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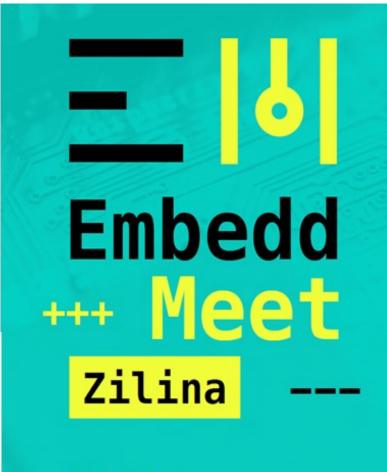
Za spoluprácu na dnešnom Meetupe ďakujeme: Matúš Chochlík (Speaker)

Za organizáciu a propagáciu eventu ďakujeme:





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# External process synchronization for fun and profit

November 19, 2019

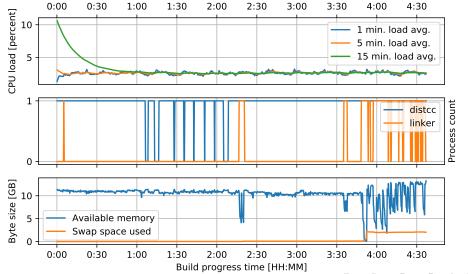
## The LLVM project

- "a collection of modular and reusable compiler and toolchain technologies"
  - 11vm code generator and optimizer for many CPUs
  - 11db debugger
  - 11d linker
  - clang C/C++/Objective-C compiler
  - clang-tidy C/C++/Objective-C static analysis tool
  - clang-format C/C++/Objective-C code format tool
  - ..
- https://llvm.org/
- https://github.com/llvm/llvm-project

## Let's build llvm and clang

- \$ git clone https://github.com/llvm/llvm-project.git
- \$ mkdir \_build
- \$ cd \_build
- \$ cmake -DCMAKE\_BUILD\_TYPE=Debug ...
- \$ make

#### make - system resource usage



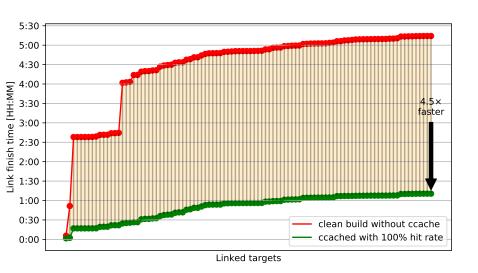
## That didn't go so well...

- System resource usage is low
- Build takes cca. 5 hours unacceptably long
- Let's try some tricks:
  - use ccache
  - use more make jobs
  - use distcc for distributed compilation

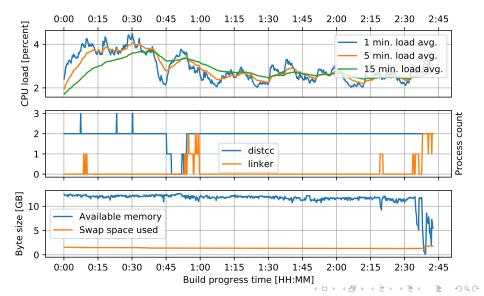
- "speeds up recompilation by caching previous compilations and detecting when the same compilation is being done again"
- https://ccache.dev/

- "distributes compilation of C or C++ code across several machines on a network"
- https://github.com/distcc/distcc

#### make - uncached clean build vs. 100% cached



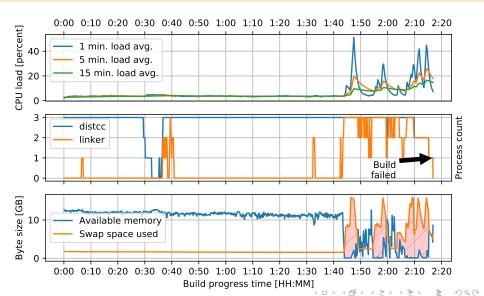
## More jobs - make clean && make -j 2



#### That's somewhat better...

- System resource usage is still low
- Build takes more than 2 and half hours still too long
- Parallelization shows potential

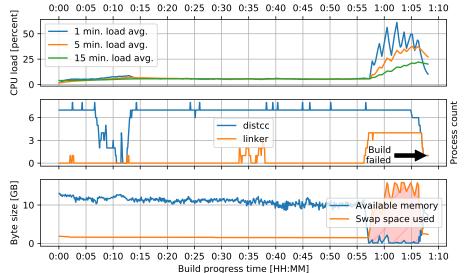
## Keep going - make clean && make -j 3



## What happened?

- System usage starts to climb
  - especially memory usage
  - signs of correlation with linker execution
- Build failed after 2 hours 20 minutes looooong
  - Linux OOM¹ process killer, kills some linker jobs
- Parallelization still shows potential

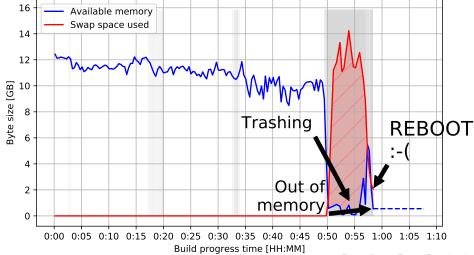
## Meh. MOAR jobs! make clean && make -j 7



#### Not great, not terrible

- System usage very high toward the end of the build
  - especially memory usage
  - still correlated with linker execution
- Build failed after 1 hours 5 minutes getting better
  - OOM process killer, kills some linker jobs
- Parallelization rules!

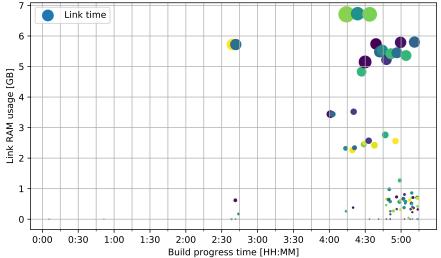
#### Still come cores left... make clean && make -j 9



#### What have we learned so far

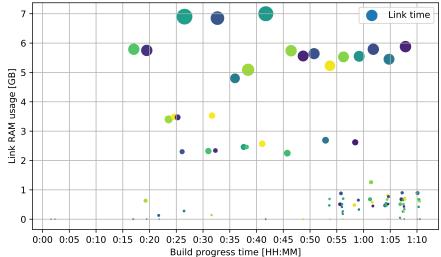
- Most of the build time is spent by compilation
- Compilation caching good.
- Compilation paralelization good.
  - We have not even really used distcc yet.
- Linking paralelization bad! Why?
  - Let's have a closer look!

## Linker memory usage vs link time – no cache

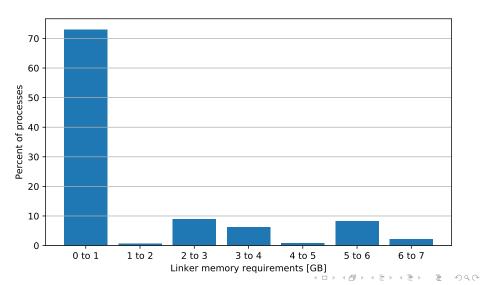




#### Linker memory usage vs link time - cached



#### Link target memory requirements – classification



## Linking llvm and clang in a nutshell

- Around 100 different link targets
  - Executables, Shared libraries
  - Many of them are small.
  - Few of the huge.
- Most of the linking is done towards the end of the build process.
  - ⇒ Using many parellel jobs, multiple instances of linker run concurrently.
    - ⇒ Many big targets are linked at the same time.

## ⇒ OOM!

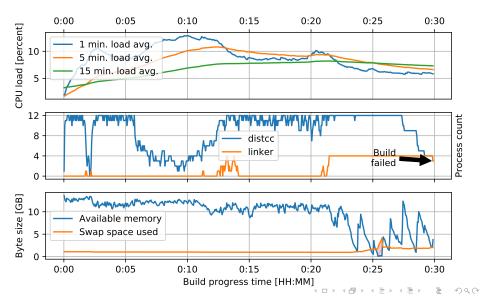
• BTW: so much swapping is bad for the SSDs.

## Making parallel linking work

- Let's try some additional tricks:
  - Use GNU gold instead of GNU ld.
    - Can link only in ELF format.
    - + Faster than 1d.
    - + Uses less memory than <a>1d</a>.
  - Use zram.
    - a Linux kernel module for creating a compressed block device in RAM
    - can be used as swap device part of swap is in compressed RAM
  - Make only some targets
    - Practical when developing and debugging.
    - Sometimes you just need to do make all.



## zram + distcc + ccache + gold + make - j 12



## If only...

#### • ... we could

- prevent so many big targets from being linked at once,
- prevent excessive swapping to disk,
- but still
  - use the hardware resources efficiently,
  - use available remote cores via distcc to full extent.

## Let's synchronize the execution of ld/gold

- atmost to the rescue!
  - Simple utility written in C.
  - Takes a command-line (executable + arguments).
  - Waits until specified conditions are met.
  - Executes the command-line.
  - https://github.com/matus-chochlik/various/atmost

#### atmost - description

"The atmost utility allows to limit concurrent execution of a single executable or a set of executables, depending on various criteria like CPU load, memory and swap usage, thermal zone temperatures, current number of I/O operations, etc., or just by a maximum number."

#### RTFM - man atmost

ATMOST(1) General Commands Manual

ATMOST(1)

NAME

atmost - tool for limiting the concurrent execution of a specified executable.

#### SYNOPSIS

```
atmost [-v|--verbose] [-f|--file file-path]

[-s|--socket socket-path] [-r|--reset]

[-i|--sleep-interval seconds] [-c|--print-current]

[-C|--print-all-current] [-n|--max-instances count]

[-l|--max-cpu-load-Im percent] [-m|--min-avail-ram percent]

[-M|--min-free-ram percent] [-S|--min-free-swap percent]

[-p|--max-total-procs count] [-tc|--max-cpu-temp temp]

[-tg|--max-gpu-temp temp] [-tb|--max-bat-temp temp]

[-io|--max-io-ops count] [-nw|--max-nw-speed speed]

[-snw|--slow-nw-speed speed] [-- executable [args...]]
```

atmost --help

#### atmost -n NUMBER -- EXECUTABLE - basic usage

- Limits the concurrent execution of **EXECUTABLE** to the specified **NUMBER** of instances.
- Uses an IPC semaphore set<sup>2</sup>.
  - Acquires<sup>3</sup> the semaphore before executing EXECUTABLE.
  - Releases the semaphore after executing EXECUTABLE.
  - realpath <sup>4</sup> of EXECUTABLE is used as the semaphore set token<sup>5</sup>.



<sup>&</sup>lt;sup>2</sup>see man 2 semget

<sup>&</sup>lt;sup>3</sup>see man 2 semop

<sup>4</sup>see man 3 realpath

<sup>&</sup>lt;sup>5</sup>see man 3 ftok

## Using atmost - wrapper scripts

- In order to synchronize an **EXECUTABLE** with **atmost**:
  - Create a shell script with the name of **EXECUTABLE**.
  - From the script call **atmost** with appropriate arguments.
  - Put the wrapper script to \$\{PATH\}

```
$ vim /opt/bin/ld
```

```
#!/bin/bash
atmost -n 2 -- /usr/bin/gold "${@}"
```

- \$ chmod u+x /opt/bin/ld
- \$ export PATH="/opt/bin:\${PATH}"

#### atmost -n 2 -- ld

- + Empirically we have found that 2 instances of ld can run safely with 16GB RAM.
- + Simple
- + Low overhead
  - Too coarse and restrictive
    - The majority of the <a href="llvm">11vm</a> targets has low RAM requirements for linking.
    - Targets from other projects also have low linker RAM requirements.
    - We could safely use more parallel link jobs.

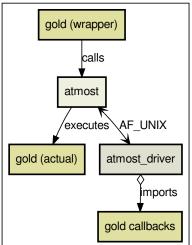
#### What if...

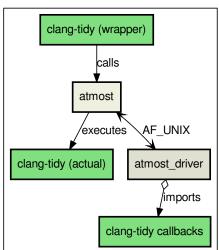
- ... we could determine if it is safe to start the linker per each invocation?
  - Enter atmost\_driver!
    - Running as a local AF\_UNIX socket server.
    - atmost is the client (with the --socket PATH option).
    - When atmost is started, it connects to the driver and sends information about the wrapped EXECUTABLE.
    - When the *driver* determines that it's safe to run **EXECUTABLE**, it responds to **atmost**.
    - Then atmost executes EXECUTABLE.

#### Let's make it even more flexible

- atmost\_driver callbacks:
  - The <a href="mailto:atmost\_driver">atmost\_driver</a> implements the reusable common server functionality.
  - The specific logic determining when it is safe to start executing, is implemented in a set of callbacks.
  - load\_callback\_data
  - save\_callback\_data
  - process\_initialized
  - let\_process\_go
  - process\_finished

#### atmost + atmost\_driver +callbacks





#### Callback - load\_callback\_data

```
def load_callback_data():
    resources = do_load_resources(...)
    return resources
```

- Called once, when the <a href="atmost\_driver">atmost\_driver</a> is being started.
- Typically does startup initialization.
- Can load resources used by the other callbacks.

#### Callback - save\_callback\_data

```
def save_callback_data(resources):
   do_cleanup_resources(resources)
```

- Called once, when the <a href="atmost\_driver">atmost\_driver</a> is being shutdown.
- The resource parameter is the return value from load\_callback\_data.
- Can cleanup and/or save resources used and data generated during the run.

## Callback - process\_initialized

```
def process_initialized(resources, proc):
    # handle new process
    proc.set_callback_data(callback_data)
```

- Called when <a href="atmost">atmost</a> sends the driver information about a new process.
- The **proc** parameter provides information about the new process.
  - PID, command-line, environment, working-directory, etc.
  - The callback can store its bookkeeping data in proc.

# Callback - let\_process\_go

```
def let_process_go(resources, procs):
   if can_process_start(resource, procs):
     return True
   return False
```

- Called repeatedly.
- Determines if a process can start executing the synchronized executable.
- The **procs** argument contains info about *all* currently managed processes, split into 3 groups.
  - active processes that are already executing.
  - waiting processes that are waiting for execution.
  - current the process to be let go.

# Callback - process\_finished

```
def process_finished(resources, proc):
   callback_data = proc.callback_data()
# cleanup info about the process
```

- Called when a process handled by the driver has finished
- The callback can retrieve its bookkeeping data from proc.

#### RTFM - man atmost\_driver

```
ATMOST(1)

NAME

atmost_driver - driver server for the atmost command.

SYNOPSIS

atmost_driver [-s|--socket socket-path] [-c|--callbacks file-path]

[-u|--update interval]

atmost_driver -h
atmost_driver --help
```

# Back to linking...

- What if the memory usage can be determined from the linker command-line arguments?
  - 1) Run a lot of builds of various projects.
  - 2) Gather a lot of data
    - command line arguments,
    - input file sizes,
    - actual linker memory usage,
    - etc.
  - 3) ???
  - 4) Profit!

# What exactly is step 3) ???

- Feed the gathered data into a machine learning model and train it.
- Use the trained ML model to predict linker memory usage from command-line arguments.
- Integrate the ML classifier into a atmost\_driver callback script.
  - Keep track of available memory.
  - Predict the memory usage of new instances of linker.
  - Let the new linkers go only if there is enough available memory.

## The ML part

• Uses the neural\_network.MLPClassifier
from the sklearn
6 Python package.

#### • Inputs:

- optimization level, the PIE<sup>7</sup> flag,
- count and combined size of input shared and static object files,
- etc.

#### Output:

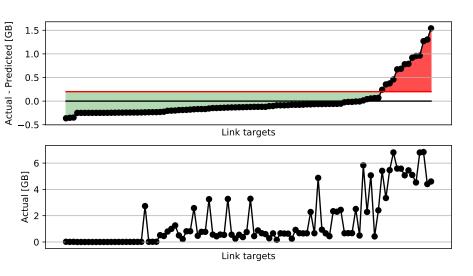
- the predicted memory requirements, as a multiple of 256MB chunks,
- i.e. if the output is 5 the predicted size is 5\*256[MB] = 1.25[GB].



<sup>&</sup>lt;sup>6</sup>https://scikit-learn.org/stable/documentation.html

<sup>&</sup>lt;sup>7</sup>position-independent executable

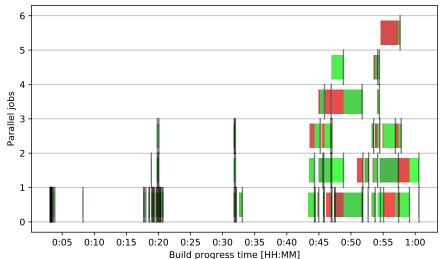
# Classifier prediction accuracy



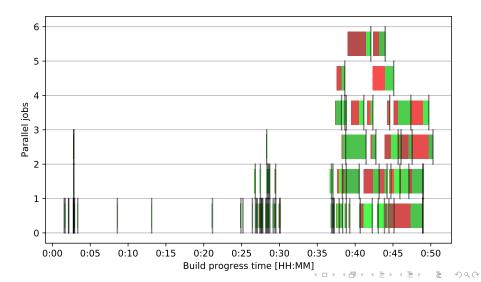
#### Integrating classifier with atmost\_driver

- load\_callback\_data load the trained ML model.
- process\_initialized use the model to predict memory usage.
- let\_process\_go only if predicted linker memory usage is less than current available memory.
- process\_finished output actual memory usage that can be stored and used for additional model training.

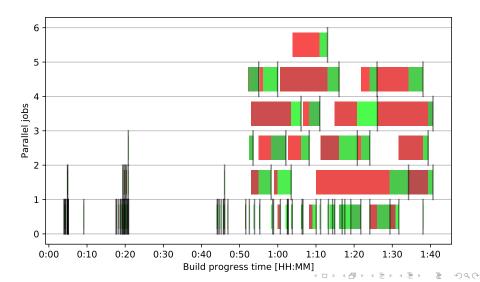
## Linker schedule (no ccache, 16GB RAM) - make -j 10



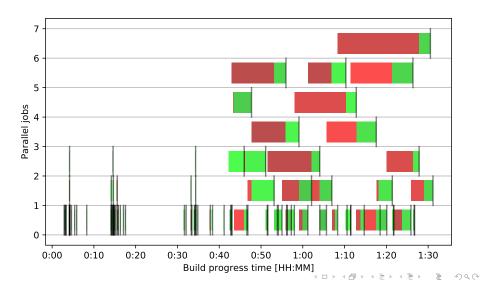
# Linker schedule (no ccache, 16GB RAM) - make -j 20



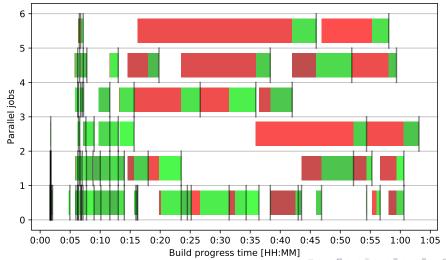
#### Linker schedule (no ccache, 8GB RAM) - make -j 10



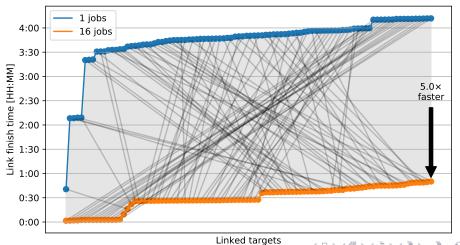
## Linker schedule (no ccache, 8GB RAM) - make -j 20



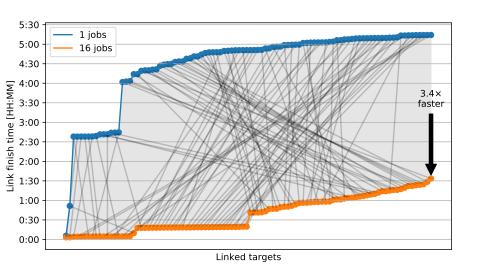
# Linker schedule (ccached, 8GB RAM) - make -j 20



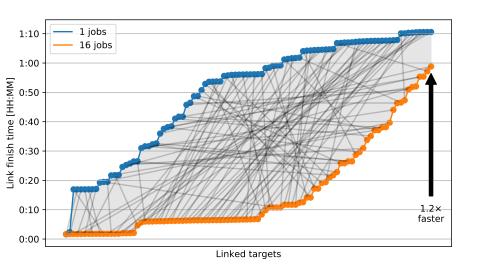
# Link order vs. time (no ccache, 16GB, 4+12 cores, 16 jobs)



# Link order vs. time (no ccache, 8GB, 8+8 cores, 16 jobs)



# Link order vs. time (ccached, 8GB, 8+8 cores, 16 jobs)



#### Parallelization statistic indicators

- Run time with j jobs
  - T<sub>j</sub>
- Speedup with j jobs

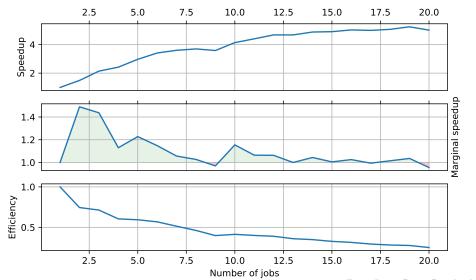
$$\bullet S_j = T_1/T_j$$

Marginal speedup with j jobs

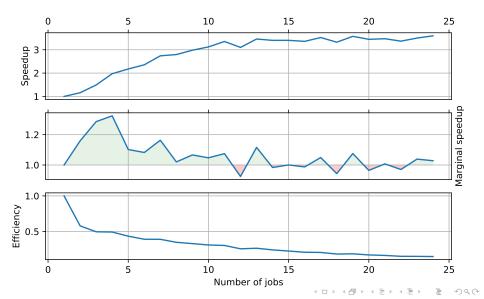
$$\bullet M_j = T_{j-1}/T_j$$

- Efficiency with j jobs
  - $E_i = S_i/j$

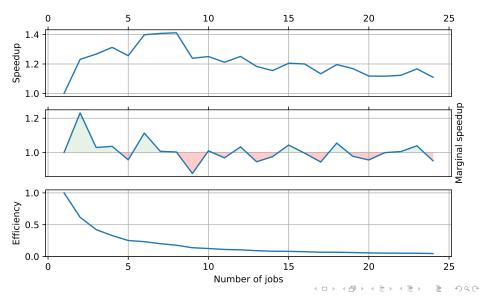
# Parallelization statistics (no ccache, 16GB, 4+12 cores)



# Parallelization statistics (no ccache, 8GB, 8+8 cores)



# Parallelization statistics (ccached, 8GB, 8+8 cores)



#### Conclusions

- atmost is a highly flexible tool for synchronizing execution of parallel processes,
  - improves build times by allowing many parallel make jobs while serializing huge link jobs,
  - prevents system stalls, crashing and reboots due to low memory conditions,
  - makes building big projects safer,
  - allows to build <u>llvm</u> on systems with low RAM capacity,
  - reduces the amount of writes to swap disk partitions.

# the fun part

- Learned new things:
  - new Python 3 features
  - Python *SciKit-Learn* package
  - Python MatPlotLib package
  - Secure distcc setup over ssh.
  - etc.

# the *profit* part

- Stable and relatively quick llvm/clang builds.
- Reasonable HW usage.
- Saving the SSDs from excessive wear.
- Not getting random reboots during development and building.

# the *moral* of the story

# Go get more RAM!<sup>8</sup>

## Further improvements

- Try to improve the prediction of memory usage based on command-line arguments.
  - Examine the arguments more deeply.
  - Use a different classifier.
  - Develop a custom model.
  - Better fitting of waiting linker processes into memory.
  - etc.
- Allow to change environment variables for the synchronized process from the <u>atmost\_driver</u> callbacks.
- . . .

https://github.com/matus-chochlik/ various/atmost/presentation/embedmeet.pdf

# Extras

#### atmost - additional options

- Besides the basic semaphore-based synchronization and the customizable driver-based synchronization, atmost supports synchronization depending on:
  - 1 and 5 minutes average load,
  - amount of available RAM,
  - amount of free swap,
  - CPU temperature,
  - network speed,
  - battery status,
  - etc.

## atmost - average load

- -l or --max-cpu-load-1m
- -L or --max-cpu-load-5m

Start executing only if 1 minute average load is less than or equal 15%:

```
atmost -1 15 -- executable arguments...
```

Start executing only if 5 minute average load is less than or equal 8.5%:

```
atmost -L 8.5 -- executable arguments...
```

#### atmost - memory / swap usage

- -m or --min-avail-ram
- -M or --min-free-ram
- -S or --min-free-swap

Start executing only if at least 20% or RAM is available:

```
atmost -m 20 -- executable arguments...
```

Start executing only if at least 70% or swap is free:

```
atmost -S 70 -- executable arguments...
```

#### atmost - total process count

• -p or --max-total-procs

Start executing only if the number of currently running processes is  $\leq 1000$ :

atmost -p 1000 -- executable arguments...

#### atmost - thermal zone temperatures

- -tc or --max-cpu-temp
- -tg or --max-gpu-temp
- -tb or --max-bat-temp

Start executing only if the CPU temperature is less than or equal to  $70^{\circ}$  C:

```
atmost -tc 70 -- executable arguments...
```

Start executing only if the battery temperature is less than or equal to  $55.5^{\circ}C$ :

```
atmost -tb 55.5 -- executable arguments\
...
```

#### atmost - modifiers

- Modifiers allow to change the limits by a specified amount under special conditions.
  - when running on battery,
  - when connected to "slow" network,
  - when disconnected from network.

#### atmost - on A/C vs. on battery

-batt or --on-battery

When on A/C, start only if 1 minute average load is less than or equal to 15%,

when on battery, start only if average load is less than or equal to 5% = (15 - 10):

```
atmost -1 15 -batt 10 -- executable \ arguments...
```

#### atmost - combining limits and modifiers

When on A/C, start only if 1 minute load is  $\leq$  20% and 5 minutes load is  $\leq$  10%, when on battery, start only if 1 minute load is  $\leq$  8% and 5 minutes load is  $\leq$  4%:

```
atmost -1 20 -batt 12 -L 10 -batt 6 -- \
executable ...
```

# Limiting concurrent instances of clang-tidy

- Is it useful to limit concurrent execution of clang-tidy?
  - 8 core CPU.
  - 2 local cores for compiling,
  - 4 local cores for compiling,
  - 6 local cores for compiling,
  - the rest used by clang-tidy.
- Measured on the test-suite of the OGLplus<sup>9</sup> project.



# atmost -n $\{N\}$ clang-tidy " $\{Q\}$ "

