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## 518. Coin Change II [↗](/problems/coin-change-2/) (/problems/coin-change-2/)

Nov. 16, 2019 | 2.8K views

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You are given coins of different denominations and a total amount of money. Write a function to compute the number of combinations that make up that amount. You may assume that you have infinite number of each kind of coin.

### Example 1:

**Input:** amount = 5, coins = [1, 2, 5]

**Output:** 4

**Explanation:** there are four ways to make up the amount:

5=5

5=2+2+1

5=2+1+1+1

5=1+1+1+1+1

### Example 2:

**Input:** amount = 3, coins = [2]

**Output:** 0

**Explanation:** the amount of 3 cannot be made up just with coins of 2.

### Example 3:

**Input:** amount = 10, coins = [10]

**Output:** 1

## Note:

You can assume that

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- $0 \leq \text{amount} \leq 5000$
- $1 \leq \text{coin} \leq 5000$
- the number of coins is less than 500
- the answer is guaranteed to fit into signed 32-bit integer

# Solution

## Approach 1: Dynamic Programming

### Template

This is a classical dynamic programming problem.

Here is a template one could use:

- Define the base cases for which the answer is obvious.
- Develop the strategy to compute more complex case from more simple one.
- Link the answer to base cases with this strategy.

### Example

Let's pick up an example: amount = 11, available coins - 2 cent, 5 cent and 10 cent. Note, that coins are unlimited.

Number of combinations that make up 11



### Base Cases: No Coins or Amount = 0

If the total amount of money is zero, there is only one combination: to take zero coins.


Another base case is no coins: zero combinations for amount > 0 and one combination for amount == 0 .


amount =	0	1	2	3	4	5	6	7	8	9	10	11
combinations using no coins =	1	0	0	0	0	0	0	0	0	0	0	0


## 2 Cent Coins


Let's do one step further and consider the situation with one kind of available coins: 2 cent.


amount =	0	1	2	3	4	5	6	7	8	9	10	11
combinations using no coins =	1	0	0	0	0	0	0	0	0	0	0	0
combinations with 2 cents =	1	0	1	0	1	0	1	0	1	0	1	0











It's quite evident that there could be 1 or 0 combinations here, 1 combination for even amount and 0 combinations for the odd one.

The same answer could be received in a recursive way, by computing the number of combinations for all amounts of money, from 0 to 11.

First, that's quite obvious that all amounts less than 2 are *not* impacted by the presence of 2 cent coins. Hence for amount = 0 and for amount = 1 one could reuse the results from the figure 2.

Starting from amount = 2 , one could use 2 cent coins in the combinations. Since the amounts are considered gradually from 2 to 11, at each given moment one could be sure to add not more than one coin to the previously known combinations.

So let's pick up 2 cent coin, and use it to make up amount = 2 . The number of combinations with this 2 cent coin is a number combinations for amount = 0 , i.e. 1.

amount =

0	1	2	3	4	5	6	7	8	9	10	11
combinations using no coins =	1	0	0	0	0	0	0	0	0	0	0
combinations with 2 cents =	1	0	1								

number of combinations for [amount = 2] =  
 number of current combinations for [amount = 2] +  
 number of current combinations for [amount - 2 = 0]  
 = 0 + 1 = 1

Now let's pick up 2 cent coin, and use it to make up amount = 3. The number of combinations with this 2 cent coin is a number combinations for amount = 1, i.e. 0.

amount =

0	1	2	3	4	5	6	7	8	9	10	11
combinations using no coins =	1	0	0	0	0	0	0	0	0	0	0
combinations with 2 cents =	1	0	1	0							

number of combinations for [amount = 3] =  
 number of current combinations for [amount = 3] +  
 number of current combinations for [amount - 2 = 1]  
 = 0 + 0 = 0

That leads to DP formula for number of combinations to make up the amount = x:  $dp[x] = dp[x] + dp[x - \text{coin}]$ , where coin = 2 cents is a value of coins we're currently adding.

amount =

0	1	2	3	4	5	6	7	8	9	10	11
combinations using no coins =	1	0	0	0	0	0	0	0	0	0	0
combinations with 2 cents =	1	0	1	0	1	0	1	0	1	0	1

$dp[x] = dp[x] + dp[x - \text{coin}]$   
 coin = 2 cents

## 2 Cent Coins + 5 Cent Coins + 10 Cent Coins

Now let's add 5 cent coins. The formula is the same, but do not forget to add  $dp[x]$ , number of combinations with 2 cent coins.

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amount =	0	1	2	3	4	5	6	7	8	9	10	11
combinations using no coins =	1	0	0	0	0	0	0	0	0	0	0	0
combinations with 2 cents =	1	0	1	0	1	0	1	0	1	0	1	0
combinations with 2 and 5 cents =	1	0	1	0	1	1	1	1	1	1	2	1

$dp[x]$  for  $x < \text{coin} = 5$  cents  
are not going to change

$dp[x] = dp[x] + dp[x - \text{coin}]$   
 $\text{coin} = 5$  cents

The story is the same for 10 cent coins.

amount =	0	1	2	3	4	5	6	7	8	9	10	11
combinations using no coins =	1	0	0	0	0	0	0	0	0	0	0	0
combinations with 2 cents =	1	0	1	0	1	0	1	0	1	0	1	0
combinations with 2 and 5 cents =	1	0	1	0	1	1	1	1	1	1	2	1
combinations with 2, 5, and 10 cents =	1	0	1	0	1	1	1	1	1	1	3	1

$dp[x]$  for  $x < \text{coin} = 10$  cents  
are not going to change

$dp[x] = dp[x] + dp[x - \text{coin}]$   
 $\text{coin} = 10$  cents

Now the strategy is here:

- Add coins one-by-one, starting from base case "no coins".
- For each added coin, compute recursively the number of combinations for each amount of money from 0 to amount .

### Algorithm

- Initiate number of combinations array with the base case "no coins":  $dp[0] = 1$  , and all the rest = 0.
- Loop over all coins:

- For each coin, loop over all amounts from 0 to amount :

- For each amount x, compute the number of combinations:  $dp[x] += dp[x - coin]$  .

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- Return  $dp[amount]$  .

## Implementation

Java

Python

Copy

```
1 class Solution:
2     def change(self, amount: int, coins: List[int]) -> int:
3         dp = [0] * (amount + 1)
4         dp[0] = 1
5
6         for coin in coins:
7             for x in range(coin, amount + 1):
8                 dp[x] += dp[x - coin]
9         return dp[amount]
```

## Complexity Analysis

- Time complexity:  $\mathcal{O}(N \times \text{amount})$ , where N is a length of coins array.
- Space complexity:  $\mathcal{O}(\text{amount})$  to keep dp array.

Analysis written by @liaison (<https://leetcode.com/liaison/>) and @andvary (<https://leetcode.com/andvary/>)

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