## Chapter 1:

Three basic components of computer security:

*Confidentiality*: The concealment of information or resources.

*Integrity*: The trustworthiness of data or resources, preventing unauthorized changes.

*Availability:* The ability to use information or resources (reliability).

Threat Model:

Classifications: Deception (fake news), Disruption (prevent operation), Usurpation (unauthorized control), Disclosure.

Characterizations: Alteration, Spoofing, Repudiation, Denial of Receipt, Delay, Denial of Service.

Policy and Mechanism:

*Security Policy*: A statement of what is, and what is not allowed.

*Security Mechanism*: A method, tool or procedure for enforcing a security policy.

Assumptions and Trust:

*Trust*: Your belief the system is trustworthy.

*Assurance*: Level at which the security mechanism implements the policy.

Let P be the set of all possible states, Q be the set of secure states, R be the set of states restricted by the security system.

A security mechanism is *secure* if R in Q, *precise* if R = Q, and *broad* if there are some states r not in Q.

*Specification*: A statement of the desired functioning of the system.

A system is said to *satisfy* a specification if the specification correctly states how the system will function.

*Design*: Translates the specifications into components that will implement them.

*Implementation*: Creates a system that satisfies the design.

A program is *correct* if its implementation performs as specified.

## Chapter 2: ACM

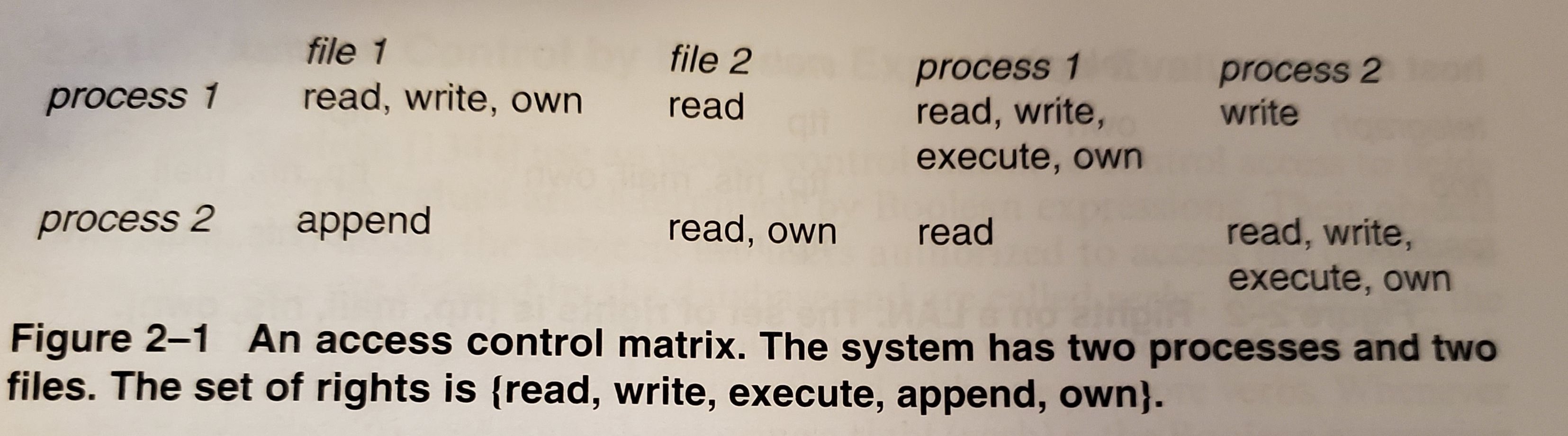
Protection State:

*State*: The collection of the current values in memory locations.

*Access Control Matrix Model*: A tool that can describe the current protection state. Which Describes the rights of subjects over all entities in the matrix.

Consider the set of protection states P. Some subset Q of P consists of exactly those states in which is the system is authorized to reside.

When a command changes the state of the system a *state transition* occurs.



Protection State Transitions: Let the intial state of the system be X\_0 = (S\_0, O\_0, A\_0) the set of state transitions is represented as a set of operations t,t2... successive states are represented as X\_1, X\_2,… where the notation X\_i (sideways T) \_(t\_i + 1) means that sate transition (t\_i + 1) moves the system from state X\_i to state X\_i + 1.

Harrison, Ruzzo, Ullman (HRU):

Primitive Commands:

1. Precondition: s not in S.
   1. Command: **create subject** s.
2. Precondition: o not in O.
   1. Command: **create object** o.
3. Precondition: s in S, o in O, r in R.
   1. Command: **enter** r **into** a[s,o].
4. Precondition: s in S, o in O, r in R.
   1. Command: **delete** r froma[s,o].
5. Precondition: s in S.
   1. Command: **destroy subject** s.
6. Precondition: o in O.
   1. Command: **destroy object** s.

Example:

**command** create\_file(p, f)

create **object** f;

**enter** *own***into** A[p,f];

**enter** *r***into** A[p,f];

**end**

**command** grant\_read\_file(p, f, q)

**if** *r* **in** A[p,f] **and** *g* **in** A[p,f]

**Then**

**enter** *r***into** A[p,f];

**end**

Note commands cannot have or negation like not in operators otherwise they would be two commands.

*Copy Right*: Often called the grant right allows the processor to grant rights to another.

*Own Right*: Enables the possessors to add or delete privileges.

Principle of Attenuation of Privilege: A subject may not increase its rights, nor grant rights that it does not possess to another subject.

## Chapter 3:

*Leaked*: When a generic right *r* is added to an element of the ACM that did not contain *r* initially, that right is said to be *leaked.*

*Safe System*: If a system can never leak the right *r*, the system is called *safe* with respect to right *r*. Alternatively is *unsafe with respect to right r.*

**Theorem 3.1:** There exists an algorithm that will determine whether a given mono-operation protection system with initial sate *s\_0* is safe with respect to generic right r.

You can omit the **delete, destroy, and create** commands beyond the first sequence.

**Theorem 3.1:** It is undecidable weather a given state of a given system is safe for a given generic right.

Proof by the Turing equivalent halting problem.

HRU in general undecidable – must try all possible command sequences.

Delete the **create** command, then the problem becomes complete in **P-SPACE.**

Delete the destroy, delete commands then the safety question is still undecidable.

Monotonic protection systems is decidable.

Mono conditional protection systems with **create, enter, delete (no destroy)** is decidable.

Take Grant Protection Model:

Subjects = Black Dot.

Objects = White Dot.

Subject or Object (doesn’t matter) = Dot with X filled in.

G |–x G' apply a rewriting rule x (witness) to G to get G'

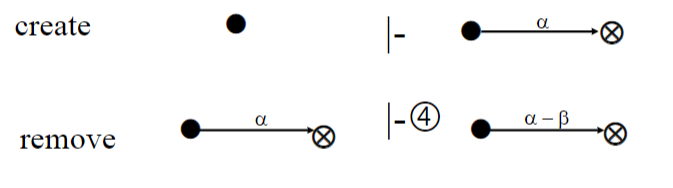
G |–\* G' apply a sequence of rewriting rules (witness) to G to get G'

R = { t, g, r, w, ... } set of rights

These four rules are called the de jure rules:

* *take rule* allows a subject to take rights of another object (add an edge originating at the subject)
  + Preconditions:
    - subject *s* has the right Take for *o*.
    - object *o* has the right *r* on *p*.
* *grant rule* allows a subject to grant own rights to another object (add an edge terminating at the subject)
  + Preconditions:
    - subject *s* has the right Grant for *o*.
    - *s* has the right *r* on *p*.
* *create rule* allows a subject to create new objects (add a vertex and an edge from the subject to the new vertex)
* *remove rule* allows a subject to remove rights it has over on another object (remove an edge originating at the subject)

## 



*Bridge:* Connects two islands, special type of path that has tg between islands.

*Can.Share(r,* ***x, y,*** *G\_0):* If and only if there is a sequence of protection graphs G\_0, … G\_n such that G\_0 |–\* G\_n using only *dejure* rules in G\_n there is an edge from **x** to **y** labeled *r.*

*Subject.Can.Share(α,* ***x, y,*** *G\_0)*

* X,y are subjects, *α right*
* And
* An edge x -*α*-> y in G\_0.
* Or
* Subject *s* in G\_0, s has x -*α*-> y
* Island\_1, … I\_n exist with x in I\_1, s in I\_n and a bridge from I\_*α to* I\_*α + 1*
  + Because there is a bridge between islands you can share rights between x and s because s has rights over y.

*Object.Can.Share(α,* ***x, y,*** *G\_0)*

Initially Span: x` = x where x` is a subject. Or x` is a subject and there is a tg path between x` and x.

Terminally Span: There is s` = subject and s` = s. Or there is a tg path between s` and s.

can•share(α, **x, y,** G\_0) if, and only if, there is an edge from **x** to **y** labeled α in G\_0, or the following hold simultaneously:

•There is an **s** in G\_0 with an s-to-y edge labeled α There is a subject **x` = x** or initially spans to **x**

•There is a subject **s`** = **s** or terminally spans to **s**

•There are islands I\_1,..., I\_k connected by bridges, and **x`** in I1 and **s`** in I\_k

## Chapter 4:

## Chapter 5:

## Chapter 6: