

Problem Sheet 7.2

Exercise 1

a) $\lambda x^o. x$

The type is $o \rightarrow o$ and the derivation is:

$$\frac{x: o \vdash x: o}{\vdash \lambda x^o. x: o \rightarrow o}$$

b) $\lambda x^o. \lambda y^o. \lambda z^{o \rightarrow o}. zx$

The type is $o \rightarrow o \rightarrow (o \rightarrow o) \rightarrow o$. Considering $\Gamma = x: o, y: o, z: o \rightarrow o$, the typing derivation is:

$$\frac{\frac{\frac{\frac{\Gamma \vdash z: o \rightarrow o}{\Gamma \vdash zx: o} \quad \Gamma \vdash x: o}{\Gamma \vdash zx: o}}{x: o, y: o, z: o \rightarrow o \vdash zx: o}}{x: o, y: o \vdash \lambda z^{o \rightarrow o}. zx: (o \rightarrow o) \rightarrow o}}{x: o \vdash \lambda x^o. \lambda y^o. \lambda z^{o \rightarrow o}. zx: o \rightarrow (o \rightarrow o) \rightarrow o}}{\vdash \lambda x^o. \lambda y^o. \lambda z^{o \rightarrow o}. zx: o \rightarrow o \rightarrow (o \rightarrow o) \rightarrow o}$$

c) $\lambda f^{(o \rightarrow o) \rightarrow o}. \lambda g^{o \rightarrow o \rightarrow o}. f(\lambda x^o. gxx)$

The type is $((o \rightarrow o) \rightarrow o) \rightarrow (o \rightarrow o \rightarrow o) \rightarrow o$. Considering $\Gamma = f: (o \rightarrow o) \rightarrow o, g: o \rightarrow o \rightarrow o$, the typing derivation is:

$$\frac{\frac{\frac{\Gamma \vdash f: (o \rightarrow o) \rightarrow o}{\Gamma \vdash f(\lambda x^o. gxx): o} \quad \frac{\frac{\frac{\Gamma, x: o \vdash g: o \rightarrow o \rightarrow o}{\Gamma, x: o \vdash gx: o \rightarrow o} \quad \Gamma, x: o \vdash x: o}{\Gamma \vdash \lambda x^o. gxx: o \rightarrow o}}{\Gamma \vdash f(\lambda x^o. gxx): o}}{f: ((o \rightarrow o) \rightarrow o), g: (o \rightarrow o \rightarrow o) \vdash f(\lambda x^o. gxx): o}}{f: ((o \rightarrow o) \rightarrow o) \vdash \lambda g^{o \rightarrow o \rightarrow o}. f(\lambda x^o. gxx): (o \rightarrow o \rightarrow o) \rightarrow o}}{\vdash \lambda f^{(o \rightarrow o) \rightarrow o}. \lambda g^{o \rightarrow o \rightarrow o}. f(\lambda x^o. gxx): ((o \rightarrow o) \rightarrow o) \rightarrow (o \rightarrow o \rightarrow o) \rightarrow o}$$

Exercise 2

The term $(\lambda x. x)(\lambda y. y)$ can be typed. According to the derivation rule for applications:

$$\frac{\Gamma \vdash M: \tau \rightarrow \sigma \quad \Gamma \vdash N: \tau}{\Gamma \vdash MN: \sigma}$$

The argument of M must have the same type as N. Therefore, if M is of type $\tau \rightarrow \tau$ the type of N must be τ . Expressions like $\lambda x. xx$ cannot be typed since it would lead to a contradiction where x must be of type $\tau \rightarrow \sigma$ and τ at the same time.

Exercise 3

$$N = (o \rightarrow o) \rightarrow o \rightarrow o$$

- a) Considering $\Gamma = n: N, f: o \rightarrow o, x: o$, the typing derivation for the successor operator is:

$$\frac{\frac{\frac{\theta = o, \eta = o}{\Gamma \vdash f: \theta \rightarrow \eta} \quad \frac{\frac{\frac{\tau = o \rightarrow o, \omega = o, \theta = o}{\Gamma \vdash n: \tau \rightarrow (\omega \rightarrow \theta)} \quad \frac{\tau = o \rightarrow o}{\Gamma \vdash f: \tau}}{\Gamma \vdash nf: \omega \rightarrow \theta} \quad \frac{\omega: o}{\Gamma \vdash x: \omega}}{\Gamma \vdash nf x: \theta} \quad \frac{n: N, f: o \rightarrow o, x: o \vdash f(nf x): \eta}{n: N, f: o \rightarrow o \vdash \lambda x^o. f(nf x): \alpha = \epsilon = o \rightarrow \eta}}{n: N \vdash \lambda f^{o \rightarrow o}. \lambda x^o. f(nf x): \alpha = \gamma: (o \rightarrow o) \rightarrow \epsilon} \quad \vdash \lambda n^N. \lambda f^{o \rightarrow o}. \lambda x^o. f(nf x): \alpha = N \rightarrow \gamma$$

$$\frac{\frac{\frac{\Gamma \vdash n: \tau \rightarrow N}{\Gamma \vdash nf: o \rightarrow o} \quad \frac{\Gamma \vdash f: o \rightarrow o}{\Gamma \vdash x: o}}{\Gamma \vdash nf x: o} \quad \frac{n: N, f: o \rightarrow o, x: o \vdash f(nf x): o}{n: N, f: o \rightarrow o \vdash \lambda x^o. f(nf x): o \rightarrow o}}{n: N \vdash \lambda f^{o \rightarrow o}. \lambda x^o. f(nf x): (o \rightarrow o) \rightarrow o \rightarrow o} \quad \vdash \lambda n^N. \lambda f^{o \rightarrow o}. \lambda x^o. f(nf x): N \rightarrow N$$

- b) $N \rightarrow N \rightarrow N$

- c) $add \triangleq \lambda m^N. \lambda n^N. \lambda f^{o \rightarrow o}. \lambda x^o. mf(nf x)$

- d) Let $\Gamma = m: N, n: N, f: o \rightarrow o, x: o$, the typing derivation for the addition term is:

$$\frac{\frac{\frac{\tau \rightarrow (\omega \rightarrow \theta) = N}{\Gamma \vdash m: \tau \rightarrow (\omega \rightarrow \theta)} \quad \frac{\tau = o \rightarrow o}{\Gamma \vdash f: \tau}}{\Gamma \vdash mf: \omega \rightarrow \theta} \quad \frac{\frac{\frac{\mu \rightarrow (\Omega \rightarrow \omega): N}{\Gamma \vdash n: \mu \rightarrow (\Omega \rightarrow \omega)} \quad \frac{\mu = o \rightarrow o}{\Gamma \vdash f: \mu}}{\Gamma \vdash nf: \Omega \rightarrow \omega} \quad \frac{\Omega = o}{\Gamma \vdash x: \Omega}}{\Gamma \vdash mf(nf x): \theta} \quad \frac{m: N, n: N, f: o \rightarrow o \vdash \lambda x^o. mf(nf x): \eta = o \rightarrow \theta}{m: N, n: N \vdash \lambda f^{o \rightarrow o}. \lambda x^o. mf(nf x): \epsilon = (o \rightarrow o) \rightarrow \eta}}{m: N \vdash \lambda n^N. \lambda f^{o \rightarrow o}. \lambda x^o. mf(nf x): \alpha = N \rightarrow \gamma} \quad \vdash \lambda m^N. \lambda n^N. \lambda f^{o \rightarrow o}. \lambda x^o. mf(nf x): \alpha = N \rightarrow \gamma$$

$$\frac{\frac{\frac{\Gamma \vdash m: N}{\Gamma \vdash mf: o \rightarrow o} \quad \frac{\Gamma \vdash f: o \rightarrow o}{\Gamma \vdash x: o}}{\Gamma \vdash mf(nf x): o} \quad \frac{n: N, n: N, f: o \rightarrow o \vdash \lambda x^o. mf(nf x): o \rightarrow o}{m: N, n: N \vdash \lambda f^{o \rightarrow o}. \lambda x^o. mf(nf x): (o \rightarrow o) \rightarrow o \rightarrow o}}{m: N \vdash \lambda n^N. \lambda f^{o \rightarrow o}. \lambda x^o. mf(nf x): N \rightarrow N} \quad \vdash \lambda m^N. \lambda n^N. \lambda f^{o \rightarrow o}. \lambda x^o. mf(nf x): N \rightarrow N \rightarrow N$$

- e) No, it cannot be similarly typed. The exponentiation term is:

$$exp \triangleq \lambda m. \lambda n. \lambda f. \lambda x. (nm)fx$$

Since n is applied to m , and n expects a function but the type of m does not match $\alpha \rightarrow \alpha$, the expression cannot be similarly typed. The first argument, m , must have a smaller type than n .