

Week 06 Quiz

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Due Apr 23 at 23:59**Points** 10**Questions** 9**Available** Apr 14 at 10:00 - Apr 23 at 23:59 10 days**Time Limit** None**Allowed Attempts** Unlimited

Instructions

You should attempt the quiz after the lecture and your tutorial.

- You may attempt the quiz multiple times (if you happen to get a question wrong, you can do it again)
- Your score on the quiz will be recorded in the grade book.
- The quiz might not display equations correctly in some browsers. If you experience problems, we recommend that you use Firefox.

Note: you must complete at least eight of the weekly quizzes to meet one of the hurdle requirements in this subject

This quiz was locked Apr 23 at 23:59.

Attempt History

Attempt**Time****Score**

	Attempt	Time	Score
LATEST	Attempt 1	1,280 minutes	8 out of 10

Score for this attempt: **8** out of 10

Submitted Apr 23 at 21:08

This attempt took 1,280 minutes.

Correct!

Question 1

1 / 1 pts

The recurrence relation for merge sort is:

- ☒ $T(n) = 2T(n/2) + O(n)$
- ☐ $T(n) = T(n-1) + O(n)$
- ☐ $T(n) = 2T(n/2) + O(1)$
- ☐ $T(n) = T(n/2) + O(1)$

Let's think about how merge sort works. We split our array or list in to two halves, and run merge sort recursively on each half (this is where the $2T(n/2)$ part arises). We then merge the two sorted halves (this is where the $O(n)$ part arises).

Question 2

1 / 1 pts

Quicksort uses Hoare partitioning. Assume an array contains ten keys: 6 3 1 7 9 5 8 2 4 0. After a first round of simple Hoare partitioning (not median-of-three), the array looks like so:

- ☐ 5 3 1 4 0 2 6 7 9 8

Correct!☐ 3 1 5 2 4 0 6 7 9 8☒ 2 3 1 0 4 5 6 8 9 7☐ 2 3 0 1 5 4 6 7 8 9☐ 5 3 1 0 4 2 6 8 9 7

Well done!

Question 3**1 / 1 pts**

Consider this recurrence relation:

$$T(1) = 1$$

$$T(n) = 2 T(n/3) + 2n + 1 \quad \text{for } n > 1$$

The Master Theorem says that

☐ $T(n) \in \Theta(n^3)$ ☐ $T(n) \in \Theta(n^2)$ ☐ $T(n) \in \Theta(n \log \log n)$ ☒ $T(n) \in \Theta(n)$ ☐ $T(n) \in \Theta(n \log n)$ **Correct!**That's right. In this case we have $a=2$, $b=3$, and $d=1$. And indeed $2 < 3$.

Question 4**1 / 1 pts**

Consider this recurrence relation:

$$T(1) = 1$$

$$T(2) = 1$$

$$T(n) = 4 T(n-2) + 2n^2 \quad \text{for } n > 2$$

The Master Theorem tells us

☐ $T(n) \in \Theta(n \log n)$

☐ $T(n) \in \Theta(n^3)$

☐ $T(n) \in \Theta(n^2 \log n)$

☐ $T(n) \in \Theta(n^2)$

☒ nothing

Correct!

That's right, the Master Theorem does not help here, as the recurrence is not of the required form.

Question 5**0 / 1 pts**

Which of the following sorting algorithm has the running time that is least dependant on the initial ordering of the input?

☐ Selection sort

☒ Merge sort

Correct Answer**You Answered**

☐ Quick sort☐ Insertion sort

Try again (see lecture slides)

Question 6**1 / 1 pts**

Suppose we have an array A with 33,554,431 elements. We want to apply binary search to look for some element k. A test of the form "is $k = A[i]$?" is a probe. How many probes will be performed in the worst case?

Correct!

Yes, the number of elements is $2^{25} - 1$. We have a worst-case instance if k is not in the array.

Correct Answers

25 (with margin: 0)

Question 7**1 / 1 pts**

What is the tight bound using Big-Theta notation for the time complexity of the following recurrence relation:

$$T(n) = 3T(n/2) + 1$$

when the Master Theorem is used?

☐ $\theta(n \log n)$

Correct!

- ☐ $\theta(n^{\log_3 2})$
- ☒ $\theta(n^{\log_2 3})$
- ☐ $\theta(\log_2 n) + O(1)$

Question 8**2 / 2 pts**

What is the tight bound using Big-Theta notation for the time complexity of the following recurrence relation:

$$T(n) = 3T\left(\frac{n}{9}\right) + \sqrt{n}$$

when the Master Theorem is used?

Correct!

- ☒ $\theta(n^{\frac{1}{2}} \log n)$
- ☐ $\theta(\sqrt{n})$
- ☐ $\theta(n^3)$
- ☐ $\theta(\log n^3)$

Question 9**0 / 1 pts**

Consider a modification to QuickSort where each time the partition function is called, the median of the partition array is always found (in constant time) and used as the pivot.

The worst-case running time for the algorithm is:

Correct Answer

- ☐ $\Theta(n \log n)$

You Answered

☐ $\Theta(n)$

☒ $\Theta(\log n)$

☐ $\Theta(n^2)$

Quiz Score: **8** out of 10