

School of Computing and Information Systems  
COMP90038 Algorithms and Complexity Tutorial Week 4

Sample answers

1. One possible way of representing a polynomial

$$p(x) = a_n x^n + a_{n-1} x^{n-1} + \cdots + a_1 x + a_0$$

is as an array  $A$  of length  $n + 1$ , with  $A[i]$  holding the coefficient  $a_i$ .

- (a) Design a brute-force algorithm for computing the value of  $p(x)$  at a given point  $x$ . Express this as a function  $\text{PEVAL}(A, n, x)$  where  $A$  is the array of coefficients,  $n$  is the degree of the polynomial, and  $x$  is the point for which we want the value of  $p$ .
- (b) If your algorithm is  $\Theta(n^2)$ , try to find a linear algorithm.
- (c) Is it possible to find an algorithm that solves the problem in sub-linear time?

**Answer.**

- (a) Working from right-to-left, the following algorithm is the natural formulation:

```
function PEVAL( $A, n, x$ )  
     $result \leftarrow 0.0$   
    for  $i \leftarrow n$  downto 0 do  
         $summand \leftarrow 1.0$   
        for  $j \leftarrow 1$  to  $i$  do  
             $summand \leftarrow x \times summand$   
         $result \leftarrow result + A[i] \times summand$   
    return  $result$ 
```

The complexity is  $\Theta(n^2)$ .

- (b) Working from left-to-right allows us to avoid many redundant calculations of  $x^i$ . It gives an algorithm that is both simpler and more efficient:

```
function PEVAL( $A, n, x$ )  
     $result \leftarrow A[0]$   
     $summand \leftarrow 1.0$   
    for  $i \leftarrow 1$  to  $n$  do  
         $summand \leftarrow x \times summand$   
         $result \leftarrow result + A[i] \times summand$   
    return  $result$ 
```

- (c) We cannot solve the problem in less than linear time, because we clearly need to access each of the  $n + 1$  coefficients.

2. Trace the brute-force string search algorithm on the following input: The path  $p$  is ‘needle’, and the text  $t$  is ‘there\_need\_not\_be\_any’. How many comparisons (successful and unsuccessful) are made?

**Answer.** 21 character comparisons are made.

3. Assume we have a text consisting of one million zeros. For each of these patterns, determine how many character comparisons the brute-force string matching algorithm will make:

(a) 010001      (b) 000101      (c) 011101

**Answer.**

- (a)  $2 \times 999995$  comparisons
  - (b)  $4 \times 999995$  comparisons
  - (c)  $2 \times 999995$  comparisons
4. Give an example of a text of length  $n$  and a pattern, which together constitute a worst-case scenario for the brute-force string matching algorithm. How many character comparisons, as a function of  $n$ , will be made for the worst-case example.

**Answer.** The worst case happens when we have a text of length  $n$  consisting of the same character  $c$  repeated  $n$  times, together with a pattern of length  $m$ , consisting of  $m - 1$  occurrences of  $c$ , followed by a single character different from  $c$ . In this case, the outer loop is traversed  $n - m + 1$  times, and each time,  $m$  character comparisons are made before failure is detected. Altogether we have  $(n - m + 1)m = (n + 1)m - m^2$  comparisons. As a function of  $m$ , this has its maximal value where  $n + 1 - 2m = 0$ , that is, when the length of the pattern is about half that of the text.

5. The *assignment problem* asks how to best assign  $n$  jobs to  $n$  contractors who have put in bids for each job. An instance of this problem is an  $n \times n$  *cost matrix*  $C$ , with  $C[i, j]$  specifying what it will cost to have contractor  $i$  do job  $j$ . The aim is to minimise the total cost. More formally, we want to find a permutation  $\langle j_1, j_2, \dots, j_n \rangle$  of  $\langle 1, 2, \dots, n \rangle$  such that  $\sum_{i=1}^n C[i, j_i]$  is minimized.

Use brute force to solve the following instance:

	Job 1	Job 2	Job 3	Job 4
Contractor 1	9	2	7	8
Contractor 2	6	4	3	7
Contractor 3	5	8	1	8
Contractor 4	7	6	9	4

**Answer.**

Permutation	Cost
1,2,3,4	$9+4+1+4 = 18$
1,2,4,3	$9+4+8+9 = 30$
1,3,2,4	$9+3+8+4 = 24$
1,3,4,2	$9+3+8+6 = 26$
1,4,2,3	$9+7+8+9 = 33$
1,4,3,2	$9+7+1+6 = 23$
2,1,3,4	$2+6+1+4 = 13$
2,1,4,3	$2+6+8+9 = 25$
$\vdots$	

and so on. The minimal cost is 13, for permutation  $\langle 2, 1, 3, 4 \rangle$ .

6. Give an instance of the assignment problem which does not include the smallest item  $C[i, j]$  of its cost matrix.

**Answer.**

	Job 1	Job 2
Contractor 1	1	2
Contractor 2	2	4