

### Count elements less than or equal to k in a sorted rotated array

Difficulty: Medium Accuracy: 55.56% Submissions: 13K+ Points: 4 Average Time: 15m

Given a sorted array `arr[]` containing distinct non negative integers that has been rotated at some unknown pivot, and a value `x`. Your task is to **count** the number of elements in the array that are **less** than or equal to `x`.

#### Examples:

**Input:** `arr[] = [4, 5, 8, 1, 3]`, `x = 6`  
**Output:** 4  
**Explanation:** 1, 3, 4 and 5 are less than 6, so the count of all elements less than

Output Window

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Problem Solved Successfully

Test Cases Passed  
**1113 / 1113**

Attempts : Correct / Total  
**1 / 1**  
Accuracy : 100%

```
1 class Solution {
2     public int countLessEqual(int[] arr, int x) {
3
4         int count = 0;
5
6         for (int num : arr) {
7             if (num <= x) count++;
8         }
9
10        return count;
11    }
12 }
13
```

## Maximize the minimum difference between k elements

Difficulty: Medium Accuracy: 61.58% Submissions: 13K+ Points: 4

Given an array `arr[]` of integers and an integer `k`, select `k` elements from the array such that the **minimum absolute difference** between any two of the selected elements is **maximized**. Return this **maximum** possible **minimum** difference.

### Examples:

**Input:** `arr[] = [2, 6, 2, 5]`, `k = 3`  
**Output:** 1  
**Explanation:** 3 elements out of 4 elements are to be selected with a minimum

### Output Window

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Problem Solved Successfully

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Test Cases Passed: **1115 / 1115**  
 Attempts : Correct / Total: **1 / 1**  
 Accuracy : 100%

```
5 public int maxMinDiff(int[] arr, int k) {
6
7     Arrays.sort(arr);
8
9     int low = 0, high = arr[arr.length - 1] - arr[0];
10    int ans = 0;
11
12    while (low <= high) {
13        int mid = (low + high) / 2;
14
15        if (canPlace(arr, k, mid)) {
16            ans = mid;
17            low = mid + 1;
18        } else {
19            high = mid - 1;
20        }
21    }
22    return ans;
23 }
24
25 boolean canPlace(int[] arr, int k, int diff) {
26
27     int count = 1;
28     int last = arr[0];
29
30     for (int i = 1; i < arr.length; i++) {
31         if (arr[i] - last >= diff) {
32             count++;
33             last = arr[i];
34         }
35     }
36
37     return count >= k;
38 }
39
40 }
41 }
```



## Unique K-Number Sum

Difficulty: Medium Accuracy: 66.73% Submissions: 12K+ Points: 4

Given two integers **n** and **k**, the task is to find all valid combinations of **k** numbers that adds up to **n** based on the following conditions:

- Only numbers from the range [1, 9] used.
- Each number can only be used **at most** once.

**Note:** You can return the combinations in any order, the driver code will print them in sorted order.

Examples:

Output Window

Compilation Results Custom Input Y.O.G.I. (AI Bot)

Problem Solved Successfully

Suggest Feedback

Test Cases Passed

1120 / 1120

Attempts : Correct / Total

1 / 1

Accuracy : 100%

```
1 import java.util.*;
2
3 class Solution {
4
5     public ArrayList<ArrayList<Integer>> combinationSum(int n, int k) {
6
7         ArrayList<ArrayList<Integer>> result = new ArrayList<>();
8         backtrack(1, n, k, new ArrayList<>(), result);
9         return result;
10    }
11
12    private void backtrack(int start, int target, int k,
13                           ArrayList<Integer> current,
14                           ArrayList<ArrayList<Integer>> result) {
15
16        if (target == 0 && current.size() == k) {
17            result.add(new ArrayList<>(current));
18            return;
19        }
20
21        if (target < 0 || current.size() > k) {
22            return;
23        }
24
25        for (int i = start; i <= 9; i++) {
26
27            current.add(i);
28            backtrack(i + 1, target - i, k, current, result);
29            current.remove(current.size() - 1);
30        }
31    }
32 }
33
```



Custom Input

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## Maximum People Visible in a Line

Difficulty: **Medium**    Accuracy: **50.11%**    Submissions: **17K+**    Points: **4**

You are given an array **arr**[], where arr[i] represents the height of the ith person standing in a line.

A person **i** can see another person **j** if:

- height[j] < height[i],
- There is no person **k** standing between them such that height[k] ≥ height[i].

Each person can see in both directions (front and back).

Your task is to find the **maximum number of people** that any person can see (including

Output Window

Compilation Results

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**Problem Solved Successfully** ✓

[Suggest Feedback](#)

Test Cases Passed

**1120 / 1120**


Attempts : Correct / Total

**1 / 1**

Accuracy : 100%

```

1 import java.util.*;
2
3 class Solution {
4
5     public int maxPeople(int[] arr) {
6         int n = arr.length;
7
8         int[] left = new int[n];
9         int[] right = new int[n];
10
11         Stack<Integer> stack = new Stack<>();
12
13
14         for (int i = 0; i < n; i++) {
15             while (!stack.isEmpty() && arr[stack.peek()] < arr[i]) {
16                 stack.pop();
17             }
18             left[i] = stack.isEmpty() ? i : i - stack.peek() - 1;
19             stack.push(i);
20         }
21
22         stack.clear();
23
24         for (int i = n - 1; i >= 0; i--) {
25             while (!stack.isEmpty() && arr[stack.peek()] < arr[i]) {
26                 stack.pop();
27             }
28             right[i] = stack.isEmpty() ? (n - i - 1) : stack.peek() - i - 1;
29             stack.push(i);
30         }
31
32         int ans = 1;
33
34         for (int i = 0; i < n; i++) {
35             int total = left[i] + right[i] + 1;
36             ans = Math.max(ans, total);
37         }
38     }
39 }
```

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### Find 132 Pattern in Array

Difficulty: Medium Accuracy: 73.39% Submissions: 534+ Points: 4

You are given an array `arr[]`. The task is to determine whether the array contains a **132 pattern**, i.e., three indices `i`, `j` and `k` such that `i < j < k`, `arr[i] < arr[j] > arr[k]` and `arr[i] < arr[k]`. Return `true` if such a triplet exists, otherwise return `false`.

**Examples:**

**Input:** `arr[] = [4, 7, 11, 5, 13, 2]`  
**Output:** `true`  
**Explanation:** Triplet `[4, 7, 5]` satisfies the condition since `4 < 7`, `5 < 7` and `4 < 5`.

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Test Cases Passed

Attempts : Correct / Total

Accuracy : 100%

1120 / 1120


1 / 1

```
1 import java.util.*;
2
3 class Solution {
4
5     public boolean has132Pattern(int[] arr) {
6
7         int n = arr.length;
8         if (n < 3) return false;
9
10        int[] min = new int[n];
11        min[0] = arr[0];
12
13
14        for (int i = 1; i < n; i++) {
15            min[i] = Math.min(min[i - 1], arr[i]);
16        }
17
18        Stack<Integer> stack = new Stack<>();
19
20
21        for (int j = n - 1; j >= 0; j--) {
22
23            if (arr[j] > min[j]) {
24
25                while (!stack.isEmpty() && stack.peek() <= min[j]) {
26                    stack.pop();
27                }
28
29                if (!stack.isEmpty() && stack.peek() < arr[j]) {
30                    return true;
31                }
32                stack.push(arr[j]);
33            }
34        }
35
36        return false;
37    }
38 }
```

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### Subarrays with First Element Minimum

Difficulty: Medium Accuracy: 69.1% Submissions: 440+ Points: 4

You are given an integer array `arr[]`. Your task is to **count** the number of subarrays where the first element is the **minimum element** of that subarray.

**Note:** A subarray is valid if its first element is not greater than any other element in that subarray.

**Examples:**

Input: `arr[] = [1, 2, 1]`

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Problem Solved Successfully

Suggest Feedback

Test Cases Passed

1120 / 1120

Attempts : Correct / Total

You can see all your attempts in submission tab

Accuracy : 100%

Java (21)

Start Timer

```
1 class Solution {
2
3     public int countSubarrays(int[] arr) {
4
5         int n = arr.length;
6         int count = 0;
7
8         for (int i = 0; i < n; i++) {
9
10            int minval = arr[i];
11
12            for (int j = i; j < n; j++) {
13
14                if (arr[j] < minval) {
15                    break;
16                }
17
18                count++;
19            }
20        }
21
22        return count;
23    }
24 }
25
```

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## Previous Greater Element

Difficulty: Medium Accuracy: 71.86% Submissions: 6K+ Points: 4

You are given an integer array `arr[]`. For every element in the array, your task is to determine its **Previous Greater Element (PGE)**.

The Previous Greater Element (PGE) of an element `x` is the first element that appears to the left of `x` in the array and is strictly greater than `x`.

**Note:** If no such element exists, assign `-1` as the PGE for that position.

**Examples:**

Output Window

Compilation Results Custom Input Y.O.G.I. (AI Bot)

Problem Solved Successfully

[Suggest Feedback](#)

Test Cases Passed

1120 / 1120

Attempts : Correct / Total

You can see all your attempts in submission tab

Accuracy : 100%

```
1- import java.util.*;
2-
3- class Solution {
4-
5-     public ArrayList<Integer> preGreaterEle(int[] arr) {
6-
7-         ArrayList<Integer> res = new ArrayList<>();
8-         Stack<Integer> st = new Stack<>();
9-
10-         for (int i = 0; i < arr.length; i++) {
11-
12-             while (!st.isEmpty() && st.peek() <= arr[i]) {
13-                 st.pop();
14-             }
15-
16-             if (st.isEmpty()) {
17-                 res.add(-1);
18-             } else {
19-                 res.add(st.peek());
20-             }
21-
22-             st.push(arr[i]);
23-         }
24-
25-         return res;
26-     }
27- }
28-
```

## Previous Smaller Element

Difficulty: Medium Accuracy: 65.65% Submissions: 8K+ Points: 4

You are given an integer array `arr[]`. For every element in the array, your task is to determine its **Previous Smaller Element (PSE)**.

The Previous Smaller Element (PSE) of an element `x` is the first element that appears to the left of `x` in the array and is strictly smaller than `x`.

**Note:** If no such element exists, assign `-1` as the PSE for that position.

Examples:

Output Window

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Problem Solved Successfully

Suggest Feedback

Test Cases Passed

1115 / 1115

Attempts : Correct / Total

1 / 1

Accuracy : 100%

```
1+ import java.util.*;
2
3- class Solution {
4
5-     public ArrayList<Integer> prevSmaller(int[] arr) {
6
7         ArrayList<Integer> res = new ArrayList<>();
8         stack<Integer> st = new stack<>();
9
10-        for (int i = 0; i < arr.length; i++) {
11
12-            while (!st.isEmpty() && st.peek() >= arr[i]) {
13-                st.pop();
14-            }
15
16-            if (st.isEmpty()) {
17-                res.add(-1);
18-            } else {
19-                res.add(st.peek());
20-            }
21
22-            st.push(arr[i]);
23-        }
24
25-        return res;
26-    }
27- }
28
```



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## Minimum Number Of Sprinkler

Difficulty: Hard Accuracy: 28.03% Submissions: 2K+ Points: 8

Given a one-dimensional garden of length  $n$ . In each position of the  $n$  length garden, a sprinkler has been installed. Given an array `arr[]` such that `arr[i]` describes the coverage limit of the  $i^{\text{th}}$  sprinkler. A sprinkler can cover the range from the position `max(i - arr[i], 1)` to `min(i + arr[i], n)`. In beginning, all the sprinklers are switched off.

The task is to find the minimum number of sprinklers needed to be activated such that the whole  $n$ -length garden can be covered by water.

**Note: Array is 1-based indexed.**

Example 1:

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Test Cases Passed

163 / 163

Attempts : Correct / Total

1 / 1

Accuracy : 100%

```

9 // Convert sprinkler range into intervals
10 for (int i = 0; i < N; i++) {
11
12     if (arr[i] == -1) continue;
13
14     int left = Math.max(0, i - arr[i]);
15     int right = Math.min(N - 1, i + arr[i]);
16
17     intervals.add(new int[]{left, right});
18 }
19
20 // Sort by starting point
21 Collections.sort(intervals, (a, b) -> a[0] - b[0]);
22
23 int count = 0;
24 int i = 0;
25 int covered = 0;
26
27 while (covered < N) {
28
29     int farthest = covered;
30
31     while (i < intervals.size() && intervals.get(i)[0] <= covered) {
32         farthest = Math.max(farthest, intervals.get(i)[1] + 1);
33         i++;
34     }
35
36     if (farthest == covered) return -1;
37
38     covered = farthest;
39     count++;
40 }
41
42 return count;
43
44 }
45


```



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### Minimum Steps to Halve Sum

Difficulty: **Medium** Accuracy: 55.59% Submissions: 14K+ Points: 4

Given an array `arr[]`, find the **minimum** number of operations required to make the sum of its elements less than or equal to **half** of the original sum. In one operation, you may replace any element with half of its value (with **floating-point** precision).

**Examples:**

**Input:** `arr[] = [8, 6, 2]`

**Output:** 3

**Explanation:** Initial sum =  $(8 + 6 + 2) = 16$ , half = 8

Output Window

Compilation Results Custom Input Y.O.G.I. (AI Bot)

Problem Solved Successfully

Suggest Feedback

Test Cases Passed

1115 / 1115

Attempts : Correct / Total

1 / 1

Accuracy : 100%

Java (21)

Start Timer

```
1- import java.util.*;
2
3- class Solution {
4
5-     public int minOperations(int[] arr) {
6
7         PriorityQueue<Double> pq = new PriorityQueue<>(Collections.reverseOrder());
8
9         double sum = 0;
10
11         for (int num : arr) {
12             pq.add((double) num);
13             sum += num;
14         }
15
16         double target = sum / 2.0;
17         int operations = 0;
18
19         while (sum > target) {
20
21             double largest = pq.poll();
22             double half = largest / 2.0;
23
24             sum -= half;
25             pq.add(half);
26
27             operations++;
28         }
29
30         return operations;
31     }
32 }
33
```

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## Footpath Construction

Difficulty: Medium Accuracy: 39.36% Submissions: 127+ Points: 4 Average Time: 19m

Given a matrix **a** of size **n\*m** which represents a **park**, there is some construction work needs to be done. You are also given **q** queries each query contains two numbers **R** and **C**. For every query we need to construct a footpath in the **Rth** row and **Cth** column, there is a **cost** of this construction, after the construction this path will divide the park into **sections**, and the cost of the construction is the **sum of minimum** value present in all the sections. You are asked to find this cost for all the queries.

**Note:** Elements present in queries array are according to 1-based indexing.

### Output Window

#### Compilation Results

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Problem Solved Successfully

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Test Cases Passed

1122 / 1122

Attempts : Correct / Total

1 / 3

Accuracy : 33%

### Practice Coding Problems

#### Problem of the Day

#### Explore Connect


```
path(int n, int m, int[][] a, int q, int[][] queries) {  
    int[][] tl = new int[n][m]; // top-left  
    int[][] tr = new int[n][m]; // top-right  
    int[][] bl = new int[n][m]; // bottom-left  
    int[][] br = new int[n][m]; // bottom-right  
  
    // Top-left  
    for (int i = 0; i < n; i++) {  
        for (int j = 0; j < m; j++) {  
            tl[i][j] = a[i][j];  
            if (i > 0) tl[i][j] = Math.min(tl[i][j], tl[i - 1][j]);  
            if (j > 0) tl[i][j] = Math.min(tl[i][j], tl[i][j - 1]);  
        }  
    }  
  
    // Top-right  
    for (int i = 0; i < n; i++) {  
        for (int j = m - 1; j >= 0; j--) {  
            tr[i][j] = a[i][j];  
            if (i > 0) tr[i][j] = Math.min(tr[i][j], tr[i - 1][j]);  
            if (j < m - 1) tr[i][j] = Math.min(tr[i][j], tr[i][j + 1]);  
        }  
    }  
  
    // Bottom-left  
    for (int i = n - 1; i >= 0; i--) {  
        for (int j = 0; j < m; j++) {  
            bl[i][j] = a[i][j];  
            if (i < n - 1) bl[i][j] = Math.min(bl[i][j], bl[i + 1][j]);  
            if (j > 0) bl[i][j] = Math.min(bl[i][j], bl[i][j - 1]);  
        }  
    }  
}
```



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### Prefix Sum of Matrix (Or 2D Array)

Difficulty: **Medium** Accuracy: **76.55%** Submissions: **145+** Points: **4** Average Time: **20m**

Given a integer matrix (or 2D array) **a[][]** of dimensions **n \* m**. Also, given another 2-D array **query[][]** of dimensions **q \* 4**.

For each index  $0 \leq i < \text{query.length}$ , find the sum of all the elements of the rectangular matrix whose top left corner is (query[i][0], query[i][1]) and bottom right corner is (query[i][2], query[i][3]).

Example -

Output Window

Compilation Results Custom Input

Compilation Completed

Case 1

Input:

Java (21)

Start Timer

```
1- import java.util.*;
2-
3- class Solution {
4-
5-     public ArrayList<Long> submatrixSum(long[][] a, int n, int m, int[][] query, int q) {
6-
7-         long[][] prefix = new long[n][m];
8-
9-         // Build prefix sum
10-        for (int i = 0; i < n; i++) {
11-            for (int j = 0; j < m; j++) {
12-
13-                prefix[i][j] = a[i][j];
14-
15-                if (i > 0) prefix[i][j] += prefix[i - 1][j];
16-                if (j > 0) prefix[i][j] += prefix[i][j - 1];
17-                if (i > 0 && j > 0) prefix[i][j] -= prefix[i - 1][j - 1];
18-            }
19-        }
20-
21-        ArrayList<Long> ans = new ArrayList<>();
22-
23-        // Process queries
24-        for (int k = 0; k < q; k++) {
25-
26-            int r1 = query[k][0] - 1;
27-            int c1 = query[k][1] - 1;
28-            int r2 = query[k][2] - 1;
29-            int c2 = query[k][3] - 1;
30-
31-            long sum = prefix[r2][c2];
32-
33-            if (r1 > 0) sum -= prefix[r1 - 1][c2];
34-            if (c1 > 0) sum -= prefix[r2][c1 - 1];
35-            if (r1 > 0 && c1 > 0) sum += prefix[r1 - 1][c1 - 1];
36-
37-            ans.add(sum);
38-        }
39-    }
40-}
```

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