

## Exercise 4 - Binary Exploitation

### General Hints

- The exploitable source file `exploitable.c` can be downloaded from Ilias here: [https://elearning.hs-albsig.de/goto.php?target=file\\_395340\\_download&client\\_id=HSALBSIG](https://elearning.hs-albsig.de/goto.php?target=file_395340_download&client_id=HSALBSIG).
- You have to deactivate most countermeasures for the exploits to work in a simple way as shown in the lecture slides.
- Furthermore, it is easier to attack a 32 bit target (x86 architecture) instead of attacking a 64 bit target (x86\_64 architecture). You can use either directly the Metasploitable VM or compile your source code using the `-m32` option of gcc.
- To be able to debug and inspect the code in `gdb` and using `objdump`, you have to compile the source code using the `-g` option to include debugging symbols in the executable.
- If you are unable to figure out the correct addresses to jump to quickly, you should try to add a simple NOP sled.

### Tools

If not already familiar with the tools, you need to get familiar with the following tools:

- From the GNU toolchain `gcc`, `gdb`, and `objdump`.
- As assembler you can use either `nasm` or the GNU assembler `gas`.



### Exercise 4.1 - A Basic Stack Overflow Attack

Exploit the error with `gets` in `exploitable.c` to jump to the function `void hidden(void)`.

### Hints

- The input to the executable can be supplied using the standard input `stdin`, e.g., using a pipe. The easiest way is to prepare your payload outside the debugger (e.g., as `exploit.bin`) and to use `r < exploit.bin` as command to run the exploit.
- Follow the approach presented on the lecture slides and adapt it to use `stdin` instead of command line arguments.
- You have to figure out the target address of the hidden function and overwrite the return address on the stack to jump to that function.

## Success Condition

You are successful, if the `printf()` called by the hidden function is invoked.



### Exercise 4.2 - Spawning A Shell With Shellcode

Exploit the error with `gets` in `exploitable.c` to spawn a shell (`/bin/sh`) using your shellcode and the system call `execve()`.

## Hints

Calling a system call on the x86 architecture works as follows:

- The `execve()` documentation is available at <https://man7.org/linux/man-pages/man2/execve.2.html>.
- The system call number is put into register `eax`.
- The official system call numbers for the x86 architecture can be found here: [https://github.com/torvalds/linux/blob/master/arch/x86/entry/syscalls/syscall\\_32.tbl](https://github.com/torvalds/linux/blob/master/arch/x86/entry/syscalls/syscall_32.tbl). Note, that the system call numbers are different for other architectures.
- The first parameter is put into `ebx`, second parameter in `ecx`, third parameter in `edx`, fourth parameter in `esi`, fifth parameter in `edi` and the sixth parameter in `ebp`. The result is returned in `eax`.
- If the all registers are filled, the system call is executed with `int $0x80`
- Details how to call a system call can be found at [https://en.wikibooks.org/wiki/X86\\_Assembly/Interfacing\\_with\\_Linux](https://en.wikibooks.org/wiki/X86_Assembly/Interfacing_with_Linux).

The assembly code must be turned into machine code using an assembler such as `nasm` or `gas`. Using `objdump` it is possible to access the machine code which is necessary to craft the final shellcode.

## Success Condition

You are successful, if a shell is started. In your report you should proof your success by executing some commands in the shell. Furthermore, you should test, if your exploit also works with activates ASLR.



### Exercise 4.3 - Spawning A Shell With `ret2libc`

Exploit the error with `gets` in `exploitable.c` to spawn a shell (`/bin/sh`) using the glibc function call `system()`.

Answer the following question: What is the major difference and benefit for the attacker between Exercise 4.2 and 4.3?

## Hints

Calling a library function works differently from executing a system call.

- The documentation for `system()` can be accessed here: <https://man7.org/linux/man-pages/man3/system.3.html>
- The function `system()` takes exactly one argument, which must be put on the stack.
- You need to either find the string `/bin/sh` somewhere in the memory space of the executable (e.g., in an environment variable), or you have to put the string on the stack yourselves.
- Finding the string `/bin/sh` works with

## Success Condition

You are successful, if a shell is started. In your report you should proof your success by executing some commands in the shell. Furthermore, you should test, if your exploit also works with activated ASLR.



### Exercise 4.4 - Spawning Meterpreter

Exploit the error with `gets` in `exploitable.c` to spawn a bind shell generated using `msfvenom` from Metasploit.

## Hints

The general approach is very similar to the shellcode injection in Exercise 4.2. However, instead of executing a system call, the payload generated with `msfvenom` must be injected and executed. You can then connect using netcat to the TCP port.

There are many ways to create a payload using `msfvenom`. This is the recommended way for this exercise:

- Use `msfvenom -help` and Internet searches to find out more about the payload generation process. It is also useful to use the `-l` option of `msfvenom` to find out more about the available payloads, encoders, etc.
- You should use a `shell_bind_tcp` payload appropriate for your platform (linux) and architecture (x86). This specific payload will open a simple TCP port, which you can connect to using netcat.
- Specify a port using, e.g., by using `LPORT=8888` as option to `msfvenom` to select the TCP port the bind shell will listen on.
- Connecting using netcat works with `nc TARGET_IP TARGET_PORT`. If successful, all commands you type (and send using enter) are executed on the target host. As proof you should list a directory such as `/bin`.

**Success Condition**

You have successfully connected to the remote shell using netcat at least once and executed some commands on the remote shell.