL SORT ALGORITHMS

Ma et al. 1997

- Runtime: $\mathcal{O}(\log^2 N)$
- Reduces to matrix-matrix multiplication problem

PROBLEM:

$\mathcal{O}(N^3)$ execution units required

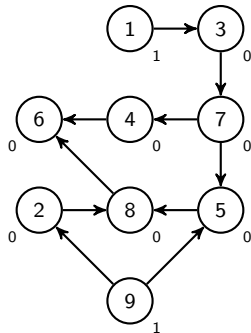
"EFFICIENT" PARALLEL AND DISTRIBUTED TOPOLOGICAL SORT ALGORITHMS

M. C. Er 1983

- Runtime: $\mathcal{O}(N)$
- More precise: longest distance between source node and sink
- Propagation of node values from all source nodes to all sink nodes

TOPOLOGICAL ORDERING - ALGORITHM

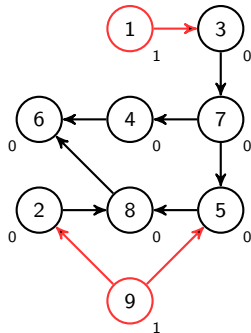
INPUT Directed acyclic graph (DAG) with N nodes
OUTPUT Topological Sortings of DAG



TOPOLOGICAL ORDERING - ALGORITHM

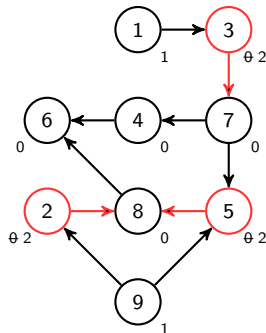
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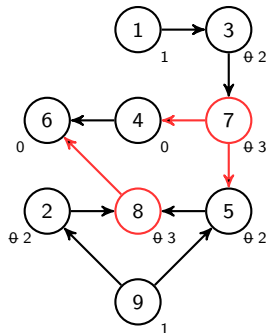
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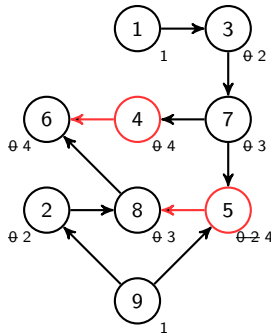
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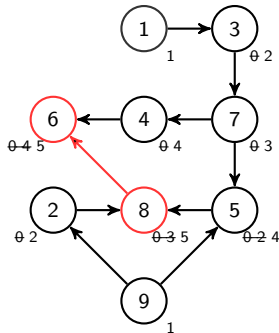
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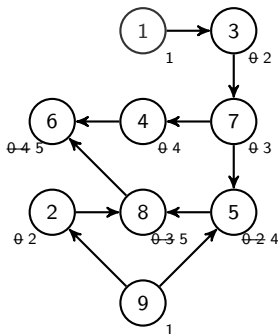


TOPOLOGICAL ORDERING - ALGORITHM

INPUT Directed acyclic graph (DAG) with N nodes
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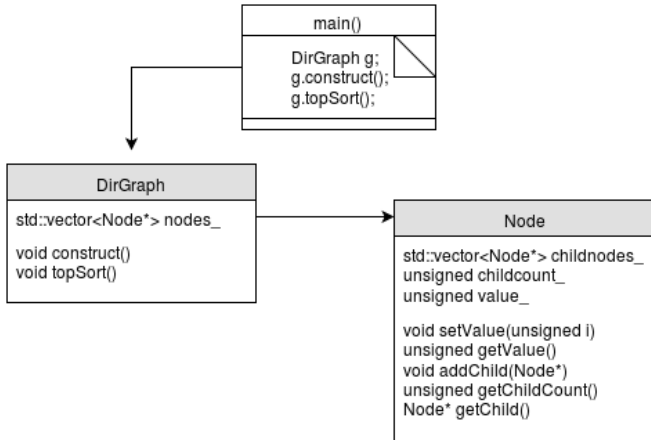


TOPOLOGICAL ORDERING - ALGORITHM



Name	1	9	2	3	7	4	5	8	6
Value	1	1	2	2	3	4	4	5	5

UML DIAGRAM



SERIAL CODE

```
std::list<Node*> currentnodes;
```

```
Node* parent;
```

```
Node* child;
```

```
unsigned childcount = 0;
```

```
unsigned currentvalue = 0;
```

SERIAL CODE

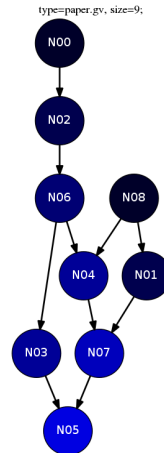
```
std::list<Node*> currentnodes;

Node* parent;
Node* child;
unsigned childcount = 0;
unsigned currentvalue = 0;

while(!currentnodes.empty()) { // stop when queue is empty
    parent = currentnodes.front();
    currentnodes.pop_front(); // remove current node - already visited
    currentvalue = parent->getValue();
    ++currentvalue; // increase value for child nodes
    childcount = parent->getChildCount();
    for(unsigned i=0; i<childcount; ++i) {
        child = parent->getChild(i);
        currentnodes.push_back(child); // add child node at end of queue
        child->setValue(currentvalue); // set value of child node to parentvalue+1
    }
}
```

PARALLELIZATION IDEAS

- Start in parallel at source nodes
- Spawn a new thread for every child node
- Synchronization needed when a node has multiple parents
- Performance dependent on topology of graph. Worst case $\mathcal{O}(N)$



- ```
type=paper.gv, size=9;
```



# HARDWARE / TOOLS

- Shared memory parallelization
- OpenMP or C++11 threads.
- On Euler: 12-core Intel Xeon E5
- Intel Xeon Phi

# CHALLENGES

- Task queue
- Find a way to cope with chain-like graphs.
- Optional goal: Find all possible topological sortings.

# QUESTIONS

- Up to how many execution units should we parallelize at least?
- From your experience, what are the main time-consuming, maybe unexpected obstacles when running on Intel Xeon Phi?