Progress Seminar

An automated triple modular redundancy EDA flow for Yosys

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Background

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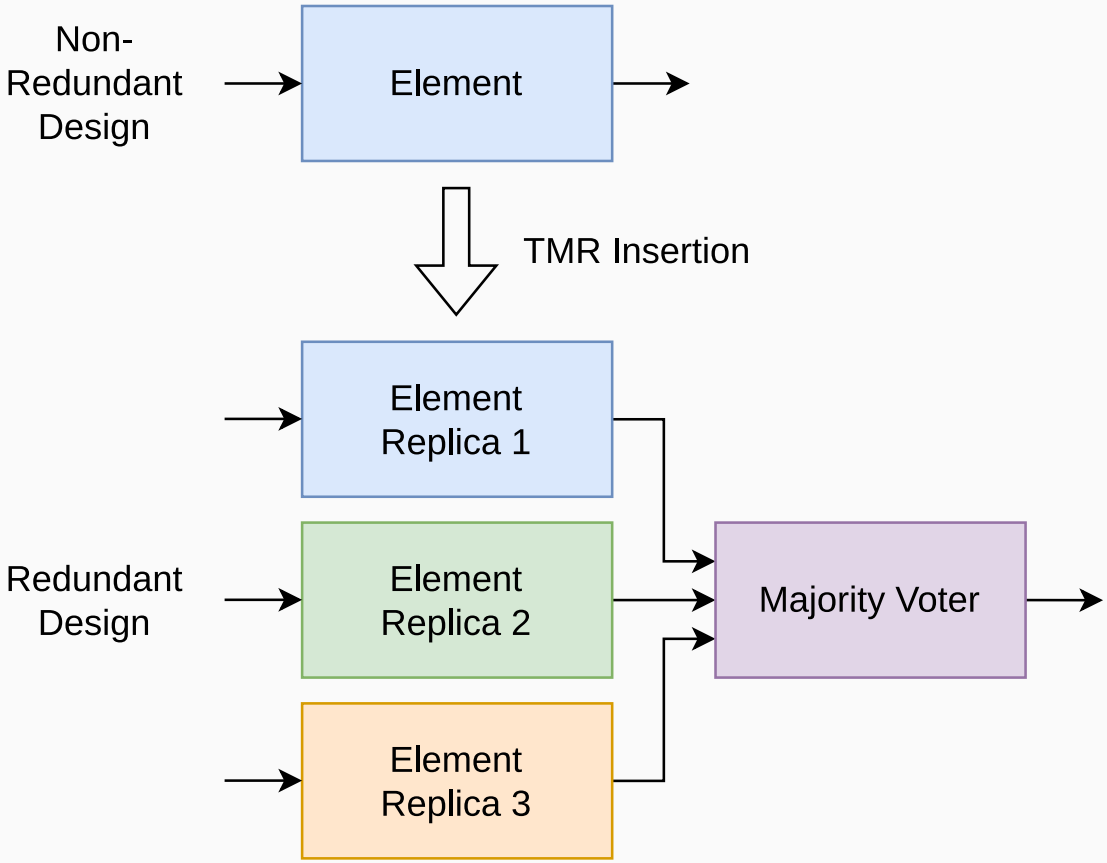
Application Specific Integrated Circuits (ASICs) and Field Programmable Gate Arrays (FPGAs) commonly deployed in space (and on Earth)... but protection from SEUs remains expensive!

RAD750 CPU [1] (James Webb Space Telescope, Curiosity rover, + many more) is commonly used, but costs **>$200,000 USD** [2]!

Triple Modular Redundancy



Triple Modular Redundancy

TMR can be added manually...

but this is **time consuming** and **error prone**.

Can we automate it?

TaMaRa

Concept

Implement TMR as a pass in an EDA synthesis tool.

- Integrated with the rest of the flow
- Easy to use
- Fully automated

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Yosys [3] is the best (and the only) open-source, research grade EDA synthesis tool.

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Yosys [3] is the best (and the only) open-source, research grade EDA synthesis tool.

- Proprietary vendor tools (Synopsys, Cadence, Xilinx, etc) immediately discarded
- Can’t be extended to add custom passes

Existing works

Two main paradigms:

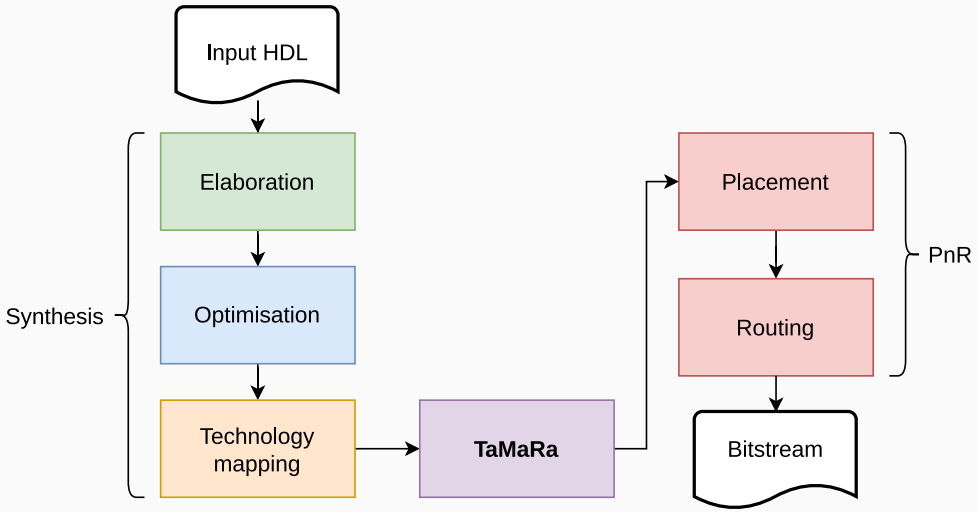
- **Design-level approaches** (“thinking in terms of HDL”)
  - Kulis [4], Lee [5]
- **Netlist-level approaches** (“thinking in terms of circuits”)
  - Johnson [6], Benites [7], Skouson [8]

The TaMaRa algorithm

TaMaRa is mainly netlist-driven. Voter insertion is inspired by Benites [7] “logic cones” concept, and parts of Johnson [6].

Also aim to propagate a `(* triplicate *)` HDL annotation to select TMR granularity (similar to Kulis [4]).

Runs after techmapping (i.e. after `abc` in Yosys)



The TaMaRa algorithm

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Why netlist driven with the `(* triplicate *)` annotation?

- Removes the possibility of Yosys optimisation eliminating redundant TMR logic
- Removes the necessity of complex blackboxing logic and trickery to bypass the normal design flow
- Cell type shouldn't matter, TaMaRa targets FPGAs and ASICs
- Still allows selecting TMR granularity - **best of both worlds**

Verification

Comprehensive verification procedure using formal methods, simulation and fuzzing.

Driven by SymbiYosys tools *eqy* and *mcy*

- In turn driven by Satisfiability Modulo Theorem (SMT) solvers (Yices [9], Boolector [10], etc)

Formal verification

Equivalence checking: Formally verify that the circuit is functionally equivalent before and after the TaMaRa pass.

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Also considering Beltrame's verification tool [11], and other literature on TMR formal verification.

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Idea:

1. Use Verismith [12] to generate random Verilog RTL.
2. Run TaMaRa synthesis end-to-end.
3. Use formal equivalence checking to verify the random circuits behave the same before/after TMR.

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Problem: Mutation

- We need valid testbenches for these random circuits, how would we generate that?
- Under active research in academia (may not be possible at the moment)

Simulation

We want to simulate an SEU environment.

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Use one of Verilator or Yosys' own cxxrtl to simulate a full design.

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- Currently considering picorv32 or Hazard3 RISC-V CPUs as the Device Under Test (DUT)

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Concept:

- Iterate over the netlist, randomly consider flipping a bit every cycle
- Write a self-checking testbench and ensure that the DUT responds correctly (e.g. RISC-V CoreMark)

Technical implementation

Implemented in C++20, using CMake.

Load into Yosys: `plugin -i libtamara.so`

TMR is implemented as two separate commands: `tamara_propagate` and `tamara_tmr`

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Run `tamara_tmr` after techmapping to perform triplication and voter insertion (add TMR).

Current status & future

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Algorithm design and planning essentially complete. Yosys internals (particularly RTLIL) understood to a satisfactory level (still learning as I go).

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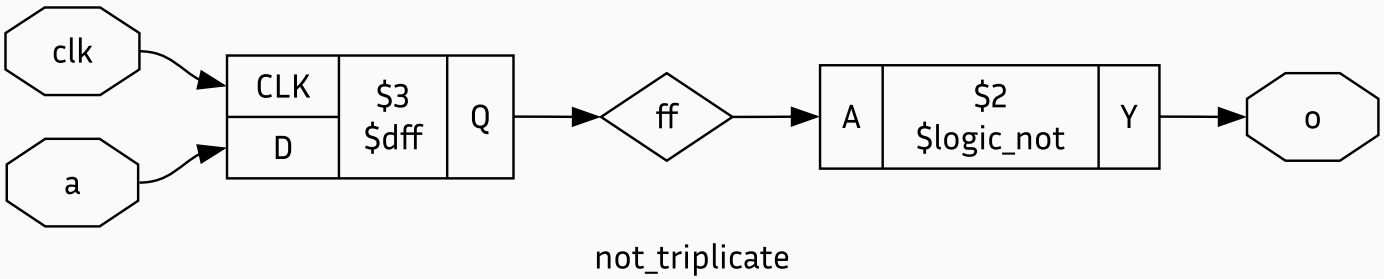
Programming hopefully finished *around* February 2025, verification by April 2025.

Progress: Automatically triplicating a NOT gate and inserting a voter

Original circuit:



not_triplicate

```
(* tamara_triplicate *)
module not_triplicate(
    input logic a,
    input logic clk,
    output logic o
);

logic ff;

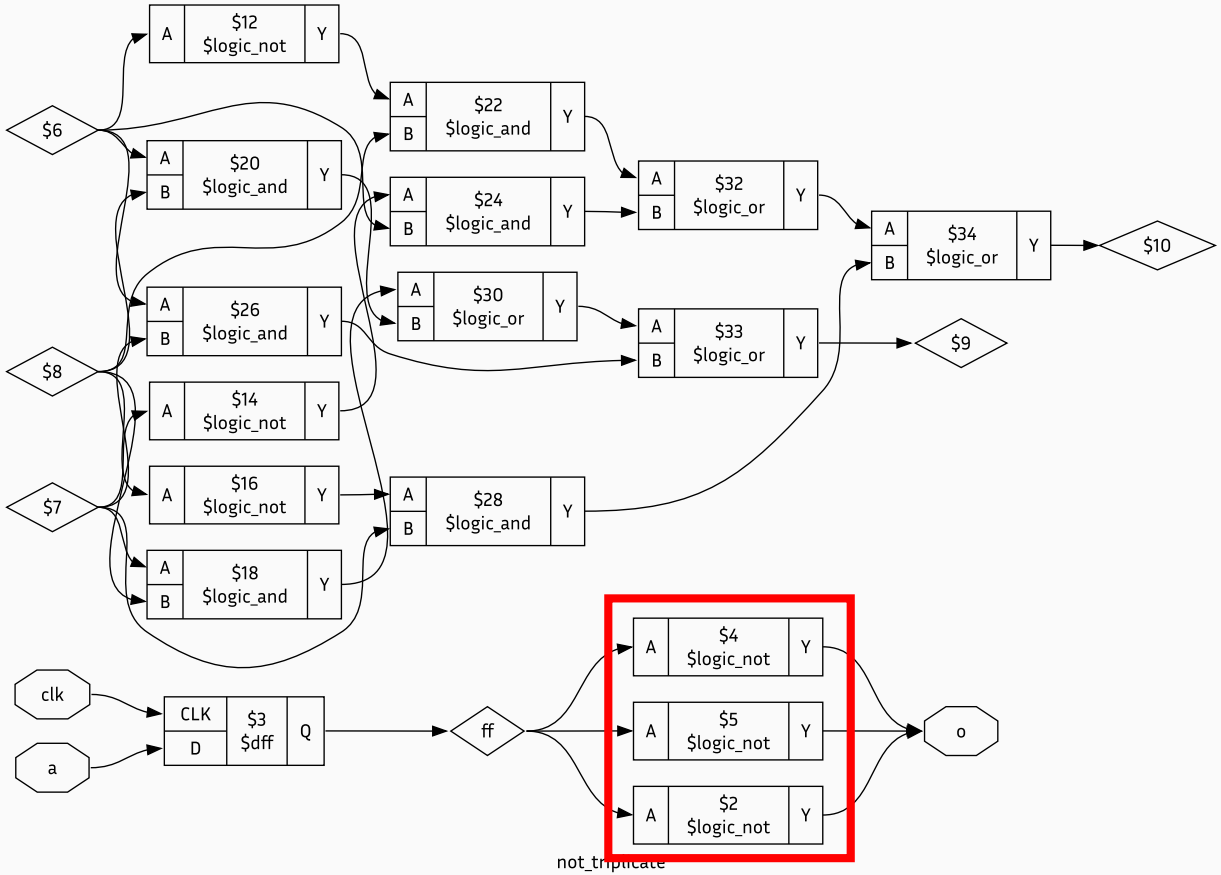
always_ff @(posedge clk) begin
    ff <= a;
end

assign o = !ff;

endmodule
```

Progress: Automatically triplicating a NOT gate and inserting a voter

After `tamara_debug replicateNot`:



Progress: Automatically triplicating a NOT gate and inserting a voter

Results:

- NOT circuit identified in `tamara::LogicGraph`
- RTLIL primitives replicated correctly
- Voter inserted using `tamara::VoterBuilder`
- Voter *not* yet wired up to main design
- Replicated components *not* yet re-wired

Progress: Equivalence checking

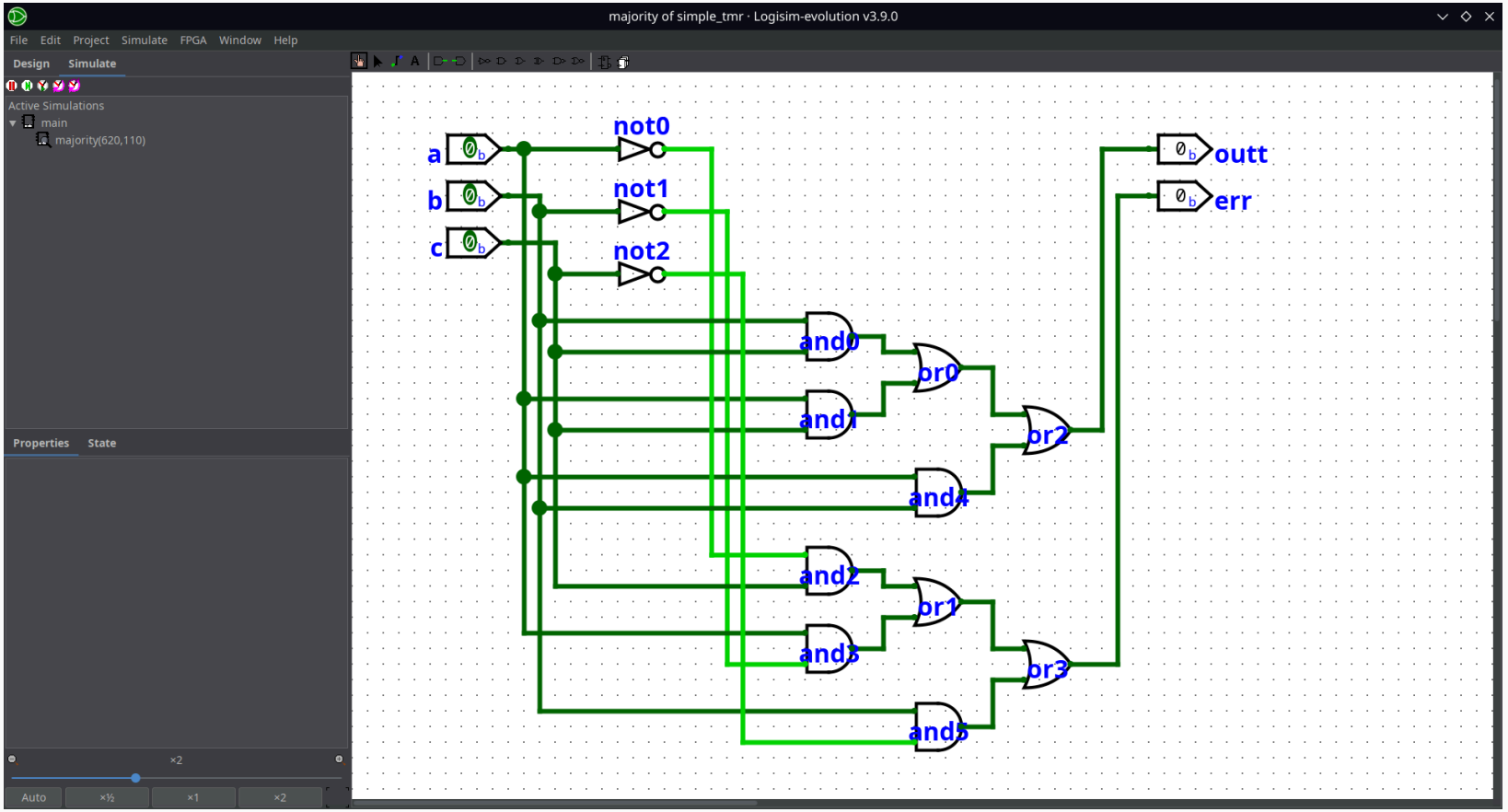
Voter circuit:

a	b	c	out	err
0	0	0	0	0
0	0	1	0	1
0	1	0	0	1
0	1	1	1	1
1	0	0	0	1
1	0	1	1	1
1	1	0	1	1
1	1	1	1	0

```
module voter(
    input logic a,
    input logic b,
    input logic c,
    output logic out,
    output logic err
);
    assign out = (a && b) || (b && c) || (a && c);
    assign err = (!a && c) || (a && !b) || (b && !c);
endmodule
```

Progress: Equivalence checking

Manual design in Logisim:



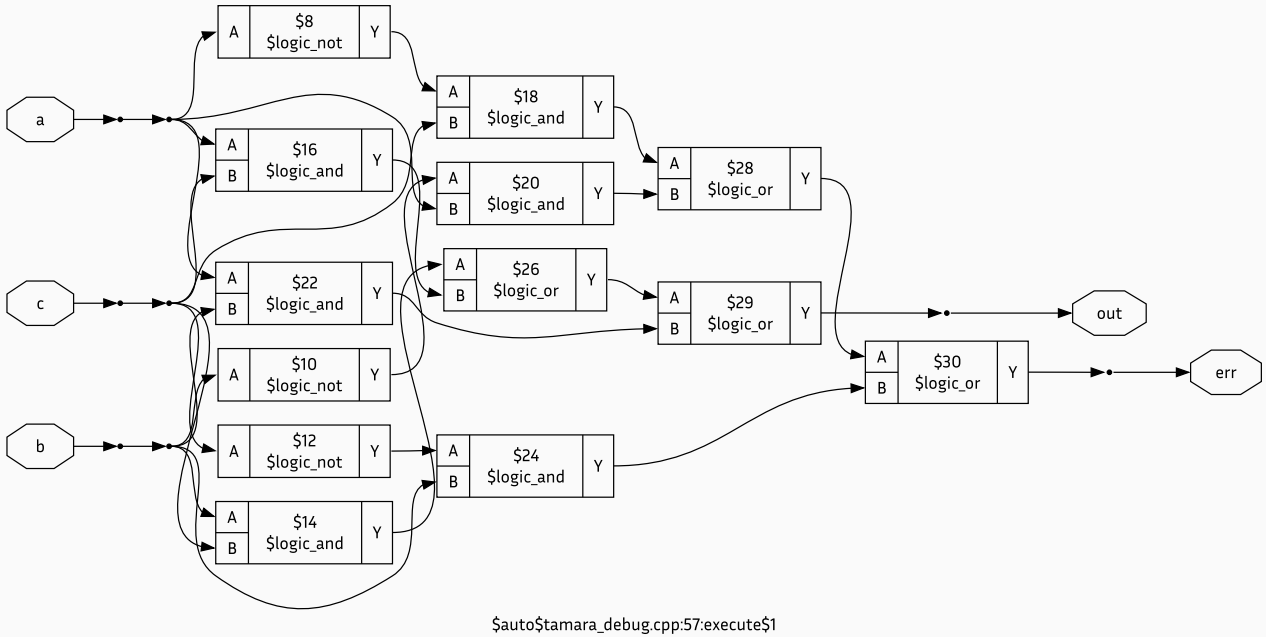
Progress: Equivalence checking

```
    Voter
tamara::VoterBuilder::build(RTLIL::Module
*module) {
        // NOT
        // a -> not0 -> and2
        WIRE(not0, and2);
        NOT(0, a, not0_and2_wire);
        ...

        // AND
        // b, c -> and0 -> or0
        WIRE(and0, or0);
        AND(0, b, c, and0_or0_wire);
        ...

        // OR
        // and0, and1 -> or0 -> or2
        WIRE(or0, or2);
        OR(0, and0_or0_wire,
and1_or0_wire, or0_or2_wire);
        ...

        return ...;
    }
```



$auto$tamara_debug.cpp:57:execute$1

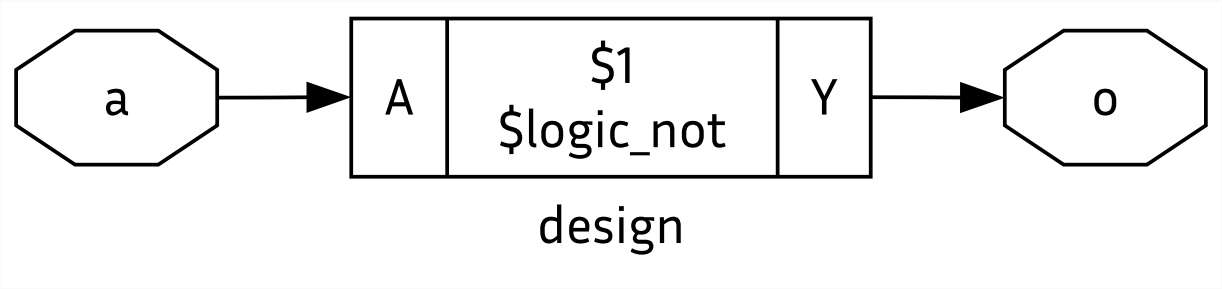
Progress: Equivalence checking

Marked equivalent by eqy in conjunction with Yices!

```
~/w/t/build (master) [n] ⤑ eqy -f ../tests/formal/equivalence/voter.eqy
EQY 22:47:32 [voter] read_gold: starting process "yosys -ql voter/gold.log voter/gold.ys"
EQY 22:47:32 [voter] read_gold: finished (returncode=0)
EQY 22:47:32 [voter] read_gate: starting process "yosys -ql voter/gate.log voter/gate.ys"
EQY 22:47:32 [voter] read_gate: finished (returncode=0)
EQY 22:47:32 [voter] combine: starting process "yosys -ql voter/combine.log voter/combine.ys"
EQY 22:47:32 [voter] combine: finished (returncode=0)
EQY 22:47:32 [voter] partition: starting process "cd voter; yosys -ql partition.log partition.ys"
EQY 22:47:32 [voter] partition: finished (returncode=0)
EQY 22:47:32 [voter] run: starting process "make -C voter -f strategies.mk"
EQY 22:47:32 [voter] run: make: Entering directory '/home/matt/workspace/tamara/build/voter'
EQY 22:47:32 [voter] run: Running strategy 'sby' on 'voter.err'..
EQY 22:47:32 [voter] run: Proved equivalence of partition 'voter.err' using strategy 'sby'
EQY 22:47:32 [voter] run: Running strategy 'sby' on 'voter.out'..
EQY 22:47:32 [voter] run: Proved equivalence of partition 'voter.out' using strategy 'sby'
EQY 22:47:32 [voter] run: make -f strategies.mk summary
EQY 22:47:32 [voter] run: make[1]: Entering directory '/home/matt/workspace/tamara/build/voter'
EQY 22:47:32 [voter] run: make[1]: Leaving directory '/home/matt/workspace/tamara/build/voter'
EQY 22:47:32 [voter] run: make: Leaving directory '/home/matt/workspace/tamara/build/voter'
EQY 22:47:32 [voter] run: finished (returncode=0)
EQY 22:47:32 [voter] Successfully proved equivalence of partition voter.out
EQY 22:47:32 [voter] Successfully proved equivalence of partition voter.err
EQY 22:47:32 [voter] Successfully proved designs equivalent
EQY 22:47:33 [voter] summary: Elapsed clock time [H:MM:SS (secs)]: 0:00:00 (0)
EQY 22:47:33 [voter] summary: Elapsed process time [H:MM:SS (secs)]: 0:00:00 (0)
EQY 22:47:33 [voter] DONE (PASS, rc=0)
```

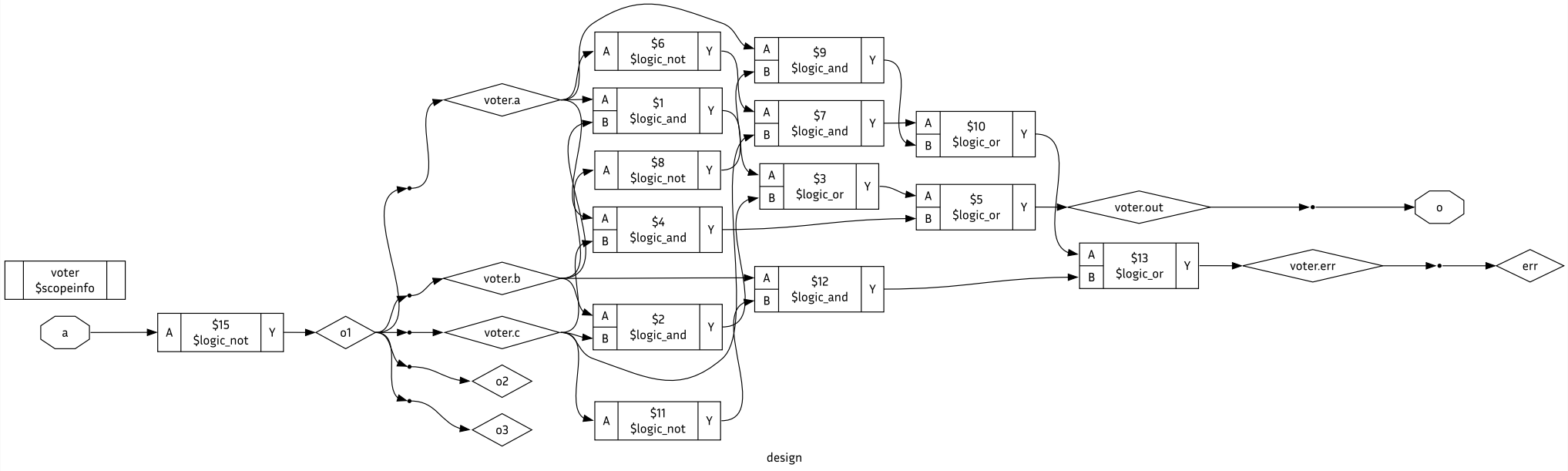
Progress: Equivalence checking (Voter insertion)

Original, very simple circuit:



Progress: Equivalence checking (Voter insertion)

After manual voter insertion (using SystemVerilog):



Progress: Equivalence checking (Voter insertion)

Are they equivalent? Yes! (Thankfully)

```
~/w/t/build (master) [n] ⤫> eqy -f ../tests/formal/equivalence/not_voter.eqy
EQY 22:10:20 [not_voter] read_gold: starting process "yosys -ql not_voter/gold.log not_voter/gold.ys"
EQY 22:10:20 [not_voter] read_gold: finished (returncode=0)
EQY 22:10:20 [not_voter] read_gate: starting process "yosys -ql not_voter/gate.log not_voter/gate.ys"
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EQY 22:10:20 [not_voter] combine: starting process "yosys -ql not_voter/combine.log not_voter/combine.ys"
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EQY 22:10:20 [not_voter] run: make: Entering directory '/home/matt/workspace/tamara/build/not_voter'
EQY 22:10:20 [not_voter] run: Running strategy 'sby' on 'design.o'..
EQY 22:10:20 [not_voter] run: Proved equivalence of partition 'design.o' using strategy 'sby'
EQY 22:10:20 [not_voter] run: make -f strategies.mk summary
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Caveat: Still need to verify circuits with more complex logic (i.e. DFFs).

The future

TODO tasks that remain needing to be done

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I have also spoken with the team at YosysHQ GmbH and Sandia National Laboratories, who are very interested in the results of this project and its applications.

Conclusion

Summary

- TaMaRa: Automated triple modular redundancy EDA flow for Yosys
- Fully integrated into Yosys suite
- Takes any circuit, helps to prevent it from experiencing SEUs by adding TMR
- Netlist-driven algorithm based on Johnson's work [6] (TODO NOT TRUE)
- **Key goal:** "Click a button" and have any circuit run in space/in high reliability environments!

I'd like to extend my gratitude to N. Engelhardt of YosysHQ, the team at Sandia National Laboratories, and my supervisor Assoc. Prof. John Williams for their support and interest during this thesis so far.

References

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Thank you! Any questions?